

② Check autocorrelation of your chain (marginally) using acf. Good chains should have decreasing acf. Faster decrease indicates fast mixing.

$$\text{acf}(1) = \frac{\frac{1}{n} \sum (x_t - \bar{x})(x_{t+1} - \bar{x})}{\frac{1}{n} \sum (x_t - \bar{x})^2}$$

③ Computing clusters allow easy implementation of multiple chains. So why not?

Simulated Annealing

Suppose we want to minimize a function

$h(\theta)$, where θ is a vector of parameters.
 $\theta \in S$

The basic idea is to simulate a MC with the stationary dist.

$$\pi_T(\theta) \propto \exp(-h(\theta)/T)$$

The parameter T is called the "temperature".

Markov chains for "simulation", annealing (cooling) for minimization.

$$\text{Let } S^* = \{\theta^* \in S : h(\theta^*) = \min_{\theta} h(\theta)\}$$

the set of global minimizers

Define a distribution π^* on S^* so that

$$\pi^*(\theta) \propto 1 \quad \text{for } \theta \in S^* \text{ and } 0 \text{ otherwise}$$

MC for simulation
"Annealing"

Then As $T \downarrow 0$ we have ~~$\pi_T \rightarrow \pi^*$~~

$$\pi_T(\theta) \rightarrow \pi(\theta) \quad \forall \theta \in S$$

FACT:

~~Then~~ As $T \downarrow 0$, we have $\pi_T \rightarrow \pi^*$

Obviously, we'd like to simulate from π^* to get the global minimum, but we can't.

But for a fixed $T > 0$, we can simulate from π_T .

Use Metropolis's algorithm with a symmetric proposal density, i.e. ~~$q(x|y) = q(y|x)$~~

$$q(\theta_n | \theta_{n-1}) = q(\theta_{n-1} | \theta_n)$$

Suppose we are in state n , θ_n ,

① Simulate $\theta \sim q(\theta | \theta_n)$

② Simulate $u \sim \text{unif}(0, 1)$

③ Set $\theta_{n+1} = \theta$ if $u \leq \min\left(\frac{\exp(-h(\theta)/T)}{\exp(-h(\theta_n)/T)}, 1\right)$

$$= \min\left(\exp\left(-\underbrace{(h(\theta) - h(\theta_n))}_{\Delta h} / T\right), 1\right)$$

otherwise set $\theta_{n+1} = \theta_n$

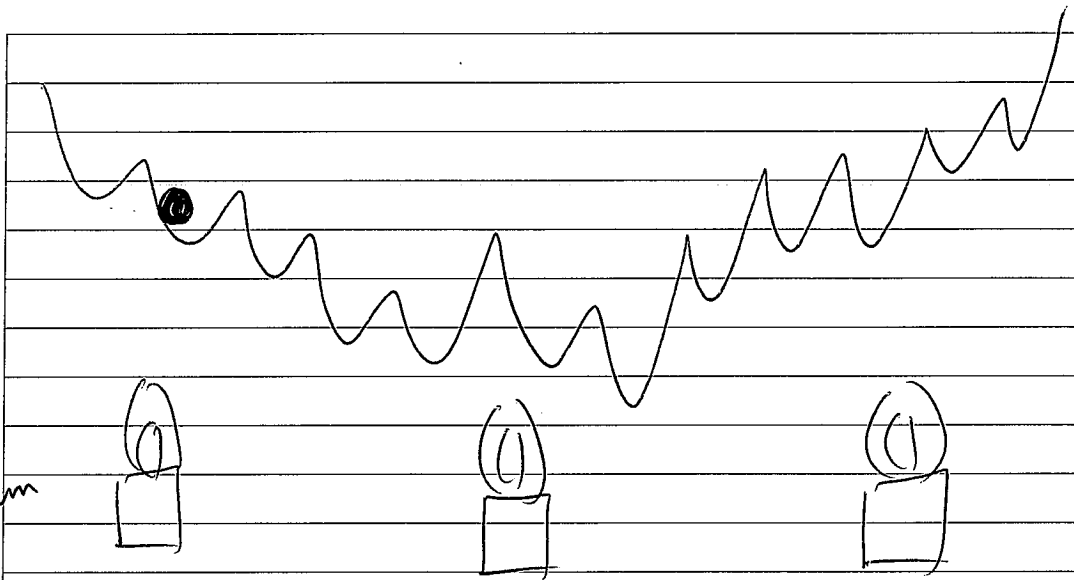
④ Decrease T .

