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# 1. algorithm

#include < algorithm > #include < numeric >

#include <algorithm> #include <numeric></numeric></algorithm>						
Algo	Params	Funcion				
sort, stable_sort	f, 1	ordena el intervalo				
$nth\_element$	f, nth, l	void ordena el n-esimo, y				
		particiona el resto				
fill, fill_n	f, l / n, elem	void llena [f, l) o [f,				
		f+n) con elem				
lower_bound, upper_bound	f, l, elem	it al primer / ultimo donde se				
		puede insertar elem para que				
		quede ordenada				
binary_search	f, l, elem	bool esta elem en [f, l)				
copy	f, l, resul	hace resul+ $i$ =f+ $i$ $\forall i$				
find, find_if, find_first_of	f, l, elem	$it$ encuentra i $\in$ [f,l) tq. i $=$ elem,				
	/ pred / f2, l2	$\operatorname{pred}(i), i \in [f2, l2)$				
count, count_if	f, l, elem/pred	cuenta elem, pred(i)				
search	f, l, f2, l2	busca $[f2,l2) \in [f,l)$				
replace, replace_if	f, l, old	cambia old / pred(i) por new				
	/ pred, new					
reverse	f, 1	da vuelta				
partition, stable_partition	f, l, pred	pred(i) ad, !pred(i) atras				
min_element, max_element	f, l, [comp]	it min, max de [f,l]				
lexicographical_compare	f1,l1,f2,l2	bool con [f1,l1];[f2,l2]				
next/prev_permutation	f,l	deja en [f,l) la perm sig, ant				
set_intersection,	f1, l1, f2, l2, res	[res,) la op. de conj				
set_difference, set_union,						
set_symmetric_difference,						
push_heap, pop_heap,	f, l, e / e /	mete/saca e en heap [f,l),				
make_heap		hace un heap de [f,l)				
is_heap	f,l	bool es [f,l) un heap				
accumulate	f,l,i,[op]	$T = \sum \text{oper de [f,l)}$				
inner_product	f1, l1, f2, i	$T = i + [f1, l1) \cdot [f2,)$				
partial_sum	f, l, r, [op]	$r+i = \sum /oper de [f,f+i] \forall i \in [f,l)$				
builtin_ffs	unsigned int	Pos. del primer 1 desde la derecha				
_builtin_clz	unsigned int	Cant. de ceros desde la izquierda.				
_builtin_ctz	unsigned int	Cant. de ceros desde la derecha.				
builtin_popcount	unsigned int	Cant. de 1's en x.				
_builtin_parity	unsigned int	1 si x es par, 0 si es impar.				
_builtin_XXXXXXII	unsigned ll	= pero para long long's.				

# 2. Estructuras

# 2.1. RMQ (static)

Dado un arreglo y una operacion asociativa *idempotente*, get(i, j) opera sobre el rango [i, j). Restriccion: LVL ≥ ceil(logn); Usar [] para llenar arreglo y luego build().

```
1 | struct RMQ{
     #define LVL 10
2
     tipo vec[LVL] [1<<(LVL+1)];
     tipo &operator[](int p){return vec[0][p];}
     tipo get(int i, int j) {//intervalo [i,j)
5
       int p = 31-_builtin_clz(j-i);
6
       return min(vec[p][i],vec[p][j-(1<<p)]);</pre>
7
     }
8
     void build(int n) {//O(nlogn)
9
       int mp = 31-__builtin_clz(n);
10
       forn(p, mp) forn(x, n-(1<<p))
11
         vec[p+1][x] = min(vec[p][x], vec[p][x+(1<<p)]);
12
     }};
13
```

# 2.2. RMQ (dynamic)

```
1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera
        sobre el rango [i, j).
   #define MAXN 100000
   #define operacion(x, y) max(x, y)
   const int neutro=0;
   struct RMQ{
     int sz;
6
     tipo t[4*MAXN];
7
     tipo &operator[](int p){return t[sz+p];}
8
     void init(int n){//O(nlgn)
9
       sz = 1 \ll (32-\_builtin\_clz(n));
10
       forn(i, 2*sz) t[i]=neutro;
11
12
     void updall(){\frac{}{0}}
13
       dforn(i, sz) t[i]=operacion(t[2*i], t[2*i+1]);}
14
     tipo get(int i, int j){return get(i,j,1,0,sz);}
15
     tipo get(int i, int j, int n, int a, int b){\frac{1}{0}}
16
       if(j<=a || i>=b) return neutro;
17
       if(i<=a && b<=j) return t[n];
18
       int c=(a+b)/2;
19
```

```
return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
20
21
     void set(int p, tipo val){//O(lgn)
22
       for(p+=sz; p>0 && t[p]!=val;){
23
         t[p]=val;
24
         p/=2;
25
         val=operacion(t[p*2], t[p*2+1]);
26
27
     }
28
   }rma;
   //Usage:
31 | cin >> n; rmg.init(n); forn(i, n) cin >> rmg[i]; rmg.updall();
                            2.3. RMQ (lazy)
1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera
        sobre el rango [i, j).
   typedef int Elem; //Elem de los elementos del arreglo
   typedef int Alt;//Elem de la alteracion
   #define operacion(x,y) x+y
   const Elem neutro=0; const Alt neutro2=0;
   #define MAXN 100000
   struct RMQ{
     int sz:
     Elem t[4*MAXN]:
     Alt dirty[4*MAXN];//las alteraciones pueden ser de distinto Elem
10
     Elem &operator[](int p){return t[sz+p];}
11
     void init(int n){//O(nlgn)
12
       sz = 1 \ll (32-\_builtin\_clz(n));
13
       forn(i, 2*sz) t[i]=neutro;
14
       forn(i, 2*sz) dirty[i]=neutro2;
15
16
     void push(int n, int a, int b){//propaga el dirty a sus hijos
17
       if(dirty[n]!=0){
18
         t[n]+=dirty[n]*(b-a);//altera el nodo
19
         if(n<sz){
20
           dirty[2*n]+=dirty[n];
21
           dirty[2*n+1]+=dirty[n];
22
23
         dirty[n]=0;
24
25
26
     }
     Elem get(int i, int j, int n, int a, int b){\frac{1}{0}}
```

```
if(j<=a || i>=b) return neutro;
28
       push(n, a, b);//corrige el valor antes de usarlo
29
       if(i<=a && b<=j) return t[n];</pre>
30
       int c=(a+b)/2;
31
       return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
32
33
     Elem get(int i, int j){return get(i,j,1,0,sz);}
34
     //altera los valores en [i, j) con una alteración de val
35
     void alterar(Alt val, int i, int j, int n, int a, int b)\frac{1}{0}
36
       push(n, a, b);
37
       if(j<=a || i>=b) return;
38
       if(i<=a && b<=j){
39
         dirty[n]+=val;
40
         push(n, a, b);
         return;
42
       }
43
       int c=(a+b)/2:
44
       alterar(val, i, j, 2*n, a, c), alterar(val, i, j, 2*n+1, c, b);
45
       t[n]=operacion(t[2*n], t[2*n+1]);//por esto es el push de arriba
46
47
     void alterar(Alt val, int i, int j){alterar(val,i,j,1,0,sz);}
48
49 |}rmq;
```

# 2.4. RMQ (persistente)

```
typedef int tipo;
   tipo oper(const tipo &a, const tipo &b){
       return a+b;
3
4
  struct node{
5
     tipo v; node *1,*r;
6
     node(tipo v):v(v), 1(NULL), r(NULL) {}
7
       node(node *1, node *r) : 1(1), r(r){
8
           if(!1) v=r->v;
9
           else if(!r) v=l->v;
10
           else v=oper(1->v, r->v);
11
       }
12
   };
13
   node *build (tipo *a, int tl, int tr) {//modificar para que tome tipo a
     if (tl+1==tr) return new node(a[tl]);
15
     int tm=(tl + tr)>>1:
16
     return new node(build(a, tl, tm), build(a, tm, tr));
17
18 | }
```

```
node *update(int pos, int new_val, node *t, int tl, int tr){
     if (tl+1==tr) return new node(new_val);
20
     int tm=(tl+tr)>>1;
21
     if(pos < tm) return new node(update(pos, new_val, t->1, tl, tm), t->r)
22
     else return new node(t->1, update(pos, new_val, t->r, tm, tr));
23
24
   tipo get(int 1, int r, node *t, int tl, int tr){
25
       if(l==tl && tr==r) return t->v;
26
     int tm=(tl + tr)>>1;
27
       if(r<=tm) return get(1, r, t->1, t1, tm);
28
       else if(1>=tm) return get(1, r, t->r, tm, tr);
29
    return oper(get(1, tm, t->1, t1, tm), get(tm, r, t->r, tm, tr));
31 }
```

#### 2.5. Fenwick Tree

```
1 //For 2D threat each column as a Fenwick tree, by adding a nested for in
        each operation
2 struct Fenwick{
     static const int sz=1000001;
     tipo t[sz];
     void adjust(int p, tipo v){//valid with p in [1, sz), O(lgn)
       for(; p<sz; p+=(p&-p)) t[p]+=v; }
6
     tipo sum(int p){//cumulative sum in [1, p], O(lgn)
       tipo s=0:
8
       for(; p; p-=(p&-p)) s+=t[p];
       return s;
10
11
     tipo sum(int a, int b){return sum(b)-sum(a-1);}
12
     //get largest value with cumulative sum less than or equal to x;
13
     //for smallest, pass x-1 and add 1 to result
14
     int getind(tipo x) {//O(lgn)
15
         int idx = 0, mask = N;
16
         while(mask && idx < N) {</pre>
17
           int t = idx + mask;
18
         if(x >= tree[t])
19
             idx = t, x -= tree[t];
20
           mask >>= 1:
21
22
         return idx:
23
    }};
24
```

#### 2.6. Union Find

```
struct UnionFind{
    vector<int> f;//the array contains the parent of each node
2
    void init(int n){f.clear(); f.insert(f.begin(), n, -1);}
3
    int comp(int x){return (f[x]=-1?x:f[x]=comp(f[x]));}//0(1)
4
    bool join(int i, int j) {
5
      bool con=comp(i)==comp(j);
6
      if(!con) f[comp(i)] = comp(j);
      return con;
8
    }};
9
```

# 2.7. Disjoint Intervals

```
|bool operator< (const ii &a, const ii &b) {return a.fst<b.fst;}
   //Stores intervals as [first, second]
   //in case of a collision it joins them in a single interval
   struct disjoint_intervals {
4
     set<ii>> segs;
5
     void insert(ii v) {//O(lgn)
6
       if(v.snd-v.fst==0.) return;//0J0
7
       set<ii>>::iterator it,at;
8
       at = it = segs.lower_bound(v);
9
       if (at!=segs.begin() && (--at)->snd >= v.fst)
10
         v.fst = at->fst, --it;
11
       for(; it!=segs.end() && it->fst <= v.snd; segs.erase(it++))</pre>
12
         v.snd=max(v.snd, it->snd);
13
       segs.insert(v);
14
15
<sub>16</sub> | };
```

