

Double Fibations



Since it is not obvious what the codomain
of an indexing functor

$$A \longrightarrow ?? \quad (A \text{ dbl cat}^?)$$

should be, we start with double fibrations.



Double Fibrations

First Idea: a double category is a pseudo cat^d in
Cat and a double fibration is a pseudo cat^d in
Fib.

This would amount to:

$$\begin{array}{ccc} \underline{E}_1 \times_{\underline{E}_0} \underline{E}_1 & \xrightarrow{\otimes^T} & \underline{E}_1 \xrightarrow{s^T} \underline{E}_0 \\ \downarrow P_1 \times_{P_0} P_1 & & \downarrow P_1 \xrightarrow{y^T} t^T \\ \underline{B}_1 \times_{\underline{B}_0} \underline{B}_1 & \xrightarrow{\otimes^\perp} & \underline{B}_1 \xrightarrow{s^\perp} \underline{B}_0 \\ & & \downarrow P_1 \xrightarrow{y^\perp} t^\perp \end{array}$$

... strict dd functors with extra properties?



Double fibrations

Problem : Fib doesn't have all 2-pullbacks required for this.

Observation:

- We require the same fibrational strictness for s and t that we require for y and \otimes .

Solution: We will require that s and t are in CFib .
(i.e., they preserve cleavages)



Double Fibrations

Definition: A double fibration is a (strict) double functor $P: \mathbb{E} \rightarrow \mathbb{B}$ between pseudo double categories:

$$\begin{array}{ccc} \underline{E}_1 \times_{\underline{E}_0} \underline{E}_1 & \xrightarrow{\otimes_E} & \underline{E}_1 \xrightleftharpoons[s_E]{t_E} \underline{E}_0 \\ P_1 \times_{P_0} P_1 \downarrow & & \downarrow P_1 \qquad \qquad \qquad \downarrow P_0 \qquad \text{s.t.} \\ \underline{B}_1 \times_{\underline{B}_0} \underline{B}_1 & \xrightarrow{\otimes_B} & \underline{B}_1 \xrightleftharpoons[s_B]{t_B} \underline{B}_0 \end{array}$$

- P_0 and P_1 are fibrations with cleavages
- s_E and t_E are cleavage preserving
- \otimes_E and \otimes_B are cartesian-morphism preserving.



Examples

1. $\text{Im} : \underline{\text{Span}}(\underline{\text{Set}}) \longrightarrow \text{Rel}$
 $(A \xleftarrow{s} S \xrightarrow{t} B) \longmapsto (\text{Im}((s,t)) \subseteq A \times B)$
is a double opfibration.

2. When $E_0 = B_0 = 1$, we get :

$$\begin{array}{ccccc} E_1 & \times_{E_0} & E_1 & \xrightarrow{\otimes_E} & E_1 \leftarrow \ddot{\gamma} \longrightarrow 1 \\ P_1 \times P_0 & \downarrow & & \downarrow P_1 & \downarrow \text{id} \\ B_1 & \times_{B_0} & B_1 & \xrightarrow{\otimes_B} & B_1 \leftarrow \gamma \longrightarrow 1 \end{array}$$

a monoidal fibration!

3. The Grothendieck constructions given in David Jaz Myers' work are also double fibrations.



Examples

4. For any 2-functor $P: \underline{E} \rightarrow \underline{B}$, P is a 2-fibration as in Buckley's work if and only if $\mathbb{Q}(P): \mathbb{Q}(E) \rightarrow \mathbb{Q}(B)$ is a double fibration.

5. if P_0 and P_1 are discrete fibrations, we recover discrete double fibrations.

6. For \mathbb{D} a double cat^y, let $\mathbb{D}^z = \begin{pmatrix} \mathbb{D}_1^z \\ \downarrow \mathbb{D}_0^z \end{pmatrix}$

dom: $\mathbb{D}^z \longrightarrow \mathbb{D}$
is a double fibration



7. The codomain fibration extends to a double codomain fibration $\text{cod}: \mathbb{D}^2 \rightarrow \mathbb{D}$
- if:
- \mathbb{D}_1 and \mathbb{D}_0 have chosen finite limits
 - these limits are preserved on the nose by s and t
 - and up to iso by y and \otimes .