If ~~$\Delta h < 0$~~ $\Delta h < 0$, we always go
If $\Delta h > 0$, we go w/ prob $\exp(-\Delta h/T)$

High temp
 \hookrightarrow ball moves
 all over

Low temp
 \hookrightarrow ball stops
 in minimum

down hill
 up hill



If $\Delta h < 0$, we always go

If $\Delta h > 0$, we go w/prob $\exp(-\Delta h/T)$

When T is closer to 0, it is less likely
 we will go up hill.

Can't cool too fast or will get stuck.

Need a cooling schedule like

$$T_n = \frac{a}{\log(n+b)} \quad (\text{for global convergence})$$

Anything else is "too fast" and will get
 stuck in a local mode.

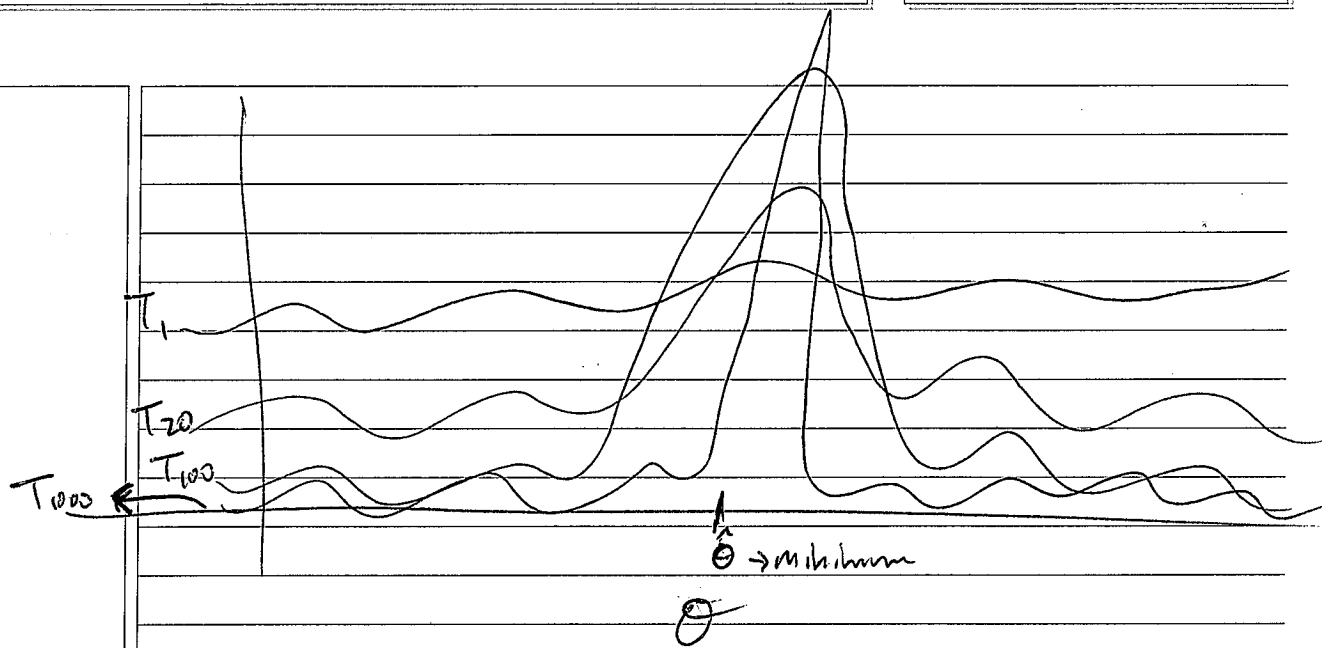
Thm: If

$$(1) T_{n+1} < T_n \quad \forall n$$

$$(2) T_n \rightarrow 0 \quad \text{as} \quad n \rightarrow \infty$$

$$(3) T_n = a / \log(n+b)$$

$$\text{Then } \|\pi_{T_n} - \pi^*\| \rightarrow 0$$



As $T_n \rightarrow 0$, the mass of π_{T_n} becomes increasingly concentrated about $\hat{\theta}$

Perfect Sampling (brief intro)

We know that under certain conditions, a MC will converge to a stationary distribution.
 \hookrightarrow at time $n \approx \infty$

But we don't know when!