# 2.8. RMQ (2D)

```
struct RMQ2D{
     static const int sz=1024;
2
     RMQ t[sz];
3
     RMQ &operator[](int p){return t[sz/2+p];}
4
     void build(int n, int m)\{//0(nm)\}
5
       forr(y, sz/2, sz/2+m)
6
         t[y].build(m);
7
       forr(y, sz/2+m, sz)
8
         forn(x, sz)
9
           t[y].t[x]=0;
10
       dforn(y, sz/2)
11
```

```
forn(x, sz)
12
           t[y].t[x]=max(t[y*2].t[x], t[y*2+1].t[x]);
13
     }
14
     void set(int x, int y, tipo v){//O(lgm.lgn)
15
       v + = sz/2;
16
       t[y].set(x, v);
17
       while (y/=2)
18
         t[y].set(x, max(t[y*2][x], t[y*2+1][x]));
19
     }
20
     //O(lgm.lgn)
21
     int get(int x1, int y1, int x2, int y2, int n=1, int a=0, int b=sz/2){
22
       if(y2<=a || y1>=b) return 0;
23
       if(y1<=a && b<=y2) return t[n].get(x1, x2);
24
       int c=(a+b)/2;
25
       return max(get(x1, y1, x2, y2, 2*n, a, c),
26
            get(x1, y1, x2, y2, 2*n+1, c, b));
27
     }
28
29
   //Example to initialize a grid of M rows and N columns:
   RMQ2D rmq;
   forn(i, M)
32
     forn(j, N)
       cin >> rmq[i][j];
34
35 rmq.build(N, M);
                               2.9. Big Int
```

```
1 #define BASEXP 6
   #define BASE 1000000
   #define LMAX 1000
   struct bint{
        int 1;
5
       11 n[LMAX];
 6
        bint(11 x=0){
7
            1=1;
8
            forn(i, LMAX){
9
                if (x) l=i+1;
10
                n[i] = x \text{BASE};
11
                x/=BASE;
12
13
            }
14
        }
15
        bint(string x){
16
```

```
l=(x.size()-1)/BASEXP+1;
17
           fill(n, n+LMAX, 0);
18
           ll r=1;
19
           forn(i, sz(x)){
20
               n[i / BASEXP] += r * (x[x.size()-1-i]-'0');
21
               r*=10; if(r==BASE)r=1;
22
           }
23
       }
24
       void out(){
25
       cout << n[1-1];
26
       dforn(i, l-1) printf("%6.61lu", n[i]);//6=BASEXP!
27
28
     void invar(){
29
       fill(n+1, n+LMAX, 0);
30
       while(1>1 && !n[1-1]) 1--;
31
     }
32
33
   bint operator+(const bint&a, const bint&b){
     bint c:
35
       c.1 = max(a.1, b.1);
36
       11 q = 0;
37
       forn(i, c.l) q += a.n[i]+b.n[i], c.n[i]=q %BASE, q/=BASE;
38
       if(q) c.n[c.l++] = q;
39
       c.invar();
40
       return c;
41
42
   pair<bint, bool> lresta(const bint& a, const bint& b) // c = a - b
44
     bint c;
45
       c.1 = max(a.1, b.1);
46
       11 q = 0;
47
       forn(i, c.l) q += a.n[i]-b.n[i], c.n[i]=(q+BASE) %BASE, q=(q+BASE)/
48
           BASE-1:
       c.invar():
49
       return make_pair(c, !q);
50
51
   bint& operator-= (bint& a, const bint& b){return a=lresta(a, b).first;}
   bint operator- (const bint&a, const bint&b) {return lresta(a, b).first;}
   bool operator< (const bint&a, const bint&b){return !lresta(a, b).second
  |bool operator<= (const bint&a, const bint&b){return lresta(b, a).second
| bool operator==(const bint&a, const bint&b){return a <= b && b <= a;}
```

```
57 | bint operator*(const bint&a, ll b){
       bint c;
       11 q = 0;
       forn(i, a.1) q += a.n[i]*b, c.n[i] = q %BASE, q/=BASE;
       while(q) c.n[c.l++] = q %BASE, q/=BASE;
       c.invar();
       return c;
   }
65
   bint operator*(const bint&a, const bint&b){
       bint c;
67
       c.1 = a.1+b.1;
       fill(c.n, c.n+b.1, 0);
       forn(i, a.1){
           11 a = 0:
           forn(j, b.l) q += a.n[i]*b.n[j]+c.n[i+j], c.n[i+j] = q \text{BASE}, q
                /=BASE:
           c.n[i+b.1] = q;
       }
74
       c.invar();
       return c;
76
77
   pair<br/><br/>bint, 11> ldiv(const bint& a, 11 b)\{// c = a / b : rm = a \% b \}
     bint c;
     11 \text{ rm} = 0;
80
     dforn(i, a.1){
               rm = rm * BASE + a.n[i];
82
               c.n[i] = rm / b;
               rm %= b;
84
       }
       c.1 = a.1;
       c.invar();
       return make_pair(c, rm);
88
89
   bint operator/(const bint&a, ll b){return ldiv(a, b).first;}
   11 operator %(const bint&a, 11 b) {return ldiv(a, b).second;}
   pair<bint, bint> ldiv(const bint& a, const bint& b){
     bint c:
93
       bint rm = 0;
94
       dforn(i, a.l){
           if (rm.l==1 && !rm.n[0])
96
                rm.n[0] = a.n[i];
97
98
           else{
```

```
dforn(j, rm.l) rm.n[j+1] = rm.n[j];
99
                rm.n[0] = a.n[i];
100
                rm.l++;
101
            }
102
            ll q = rm.n[b.1] * BASE + rm.n[b.1-1];
103
            ll u = q / (b.n[b.l-1] + 1);
104
            ll v = q / b.n[b.l-1] + 1;
105
            while (u < v-1){
106
                11 m = (u+v)/2;
107
                if (b*m <= rm) u = m;
108
                else v = m;
109
110
            c.n[i]=u;
111
            rm-=b*u:
112
       }
113
      c.l=a.l;
114
       c.invar();
115
       return make_pair(c, rm);
116
117
    bint operator/(const bint&a, const bint&b){return ldiv(a, b).first;}
   bint operator %(const bint&a, const bint&b) {return ldiv(a, b).second;}
                             2.10. HashTables
    //Compilar: g++ --std=c++11
   struct Hash{
 2
     size_t operator()(const ii &a)const{
 3
        size_t s=hash<int>()(a.fst);
 4
       return hash<int>()(a.snd)+0x9e3779b9+(s<<6)+(s>>2);
 5
 6
      size_t operator()(const vector<int> &v)const{
 7
       size_t s=0;
 8
       for(auto &e : v)
 9
          s = hash < int > ()(e) + 0x9e3779b9 + (s < 6) + (s > 2);
10
       return s;
11
     }
12
13
   unordered_set<ii, Hash> s;
   unordered_map<ii, int, Hash> m;//map<key, value, hasher>
                              2.11. Modnum
  struct mnum{
     static const tipo mod=12582917;
```

```
tipo v:
3
     mnum(tipo v=0): v(v mod) {}
4
     mnum operator+(mnum b){return v+b.v;}
     mnum operator-(mnum b){return v>=b.v? v-b.v : mod-b.v+v;}
     mnum operator*(mnum b){return v*b.v;}
     mnum operator^(int n){
       if(!n) return 1;
       return n %2? (*this)^(n/2)*(this) : (*this)^(n/2);}
10
11 };
                           2.12. Treap para set
typedef int Key;
  typedef struct node *pnode;
   struct node{
       Kev kev;
       int prior, size;
5
       pnode l,r;
6
       node(Key key=0): key(key), prior(rand()), size(1), 1(0), r(0) {}
7
   }:
8
   static int size(pnode p) { return p ? p->size : 0; }
   void push(pnode p) {
     // modificar y propagar el dirty a los hijos aca(para lazy)
12
   // Update function and size from children's Value
   void pull(pnode p) {//recalcular valor del nodo aca (para rmg)
     p->size = 1 + size(p->1) + size(p->r);
15
16
   //junta dos arreglos
   pnode merge(pnode 1, pnode r) {
     if (!1 || !r) return 1 ? 1 : r;
     push(1), push(r);
20
     pnode t;
21
     if (1->prior < r->prior) 1->r=merge(1->r, r), t = 1;
22
     else r\rightarrow l=merge(1, r\rightarrow 1), t = r;
23
     pull(t);
24
     return t;
25
26
   //parte el arreglo en dos, l<key<=r
   void split(pnode t, Key key, pnode &1, pnode &r) {
       if (!t) return void(1 = r = 0);
29
       push(t);
30
       if (\text{key} \leftarrow \text{t->key}) split(\text{t->l}, \text{key}, l, \text{t->l}), r = t;
31
```

```
else split(t->r, key, t->r, r), l = t;
32
       pull(t);
33
   }
34
35
   void erase(pnode &t, Key key) {
       if (!t) return;
37
       push(t);
38
       if (key == t->key) t=merge(t->l, t->r);
39
       else if (key < t->key) erase(t->1, key);
40
       else erase(t->r, key);
41
       if(t) pull(t);
42
43
44
   ostream& operator<<(ostream &out, const pnode &t) {
     if(!t) return out;
46
       return out << t->l << t->key << ''' << t->r;
47
48
   pnode find(pnode t, Key key) {
49
       if (!t) return 0:
50
       if (key == t->key) return t;
51
       if (key < t->key) return find(t->1, key);
52
       return find(t->r, key);
53
54
   struct treap {
55
       pnode root;
56
       treap(pnode root=0): root(root) {}
57
       int size() { return ::size(root); }
58
       void insert(Key key) {
59
           pnode t1, t2; split(root, key, t1, t2);
60
           t1=::merge(t1,new node(key));
61
           root=::merge(t1,t2);
62
63
       void erase(Key key1, Key key2) {
64
           pnode t1,t2,t3;
65
           split(root,key1,t1,t2);
66
           split(t2,key2, t2, t3);
67
           root=merge(t1,t3);
68
       }
69
       void erase(Key key) {::erase(root, key);}
70
       pnode find(Key key) { return ::find(root, key); }
71
       Key &operator[](int pos){return find(pos)->key;}//ojito
72
73
  treap merge(treap a, treap b) {return treap(merge(a.root, b.root));}
```

# 2.13. Treap para arreglo

```
1 typedef int Value;
   typedef struct node *pnode;
   struct node{
       Value val:
       int prior, size;
       pnode 1,r;
       node(Value val=0): val(val), prior(rand()), size(1), 1(0), r(0) {}
   };
   static int size(pnode p) { return p ? p->size : 0; }
   void push(pnode p) {
     // modificar y propagar el dirty a los hijos aca(para lazy)
12
   // Update function and size from children's Value
   void pull(pnode p) {//recalcular valor del nodo aca (para rmg)
     p->size = 1 + size(p->1) + size(p->r);
16
   //junta dos arreglos
   pnode merge(pnode 1, pnode r) {
     if (!1 || !r) return 1 ? 1 : r;
     push(1), push(r);
     pnode t;
     if (1-\text{prior} < r-\text{prior}) 1-\text{r-merge}(1-\text{r}, r), t = 1;
     else r\rightarrow l=merge(1, r\rightarrow 1), t = r;
     pull(t);
24
     return t;
25
26
   //parte el arreglo en dos, sz(1)==tam
   void split(pnode t, int tam, pnode &1, pnode &r) {
     if (!t) return void(1 = r = 0);
     push(t);
     if (tam \le size(t->1)) split(t->1, tam, 1, t->1), r = t;
31
     else split(t->r, tam - 1 - size(t->l), t->r, r), l = t;
32
     pull(t);
33
34
   pnode at(pnode t, int pos) {
     if(!t) exit(1):
     if(pos == size(t->1)) return t;
37
     if(pos < size(t->1)) return at(t->1, pos);
     return at(t->r, pos - 1 - size(t->l));
39
   }
40
41
```

```
42 ostream& operator<<(ostream &out, const pnode &t) {
     if(!t) return out;
43
       return out << t->1 << t->val << ''' << t->r;
44
45
46
47
   struct arr {//para usar el treap como un arreglo
       pnode root;
49
       arr(Value val): root(new node(val)) {}
50
       arr(pnode root=0): root(root) {}
51
       void insert(int pos, Value val) {//agrega un elemento
52
           pnode t1,t2; ::split(root, pos, t1, t2);
53
           t1=merge(t1, new node(val));
54
           root=merge(t1,t2);
55
56
       void erase(int i, int j=-1) {//borra un rango
57
       if(j==-1) j=i+1;
58
           pnode t1,t2,t3;
59
           ::split(root,i,t1,t2);
60
           ::split(t2, j-i, t2, t3);
61
           root=merge(t1, t3);
62
       }
63
       void push_back(Value val) {root=merge(root, new node(val));}
64
       Value &operator[](int pos){return at(root, pos)->val;}//ojito
65
       //parte el arreglo en dos con sz(1)==tam
66
       void split(int tam, arr &1, arr &r){::split(root, tam, 1.root, r.
67
           root):}
       //dado un rango, parte el arreglo en tres. m = rango [i, j)
68
       void split(int i, int j, arr &l, arr &m, arr &r){
69
       ::split(root, i, l.root, m.root);
70
       ::split(m.root, j-i, m.root, r.root);
71
     }
72
73
   //concatena dos arreglos
  arr merge(arr a, arr b){return arr(merge(a.root, b.root));}
                       2.14. Convex Hull Trick
const ll is_query = -(1LL<<62);
  struct Line {
       ll m, b;
```

```
mutable multiset<Line>::iterator it:
4
       const Line *succ(multiset<Line>::iterator it) const;
5
       bool operator<(const Line& rhs) const {</pre>
6
            if (rhs.b != is_query) return m < rhs.m;</pre>
            const Line *s=succ(it);
8
           if(!s) return 0;
           11 x = rhs.m;
            return b - s -> b < (s -> m - m) * x;
11
       }
12
   };
13
   struct HullDynamic : public multiset<Line>{ // will maintain upper hull
14
       for maximum
       bool bad(iterator v) {
15
            iterator z = next(y);
16
           if (y == begin()) {
                if (z == end()) return 0;
                return y->m == z->m && y->b <= z->b;
19
            iterator x = prev(y);
21
            if (z == end()) return y \rightarrow m == x \rightarrow m && y \rightarrow b <= x \rightarrow b;
           return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - x-m)
23
                );
24
       iterator next(iterator y){return ++y;}
25
       iterator prev(iterator y){return --y;}
26
       void insert_line(ll m, ll b) {
27
           iterator y = insert((Line) { m, b });
28
29
            y->it=y;
            if (bad(y)) { erase(y); return; }
30
            while (next(y) != end() && bad(next(y))) erase(next(y));
31
           while (y != begin() && bad(prev(y))) erase(prev(y));
32
       }
33
       ll eval(ll x) {
34
            Line l = *lower_bound((Line) { x, is_query });
35
            return 1.m * x + 1.b;
36
       }
37
   }h:
38
   const Line *Line::succ(multiset<Line>::iterator it) const{
       return (++it==h.end()? NULL : &*it);}
```

#### 2.15. Gain-Cost Set

```
1 //esta estructura mantiene pairs(beneficio, costo)
  //de tal manera que en el set quedan ordenados
   //por beneficio Y COSTO creciente. (va borrando los que no son optimos)
   struct V{
     int gain, cost;
     bool operator<(const V &b)const{return gain<b.gain;}</pre>
7
   set<V> s;
8
   void add(V x){
     set<V>::iterator p=s.lower_bound(x);//primer elemento mayor o igual
10
     if(p!=s.end() && p->cost <= x.cost) return;//ya hay uno mejor
11
     p=s.upper_bound(x);//primer elemento mayor
12
     if(p!=s.begin()){//borro todos los peores (<=beneficio y >=costo)
13
       --p;//ahora es ultimo elemento menor o igual
14
       while(p->cost >= x.cost){
15
         if(p==s.begin()){s.erase(p); break;}
16
         s.erase(p--);
17
       }
18
     }
19
     s.insert(x);
20
21
   int get(int gain){//minimo costo de obtener tal ganancia
22
     set<V>::iterator p=s.lower_bound((V){gain, 0});
23
     return p==s.end()? INF : p->cost;}
                      2.16. Set con busq binaria
  #include <ext/pb_ds/assoc_container.hpp>
  #include <ext/pb_ds/tree_policy.hpp>
   using namespace __gnu_pbds;
```