8. For \mathcal{C} a small cat $^{\mathbb{D}}$, $\text{Fun}(\mathcal{C})$ has:

obj: $f: I \rightarrow \mathcal{C}$, or $(I, \{C_i\}_{i \in I})$

arrows: $\frac{(h, \alpha): f \rightarrow g}{h: I \rightarrow J, \alpha: f \Rightarrow g h} \quad \begin{array}{c} I \xrightarrow{h} J \\ f \searrow \alpha \swarrow g \end{array}$

or: $(I, \{C_i\}_{i \in I})$

\downarrow $(h: I \rightarrow J, \{\alpha_i: C_i \rightarrow C_{h(i)}\}_{i \in I})$

$(J, \{C_j\}_{j \in J})$



Pro arrows : $f \xrightarrow{(S, d_0, d_1, \theta)}$

for natural transformations

$$d_0 \downarrow \begin{matrix} S \\ \xrightarrow{\theta} \\ \downarrow p \end{matrix} K \quad \text{for a span of functions}$$

$$I \xrightarrow{f} e \quad I \xrightarrow{d_0} S \xrightarrow{d_1} K$$

Or :

$$(I, \{C_i\}_{i \in I}) \xrightarrow{((d_0, S, d_1), \theta_s)} (K, \{C_k\})$$

for a span of functions $I \xrightarrow{d_0} S \xrightarrow{d_1} K$

+ a family of arrows

$$\theta_s : C_{d_0(s)} \longrightarrow C_{d_1(s)}$$



cells:

a cell from

$$\begin{array}{ccc} S & \xrightarrow{d_1} & K \\ d_0 \downarrow & \cong & \downarrow p \\ I & \xrightarrow{f} & e \end{array} \quad \text{to} \quad \begin{array}{ccc} T & \xrightarrow{d'_1} & L \\ d'_0 \downarrow & \cong & \downarrow q \\ J & \xrightarrow{g} & e \end{array}$$

is given by a morphism of spans:

$$\begin{array}{ccccc} I & \xleftarrow{d_0} & S & \xrightarrow{d_1} & K \\ h \downarrow & & \downarrow m & & \downarrow r \\ J & \xleftarrow{d'_0} & T & \xrightarrow{d'_1} & L \end{array}$$

with 2-cells:

$$\begin{array}{ccc} e & \begin{matrix} \nearrow I \\ \nearrow \alpha \\ \searrow J \end{matrix} & h \\ & \nearrow g & \end{array} \quad \begin{array}{ccc} e & \begin{matrix} \nearrow K \\ \nearrow \beta \\ \searrow L \end{matrix} & r \\ & \nearrow q & \end{array} \quad \begin{matrix} \text{st:} \\ (\beta * d_1) \theta = \\ (\delta * g) (\alpha * d_0) \end{matrix}$$



or: a family of cells:

$$(I, \{C_i\}_{i \in \Sigma}) \xrightarrow{((d_0, S, d_1), \theta)} (K, \{C_k\}_{k \in K})$$

$$(h, (\alpha_i)) \downarrow \quad m \quad \downarrow (r, (\beta_k))$$

$$(J, \{C_j\}_{j \in J}) \xrightarrow{((d'_0, T, d'_1), \theta')} (L, \{C_l\}_{l \in L})$$

where $m : S \rightarrow T$ fits in

$$I \xleftarrow{d_0} S \xrightarrow{d_1} K$$

$$h \downarrow \quad \quad \quad \downarrow m \quad \quad \quad \downarrow r$$

$$J \xleftarrow{d'_0} T \xrightarrow{d'_1} L$$

and we require that for each $s \in S$:



$$\begin{array}{ccc}
 C_{d_0(s)} & \xrightarrow{\Theta_s} & C_{d_1(s)} \\
 \alpha_{d_0(s)} \downarrow & & \downarrow \beta_{d_1(s)} \\
 C_{h(d_0 s)} & \xrightarrow{\Theta'_{m(s)}} & C_{r(d_1 s)}
 \end{array}$$

this square is well defined because

$$d_0^l(m(s)) = h d_0(s)$$

$$\text{and } d_1^l(m(s)) = r d_1(s)$$

We only need to require that

$$\beta_{d_1(s)} \circ \Theta_s = \Theta'_{m(s)} \circ \alpha_{d_0(s)},$$

because C has no 2-cells.

