

Advanced Techniques for Combinatorial Algorithms: Introduction to Information Theory

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Entropy

Entropy

Let A be a discrete random variable over universe set U .
Then its entropy is

$$H(A) = \sum_{x \in U} p_x \log_2 \frac{1}{p_x} = - \sum_{x \in U} p_x \log_2 p_x$$

where p_x is the probability of the event x .

Entropy measures information content.

Entropy 2

Entropy of a text T

Let $f(\sigma)$ be the relative frequency of symbol σ in T . Then its entropy is

$$H(T) = - \sum_{\sigma \in \Sigma} f(\sigma) \log_2 f(\sigma)$$

Jensen Inequality

Jensen Inequality

Let A be a random variable, and let ψ be a convex function. Then

$$\psi(E(A)) \leq E(\psi(A))$$

where $E()$ is the expectation.

Prefix code

Code

$$\phi : A \subset \mathbb{N} \mapsto \{0, 1\}^*$$

Variable length codewords

Example

$0 \mapsto 1010$

$1 \mapsto 0010$

$2 \mapsto 00$

Not prefix code

Prefix code

Prefix Code

No codeword is the prefix of another codeword.

Examples

$0 \mapsto 1$

$1 \mapsto 01$

$2 \mapsto 001$

Lemma

All prefix codes are uniquely decodable.

Elias gamma encoding

Encoding $n \geq 1$

- 1 $N \leftarrow \lfloor \log_2 n \rfloor$
- 2 N zeroes followed by n in binary

Requires $2\lfloor \log_2 n \rfloor + 1$ bits

Examples

$9 = 1001_2 \mapsto 0001001$

$15 = 1111_2 \mapsto 0001111$

Elias delta encoding

Encoding $n \geq 1$

- 1 $N \leftarrow \lfloor \log_2 n \rfloor$
- 2 $L \leftarrow \lfloor \log_2 N + 1 \rfloor$
- 3 Elias γ encoding of $N + 1$, followed by $n - 2^N$ as N -bits

Requires $2\lfloor \log_2 (\lfloor \log_2 n \rfloor + 1) \rfloor + \lfloor \log_2 n \rfloor + 1$ bits

Examples

$$9 = 1001_2 \mapsto 011001$$

$$15 = 1111_2 \mapsto 011111$$

Huffman encoding

All elements of U are weighted.

Algorithm 1: Huffman encoding

```
1  $a, b \leftarrow$  the two lightest elements of  $U$ ;  
2  $c \leftarrow$  new element  $w(c) \leftarrow w(a) + w(b)$ ;  
3 if  $|U| > 1$  then  
4    $\phi_1 \leftarrow \text{Huffman}(U \setminus \{a, b\} \cup \{c\})$ ;  
5    $\phi \leftarrow \phi_1$ ;  
6    $\phi(a) \leftarrow \phi_1(c)0$ ;  $\phi(b) \leftarrow \phi_1(c)1$ ;  
7   Remove  $c, \phi(c)$ ;  
8 else  
9    $\phi(u) \leftarrow$  empty codeword;  
10 Return( $\phi$ );
```