# using namespace \_\_gnu\_pbds; typedef tree<int,null\_type,less<int>,//key,mapped type, comparator rb\_tree\_tag,tree\_order\_statistics\_node\_update> set\_t; //find\_by\_order(i) devuelve iterador al i-esimo elemento

# 3. Algos

# 3.1. Longest Increasing Subsecuence

```
1 //Para non-increasing, cambiar comparaciones y revisar busq binaria
```

```
2 //Given an array, paint it in the least number of colors so that each
       color turns to a non-increasing subsequence.
3 //Solution:Min number of colors=Length of the longest increasing
       subsequence
4 | int N, a[MAXN];//secuencia y su longitud
  ii d[MAXN+1];//d[i]=ultimo valor de la subsecuencia de tamanio i
   int p[MAXN];//padres
   vector<int> R;//respuesta
   void rec(int i){
     if(i==-1) return;
     R.push_back(a[i]);
10
     rec(p[i]);
11
12
   }
   int lis(){//O(nlogn)
     d[0] = ii(-INF, -1); forn(i, N) d[i+1]=ii(INF, -1);
14
     forn(i, N){
       int j = upper_bound(d, d+N+1, ii(a[i], INF))-d;
16
       if (d[j-1].first < a[i]&&a[i] < d[j].first){</pre>
         p[i]=d[j-1].second;
18
         d[j] = ii(a[i], i);
19
       }
20
     }
21
     R.clear();
22
     dforn(i, N+1) if(d[i].first!=INF){
23
       rec(d[i].second);//reconstruir
24
       reverse(R.begin(), R.end());
25
       return i;//longitud
26
     }
27
     return 0;
28
29 }
```

#### 3.2. Manacher

```
int d1[MAXN];//d1[i]=long del maximo palindromo impar con centro en i
int d2[MAXN];//d2[i]=analogo pero para longitud par

//0 1 2 3 4
//a a b c c <--d1[2]=3
//a a b b <--d2[2]=2 (estan uno antes)

void manacher(){
  int l=0, r=-1, n=sz(s);
  forn(i, n){
    int k=(i>r? 1 : min(d1[l+r-i], r-i));
    while(i+k<n && i-k>=0 && s[i+k]==s[i-k]) ++k;
```

```
d1[i] = k--;
11
       if(i+k > r) l=i-k, r=i+k;
12
     }
13
     1=0, r=-1;
14
     forn(i, n){
15
       int k=(i>r? 0 : min(d2[1+r-i+1], r-i+1))+1;
16
       while(i+k-1 \le k = 0 \ k \le [i+k-1] == s[i-k]) k++;
17
       d2[i] = --k;
18
       if(i+k-1 > r) l=i-k, r=i+k-1;
19
20
```

# 3.3. Alpha-Beta prunning

```
| 11 alphabeta(State &s, bool player = true, int depth = 1e9, 11 alpha = -
       INF, 11 beta = INF) { //player = true -> Maximiza
       if(s.isFinal()) return s.score;
2
     //~ if (!depth) return s.heuristic();
3
       vector<State> children;
       s.expand(player, children);
5
       int n = children.size();
6
       forn(i, n) {
           ll v = alphabeta(children[i], !player, depth-1, alpha, beta);
8
           if(!player) alpha = max(alpha, v);
9
           else beta = min(beta, v);
10
           if(beta <= alpha) break;</pre>
11
       }
12
       return !player ? alpha : beta;}
13
```

# 4. Strings

#### 4.1. KMP

```
string T;//cadena donde buscar(where)
   string P;//cadena a buscar(what)
   int b[MAXLEN];//back table
   void kmppre(){//by gabina with love
4
       int i =0, j=-1; b[0]=-1;
5
       while(i<sz(P)){</pre>
6
           while(j>=0 && P[i] != P[j]) j=b[j];
7
           i++, j++;
           b[i] = j;
9
       }
10
11 }
```

```
12
   void kmp(){
       int i=0, j=0;
14
       while(i<sz(T)){</pre>
15
            while(j>=0 && T[i]!=P[j]) j=b[j];
16
            i++, j++;
            if(j==sz(P)){
18
                printf("P_is_found_at_index_\%d_in_T\n", i-j);
19
                j=b[j];
20
            }
21
       }
22
23 }
```

#### 4.2. Trie

```
1 | struct trie{
     map<char, trie> m;
     void add(const string &s, int p=0){
       if(s[p]) m[s[p]].add(s, p+1);
4
    }
5
     void dfs(){
6
       //Do stuff
       forall(it, m)
8
         it->second.dfs();
9
    }
10
11 };
```

#### 4.3. Suffix Array (largo, nlogn)

```
1 #define MAX_N 1000
   #define rBOUND(x) (x<n? r[x] : 0)
   //sa will hold the suffixes in order.
   int sa[MAX_N], r[MAX_N], n;
   string s; //input string, n=sz(s)
   int f[MAX_N], tmpsa[MAX_N];
   void countingSort(int k){
     zero(f);
     forn(i, n) f[rBOUND(i+k)]++;
     int sum=0:
11
     forn(i, max(255, n)){
       int t=f[i]; f[i]=sum; sum+=t;}
13
     forn(i, n)
14
       tmpsa[f[rBOUND(sa[i]+k)]++]=sa[i];
15
```

```
memcpy(sa, tmpsa, sizeof(sa));
16
   }
17
   void constructsa(){\frac{}{0} n log n)
18
     n=sz(s);
19
     forn(i, n) sa[i]=i, r[i]=s[i];
20
     for(int k=1; k<n; k<<=1){
21
       countingSort(k), countingSort(0);
22
       int rank, tmpr[MAX_N];
23
       tmpr[sa[0]]=rank=0;
24
       forr(i, 1, n)
25
         tmpr[sa[i]] = r[sa[i-1]] & r[sa[i]+k] = r[sa[i-1]+k])?
26
             rank: ++rank:
       memcpy(r, tmpr, sizeof(r));
27
       if(r[sa[n-1]]==n-1) break:
28
     }
29
30
   void print(){//for debug
31
     forn(i, n)
32
       cout << i << ''' <<
33
       s.substr(sa[i], s.find( '$', sa[i])-sa[i]) << endl;}
34
```

#### 4.4. String Matching With Suffix Array

```
//returns (lowerbound, upperbound) of the search
   ii stringMatching(string P){ //O(sz(P)lgn)
     int lo=0, hi=n-1, mid=lo;
3
     while(lo<hi){
4
       mid=(lo+hi)/2;
5
       int res=s.compare(sa[mid], sz(P), P);
6
       if(res>=0) hi=mid;
7
       else lo=mid+1;
8
9
     if(s.compare(sa[lo], sz(P), P)!=0) return ii(-1, -1);
10
     ii ans; ans.fst=lo;
11
     lo=0, hi=n-1, mid;
12
     while(lo<hi){</pre>
13
       mid=(lo+hi)/2;
14
       int res=s.compare(sa[mid], sz(P), P);
15
       if(res>0) hi=mid:
16
       else lo=mid+1:
17
18
     if(s.compare(sa[hi], sz(P), P)!=0) hi--;
19
     ans.snd=hi;
20
```

```
return ans:
22 }
                4.5. LCP (Longest Common Prefix)
1 //Calculates the LCP between consecutives suffixes in the Suffix Array.
   //LCP[i] is the length of the LCP between sa[i] and sa[i-1]
   int LCP[MAX_N], phi[MAX_N], PLCP[MAX_N];
   void computeLCP(){//0(n)}
     phi[sa[0]]=-1;
     forr(i, 1, n) phi[sa[i]]=sa[i-1];
     int L=0:
     forn(i, n){
8
       if(phi[i]==-1) {PLCP[i]=0; continue;}
9
       while(s[i+L]==s[phi[i]+L]) L++;
10
       PLCP[i]=L:
11
       L=\max(L-1, 0):
12
13
     forn(i, n) LCP[i]=PLCP[sa[i]];
14
15 | }
                              4.6. Corasick
1
   struct trie{
     map<char, trie> next;
     trie* tran[256];//transiciones del automata
     int idhoja, szhoja;//id de la hoja o 0 si no lo es
     //link lleva al sufijo mas largo, nxthoja lleva al mas largo pero que
         es hoja
     trie *padre, *link, *nxthoja;
7
     char pch://caracter que conecta con padre
8
     trie(): tran(), idhoja(), padre(), link() {}
9
     void insert(const string &s, int id=1, int p=0){//id>0!!!
10
       if(p<sz(s)){</pre>
11
         trie &ch=next[s[p]];
12
         tran[(int)s[p]]=&ch;
13
         ch.padre=this, ch.pch=s[p];
14
         ch.insert(s, id, p+1);
15
16
       else idhoja=id, szhoja=sz(s);
17
18
     trie* get_link() {
19
```

if(!link){

20

10

11 }

```
if(!padre) link=this;//es la raiz
21
         else if(!padre->padre) link=padre;//hijo de la raiz
^{22}
         else link=padre->get_link()->get_tran(pch);
23
24
       return link;
25
26
     trie* get_tran(int c) {
27
       if(!tran[c])
28
         tran[c] = !padre? this : this->get_link()->get_tran(c);
29
       return tran[c];
30
31
     trie *get_nxthoja(){
32
       if(!nxthoja) nxthoja = get_link()->idhoja? link : link->nxthoja;
33
       return nxthoja;
34
35
     void print(int p){
36
       if(idhoja)
37
         cout << "found," << idhoja << ", at position," << p-szhoja << endl
38
       if(get_nxthoja()) get_nxthoja()->print(p);
39
40
     void matching(const string &s, int p=0){
41
       print(p);
42
       if(p<sz(s)) get_tran(s[p])->matching(s, p+1);
43
```

# Geometria

#### 5.1. Punto

```
1 | struct pto{
```

#### 5.2. Line

```
int sgn(ll x){return x<0? -1 : !!x;}
  struct line{
2
    line() {}
3
    double a,b,c;//Ax+By=C
4
  //pto MUST store float coordinates!
    line(double a, double b, double c):a(a),b(b),c(c){}
    line(pto p, pto q): a(q.y-p.y), b(p.x-q.x), c(a*p.x+b*p.y) {}
    int side(pto p){return sgn(ll(a) * p.x + ll(b) * p.y - c);}
8
9
  |bool parallels(line 11, line 12){return abs(11.a*12.b-12.a*11.b)<EPS;}
```

```
pto inter(line 11, line 12){//intersection
     double det=11.a*12.b-12.a*11.b;
12
     if(abs(det) < EPS) return pto(INF, INF); //parallels</pre>
13
     return pto(12.b*11.c-11.b*12.c, 11.a*12.c-12.a*11.c)/det;
14
15 }
                              5.3. Segment
1 struct segm{
     pto s,f;
     segm(pto s, pto f):s(s), f(f) {}
     pto closest(pto p) {//use for dist to point
        double 12 = dist_sq(s, f);
5
        if(12==0.) return s;
6
        double t = ((p-s)*(f-s))/12/12;
7
        if (t<0.) return s;//not write if is a line
8
        else if(t>1.)return f;//not write if is a line
9
        return s+((f-s)*t);
10
    }
11
       bool inside(pto p){return abs(dist(s, p)+dist(p, f)-dist(s, f))<EPS
12
           ;}
13
   };
14
   pto inter(segm s1, segm s2){
     pto r=inter(line(s1.s, s1.f), line(s2.s, s2.f));
       if(s1.inside(r) && s2.inside(r))
17
           return r;
18
     return pto(INF, INF);
19
20 }
                             5.4. Rectangle
1 | struct rect{
     //lower-left and upper-right corners
    pto lw, up;
3
   };
4
   //returns if there's an intersection and stores it in r
   bool inter(rect a, rect b, rect &r){
    r.lw=pto(max(a.lw.x, b.lw.x), max(a.lw.y, b.lw.y));
    r.up=pto(min(a.up.x, b.up.x), min(a.up.y, b.up.y));
   //check case when only a edge is common
9
```

return r.lw.x<r.up.x && r.lw.y<r.up.y;