Note: This can be extended to C a dbl. cat^y



$$\Pi_0 : \text{TFam}(\mathcal{C})_0 \rightarrow \underline{\text{Set}}$$

is a split fibration.

We extend this to

$$\Pi : \text{TFam}(\mathcal{C}) \rightarrow \text{Span}(\underline{\text{Set}}).$$

(send proarrows to their underlying spans and cells to open morphisms)

Claim: this is a split double fibration.



Relation with
Street's Internal Vibrations
(a look-off trail we may not take)



Internal Fibrations

- Internal fibrations in a 2-category were introduced by Street in 1974.
- We will use the following three 2-cats:

$$\frac{\mathcal{Dbl}_s}{\begin{array}{c} \uparrow \\ \text{strict dbl} \\ \text{functors} \end{array}} \subseteq \frac{\mathcal{Dbl}}{\begin{array}{c} \uparrow \\ \text{pseudo} \\ \text{dbl functors} \end{array}} \subseteq \frac{\mathcal{Dbl}_l}{\begin{array}{c} \uparrow \\ \text{lax dbl} \\ \text{functors} \end{array}}$$

say more
about Street's
definition?

- In all 3 cases we use vertical transformations

$\alpha: F \Rightarrow G: \mathbb{D} \rightrightarrows \mathbb{E}$:

* for each X in \mathbb{D} : an arrow $\alpha_X: Fx \rightarrow Gx$ in \mathbb{E}

* for each prob arrow $X \xrightarrow{\sim} Y$ a dbl

cell $\alpha_m: Fm \rightarrow Gm$ in \mathbb{E}_1 , (natural and functorial)



Theorem (Cruttwell, Lambert, P, Szyl) 

A strict double functor is an internal fibration in DblCat if and only if it is a double fibration

In addition, a pseudo double functor P

- is an internal fibration in DblCat,
iff P_0 and P_1 admit cleavages that
are preserved by s_E and t_E .
- is an internal fibration in DblCat
iff. in addition, y_E and \otimes_E preserve
Cartesian morphisms.

Furthermore, a strict double functor P is an internal fibration in \mathbf{DblCat}_S if and only if P_0 and P_1 are fibrations that admit cleavages that are preserved by all of s_E , t_E , y_E and \otimes_E



Double Indexing Functors (Take 1)



Double Indexing Functors

Exercise!

Note: ① Categories are monoids in $\text{Span}(\text{Set})$

$$\begin{array}{ccc}
 C_0 & \xleftarrow{\text{id}} & C_0 \xrightarrow{\text{id}} C_0 \\
 \parallel & & \downarrow y \\
 C_0 & \xleftarrow{s} & C_1 \xrightarrow{t} C_0
 \end{array}
 \quad
 \begin{array}{ccc}
 C_0 & \xleftarrow{s} & C_1 \xrightarrow{t} C_0 \\
 \parallel & & \downarrow \mu \\
 C_0 & \xleftarrow{s} & C_1 \xrightarrow{t} C_0
 \end{array}$$

Unit \hookrightarrow identities
for the cat \mathcal{Y}

multiplication \hookrightarrow
composition for the
cat \mathcal{Y} .

② Moeller and Vasilakopoulou used:

$$\underline{\text{Fib}} \simeq \underline{\text{ICat}} \implies \underline{\text{Ps Mon}}(\underline{\text{Fib}}, \times) \simeq \underline{\text{Ps Mon}}(\underline{\text{ICat}}, \times)$$



To generalize this further we need :

- * double 2-categories (pseudo category objects in 2-Cat)
- * pseudo monoids in double 2-categories

Result : pseudo categories in a 2-cat^y \mathcal{C}
correspond to pseudo monoids in $\text{Span}(\mathcal{C})$.

Recall that we want to take the source and target
from a more restricted class of arrows, say Σ :

Result: pseudo cat^s in \mathcal{C} with s, t in Σ correspond

to pseudo monoids in $\text{Span}_\Sigma(\mathcal{C})$.



