It suffices to count the number of pairs (i, j) that satisfies $a[i]$ xor $a[j] \leq x$.
1. Use trie. $O(nW)$.
2. First sort the integers, then build the compressed trie in $O(n)$ time, which has $O(n)$ nodes and edges. Perform a dfs on the trie to enumerate all possible $a[i]$'s, and the value $a[i]$ xor x will also traverse the trie once (we can simulate the traversal using bit operations), in the meantime we maintain the number of possible j's, so the running time is $O(n)$. $O(\text{sort}(n))$.

References