# 5.5. Polygon Area

```
|double area(vector<pto> &p){//0(sz(p))
     double area=0;
2
    forn(i, sz(p)) area+=p[i]^p[(i+1) %z(p)];
    //if points are in clockwise order then area is negative
     return abs(area)/2;
5
6
   //Area ellipse = M_PI*a*b where a and b are the semi axis lengths
  //Area triangle = sqrt(s*(s-a)(s-b)(s-c)) where s=(a+b+c)/2
                                5.6. Circle
   vec perp(vec v){return vec(-v.y, v.x);}
  line bisector(pto x, pto y){
     line l=line(x, y); pto m=(x+y)/2;
     return line(-1.b, 1.a, -1.b*m.x+1.a*m.y);
4
5
   struct Circle{
     pto o;
7
     double r:
8
     Circle(pto x, pto y, pto z){
9
       o=inter(bisector(x, y), bisector(y, z));
10
       r=dist(o, x);
11
12
     pair<pto, pto> ptosTang(pto p){
13
       pto m=(p+o)/2;
14
       tipo d=dist(o, m);
15
       tipo a=r*r/(2*d);
16
       tipo h=sqrt(r*r-a*a);
17
       pto m2=o+(m-o)*a/d;
18
       vec per=perp(m-o)/d;
19
       return make_pair(m2-per*h, m2+per*h);
20
21
^{22}
   //finds the center of the circle containing p1 and p2 with radius r
   //as there may be two solutions swap p1, p2 to get the other
   bool circle2PtsRad(pto p1, pto p2, double r, pto &c){
25
           double d2=(p1-p2).norm_sq(), det=r*r/d2-0.25;
26
           if(det<0) return false:
27
           c=(p1+p2)/2+perp(p2-p1)*sqrt(det);
28
           return true:
29
  |}
30
```

```
#define sqr(a) ((a)*(a))
   #define feq(a,b) (fabs((a)-(b))<EPS)</pre>
   pair<tipo, tipo > ecCuad(tipo a, tipo b, tipo c){//a*x*x+b*x+c=0
     tipo dx = sqrt(b*b-4.0*a*c);
34
     return make_pair((-b + dx)/(2.0*a), (-b - dx)/(2.0*a));
35
36
   pair<pto, pto> interCL(Circle c, line 1){
37
     bool sw=false;
     if((sw=feq(0,1.b))){}
39
     swap(l.a, l.b);
     swap(c.o.x, c.o.y);
41
42
     pair<tipo, tipo> rc = ecCuad(
43
     sqr(1.a)+sqr(1.b),
44
     2.0*1.a*1.b*c.o.y-2.0*(sqr(1.b)*c.o.x+1.c*1.a),
45
     sqr(1.b)*(sqr(c.o.x)+sqr(c.o.y)-sqr(c.r))+sqr(1.c)-2.0*1.c*1.b*c.o.y
     );
47
     pair<pto, pto> p( pto(rc.first, (l.c - l.a * rc.first) / l.b),
               pto(rc.second, (1.c - 1.a * rc.second) / 1.b) );
49
     if(sw){
     swap(p.first.x, p.first.y);
51
     swap(p.second.x, p.second.y);
52
53
     return p;
54
55
   pair<pto, pto> interCC(Circle c1, Circle c2){
    line 1:
57
     1.a = c1.o.x-c2.o.x;
     1.b = c1.o.y-c2.o.y;
     1.c = (sqr(c2.r) - sqr(c1.r) + sqr(c1.o.x) - sqr(c2.o.x) + sqr(c1.o.y)
     -sqr(c2.o.y))/2.0;
61
     return interCL(c1, 1);
62
63 }
                           5.7. Point in Poly
1 //checks if v is inside of P, using ray casting
2 //works with convex and concave.
   //excludes boundaries, handle it separately using segment.inside()
  |bool inPolygon(pto v, vector<pto>& P) {
     bool c = false:
     forn(i, sz(P)){
6
       int j=(i+1) \%z(P);
```

#### 5.8. Convex Check CHECK

```
bool isConvex(vector<int> &p){//O(N)
int N=sz(p);
if(N<3) return false;
bool isLeft=p[0].left(p[1], p[2]);
forr(i, 1, N)
if(p[i].left(p[(i+1) %], p[(i+2) %])!=isLeft)
return false;
return true; }</pre>
```

# 5.9. Convex Hull

```
//stores convex hull of P in S, CCW order
   void CH(vector<pto>& P, vector<pto> &S){
     S.clear();
     sort(P.begin(), P.end());
     forn(i, sz(P)){
5
       while(sz(S) \ge 2 \&\& S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back();
       S.pb(P[i]);
7
8
     S.pop_back();
     int k=sz(S):
     dforn(i, sz(P)){
       while(sz(S) \ge k+2 \&\& S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back
           ();
       S.pb(P[i]);
13
14
     S.pop_back();
15
16 }
```

# 5.10. Cut Polygon

```
//cuts polygon Q along the line ab
//stores the left side (swap a, b for the right one) in P
void cutPolygon(pto a, pto b, vector<pto> Q, vector<pto> &P){
P.clear();
```

```
forn(i, sz(Q)){
   double left1=(b-a)^(Q[i]-a), left2=(b-a)^(Q[(i+1)%z(Q)]-a);
   if(left1>=0) P.pb(Q[i]);
   if(left1*left2<0)
        P.pb(inter(line(Q[i], Q[(i+1)%z(Q)]), line(a, b)));
   }
}</pre>
```

#### 5.11. Bresenham

```
1 //plot a line approximation in a 2d map
   void bresenham(pto a, pto b){
    pto d=b-a; d.x=abs(d.x), d.y=abs(d.y);
    pto s(a.x<b.x? 1: -1, a.y<b.y? 1: -1);
    int err=d.x-d.y;
5
     while(1){
6
       m[a.x][a.y]=1;//plot
      if(a==b) break;
      int e2=2*err;
      if(e2 > -d.y){
         err-=d.y, a.x+=s.x;
       if(e2 < d.x)
         err+= d.x, a.y+= s.y;
13
    }
14
15 }
```

#### 5.12. Rotate Matrix

```
//rotates matrix t 90 degrees clockwise
//using auxiliary matrix t2(faster)
void rotate(){
forn(x, n) forn(y, n)
t2[n-y-1][x]=t[x][y];
memcpy(t, t2, sizeof(t));
}
```

# 5.13. Interseccion de Circulos en n3log(n)

```
struct event {
    double x; int t;
    event(double xx, int tt) : x(xx), t(tt) {}
    bool operator <(const event &o) const { return x < o.x; }
};
typedef vector<Circle> VC;
```

```
typedef vector<event> VE;
8
   int n;
   double cuenta(VE &v, double A,double B) {
9
       sort(v.begin(), v.end());
10
       double res = 0.0, lx = ((v.empty())?0.0:v[0].x);
11
       int contador = 0;
12
       forn(i,sz(v)) {
13
           // interseccion de todos (contador == n), union de todos (
14
               contador > 0),
           // conjunto de puntos cubierto por exacta k Circulos (contador
15
           if (contador == n) res += v[i].x - lx;
16
           contador += v[i].t:
17
           lx = v[i].x;
18
       }
19
       return res;
20
21
    // Primitiva de sqrt(r*r - x*x) como funcion double de una variable x.
   inline double primitiva(double x,double r) {
23
       if (x \ge r) return r*r*M_PI/4.0;
24
       if (x \le -r) return -r*r*M_PI/4.0;
25
       double raiz = sqrt(r*r-x*x);
26
       return 0.5 * (x * raiz + r*r*atan(x/raiz));
27
28
   double interCircle(VC &v) {
29
       vector<double> p; p.reserve(v.size() * (v.size() + 2));
30
       forn(i,sz(v)) {
31
           p.push_back(v[i].c.x + v[i].r);
32
           p.push_back(v[i].c.x - v[i].r);
33
       }
34
       forn(i,sz(v)) forn(j,i) {
35
           Circle &a = v[i], b = v[j];
36
           double d = (a.c - b.c).norm():
37
           if (fabs(a.r - b.r) < d \&\& d < a.r + b.r) {
38
               double alfa = acos((sqr(a.r) + sqr(d) - sqr(b.r)) / (2.0 * d)
39
                     * a.r)):
               pto vec = (b.c - a.c) * (a.r / d);
40
               p.pb((a.c + rotate(vec, alfa)).x);
41
               p.pb((a.c + rotate(vec, -alfa)).x);
42
           }
43
       }
44
       sort(p.begin(), p.end());
45
       double res = 0.0;
46
```

```
forn(i,sz(p)-1) {
47
           const double A = p[i], B = p[i+1];
48
           VE ve; ve.reserve(2 * v.size());
49
           forn(j,sz(v)) {
50
               const Circle &c = v[j];
               double arco = primitiva(B-c.c.x,c.r) - primitiva(A-c.c.x,c.r
52
                    );
               double base = c.c.y * (B-A);
53
               ve.push_back(event(base + arco,-1));
54
               ve.push_back(event(base - arco, 1));
55
56
           res += cuenta(ve,A,B);
57
       }
58
       return res;
59
60 }
```

# 6. Math

#### 6.1. Identidades

$$\sum_{i=0}^{n} \binom{n}{i} = 2^{n}$$

$$\sum_{i=0}^{n} i \binom{n}{i} = n * 2^{n-1}$$

$$\sum_{i=m}^{n} i = \frac{n(n+1)}{2} - \frac{m(m-1)}{2} = \frac{(n+1-m)(n+m)}{2}$$

$$\sum_{i=0}^{n} i = \sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i=0}^{n} i^{2} = \frac{n(n+1)(2n+1)}{6} = \frac{n^{3}}{3} + \frac{n^{2}}{2} + \frac{n}{6}$$

$$\sum_{i=0}^{n} i (i-1) = \frac{8}{6} (\frac{n}{2}) (\frac{n}{2} + 1)(n+1) \text{ (doubles)} \rightarrow \text{Sino ver caso impar y par}$$

$$\sum_{i=0}^{n} i^{3} = \left(\frac{n(n+1)}{2}\right)^{2} = \frac{n^{4}}{4} + \frac{n^{3}}{2} + \frac{n^{2}}{4} = \left[\sum_{i=1}^{n} i\right]^{2}$$

$$\sum_{i=0}^{n} i^{4} = \frac{n(n+1)(2n+1)(3n^{2}+3n-1)}{30} = \frac{n^{5}}{5} + \frac{n^{4}}{2} + \frac{n^{3}}{3} - \frac{n}{30}$$

$$\sum_{i=0}^{n} i^{p} = \frac{(n+1)^{p+1}}{p+1} + \sum_{k=1}^{p} \frac{B_{k}}{p-k+1} \binom{p}{k} (n+1)^{p-k+1}$$

$$r = e - v + k + 1$$

Teorema de Pick: (Area, puntos interiores y puntos en el borde)  $A = I + \frac{B}{2} - 1$ 

#### 6.2. Ec. Caracteristica

$$a_0T(n) + a_1T(n-1) + \ldots + a_kT(n-k) = 0$$

$$p(x) = a_0x^k + a_1x^{k-1} + \ldots + a_k$$
Sean  $r_1, r_2, \ldots, r_q$  las raíces distintas, de mult.  $m_1, m_2, \ldots, m_q$ 

$$T(n) = \sum_{i=1}^q \sum_{j=0}^{m_i-1} c_{ij} n^j r_i^n$$
Las constantes  $c_{ij}$  se determinan por los casos base.

#### 6.3. Combinatorio

```
forn(i, MAXN+1){//comb[i][k]=i tomados de a k
    comb[i][0]=comb[i][i]=1;
    forr(k, 1, i) comb[i][k]=(comb[i-1][k]+comb[i-1][k-1]) MOD;
}

ll lucas (ll n, ll k, int p){ //Calcula (n,k) %p teniendo comb[p][p]
    precalculado.

ll aux = 1;
    while (n + k){
    aux = (aux * comb[n %p][k %p]) %p;
    n/=p, k/=p;
}
return aux;
}
```

# 6.4. Exp. de Numeros Mod.

# 6.5. Exp. de Matrices y Fibonacci en log(n)

```
#define SIZE 350
   int NN;
   void mul(double a[SIZE][SIZE], double b[SIZE][SIZE])
4
       double res[SIZE] [SIZE] = {{0}};
5
      forn(i, NN) forn(j, NN) forn(k, NN) res[i][j]+=a[i][k]*b[k][j];
6
       forn(i, NN) forn(j, NN) a[i][j]=res[i][j];
7
8
   void powmat(double a[SIZE][SIZE], int n, double res[SIZE][SIZE])
10
      forn(i, NN) forn(j, NN) res[i][j]=(i==j);
11
       while(n){
12
           if(n&1) mul(res, a), n--;
13
           else mul(a, a), n/=2;
14
       }
15
  }
16
17
```

```
18 struct M22f
                    // la bl
     tipo a,b,c,d;// |c d|
     M22 operator*(const M22 &p) const {
       return (M22){a*p.a+b*p.c, a*p.b+b*p.d, c*p.a+d*p.c,c*p.b+d*p.d};}
21
   };
22
   M22 operator (const M22 &p, int n){
     if(!n) return (M22){1, 0, 0, 1};//identidad
24
     M22 q=p^(n/2); q=q*q;
25
     return n%2? p * q : q;}
26
27
   ll fibo(ll n){//calcula el fibonacci enesimo
28
     M22 \text{ mat}=(M22)\{0, 1, 1, 1\}^n;
     return mat.a*f0+mat.b*f1;//f0 y f1 son los valores iniciales
31 | }
                6.6. Matrices y determinante O(n^3)
   struct Mat {
       vector<vector<double> > vec:
       Mat(int n): vec(n, vector<double>(n) ) {}
       Mat(int n, int m): vec(n, vector<double>(m) ) {}
       vector<double> &operator[](int f){return vec[f];}
5
       const vector<double> &operator[](int f) const {return vec[f];}
6
       int size() const {return sz(vec);}
7
       Mat operator+(Mat &b) { ///this de n x m entonces b de n x m
8
           Mat m(sz(b), sz(b[0]));
           forn(i,sz(vec)) forn(j,sz(vec[0])) m[i][j] = vec[i][j] + b[i][j
10
               ];
           return m;
11
12
       Mat operator*(const Mat &b) { ///this de n x m entonces b de m x t
13
           int n = sz(vec), m = sz(vec[0]), t = sz(b[0]);
14
           Mat mat(n,t);
15
           forn(i,n) forn(j,t) {
16
               forn(k,m)
17
                   mat[i][i] += vec[i][k] * b[k][i];
18
           }
19
           return mat:
20
21
       double determinant(){//sacado de e maxx ru
22
           double det = 1:
23
           int n = sz(vec);
24
```

Mat m(\*this);