Double 2-Categories

A double 2-category is a pseudo cat^d in 2Cat

$$\mathcal{E}_1 \times_{\mathcal{E}_0} \mathcal{E}_1 \xrightarrow{\otimes} \mathcal{E}_1 \xrightarrow{s} \mathcal{E}_0$$

so we have

$$\begin{array}{ccc} \bullet & \xrightarrow{\quad} & \bullet \\ & \downarrow \alpha & \downarrow \\ \bullet & \xrightarrow{\quad} & \bullet \end{array}$$

objects arrows

$$\Gamma : \alpha \Rightarrow \beta \left(\begin{array}{c} \overbrace{\Gamma \quad \beta}^{\text{2-cells}} \\ \alpha \end{array} \right) \quad \left(\begin{array}{c} \vdots \\ \Gamma \end{array} \right)$$

2-cells



The Double 2-Category $\mathbf{Span}(\underline{\mathbf{Cat}})$

Objects categories A, B, C, \dots

Arrows functors

Proarrows spans $\mathcal{S} \xleftarrow{S} A \xrightarrow{T} J$

Dbl cells commutative diagrams

$$\begin{array}{ccccc} \mathcal{S} & \xleftarrow{S} & A & \xrightarrow{T} & J \\ g \downarrow & & \downarrow F & & \downarrow H \\ \mathcal{S}' & \xleftarrow{S'} & A' & \xrightarrow{T'} & J' \end{array}$$



Dbl 3-cells

$$\begin{array}{ccc} \mathcal{S} & \xleftarrow{\quad S \quad} & \mathcal{A} & \xrightarrow{\quad T \quad} & \mathcal{T} \\ \downarrow g & & \downarrow f & & \downarrow h \\ \mathcal{X} & \xleftarrow{\quad X \quad} & \mathcal{B} & \xrightarrow{\quad Y \quad} & \mathcal{Y} \end{array} \quad (\gamma, \varphi, \theta) \Rightarrow$$

$$\begin{array}{ccc} \mathcal{S} & \xleftarrow{\quad S \quad} & \mathcal{A} & \xrightarrow{\quad T \quad} & \mathcal{T} \\ \downarrow g' & & \downarrow f' & & \downarrow h' \\ \mathcal{X} & \xleftarrow{\quad X \quad} & \mathcal{B} & \xrightarrow{\quad Y \quad} & \mathcal{Y} \end{array}$$



Consists of

$$\begin{array}{ccccc} \mathcal{G} & \xleftarrow{s} & A & \xrightarrow{T} & \mathcal{T} \\ \downarrow \varphi & \quad F \downarrow \varphi' & \downarrow F' & \quad H \downarrow \theta' & \downarrow \theta \\ \mathcal{X} & \xleftarrow{x} & B & \xrightarrow{Y} & \mathcal{Y} \end{array}$$

such that the two cylinders commute.



Double Indexing Functors

$\mathbf{I} \underline{\text{Span}}(\underline{\text{Cat}})$ is the slice category

with

• objects: Contravariant lax double pseudo functors $F: A^{\text{op}} \longrightarrow \underline{\text{Span}}(\underline{\text{Cat}})$

(A -indexed ps. dbl. categories)

• morphisms: $(H, \tau): F \rightarrow G:$

$$\begin{array}{ccc} H \text{ lax dbl.} & & A^{\text{op}} \xrightarrow{H^{\text{op}}} B^{\text{op}} \\ \text{pseudo functor} & & \\ T: \text{lax dbl.} & \Downarrow \tau & \\ \text{ps. ntl. trato} & & \underline{\text{Span}}(\underline{\text{Cat}}) \end{array}$$



Lax Double Pseudo Functors

Definition

A lax double pseudo functor

$$F : \mathbb{D} \rightarrow \mathbb{E}$$

between dbl 2-categories \mathbb{D}, \mathbb{E} , consists of:

