

1. process for each digit.  $O(\log W)$ .
2. in  $O(\log^\epsilon W)$  time we can preprocess the reverse of  $O(\log \log W)$  bits and query in  $O(1)$ .  $O(\frac{\log W}{\log \log W})$ .

note. for binary integer, reverse integer can be solved in  $O(\log \log W)$  time by divide and conquer. otherwise we may need number base conversion.

after preprocessing in  $O(W^\epsilon)$ , we can reverse integer in  $O(1)$ .

in general, for binary integer we can shuffle the bits in  $O(\log \log W)$  time [1, Sec. 7.1.3]. we also have a matching  $\Omega(\log \log W)$  lower bound if multiplication is not allowed.

## References

- [1] Donald E Knuth. *The art of computer programming, volume 4A: combinatorial algorithms, part 1*. Pearson Education India, 2011.