25

```
forn(i, n){//para cada columna
26
                int k = i;
27
               forr(j, i+1, n)//busco la fila con mayor val abs
28
                    if(abs(m[j][i])>abs(m[k][i]))
29
                        k = j;
30
                if(abs(m[k][i])<1e-9) return 0;
31
                m[i].swap(m[k]);//la swapeo
32
                if(i!=k) det = -det;
33
                det *= m[i][i];
34
                forr(j, i+1, n) m[i][j] /= m[i][i];
35
                //hago 0 todas las otras filas
36
                forn(j, n) if (j!= i && abs(m[j][i])>1e-9)
37
                    forr(k, i+1, n) m[j][k]-=m[i][k]*m[j][i];
38
           }
39
           return det;
40
       }
41
       Mat identidad(int n) {
42
           Mat m(n);
43
           forn(i,n) m[i][i] = 1;
44
           return m;
45
       }
46
       Mat transpuesta() {
47
           Mat m(sz(vec[0]),sz(vec));
48
           forn(i,sz(vec[0])) forn(j,sz(vec))
49
                m[i][j] = vec[j][i];
50
           return m;
51
       }
52
       void print() {
53
           int n = sz(vec), m = sz(vec[0]);
54
           cout << "********* << endl:
55
           forn(i,n){
56
                forn(j,m) cout << ""+!j << vec[i][j];
57
                cout << endl:</pre>
58
           }
59
       }
60
61 | };
```

#### 6.7. Teorema Chino del Resto

$$y = \sum_{j=1}^{n} (x_j * (\prod_{i=1, i \neq j}^{n} m_i)_{m_j}^{-1} * \prod_{i=1, i \neq j}^{n} m_i)$$

#### 6.8. Criba

```
1 #define MAXP 100000 //no necesariamente primo
   int criba[MAXP+1];
   void crearcriba(){
     int w[] = \{4,2,4,2,4,6,2,6\};
     for(int p=25;p<=MAXP;p+=10) criba[p]=5;</pre>
     for(int p=9;p<=MAXP;p+=6) criba[p]=3;</pre>
     for(int p=4;p<=MAXP;p+=2) criba[p]=2;</pre>
7
     for(int p=7,cur=0;p*p<=MAXP;p+=w[cur++&7]) if (!criba[p])
       for(int j=p*p;j<=MAXP;j+=(p<<1)) if(!criba[j]) criba[j]=p;</pre>
10
11
   vector<int> primos;
   void buscarprimos(){
     crearcriba():
     forr (i,2,MAXP+1) if (!criba[i]) primos.push_back(i);
15
16
17
   //~ Useful for bit trick:
   //~ #define SET(i) ( criba[(i)>>5] |=1<<((i)&31) )
  //~ #define INDEX(i) ( (criba[i>>5]>>((i)&31))&1 )
21 //~ unsigned int criba[MAXP/32+1];
                       6.9. Funciones de primos
       Sea n = \prod p_i^{k_i}, fact(n) genera un map donde a cada p_i le asocia su k_i
1 //factoriza bien numeros hasta MAXP^2
   map<11,11> fact(11 n){ //0 (cant primos)
     map<ll,ll> ret;
     forall(p, primos){
       while(!(n %*p)){
         ret[*p]++;//divisor found
6
         n/=*p;
8
9
     if(n>1) ret[n]++;
10
     return ret:
11
12
13
   //factoriza bien numeros hasta MAXP
   |map<11,11> fact2(11 n){ //0 (lg n)}
15
     map<ll,ll> ret;
```

```
while (criba[n]){
                                                                                      map<ll,ll> f=fact(n);
17
                                                                                     forall(it, f) rta *= (it->second + 1);
       ret[criba[n]]++;
                                                                                60
18
       n/=criba[n];
                                                                                     return rta;
                                                                                 61
19
                                                                                 62
20
     if(n>1) ret[n]++;
21
                                                                                 63
     return ret;
                                                                                    ll sumDiv (ll n){
22
                                                                                     ll rta = 1;
23
                                                                                     map<ll,11> f=fact(n);
24
   //Usar asi: divisores(fac, divs, fac.begin()); NO ESTA ORDENADO
                                                                                     forall(it, f) rta *= ((ll)pow((double)it->first, it->second + 1.0)-1)
   void divisores(const map<11,11> &f, vector<11> &divs, map<11,11>::
                                                                                          / (it->first-1);
       iterator it, ll n=1){
                                                                                     return rta;
                                                                                 68
       if(it==f.begin()) divs.clear();
                                                                                    }
27
                                                                                 69
       if(it==f.end()) {
28
                                                                                 70
           divs.pb(n);
                                                                                    ll eulerPhi (ll n){ // con criba: O(lg n)
29
                                                                                     11 \text{ rta} = n;
           return;
30
       }
                                                                                     map<11,11> f=fact(n);
31
       11 p=it->fst, k=it->snd; ++it;
                                                                                     forall(it, f) rta -= rta / it->first;
32
       forn(_, k+1) divisores(f, divs, it, n), n*=p;
                                                                                 75
                                                                                      return rta;
33
                                                                                    }
                                                                                 76
34
                                                                                 77
35
   11 numPrimeFactors (11 n){
                                                                                   36
     ll rta = 0;
                                                                                     11 r = n;
37
                                                                                79
     map<11,11> f=fact(n);
                                                                                     forr (i,2,n+1){
38
     forall(it, f) rta += it->second;
                                                                                     if ((l1)i*i > n) break;
39
     return rta;
                                                                                       if (n \% i == 0){
40
                                                                                          while (n\% == 0) n/=i;
41
                                                                                          r = r/i;
                                                                                84
42
   11 numDiffPrimeFactors (11 n){
                                                                                        }
                                                                                 85
43
     11 rta = 0:
                                                                                     }
                                                                                 86
44
     map<11,11> f=fact(n);
                                                                                     if (n != 1) r= r/n;
45
     forall(it, f) rta += 1;
                                                                                     return r;
                                                                                 89 }
     return rta;
47
48
49
                                                                                                    6.10. Phollard's Rho (rolando)
   11 sumPrimeFactors (11 n){
50
     ll rta = 0:
51
    map<ll,ll> f=fact(n);
                                                                                 1 | ll gcd(ll a, ll b){return a?gcd(b %a, a):b;}
52
    forall(it, f) rta += it->first;
53
     return rta;
                                                                                   | 11 mulmod (11 a, 11 b, 11 c) { //returns (a*b) %c, and minimize overfloor
54
                                                                                     11 x = 0, y = a\%;
55
                                                                                     while (b > 0){
56
   11 numDiv (ll n){
                                                                                        if (b \% 2 == 1) x = (x+y) \% c;
    ll rta = 1;
                                                                                        y = (y*2) \% c;
```

```
b /= 2:
                                                                                            y = (mulmod(y, y, n) + c) n;
                                                                                 51
                                                                                            y = (mulmod(y, y, n) + c) n;
                                                                                 52
9
                                                                                            if(x - y \ge 0) d = gcd(x - y, n);
     return x % c;
10
                                                                                            else d = gcd(y - x, n);
11
                                                                                        }
                                                                                 55
12
   ll expmod (ll b, ll e, ll m){\frac{1}{0}}
                                                                                        return d;
                                                                                 56
                                                                                    }
     if(!e) return 1;
                                                                                 57
     11 q = expmod(b, e/2, m); q = mulmod(q, q, m);
                                                                                 58
    return e %2? mulmod(b,q,m) : q;
                                                                                    map<ll,ll> prim;
16
                                                                                    void factRho (ll n){ //O (lg n)^3. un solo numero
17
                                                                                      if (n == 1) return;
18
   bool es_primo_prob (ll n, int a)
                                                                                     if (rabin(n)){
                                                                                       prim[n]++;
20
     if (n == a) return true:
                                                                                       return;
21
     11 s = 0, d = n-1;
     while (d \% 2 == 0) s++, d/=2;
                                                                                     11 factor = rho(n);
23
                                                                                     factRho(factor):
24
     11 x = expmod(a,d,n);
                                                                                     factRho(n/factor);
25
     if ((x == 1) \mid | (x+1 == n)) return true:
                                                                                 69 }
26
27
                                                                                                                6.11. GCD
     forn (i, s-1){
28
      x = mulmod(x, x, n);
29
                                                                                 tipo gcd(tipo a, tipo b){return a?gcd(b %a, a):b;}
       if (x == 1) return false;
30
       if (x+1 == n) return true;
31
                                                                                                         6.12. Extended Euclid
     }
32
     return false;
33
                                                                                 void extendedEuclid (ll a, ll b) \{ //a * x + b * y = d \}
34
                                                                                     if (!b) { x = 1; y = 0; d = a; return;}
35
                                                                                     extendedEuclid (b, a%);
   bool rabin (ll n){ //devuelve true si n es primo
36
                                                                                     11 x1 = y;
     if (n == 1) return false;
37
                                                                                     11 y1 = x - (a/b) * y;
     const int ar[] = \{2,3,5,7,11,13,17,19,23\};
38
                                                                                      x = x1; y = y1;
     forn (j,9)
39
                                                                                 7 | }
       if (!es_primo_prob(n,ar[j]))
40
         return false:
                                                                                                                6.13. LCM
41
     return true;
42
43
                                                                                 1 | tipo lcm(tipo a, tipo b){return a / gcd(a,b) * b;}
44
                                                                                                              6.14. Inversos
   ll rho(ll n){
       if( (n & 1) == 0 ) return 2;
46
      11 x = 2, y = 2, d = 1;
                                                                                 1 #define MAXMOD 15485867
47
      ll c = rand() % n + 1;
                                                                                 1 ll inv[MAXMOD];//inv[i]*i=1 mod MOD
48
       while(d == 1){
                                                                                 3 void calc(int p){//O(p)
49
           x = (mulmod(x, x, n) + c) n;
50
                                                                                     inv[1]=1;
```

```
forr(i, 2, p) inv[i] = p-((p/i)*inv[p%i])%p;

int inverso(int x){//O(log x)

return expmod(x, eulerphi(MOD)-2);//si mod no es primo(sacar a mano)
return expmod(x, MOD-2);//si mod es primo
}
```

# 6.15. Simpson

```
double integral(double a, double b, int n=10000) {//O(n), n=cantdiv
double area=0, h=(b-a)/n, fa=f(a), fb;
forn(i, n){
    fb=f(a+h*(i+1));
    area+=fa+ 4*f(a+h*(i+0.5)) +fb, fa=fb;
}
return area*h/6.:}
```

#### 6.16. Fraction

```
tipo mcd(tipo a, tipo b){return a?mcd(b%, a):b;}
   struct frac{
     tipo p,q;
3
     frac(tipo p=0, tipo q=1):p(p),q(q) {norm();}
     void norm(){
       tipo a = mcd(p,q):
       if(a) p/=a, q/=a;
       else q=1;
8
       if (q<0) q=-q, p=-p;}
9
     frac operator+(const frac& o){
10
       tipo a = mcd(q, o.q);
11
       return frac(p*(o.q/a)+o.p*(q/a), q*(o.q/a));}
12
     frac operator-(const frac& o){
13
       tipo a = mcd(q, o.q);
14
       return frac(p*(o.q/a)-o.p*(q/a), q*(o.q/a));}
15
     frac operator*(frac o){
16
       tipo a = mcd(q,o.p), b = mcd(o.q,p);
17
       return frac((p/b)*(o.p/a), (q/a)*(o.q/b));}
18
     frac operator/(frac o){
19
       tipo a = mcd(q,o.q), b = mcd(o.p,p);
20
       return frac((p/b)*(o.q/a),(q/a)*(o.p/b));}
21
     bool operator<(const frac &o) const{return p*o.q < o.p*q;}</pre>
22
     bool operator==(frac o){return p==o.p&kq==o.q;}
23
24 | };
```