- Pseudo functors: $F_0 : \mathbb{D}_0 \rightarrow \mathbb{E}_0$ (now 2-categories!)
 $F_1 : \mathbb{D}_1 \rightarrow \mathbb{E}_1$ pseudo in the arrow direction!
- Comparison pseudo natural transformations:

$$\begin{array}{ccc} \mathbb{D}_0 \times \mathbb{D}_1 & \xrightarrow{\otimes} & \mathbb{D}_1 \\ F_0 \times F_1 \downarrow \quad \quad \quad \Downarrow \phi \quad \quad \quad \downarrow F_1 \\ \mathbb{E}_0 \times_{\mathbb{E}_0} \mathbb{E}_1 & \xrightarrow{\otimes} & \mathbb{E}_1 \end{array} \qquad \begin{array}{ccc} \mathbb{D}_0 & \xrightarrow{x} & \mathbb{D}_1 \\ F_0 \downarrow \quad \quad \quad \Downarrow \gamma \quad \quad \quad \downarrow F_1 \\ \mathbb{E}_0 & \xrightarrow{y} & \mathbb{E}_1 \end{array}$$

lax in the proarrow direction



- Invertible associativity and unitor modifications

$$\begin{array}{ccccc}
 & \mathbb{D}_i^{(3)} & & \mathbb{D}_i^{(2)} & \\
 \mathbb{E}_i^{(3)} & \xleftarrow{\tau_i^{(3)}} & \mathbb{D}_i^{(3)} & \xrightarrow{\otimes x_1} & \mathbb{D}_i^{(2)} \\
 & \downarrow 1 \times \phi & \downarrow \alpha \cong & & \downarrow \otimes \\
 & \mathbb{D}_i^{(2)} & & \mathbb{D}_i^{(2)} & \\
 & \downarrow 1 \times \otimes & & \downarrow \phi & \\
 \mathbb{E}_i^{(2)} & & \mathbb{D}_i & & \mathbb{D}_i \\
 & \searrow \otimes & \swarrow \mathbb{F}_i & & \\
 & & \mathbb{E}_i & &
 \end{array}$$

$\vdash \dashv$

$$\begin{array}{ccccc}
 & \mathbb{D}_i^{(3)} & & \mathbb{D}_i^{(2)} & \\
 \mathbb{E}_i^{(3)} & \xleftarrow{\tau_i^{(3)}} & \mathbb{D}_i^{(3)} & \xrightarrow{\phi x_1} & \mathbb{D}_i^{(2)} \\
 & \downarrow 1 \times \phi & \downarrow \alpha \cong & \downarrow \phi & \downarrow \otimes \\
 & \mathbb{E}_i^{(2)} & & \mathbb{D}_i^{(2)} & \\
 & \downarrow 1 \times \otimes & & \downarrow \phi & \\
 \mathbb{E}_i^{(2)} & & \mathbb{D}_i & & \mathbb{D}_i \\
 & \searrow \otimes & \swarrow \mathbb{F}_i & & \\
 & & \mathbb{E}_i & &
 \end{array}$$

e.t.c.

satisfying well-definedness and coherence conditions.



Lax Double Pseudo Natural Transformations

A lax dbl ps. ntl transformation:

$$\tau : \mathcal{F} \Rightarrow \mathcal{G} : \mathbb{D} \rightrightarrows \mathbb{E}$$

consists of:

- pseudo ntl transformations

$$\tau_0 : \mathcal{F}_0 \Rightarrow \mathcal{G}_0, \tau_1 : \mathcal{F}_1 \Rightarrow \mathcal{G}_1$$

(arrow-component transformations)

- modifications:

$$\begin{array}{ccc} \mathbb{D}_0 \times_{\mathbb{D}_0} \mathbb{D}_1 & = & \mathbb{D}_0 \times_{\mathbb{D}_0} \mathbb{D}_1 \xrightarrow{\otimes} \mathbb{D}_1 \\ \downarrow \mathcal{F}_0 \times \mathcal{F}_1 & \xrightarrow{\tau_0 \times \tau_1} & \downarrow \mathcal{G}_1 \\ \mathbb{E}_0 \times_{\mathbb{E}_0} \mathbb{E}_1 & = & \mathbb{E}_0 \times_{\mathbb{E}_0} \mathbb{E}_1 \xrightarrow{\otimes} \mathbb{E}_1 \end{array} \quad \Rightarrow \quad \begin{array}{ccc} \mathbb{D}_0 \times_{\mathbb{D}_0} \mathbb{D}_1 & \xrightarrow{\otimes} & \mathbb{D}_1 \\ \downarrow \mathcal{F}_0 \times \mathcal{F}_1 & \xrightarrow{\phi} & \downarrow \mathcal{F}_1 \\ \mathbb{E}_0 \times_{\mathbb{E}_0} \mathbb{E}_1 & \xrightarrow{\otimes} & \mathbb{E}_1 \end{array}$$