We know a burn-in is a good idea, but how much?

When can we declare "convergence"?

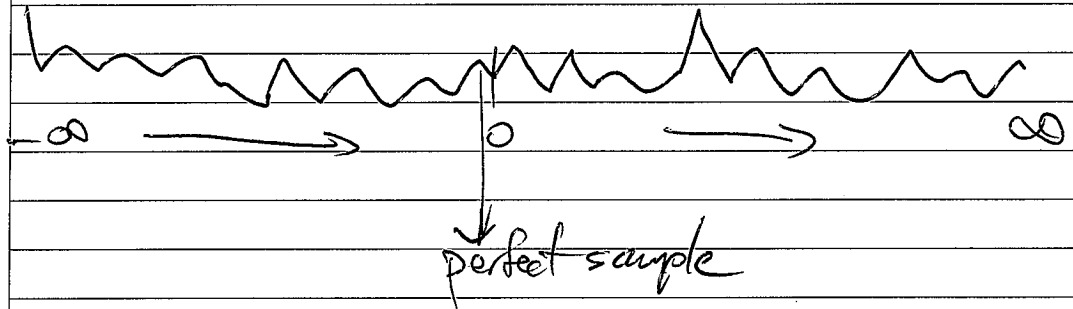
Propp + Wilson (1996) discovered that perfect random samples can be obtained in finite (but stochastic) time.

\Rightarrow Coupling from the past algorithm

Basic idea
of CFTP

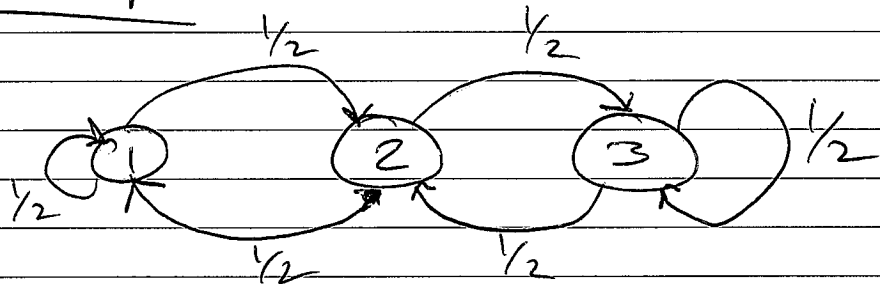
① If we start at $n=0$, then a sample from state $n=\infty$ will be a "perfect" sample from π , the stationary distribution.

② So why don't we start the chain at state $n=-\infty$ and ~~take a sample from~~ use the value drawn at state $n=0$?



Both chains run for ∞ amount of time so must be stationary.

Simple example



Discrete
state space

Transition matrix

$$P = \begin{bmatrix} 1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 \\ 0 & 1/2 & 1/2 \end{bmatrix}$$

Stationary dist: $\pi = (1/3, 1/3, 1/3)$

The stationary dist. is uniform on all 3 states.

How do we simulate this Markov Chain?

Suppose ~~we~~ X_n is the value in state n .

Draw a $u \sim \text{uniform}(0,1)$.

If $X_n = 1$ then goto $\begin{cases} X_{n+1} = 1 & u \leq \frac{1}{2} \\ X_{n+1} = 2 & u > \frac{1}{2} \end{cases}$

If $X_n = 2$ then goto $\begin{cases} X_{n+1} = 1 & u \leq \frac{1}{2} \\ X_{n+1} = 3 & u > \frac{1}{2} \end{cases}$

If $X_n = 3$ then goto $\begin{cases} X_{n+1} = 2 & u \leq \frac{1}{2} \\ X_{n+1} = 3 & u > \frac{1}{2} \end{cases}$

Imagine running the chain in this way starting at $n = -\infty$ and ending at $n = 0$.