#### 6.17. Polinomio

```
struct poly {
       vector<tipo> c;//guarda los coeficientes del polinomio
2
       poly(const vector<tipo> &c): c(c) {}
3
       poly() {}
4
     int gr(){//calculates grade of the polynomial
       return sz(c); }
6
     bool isnull() {return c.empty();}
7
       poly operator+(const poly &o) const {
           int m = sz(c), n = sz(o.c);
9
           vector<tipo> res(max(m,n));
10
           forn(i, m) res[i] += c[i];
11
           forn(i, n) res[i] += o.c[i];
12
           return poly(res);
13
14
       polv operator*(const polv &o) const {
15
           int m = sz(c), n = sz(o.c);
16
           vector<tipo> res(m+n-1);
17
           forn(i, m) forn(j, n) res[i+j]+=c[i]*o.c[j];
18
           return poly(res);
20
     tipo eval(tipo v) {
21
       tipo sum = 0:
       dforn(i, sz(c)) sum=sum*v + c[i];
23
       return sum:
24
    }
25
       //poly contains only a vector<int> c (the coeficients)
26
     //the following function generates the roots of the polynomial
27
   //it can be easily modified to return float roots
     set<tipo> roots(){
29
       set<tipo> roots;
30
       tipo a0 = abs(c[0]), an = abs(c[sz(c)-1]);
31
       vector<tipo> ps,qs;
32
       forr(p,1,sqrt(a0)+1) if (a0%p==0) ps.pb(p),ps.pb(a0/p);
33
       forr(q,1,sqrt(an)+1) if (an \% ==0) qs.pb(q),qs.pb(an/q);
34
       forall(pt,ps)
35
         forall(qt,qs) if ( (*pt) % (*qt)==0 ) {
36
           tipo root = abs((*pt) / (*qt));
37
           if (eval(root)==0) roots.insert(root);
38
39
       return roots;
40
41
```

```
42 };
                                                                                        int uc=i. uf=i:
                                                                                 5
  poly interpolate(const vector<tipo> &x, const vector<tipo> &y) {
                                                                                        forr(f, i, n) forr(c, i, m) if(fabs(a[f][c])>fabs(a[uf][uc])) {uf=f;
                                                                                 6
       int n = sz(x);
44
                                                                                        if (feq(a[uf][uc], 0)) { rw = i; break; }
       poly p;
45
       vector<tipo> aux(2);
                                                                                        forn(j, n) swap(a[j][i], a[j][uc]);
46
                                                                                        swap(a[i], a[uf]); swap(y[i], y[uf]); swap(p[i], p[uc]);
       forn(i, n) {
47
                                                                                        tipo inv = 1 / a[i][i]; //aca divide
          double a = y[i] - p.eval(x[i]);
48
          forn(j, i) a /= x[i] - x[j];
                                                                                        forr(j, i+1, n) {
49
          poly add(vector<tipo>(1, a));
                                                                                          tipo v = a[j][i] * inv;
                                                                                 12
50
          forn(j, i) aux[0]=-x[j], aux[1]=1, add = add*aux;
                                                                                          forr(k, i, m) a[i][k]-=v * a[i][k];
51
                                                                                          y[j] -= v*y[i];
          p = p + add;
52
                                                                                 14
      }
53
                                                                                 15
                                                                                      } // rw = rango(a), aca la matriz esta triangulada
       return p;
54
                                                                                      forr(i, rw, n) if (!feq(y[i],0)) return false; // checkeo de
                                                                                 17
55
   //the following functions allows parsing an expression like
                                                                                          compatibilidad
   //34+150+4*45
                                                                                      x = vector < tipo > (m, 0);
   //into a polynomial(el numero en funcion de la base)
                                                                                      dforn(i, rw){
   #define LAST(s) (sz(s)? s[sz(s)-1]:0)
                                                                                        tipo s = y[i];
   #define POP(s) s.erase(--s.end()):
                                                                                        forr(j, i+1, rw) s -= a[i][j]*x[p[j]];
   poly D(string &s) {
                                                                                        x[p[i]] = s / a[i][i]; //aca divide
    poly d;
                                                                                      }
                                                                                 23
62
     for(int i=0; isdigit(LAST(s)); i++) d.c.push_back(LAST(s)-'0'), POP(s)
                                                                                      ev = Mat(m-rw, Vec(m, 0)); // Esta parte va SOLO si se necesita el ev
63
                                                                                      forn(k, m-rw) {
                                                                                 25
                                                                                        ev[k][p[k+rw]] = 1;
     return d;}
                                                                                 26
64
                                                                                        dforn(i, rw){
                                                                                 27
65
                                                                                          tipo s = -a[i][k+rw];
   poly T(string &s) {
66
                                                                                          forr(j, i+1, rw) s -= a[i][j]*ev[k][p[j]];
    poly t=D(s);
67
     if (LAST(s)="'*"){POP(s); return T(s)*t;}
                                                                                          ev[k][p[i]] = s / a[i][i]; //aca divide
68
                                                                                        }
     return t;
                                                                                 31
69
                                                                                      }
70
   //main function, call this to parse
                                                                                      return true;
   poly E(string &s) {
    poly e=T(s);
73
                                                                                                                 6.19. FFT
    if (LAST(s)=='+')\{POP(s): return E(s)+e:\}
74
    return e:
75
76 }
                                                                                 1 //~ typedef complex<double> base; //menos codigo, pero mas lento
                                                                                 2 //elegir si usar complejos de c (lento) o estos
                           6.18. Ec. Lineales
                                                                                 3 struct base{
                                                                                        double r,i;
  | bool resolver_ev(Mat a, Vec y, Vec &x, Mat &ev){
                                                                                        base(double r=0, double i=0):r(r), i(i){}
     int n = a.size(), m = n?a[0].size():0, rw = min(n, m);
                                                                                        double real()const{return r;}
                                                                                 6
    vector<int> p; forn(i,m) p.push_back(i);
                                                                                        void operator/=(const int c){r/=c, i/=c;}
                                                                                 7
3
    forn(i, rw) {
                                                                                 8 };
```

```
base operator*(const base &a, const base &b){
       return base(a.r*b.r-a.i*b.i, a.r*b.i+a.i*b.r);}
10
   base operator+(const base &a, const base &b){
11
       return base(a.r+b.r, a.i+b.i);}
^{12}
   base operator-(const base &a, const base &b){
       return base(a.r-b.r, a.i-b.i);}
14
   vector<int> rev;
15
   vector<base> wlen_pw;
   inline static void fft(base a[], int n, bool invert) {
17
       forn(i, n) if(i<rev[i]) swap(a[i], a[rev[i]]);</pre>
18
     for (int len=2; len<=n; len<<=1) {
19
       double ang = 2*M_PI/len * (invert?-1:+1);
20
       int len2 = len >> 1:
21
       base wlen (cos(ang), sin(ang)):
22
       wlen_pw[0] = base(1, 0);
23
           forr(i, 1, len2) wlen_pw[i] = wlen_pw[i-1] * wlen;
24
       for (int i=0; i<n; i+=len) {</pre>
25
         base t,
26
           *pu = a+i.
27
           *pv = a+i+len2,
28
           *pu_end = a+i+len2,
29
           *pw = &wlen_pw[0];
30
         for (; pu!=pu_end; ++pu, ++pv, ++pw) {
31
           t = *pv * *pw;
32
           *pv = *pu - t;
33
            *pu = *pu + t;
34
         }
35
       }
36
     }
37
     if (invert) forn(i, n) a[i]/= n;
38
39
   inline static void calc_rev(int n){//precalculo: llamar antes de fft!!
40
       wlen_pw.resize(n);
41
       rev.resize(n):
42
       int lg=31-__builtin_clz(n);
43
       forn(i, n){
44
       rev[i] = 0:
45
           forn(k, lg) if(i&(1<<k))
46
                rev[i] |=1<<(lg-1-k);
47
       }
48
49
   inline static void multiply(const vector<int> &a, const vector<int> &b,
       vector<int> &res) {
```

```
vector<base> fa (a.begin(), a.end()), fb (b.begin(), b.end());
51
       int n=1:
52
     while(n < max(sz(a), sz(b))) n <<= 1;
53
     n <<= 1;
54
       calc_rev(n);
55
     fa.resize (n), fb.resize (n);
56
     fft (&fa[0], n, false), fft (&fb[0], n, false);
     forn(i, n) fa[i] = fa[i] * fb[i];
58
     fft (&fa[0], n, true);
59
     res.resize(n);
60
       forn(i, n) res[i] = int (fa[i].real() + 0.5);
61
62
   void toPoly(const string &s. vector<int> &P){//convierte un numero a
       polinomio
       P.clear():
64
       dform(i, sz(s)) P.pb(s[i]-'0');
65
66 }
```

# 6.20. Tablas y cotas (Primos, Divisores, Factoriales, etc)

```
Factoriales
0! = 1
                  11! = 39.916.800
1! = 1
                  12! = 479.001.600 \ (\in int)
2! = 2
                  13! = 6.227.020.800
3! = 6
                  14! = 87.178.291.200
4! = 24
                  15! = 1.307.674.368.000
5! = 120
                  16! = 20.922.789.888.000
6! = 720
                  17! = 355.687.428.096.000
7! = 5.040
                  18! = 6.402.373.705.728.000
8! = 40.320
                  19! = 121.645.100.408.832.000
9! = 362.880
                  20! = 2.432.902.008.176.640.000 (\in tint)
10! = 3.628.800 \mid 21! = 51.090.942.171.709.400.000
       max signed tint = 9.223.372.036.854.775.807
     max unsigned tint = 18.446.744.073.709.551.615
```

#### Primos

#### Primos cercanos a $10^n$

9941 9949 9967 9973 10007 10009 10037 10039 10061 10067 10069 10079 99961 99971 99989 99991 100003 100019 100043 100049 100057 100069 999959 999961 999979 999983 1000003 1000033 1000037 1000039 9999943 9999971 9999973 9999991 10000019 10000079 10000103 10000121 99999941 99999959 9999971 99999989 100000007 100000037 100000039 100000049 99999893 99999929 99999937 1000000007 1000000009 1000000021 1000000033

#### Cantidad de primos menores que $10^n$

```
\pi(10^1) = 4; \pi(10^2) = 25; \pi(10^3) = 168; \pi(10^4) = 1229; \pi(10^5) = 9592

\pi(10^6) = 78.498; \pi(10^7) = 664.579; \pi(10^8) = 5.761.455; \pi(10^9) = 50.847.534

\pi(10^{10}) = 455.052,511; \pi(10^{11}) = 4.118.054.813; \pi(10^{12}) = 37.607.912.018
```

#### Divisores

```
Cantidad de divisores (\sigma_0) para algunos n/\neg \exists n' < n, \sigma_0(n') \ge \sigma_0(n)
        \sigma_0(60) = 12; \sigma_0(120) = 16; \sigma_0(180) = 18; \sigma_0(240) = 20; \sigma_0(360) = 24
    \sigma_0(720) = 30; \sigma_0(840) = 32; \sigma_0(1260) = 36; \sigma_0(1680) = 40; \sigma_0(10080) = 72
        \sigma_0(15120) = 80; \sigma_0(50400) = 108; \sigma_0(83160) = 128; \sigma_0(110880) = 144
    \sigma_0(498960) = 200; \sigma_0(554400) = 216; \sigma_0(1081080) = 256; \sigma_0(1441440) = 288
                            \sigma_0(4324320) = 384 : \sigma_0(8648640) = 448
             Suma de divisores (\sigma_1) para algunos n/\neg \exists n' < n, \sigma_1(n') \ge \sigma_1(n)
    \sigma_1(96) = 252; \sigma_1(108) = 280; \sigma_1(120) = 360; \sigma_1(144) = 403; \sigma_1(168) = 480
        \sigma_1(960) = 3048; \sigma_1(1008) = 3224; \sigma_1(1080) = 3600; \sigma_1(1200) = 3844
     \sigma_1(4620) = 16128; \sigma_1(4680) = 16380; \sigma_1(5040) = 19344; \sigma_1(5760) = 19890
    \sigma_1(8820) = 31122; \sigma_1(9240) = 34560; \sigma_1(10080) = 39312; \sigma_1(10920) = 40320
\sigma_1(32760) = 131040; \sigma_1(35280) = 137826; \sigma_1(36960) = 145152; \sigma_1(37800) = 148800
\sigma_1(60480) = 243840; \sigma_1(64680) = 246240; \sigma_1(65520) = 270816; \sigma_1(70560) = 280098
            \sigma_1(95760) = 386880; \sigma_1(98280) = 403200; \sigma_1(100800) = 409448
        \sigma_1(491400) = 2083200; \sigma_1(498960) = 2160576; \sigma_1(514080) = 2177280
        \sigma_1(982800) = 4305280; \sigma_1(997920) = 4390848; \sigma_1(1048320) = 4464096
    \sigma_1(4979520) = 22189440; \sigma_1(4989600) = 22686048; \sigma_1(5045040) = 23154768
    \sigma_1(9896040) = 44323200; \sigma_1(9959040) = 44553600; \sigma_1(9979200) = 45732192
```

# 7. Grafos

# 7.1. Dijkstra

```
1 #define INF 1e9
2 int N;
   #define MAX_V 250001
   vector<ii> G[MAX V]:
   //To add an edge use
   #define add(a, b, w) G[a].pb(make_pair(w, b))
7
   ll dijkstra(int s, int t){\frac{}{|0(|E| \log |V|)}}
     priority_queue<ii, vector<ii>, greater<ii> > Q;
     vector<ll> dist(N, INF); vector<int> dad(N, -1);
10
     Q.push(make_pair(0, s)); dist[s] = 0;
11
     while(sz(Q)){
12
       ii p = Q.top(); Q.pop();
13
       if(p.snd == t) break;
14
       forall(it, G[p.snd])
15
         if(dist[p.snd]+it->first < dist[it->snd]){
16
           dist[it->snd] = dist[p.snd] + it->fst:
17
           dad[it->snd] = p.snd:
18
           Q.push(make_pair(dist[it->snd], it->snd));
19
         }
20
     }
21
22
     return dist[t]:
     if(dist[t]<INF)//path generator</pre>
23
       for(int i=t; i!=-1; i=dad[i])
24
         printf("%d%c", i, (i==s?'\n':'\_'));
25
26 }
```

#### 7.2. Bellman-Ford

```
vector<ii>G[MAX_N];//ady. list with pairs (weight, dst)
int dist[MAX_N];
void bford(int src){//O(VE)

dist[src]=0;
forn(i, N-1) forn(j, N) if(dist[j]!=INF) forall(it, G[j])

dist[it->snd]=min(dist[it->snd], dist[j]+it->fst);

bool hasNegCycle(){
forn(j, N) if(dist[j]!=INF) forall(it, G[j])
```

# 7.3. Floyd-Warshall

```
//G[i][j] contains weight of edge (i, j) or INF
   //G[i][i]=0
  int G[MAX_N] [MAX_N];
   void floyd(){//0(N^3)}
  |forn(k, N) forn(i, N) if(G[i][k]!=INF) forn(j, N) if(G[k][j]!=INF)
     G[i][j]=min(G[i][j], G[i][k]+G[k][j]);
6
7
   bool inNegCycle(int v){
8
     return G[v][v]<0;}
9
   //checks if there's a neg. cycle in path from a to b
   bool hasNegCycle(int a, int b){
11
    forn(i, N) if(G[a][i]!=INF && G[i][i]<0 && G[i][b]!=INF)
12
       return true;
13
     return false;
14
15 }
```