and

$$\begin{array}{ccc} \mathbb{D}_0 = \mathbb{D}_0 & \xrightarrow{\alpha} & \mathbb{D}_1 \\ \downarrow \tau_0 & \Rightarrow & \downarrow \beta_0 \\ \mathbb{E}_0 = \mathbb{E}_0 & \xrightarrow{\gamma} & \mathbb{E}_1 \end{array} \quad \begin{array}{ccc} \mathbb{D}_0 & \xrightarrow{\delta} & \mathbb{D}_1 = \mathbb{D}_1 \\ \downarrow \tau_1 & \Rightarrow & \downarrow \beta_1 \\ \mathbb{E}_0 & \xrightarrow{\gamma} & \mathbb{E}_1 = \mathbb{E}_1 \end{array}$$

satisfying multiplicativity and
unitality conditions.

We write DblzCat (\mathbb{D}, \mathbb{E}) for the Cat $^{\mathcal{T}}$
of lax dbl pseudo functors and lax dbl. ps. n.t.
transformations.



The Representation Theorem



The Representation Theorem

Theorem (Cruttwell, Lambert, P., Szajd)

There is an equivalence of categories

$$\underline{\text{DblFib}} \simeq \underline{\text{I Span}}(\underline{\text{Cat}})$$

Idea for the proof: use pseudo monoids in
double 2-categories.

$$\underline{\text{Fib}} \simeq \underline{\text{I Cat}} \text{ and } \underline{\text{Cfib}} \simeq \underline{\text{I Cat}}^t$$

$$\text{so: } \underline{\text{Span}}_c(\underline{\text{Fib}}) \simeq \underline{\text{Span}}_t(\underline{\text{I Cat}})$$

(Convince yourself!)



$\mathbb{S}\text{pan}_c(\underline{\text{Fib}})$ has objects:

\underline{E} fibrations with cleavage
 $\downarrow p$
 \underline{B}

Notation:
 p

$\mathbb{S}\text{pan}_t(\underline{\text{ICat}})$ has objects: $\underline{B}^{\text{op}} \xrightarrow{F} \underline{\text{Cat}}$ ps.functor F

arrows:

$$\begin{array}{ccc} \underline{E} & \xrightarrow{f^T} & \underline{E}' \\ \downarrow p & & \downarrow p' \\ \underline{B} & \xrightarrow{f^\perp} & \underline{B}' \end{array} \quad \begin{array}{l} \text{Cartesian arrow} \\ \text{preserving} \end{array} \quad \begin{array}{l} p \\ \downarrow f \\ p' \end{array}$$

arrows:

$$\begin{array}{ccc} \underline{B}^{\text{op}} & \xrightarrow[\theta]{H} & (\underline{B}')^{\text{op}} \\ \downarrow F & \Rightarrow & \downarrow F' \\ \underline{\text{Cat}} & & \end{array} \quad \begin{array}{l} \theta \text{ pseudo} \\ \text{trafo } (H, \theta) \end{array} \quad \begin{array}{l} F \\ \downarrow \\ F' \end{array}$$



proarrows : $P \xleftarrow{l} Q \xrightarrow{r} R$ l^T, r^T cleavage preserving

proarrows : $F \xleftarrow{(L,\alpha)} G \xrightarrow{(R,\rho)} K$ α, ρ strict natural transfs

cells : $P \xleftarrow{l} Q \xrightarrow{r} R$ $l^T, r^T, l^{\dagger}, r^{\dagger}$
 $f \downarrow \quad \quad \quad g \downarrow \quad \quad \quad h \downarrow$
 $P' \xleftarrow{l'} Q' \xrightarrow{r'} R'$ cleavage preserving

cells : $F \xleftarrow{(L,\alpha)} G \xrightarrow{(R,\rho)} K$ f^T, g^T, h^T
 $(H,\theta) \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow$
 $F' \xleftarrow{(L',\alpha')} G' \xrightarrow{(R',\rho')} K'$ Cartesian arrow pres.



