

Let $k = 3$ denote the number of booking limit.

1. Balanced search tree. $O(n \log n)$.

2. Maintain the depth of each interval, using predecessor search in $O(\text{pred}(n) + \Delta)$ time per operation, where we increase the depth of Δ intervals. Also maintain a predecessor search data structure for all intervals with depth $k - 1$. The total depth of the intervals is $O(n \cdot k)$, so by amortized analysis, the total running time is $O(nk + n \cdot \text{pred}(n))$.

My Calendar II

Submission Detail

98 / 98 test cases passed.

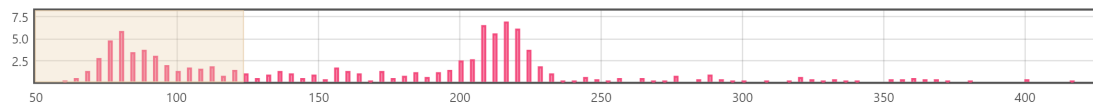
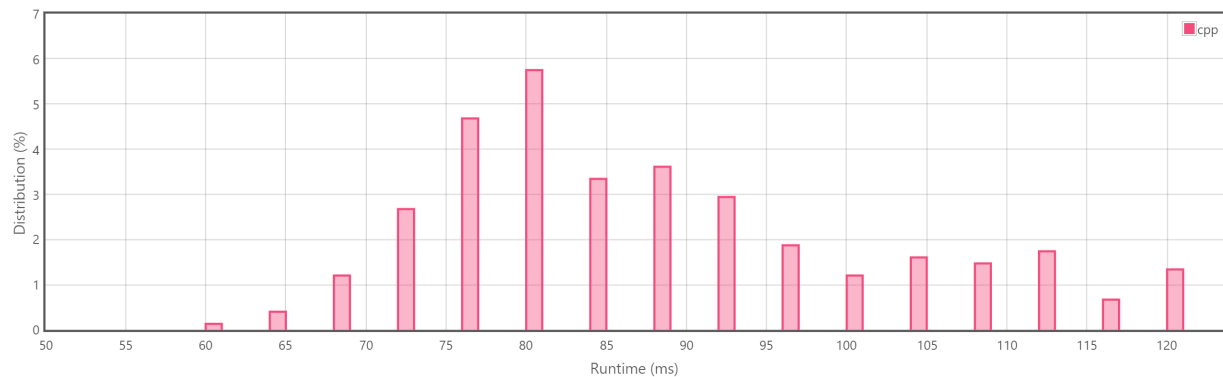
Runtime: 44 ms

Memory Usage: 33.4 MB

Status: Accepted

Submitted: 0 minutes ago

Accepted Solutions Runtime Distribution



Zoom area by dragging across this chart

Runtime: 44 ms, faster than 100.00% of C++ online submissions for My Calendar II.

Memory Usage: 33.4 MB, less than 100.00% of C++ online submissions for My Calendar II.

References