#### 7.4. Kruskal

```
struct Ar{int a,b,w;};
   bool operator (const Ar& a, const Ar &b) {return a.w <b.w;}
   vector<Ar> E;
   11 kruskal(){
       11 cost=0;
5
       sort(E.begin(), E.end());//ordenar aristas de menor a mayor
6
       uf.init(n);
       forall(it, E){
8
           if(uf.comp(it->a)!=uf.comp(it->b)){//si no estan conectados
9
               uf.unir(it->a, it->b);//conectar
10
               cost+=it->w:
11
           }
12
       }
13
       return cost;
14
15 }
```

# 7.5. Prim

```
bool taken[MAXN];
priority_queue<ii, vector<ii>, greater<ii>> pq;//min heap
   void process(int v){
       taken[v]=true;
       forall(e, G[v])
5
           if(!taken[e->second]) pq.push(*e);
6
  }
7
   11 prim(){
       zero(taken);
       process(0);
11
       11 cost=0;
12
       while(sz(pq)){
13
           ii e=pq.top(); pq.pop();
14
           if(!taken[e.second]) cost+=e.first, process(e.second);
15
       }
16
       return cost;
17
18 }
```

#### 7.6. 2-SAT + Tarjan SCC

```
1 //We have a vertex representing a var and other for his negation.
2 //Every edge stored in G represents an implication. To add an equation
       of the form allb, use addor(a, b)
   //MAX=max cant var, n=cant var
   #define addor(a, b) (G[neg(a)].pb(b), G[neg(b)].pb(a))
   vector<int> G[MAX*2];
   //idx[i]=index assigned in the dfs
  //lw[i]=lowest index(closer from the root) reachable from i
   int lw[MAX*2], idx[MAX*2], qidx;
   stack<int> q;
   int qcmp, cmp[MAX*2];
   //verdad[cmp[i]]=valor de la variable i
   bool verdad[MAX*2+1];
12
13
   int neg(int x) { return x>=n? x-n : x+n;}
14
   void tin(int v){
     lw[v]=idx[v]=++qidx;
16
    q.push(v), cmp[v]=-2;
17
    forall(it, G[v]){
18
       if(!idx[*it] || cmp[*it]==-2){
19
         if(!idx[*it]) tjn(*it);
20
         lw[v]=min(lw[v], lw[*it]);
21
```

```
}
22
     }
23
     if(lw[v]==idx[v]){
^{24}
       qcmp++;
25
        int x;
26
       do{x=q.top(); q.pop(); cmp[x]=qcmp;}while(x!=v);
27
       verdad[qcmp] = (cmp[neg(v)] < 0);</pre>
28
29
30
    //remember to CLEAR G!!!
   bool satisf(){\frac{}{0(n)}}
     memset(idx, 0, sizeof(idx)), qidx=0;
     memset(cmp, -1, sizeof(cmp)), qcmp=0;
     forn(i, n){
35
       if(!idx[i]) tjn(i);
36
        if(!idx[neg(i)]) tjn(neg(i));
37
     }
38
     forn(i, n) if(cmp[i] == cmp[neg(i)]) return false;
39
     return true:
40
41 }
```

#### Articulation Points

```
int N;
   vector<int> G[1000000]:
   //V[i]=node number(if visited), L[i]= lowest V[i] reachable from i
   int qV, V[1000000], L[1000000], P[1000000];
   void dfs(int v, int f){
    L[v]=V[v]=++qV;
6
     forall(it, G[v])
       if(!V[*it]){
8
         dfs(*it, v);
9
         L[v] = min(L[v], L[*it]);
10
         P[v] += L[*it] >= V[v];
11
       }
12
       else if(*it!=f)
13
         L[v]=min(L[v], V[*it]);
14
15
   int cantart() { //0(n)
16
     qV=0;
17
     zero(V), zero(P);
18
     dfs(1, 0); P[1]--;
19
     int q=0;
```

```
22 return q;
23 }
```

forn(i, N) if(P[i]) q++;

7 GRAFOS - 7.7 Articulation Points

# 7.8. Comp. Biconexas y Puentas

```
struct edge {
     int u,v, comp;
     bool bridge;
   };
4
   vector<edge> e;
   void addEdge(int u, int v) {
     G[u].pb(sz(e)), G[v].pb(sz(e));
     e.pb((edge){u,v,-1,false});
9
   //d[i]=id de la dfs
   //b[i]=lowest id reachable from i
   int d[MAXN], b[MAXN], t;
   int nbc;//cant componentes
   int comp[MAXN];//comp[i]=cant comp biconexas a la cual pertenece i
   void initDfs(int n) {
     zero(G), zero(comp);
     e.clear();
     forn(i,n) d[i]=-1;
     nbc = t = 0:
19
   }
20
   stack<int> st;
   void dfs(int u, int pe) {\frac{}{0(n + m)}}
     b[u] = d[u] = t++;
     comp[u] = (pe != -1);
24
     forall(ne, G[u]) if (*ne != pe){
25
       int v = e[*ne].u ^e[*ne].v ^u;
26
       if (d[v] == -1) {
27
         st.push(*ne);
28
         dfs(v,*ne);
29
         if (b[v] > d[u]){
30
           e[*ne].bridge = true; // bridge
31
32
         if (b[v] >= d[u]) \{ // art \}
33
           int last;
34
           do {
35
             last = st.top(); st.pop();
36
             e[last].comp = nbc;
37
```

```
} while (last != *ne);
38
            nbc++;
39
            comp[u]++;
40
41
          b[u] = min(b[u], b[v]);
42
43
       else if (d[v] < d[u]) \{ // back edge
44
          st.push(*ne);
45
         b[u] = min(b[u], d[v]);
46
47
     }
48
   |}
49
```

#### 7.9. LCA + Climb

```
const int MAXN=100001;
   const int LOGN=20;
   //f[v][k] holds the 2^k father of v
   //L[v] holds the level of v
   int N, f[MAXN][LOGN], L[MAXN];
   //call before build:
   void dfs(int v, int fa=-1, int lvl=0){//generate required data
     f[v][0]=fa, L[v]=lvl;
     forall(it, G[v])if(*it!=fa)
       dfs(*it, v, lvl+1);
10
11
   void build(){//f[i][0] must be filled previously, O(nlgn)
     forn(k, LOGN-1) forn(i, N) f[i][k+1]=f[f[i][k]][k];}
13
14
   #define lg(x) (31-__builtin_clz(x))//=floor(log2(x))
15
16
   int climb(int a, int d){\frac{}{0(lgn)}}
17
     if(!d) return a;
18
     dforn(i, lg(L[a])+1)
19
       if(1<<i<=d)
20
         a=f[a][i], d-=1<<i;
21
       return a;
^{22}
23
   int lca(int a, int b){\frac{1}{0}}
     if(L[a]<L[b]) swap(a, b);</pre>
25
     a=climb(a, L[a]-L[b]);
26
     if(a==b) return a;
27
     dforn(i, lg(L[a])+1)
28
```

```
if(f[a][i]!=f[b][i])
29
         a=f[a][i], b=f[b][i];
30
     return f[a][0];
31
32
   int dist(int a, int b) {//returns distance between nodes
     return L[a]+L[b]-2*L[lca(a, b)];}
```

#### 7.10. Heavy Light Decomposition

```
int treesz[MAXN];//cantidad de nodos en el subarbol del nodo v
  int dad[MAXN];//dad[v]=padre del nodo v
   void dfs1(int v, int p=-1){//pre-dfs
     dad[v]=p;
     treesz[v]=1;
    forall(it, G[v]) if(*it!=p){
       dfs1(*it, v);
       treesz[v]+=treesz[*it];
8
9
10
   //PONER Q EN O !!!!!
   int pos[MAXN], q;//pos[v]=posicion del nodo v en el recorrido de la dfs
   //Las cadenas aparecen continuas en el recorrido!
   int cantcad;
   int homecad[MAXN]://dada una cadena devuelve su nodo inicial
   int cad[MAXN];//cad[v]=cadena a la que pertenece el nodo
   void heavylight(int v, int cur=-1){
     if(cur==-1) homecad[cur=cantcad++]=v;
18
     pos[v]=q++;
19
     cad[v]=cur;
20
     int mx=-1;
21
     forn(i, sz(G[v])) if(G[v][i]!=dad[v])
22
      if(mx==-1 || treesz[G[v][mx]]<treesz[G[v][i]]) mx=i;</pre>
23
     if(mx!=-1) heavylight(G[v][mx], cur);
24
     forn(i, sz(G[v])) if(i!=mx && G[v][i]!=dad[v])
25
       heavylight(G[v][i], -1);
26
27
   //ejemplo de obtener el maximo numero en el camino entre dos nodos
   //RTA: max(query(low, u), query(low, v)), con low=lca(u, v)
   //esta funcion va trepando por las cadenas
   int query(int an, int v){\frac{1}{0}(\log n)}
    //si estan en la misma cadena:
     if(cad[an]==cad[v]) return rmq.get(pos[an], pos[v]+1);
33
     return max(query(an, dad[homecad[cad[v]]]),
```

if(u!=r) explore(u, r, it2);

if(get(v)<sz(G[v])) q.push(it);</pre>

15

```
rmq.get(pos[homecad[cad[v]]], pos[v]+1));
                                                                                   17 | }
35
36 }
                                                                                   19
                    7.11. Centroid Decomposition
                                                                                  20
                                                                                  21
   int n;
   vector<int> G[MAXN];
                                                                                  23
   bool taken[MAXN];//poner todos en FALSE al principio!!
                                                                                   24
   int padre [MAXN]; //padre de cada nodo en el centroid tree
5
                                                                                   26
   int szt[MAXN];
                                                                                  27
   void calcsz(int v, int p) {
                                                                                   28
     szt[v] = 1;
8
                                                                                   29
     forall(it,G[v]) if (*it!=p && !taken[*it])
9
       calcsz(*it,v), szt[v]+=szt[*it];
10
                                                                                   31
11
                                                                                   32 }
   void centroid(int v=0, int f=-1, int lvl=0, int tam=-1) {//0(nlogn)
     if(tam==-1) calcsz(v, -1), tam=szt[v];
13
     forall(it, G[v]) if(!taken[*it] && szt[*it]>=tam/2)
14
       {szt[v]=0; centroid(*it, f, lvl, tam); return;}
15
     taken[v]=true;
16
     padre[v]=f:
17
     forall(it, G[v]) if(!taken[*it])
18
       centroid(*it, v, lvl+1, -1);
19
20 }
                            7.12. Euler Cycle
                                                                                   8
   int n,m,ars[MAXE], eq;
   vector<int> G[MAXN];//fill G,n,m,ars,eq
                                                                                   10
   list<int> path;
                                                                                  11
   int used[MAXN];
                                                                                  12
   bool usede[MAXE];
                                                                                   13
   queue<list<int>::iterator> q;
                                                                                        }
                                                                                   14
   int get(int v){
                                                                                   15
     while(used[v]<sz(G[v]) && usede[ G[v][used[v]] ]) used[v]++;</pre>
                                                                                   16
     return used[v];
                                                                                   17
9
10
   void explore(int v, int r, list<int>::iterator it){
11
     int ar=G[v][get(v)]; int u=v^ars[ar];
12
     usede[ar]=true:
                                                                                  21
     list<int>::iterator it2=path.insert(it, u);
14
```

```
void euler(){
     zero(used), zero(usede);
     path.clear();
     q=queue<list<int>::iterator>();
     path.push_back(0); q.push(path.begin());
     while(sz(q)){
       list<int>::iterator it=q.front(); q.pop();
       if(used[*it] < sz(G[*it])) explore(*it, *it, it);</pre>
     reverse(path.begin(), path.end());
   void addEdge(int u, int v){
     G[u].pb(eq), G[v].pb(eq);
     ars[eq++]=u^v;
                         7.13. Diametro árbol
vector<int> G[MAXN];
   int n,m;
   int p[MAXN],d[MAXN],d2[MAXN];
   int bfs(int r. int *d) {
     queue<int> q;
     d[r]=0; q.push(r);
     int v;
     while(sz(q)) { v=q.front(); q.pop();
      forall(it,G[v]) if (d[*it]==-1) {
         d[*it]=d[v]+1, p[*it]=v;
         q.push(*it);
     return v;//ultimo nodo visitado
   vector<int> diams;
   vector<ii> centros:
   void diametros(){
    memset(d,-1,sizeof(d));
     memset(d2,-1,sizeof(d2));
     diams.clear(), centros.clear();
23
    forn(i, n) if(d[i]==-1){
```

```
int v,c;
25
       c=v=bfs(bfs(i, d2), d);
26
       forn(_,d[v]/2) c=p[c];
27
       diams.pb(d[v]);
28
       if(d[v]&1) centros.pb(ii(c, p[c]));
29
       else centros.pb(ii(c, c));
30
31
32 }
```