Then Lift :

$$\begin{aligned}\underline{\mathbf{Dbl}\text{-}\mathbf{Fib}} &:= \underline{\mathbf{Ps}\text{-}\mathbf{Mon}}\left(\mathbf{Span}_c(\mathbf{Fib})\right) \simeq \\ &\simeq \underline{\mathbf{Ps}\text{-}\mathbf{Mon}}\left(\mathbf{Span}_t(\mathbf{I}\underline{\mathbf{cat}})\right) \simeq \mathbf{I}\mathbf{Span}(\underline{\mathbf{Cat}}).\end{aligned}$$



Connections with Known Constructions

Monoidal Fibrations

monoids
are cats

2-Fibrations

quintet constr.

Double Fibrations

Discrete Dbl Fib.

Double Gr. Constr.

All of these have an indexed notion and
a Grothendieck category of elements.

In our case the composition of the
functors used in the proof gives us also a
category of elements construction.



The Double Category of Elements



The Double Grothendieck Construction

Start with $F : \mathbb{D}^{\circ\text{op}} \longrightarrow \text{Span}(\underline{\text{Cat}})$:

$$F_0 : \mathbb{D}_0^{\circ\text{op}} \longrightarrow \text{Span}(\underline{\text{Cat}})_0 = \underline{\text{Cat}} \quad (1)$$

$$F_1 : \mathbb{D}_1^{\circ\text{op}} \longrightarrow \text{Span}(\underline{\text{Cat}})_1$$

and a further induced functor:

$$\mathbb{D}_1^{\circ\text{op}} \xrightarrow{F_1} \text{Span}(\underline{\text{Cat}})_1 \xrightarrow{\text{apx}} \underline{\text{Cat}} \quad (2)$$

Apply the ordinary elements construction to (1) and (2):

$$\mathbb{E}\mathbb{I}(F)_0 \longrightarrow \mathbb{D}_0 \quad \mathbb{E}\mathbb{I}(F)_1 \longrightarrow \mathbb{D}_1$$

Cloven fibrations.



Details of $\mathbb{E}l(\mathcal{F})$

Some notation:

- for $m: A \rightarrow B$ in \mathbb{D} , write Fm in $\text{Span}(\underline{\text{Cat}})$

as :

$$\begin{array}{ccc} & Fm & \\ Lm \swarrow & & \searrow Rm \\ FA & & FB \end{array}$$

- for $A \xrightarrow{m} B \xrightarrow{n} C$ in \mathbb{D} , we have the laxity

comparison cell

$$\begin{array}{ccccc} FA & \xleftarrow{\quad} & Fm \times_{FB}^n & \xrightarrow{\quad} & FC \\ \parallel & & \downarrow \phi_{m,n} & & \parallel \\ FA & \xleftarrow{L_{m \otimes n}} & F(m \otimes n) & \xrightarrow{R_{m \otimes n}} & FC \end{array}$$



Now $\text{El}(F)$ is given by:

- objects: (C, x) $C \text{ in } \mathbb{D}, x \text{ in } FC.$
- arrows: $(f, \bar{f}): (C, x) \longrightarrow (D, y)$
with $f: C \rightarrow D$ in \mathbb{D}
and $\bar{f}: x \rightarrow f^*y$ ($= F(f)y$) in $FC.$
- proarrows: $(m, \bar{m}): (C, x) \dashrightarrow (D, y)$
with $C \xrightarrow{\quad m \quad} D$ in \mathbb{D}
 $\bar{m} \in Fm \quad \left(FC \xleftarrow{Lm} Fm \xrightarrow{Rm} FD \right)$
s.t. $L_m(\bar{m}) = x$
 $R_m(\bar{m}) = y$