Can we emulate this process?

Note:

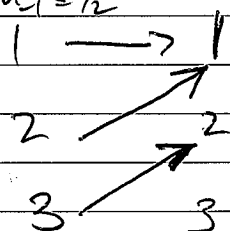
- ① We can generate $u_n, u_{n+1}, u_{n+2}, \dots$ because they are independent of everything.
- ② $X_{n+1} = \phi(u_n, X_n)$

CFTP Algorithm for 1 perfect sample

~~Suppose we draw $U \leq \frac{1}{2}$~~

Start 3 parallel MCs each starting in state 1, 2, or 3 respectively.

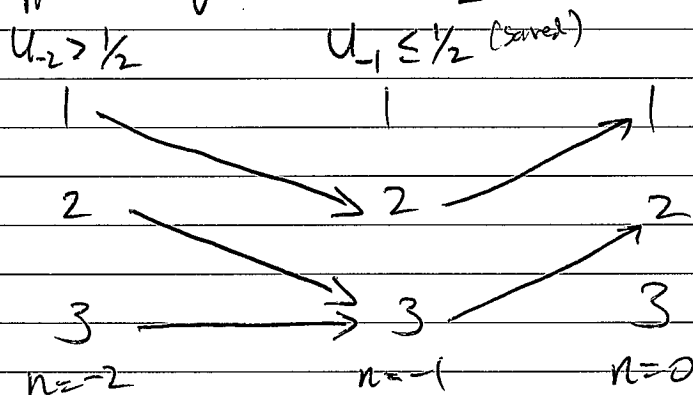
Suppose we draw $U \leq \frac{1}{2}$ and start at $n = -1$
 $U_1 \leq \frac{1}{2}$



$n = -1$ $n = 0$

If we start the chains at $n = -1$ and draw $U \leq \frac{1}{2}$, then we are at ~~state~~ state 1 or 2. (not 3)

Suppose we generate $U_{-2} > \frac{1}{2}$ and start at $n = -2$.



$n = -2$

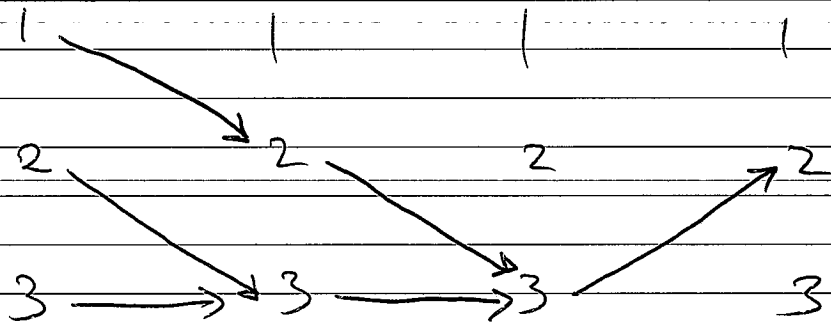
$n = -1$

$n = 0$

We "lost" the second chain, or the 2nd and 3rd chains coalesced.

Now suppose generate $U_3 > \frac{1}{2}$ and start at $n = -3$.

$U_{-3} > \frac{1}{2}$ $U_{-2} > \frac{1}{2}$ $U_1 \leq \frac{1}{2}$



$n = -3$

$n = -2$

$n = -1$

$n = 0$

When we start from $n = -3$, all of the chains coalesce at state $n = -1$.

~~Regard~~

\Rightarrow Regardless of what happens b/w ~~states~~ times $n = -\infty$ and $n = -3$ (we don't know), we know that at ~~state~~^{time} $n = -1$ they will all coalesce at ~~the~~ state (3).

\Rightarrow Regardless of starting value, at time $n = 0$, we will be at state (2).

\Rightarrow (2) is a perfect sample from stationary dist.

\Rightarrow Coalescence "breaks" dependence on the starting value ~~or initial~~

Note: In this chain, the arrows don't cross, so we only need to keep track of "top" chain and "bottom" chain

① Why can't run forward & until coaply?

\hookrightarrow only coaples at ① and ③

② Why can't run until coaply and then a few more?

How many more?