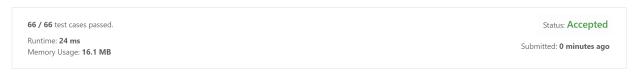
For simplicity assume m = n.

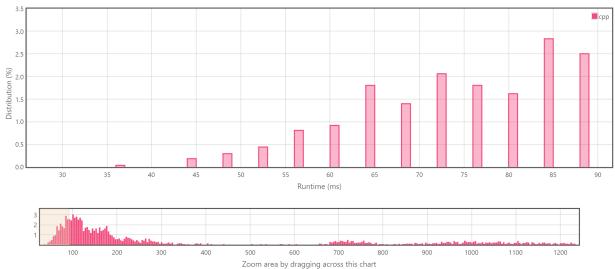
 $O^*(n^{2-\frac{2}{d+1}}) = O^*(n^{\frac{3}{2}})$  by reducing to halfspace range counting in 3D and combining partition and cutting [2], or  $O(n\sqrt{n}\log^2 n)$  [1] with near-linear space, which is optimal within polylog factors.

## **Queries on Number of Points Inside a Circle**

## **Submission Detail**



## **Accepted Solutions Runtime Distribution**



 $Runtime: 24 \ ms, faster than \ 100.00\% \ of \ C++ \ online \ submissions \ for \ Queries \ on \ Number \ of \ Points \ Inside \ a \ Circle.$ 

Memory Usage:  $16.1\,MB$ , less than 90.61% of C++ online submissions for Queries on Number of Points Inside a Circle

## References

- [1] Bernard Chazelle and Emo Welzl. Quasi-optimal range searching in spaces of finite vc-dimension. Discrete & Computational Geometry, 4(5):467-489, 1989.
- [2] Jiří Matoušek. Efficient partition trees. Discrete & Computational Geometry, 8(3):315–334, 1992.