• double cells :

$$\begin{array}{ccc} (A, x) & \xrightarrow{(m, \bar{m})} & (B, y) \\ (f, f) \downarrow & & \downarrow (g, \bar{g}) \\ (C, z) & \xrightarrow{(n, \bar{n})} & (D, w) \end{array}$$

with

$$\begin{array}{ccc} A & \xrightarrow{m} & B \\ f \downarrow & \theta & \downarrow g \\ C & \xrightarrow{n} & D \end{array}$$

$$\text{in } \mathbb{D} \quad \begin{array}{ccccc} FA & \xleftarrow{L_m} & Fm & \xrightarrow{R_m} & FB \\ \downarrow f & & \downarrow F\theta & & \downarrow \bar{f} \\ FC & \xleftarrow{L_n} & Fn & \xrightarrow{R_n} & FD \end{array}$$

and $\bar{m} \xrightarrow{\bar{\theta}} \theta^* \bar{n}$ an arrow in Fm s.t.

$$L_m(\bar{\theta}) = \bar{f} \quad \text{and} \quad R_m(\bar{\theta}) = \bar{f}.$$



Composition in the "arrow direction" is as expected:

- for arrows $(A, x) \xrightarrow{(f, \bar{f})} (B, y) \xrightarrow{(g, \bar{g})} (C, z)$
the composite is $(gf, \phi_{f,g} f^*(\bar{g}) \bar{f}): (A, x) \rightarrow (C, z)$
- for cells $(m, \bar{m}) \xrightarrow{(\theta, \bar{\theta})} (n, \bar{n}) \xrightarrow{(\delta, \bar{\delta})} (p, \bar{p})$
the composite is $(\delta\theta, \phi_{\theta,\delta} \theta^*(\bar{\delta}) \bar{\theta}): (m, \bar{m}) \rightarrow (p, \bar{p})$.
- Units: $(1_{C,x} (\varphi_C)): (C, x) \rightarrow (C, x)$
 $(1_m, (\varphi_m)_{\bar{m}}): (m, \bar{m}) \rightarrow (m, \bar{m})$.



$\mathbb{E}I(F), \frac{s}{t}, \mathbb{E}I(F)_o$ are defined by

$$s(\theta, \bar{\theta}) = (f, \bar{f})$$

$$t(\theta, \bar{\theta}) = (g, \bar{g})$$

- Proarrow composition:

for $(A, x) \xrightarrow{(m, \bar{m})} (B, y) \xrightarrow{(n, \bar{n})} (C, z)$

the composite is:

$$(m \otimes n, \phi_{m,n}(\bar{m}, \bar{n})) : (A, x) \longrightarrow (C, z)$$

- Composition for cells and proarrow units
are given using appropriate components of the
structure nos related to pseudo naturality.



Results

- $\overline{\Pi}: \mathbb{E}I(F) \rightarrow \mathbb{D}$ is a double fibration.
- This is the object part of an equivalence
of categories

$$\underline{\text{DblFib}} \simeq \mathbf{ISpan}(\underline{\text{Cat}})$$

which specializes to

$$\underline{\text{DblFib}}(\mathbb{B}) \simeq \underline{\text{DblzCat}}(\mathbb{B}^{\otimes}, \text{Span}(\underline{\text{Cat}}))$$

for each dbl cat^d \mathbb{B} .



Examples

- Let \underline{A} be a category with pushouts and $\mathbb{C}\text{sp}(\underline{A})$ the double cat^g with cells:

$$\begin{array}{ccccc} X & \longrightarrow & Z & \leftarrow & Y \\ u \downarrow & & \downarrow w & & \downarrow v \\ X' & \longrightarrow & Z' & \leftarrow & Y' \end{array}$$

- For any lax double functor

$$F: \mathbb{C}\text{sp}(\underline{A}) \rightarrow \mathbb{Span}(\underline{\text{Cat}})$$

the dbl cat^g of elements $\text{IEL}(F)$ in $F\text{-}\mathbb{C}\text{sp}$,
the double category of F -decorated cospan.
[Patterson, 2023].



This slightly generalizes the decorated cospan
from Baez-Cousens - Vasilakopoulou.