#### 7.14. Chu-liu

```
void visit(graph &h, int v, int s, int r,
     vector<int> &no, vector< vector<int> > &comp,
     vector<int> &prev, vector< vector<int> > &next, vector<weight> &mcost,
     vector<int> &mark, weight &cost, bool &found) {
     if (mark[v]) {
5
       vector<int> temp = no;
6
       found = true;
7
       do {
8
         cost += mcost[v];
         v = prev[v];
10
         if (v != s) {
11
           while (comp[v].size() > 0) {
12
             no[comp[v].back()] = s;
13
             comp[s].push_back(comp[v].back());
14
             comp[v].pop_back();
15
16
17
       } while (v != s);
18
       forall(j,comp[s]) if (*j != r) forall(e,h[*j])
19
         if (no[e->src] != s) e->w -= mcost[ temp[*i] ];
20
     }
^{21}
     mark[v] = true;
^{22}
     forall(i,next[v]) if (no[*i] != no[v] && prev[no[*i]] == v)
23
       if (!mark[no[*i]] || *i == s)
24
         visit(h, *i, s, r, no, comp, prev, next, mcost, mark, cost, found)
25
26
   weight minimumSpanningArborescence(const graph &g, int r) {
       const int n=sz(g);
28
     graph h(n);
29
     forn(u,n) forall(e,g[u]) h[e->dst].pb(*e);
30
     vector<int> no(n);
31
```

```
vector<vector<int> > comp(n);
     forn(u, n) comp[u].pb(no[u] = u);
33
     for (weight cost = 0; ;) {
34
       vector<int> prev(n, -1);
35
       vector<weight> mcost(n, INF);
36
       forn(j,n) if (j != r) forall(e,h[j])
         if (no[e->src] != no[i])
38
           if (e->w < mcost[ no[j] ])</pre>
             mcost[ no[j] ] = e->w, prev[ no[j] ] = no[e->src];
       vector< vector<int> > next(n);
       forn(u,n) if (prev[u] >= 0)
42
         next[ prev[u] ].push_back(u);
43
       bool stop = true;
44
       vector<int> mark(n):
       forn(u,n) if (u != r && !mark[u] && !comp[u].empty()) {
46
         bool found = false:
         visit(h, u, u, r, no, comp, prev, next, mcost, mark, cost, found);
         if (found) stop = false;
       }
50
       if (stop) {
         forn(u,n) if (prev[u] >= 0) cost += mcost[u];
         return cost;
53
       }
54
    }
55
56 }
```

# 7.15. Hungarian

```
1 #define MAXN 256
   #define INFTO 0x7f7f7f7f
   int n;
   int mt[MAXN] [MAXN]; // Matriz de costos (X * Y)
   int xy[MAXN], yx[MAXN]; // Matching resultante (X->Y, Y->X)
   int lx[MAXN], ly[MAXN], slk[MAXN], slkx[MAXN], prv[MAXN];
   char S[MAXN], T[MAXN];
   void updtree(int x) {
    forn(y, n) if (lx[x] + ly[y] - mt[x][y] < slk[y]) {
       slk[y] = lx[x] + ly[y] - mt[x][y];
10
       slkx[y] = x;
11
  } }
12
   int hungar(){//Matching maximo de mayor costo en grafos dirigidos (N^3)
     forn(i, n) {
14
       ly[i] = 0;
15
```

```
lx[i] = *max_element(mt[i], mt[i]+n); }
16
     memset(xy, -1, sizeof(xy));
17
     memset(yx, -1, sizeof(yx));
18
     forn(m, n) {
19
       memset(S, 0, sizeof(S));
20
       memset(T, 0, sizeof(T));
21
       memset(prv, -1, sizeof(prv));
^{22}
       memset(slk, 0x7f, sizeof(slk));
23
       queue<int> q;
^{24}
   #define bpone(e, p) { q.push(e); prv[e] = p; S[e] = 1; updtree(e); }
25
       forn(i, n) if (xy[i] == -1) { bpone(i, -2); break; }
26
       int x=0, y=-1;
27
       while (v==-1) {
28
         while (!q.empty() && y==-1) {
29
           x = q.front(); q.pop();
30
           forn(j, n) if (mt[x][j] == lx[x] + ly[j] && !T[j]) {
31
             if (yx[j] == -1) \{ y = j; break; \}
32
             T[i] = 1;
33
             bpone(yx[j], x);
34
35
         }
36
         if (v!=-1) break;
37
         int dlt = INFTO;
38
         forn(j, n) if (!T[j]) dlt = min(dlt, slk[j]);
39
         forn(k, n) {
40
           if (S[k]) lx[k] = dlt;
41
           if (T[k]) ly [k] += dlt;
42
           if (!T[k]) slk[k] -= dlt;
43
         }
44
         forn(j, n) if (!T[j] && !slk[j]) {
45
           if (yx[i] == -1) {
46
             x = slkx[j]; y = j; break;
47
           } else {
48
             T[i] = 1;
49
             if (!S[yx[j]]) bpone(yx[j], slkx[j]);
50
           }
51
         }
52
       }
53
       if (y!=-1) {
54
         for(int p = x; p != -2; p = prv[p]) {
55
           yx[y] = p;
56
           int ty = xy[p]; xy[p] = y; y = ty;
57
58
```

# 8. Network Flow

#### 8.1. Dinic

```
const int MAX = 300;
   int nodes, src, dst;
   int dist[MAX], q[MAX], work[MAX];
   struct Edge {
6
       int to, rev;
       11 f, cap;
       Edge(int to, int rev, 11 f, 11 cap) : to(to), rev(rev), f(f), cap(
           cap) {}
   };
10
   vector<Edge> G[MAX];
   // Adds bidirectional edge
   void addEdge(int s, int t, ll cap){
       G[s].pb(Edge(t, sz(G[t]), 0, cap));
15
       G[t].pb(Edge(s, sz(G[s])-1, 0, 0));
16
   }
17
18
   bool dinic_bfs(){
       fill(dist, dist+nodes, -1);
20
       dist[src]=0;
21
       int qt=0;
22
       a[at++]=src;
23
       for(int qh=0; qh<qt; qh++){</pre>
24
           int u =q[qh];
25
           forall(e, G[u]){
26
                int v=e->to;
27
                if(dist[v]<0 \&\& e>f < e>cap){
28
                    dist[v]=dist[u]+1;
29
                    q[qt++]=v;
30
                }
31
```

```
}
32
       }
33
       return dist[dst]>=0;
34
35
36
   ll dinic_dfs(int u, ll f){
37
       if(u==dst) return f;
38
       for(int &i=work[u]; i<sz(G[u]); i++){</pre>
39
            Edge &e = G[u][i];
40
           if(e.cap<=e.f) continue;</pre>
41
            int v=e.to;
42
           if(dist[v]==dist[u]+1){
43
                    11 df=dinic_dfs(v, min(f, e.cap-e.f));
                    if(df>0){
45
                             e.f+=df:
46
                             G[v][e.rev].f-= df;
47
                             return df;
48
                    }
49
           }
50
       }
51
       return 0;
52
53
54
   ll maxFlow(int _src, int _dst){
55
       src=_src;
56
       dst=_dst;
57
       11 result=0;
58
       while(dinic_bfs()){
59
           fill(work, work+nodes, 0);
60
            while(ll delta=dinic_dfs(src,INF))
61
                result+=delta;
62
       }
63
       // todos los nodos con dist[v]!=-1 vs los que tienen dist[v]==-1
64
            forman el min-cut
       return result:
65
66 }
                                 8.2. Konig
```

```
int s[maxnodes]; // numero de la bfs del koning
5 | queue<int> kq;
6 // s[e] %2==1 o si e esta en V1 y s[e] ==-1-> lo agarras
   void koning() {//O(n)
     forn(v,nodes-2) s[v] = match[v] = -1;
     forn(v,nodes-2) forall(it,g[v]) if (it->to < nodes-2 && it->f>0)
       { match[v]=it->to; match[it->to]=v;}
     form(v,nodes-2) if (match[v]==-1) {s[v]=0;kq.push(v);}
11
     while(!kq.empty()) {
12
       int e = kq.front(); kq.pop();
13
       if (s[e] %2==1) {
14
         s[match[e]] = s[e]+1;
15
         kq.push(match[e]);
16
       } else {
17
18
         forall(it,g[e]) if (it->to < nodes-2 && s[it->to]==-1) {
19
           s[it->to] = s[e]+1;
20
           kq.push(it->to);
21
         }
22
23
    }
24
25 }
```

# 8.3. Edmonds Karp's

```
1 #define MAX V 1000
   #define INF 1e9
   //special nodes
   #define SRC 0
   #define SNK 1
   map<int, int> G[MAX_V];//limpiar esto
   //To add an edge use
   #define add(a, b, w) G[a][b]=w
   int f, p[MAX_V];
   void augment(int v, int minE){
    if(v==SRC) f=minE;
11
     else if(p[v]!=-1){
12
       augment(p[v], min(minE, G[p[v]][v]));
13
       G[p[v]][v]=f, G[v][p[v]]+=f;
14
    }
15
   }
16
17 | ll maxflow(){//0(VE^2)
    11 Mf=0;
18
```

```
do{
19
       f=0;
20
       char used[MAX_V]; queue<int> q; q.push(SRC);
21
       zero(used), memset(p, -1, sizeof(p));
^{22}
       while(sz(q)){
23
         int u=q.front(); q.pop();
24
         if(u==SNK) break;
25
         forall(it, G[u])
26
           if(it->snd>0 && !used[it->fst])
27
              used[it->fst]=true, q.push(it->fst), p[it->fst]=u;
28
       }
29
       augment(SNK, INF);
30
       Mf+=f:
31
     }while(f);
32
     return Mf;
33
34 }
```

# 8.4. Push-Relabel O(N3)

```
#define MAX V 1000
  int N://valid nodes are [0...N-1]
   #define INF 1e9
   //special nodes
   #define SRC 0
   #define SNK 1
   map<int, int> G[MAX_V];
   //To add an edge use
   #define add(a, b, w) G[a][b]=w
   11 excess[MAX_V];
   int height[MAX_V], active[MAX_V], count[2*MAX_V+1];
   queue<int> Q;
12
   void enqueue(int v) {
13
     if (!active[v] && excess[v] > 0) active[v]=true, Q.push(v); }
14
   void push(int a, int b) {
15
     int amt = min(excess[a], ll(G[a][b]));
16
     if(height[a] <= height[b] || amt == 0) return;</pre>
17
     G[a][b]-=amt, G[b][a]+=amt;
18
     excess[b] += amt. excess[a] -= amt:
19
     enqueue(b);
20
21
   void gap(int k) {
22
     forn(v, N){
23
       if (height[v] < k) continue;
24
```

```
count[height[v]]--;
25
       height[v] = max(height[v], N+1);
26
       count[height[v]]++;
27
       enqueue(v);
28
     }
29
30
   void relabel(int v) {
31
     count[height[v]]--;
     height[v] = 2*N;
33
     forall(it, G[v])
34
       if(it->snd)
35
         height[v] = min(height[v], height[it->fst] + 1);
36
     count[height[v]]++;
37
     enqueue(v);
38
39
   ll maxflow() \{//0(V^3)
     zero(height), zero(active), zero(count), zero(excess);
     count[0] = N-1;
     count[N] = 1:
43
     height[SRC] = N;
     active[SRC] = active[SNK] = true;
45
     forall(it, G[SRC]){
46
       excess[SRC] += it->snd;
47
       push(SRC, it->fst);
48
     }
49
     while(sz(Q)) {
50
       int v = Q.front(); Q.pop();
51
       active[v]=false;
52
     forall(it, G[v]) push(v, it->fst);
53
     if(excess[v] > 0)
54
       count[height[v]] == 1? gap(height[v]):relabel(v);
55
     }
56
     ll mf=0:
57
     forall(it, G[SRC]) mf+=G[it->fst][SRC];
     return mf:
59
60 }
                        8.5. Min-cost Max-flow
```

```
const int MAXN=10000;
const ll INF = 1e14;
struct edge {
```

```
int u, v;
5
     ll cap, cost, flow;
     11 rem() { return cap - flow; }
   };
8
   int nodes;//numero de nodos
   vector<int> G[MAXN];
   vector<edge> e;
   void addEdge(int u, int v, ll cap, ll cost) {
     G[u].pb(sz(e)); e.pb((edge){u,v,cap,cost,0});
13
     G[v].pb(sz(e)); e.pb((edge){v,u,0,-cost,0});
14
15
   11 pot[MAXN], dist[MAXN], pre[MAXN], cap[MAXN];
   11 mxFlow, mnCost;
   void flow(int s, int t) {
     fill(pot, pot+nodes, 0);
     mxFlow=mnCost=0;
20
     while(1){
21
       fill(dist, dist+nodes, INF); dist[s] = 0;
22
       fill(pre, pre+nodes, -1); pre[s]=0;
23
       fill(cap, cap+nodes, 0); cap[s] = INF;
24
       priority_queue<pair<11,int> > q; q.push(make_pair(0,s));
25
       //~ Bellman Ford
26
       //~ forn(i,nodes) {
27
         //~ forn(j,sz(e)) if (e[j].rem()) {
28
          //~ ll c = e[j].cost + pot[e[j].u] - pot[e[j].v];
29
           //~ if (dist[e[j].v] > dist[e[j].u] + c) {
30
            //~ dist[e[j].v] = dist[e[j].u] + c;
31
             //~ pre[e[j].v] = j;
32
             //~ cap[e[j].v] = min(cap[e[j].u], e[j].rem());
33
         //~ } }
34
       //~ }
35
       //~ Dijkstra
36
       while (!q.empty()) {
37
         pair<ll,int> top = q.top(); q.pop();
38
         int u = top.second;
39
         11 d = -top.first;
40
         if (u == t) break;
41
         if (d > dist[u]) continue:
42
         forn(i,sz(G[u])) {
43
           edge E = e[G[u][i]];
44
           11 c = E.cost + pot[u] - pot[E.v];
45
           if (E.rem() && dist[E.v] > dist[u] + c) {
46
             dist[E.v] = dist[u] + c;
47
```

```
pre[E.v] = G[u][i];
48
             cap[E.v] = min(cap[u], E.rem());
49
             q.push(make_pair(-dist[E.v], E.v));
50
51
         }
52
53
       if (pre[t] == -1) break;
54
       forn(u,nodes)
         if (dist[u] == INF) pot[u] = INF;
         else pot[u] += dist[u];
       mxFlow +=cap[t];
58
       mnCost +=cap[t]*pot[t];
59
       for (int v = t; v != s; v = e[pre[v]].u) {
60
         e[pre[v]].flow += cap[t];
61
         e[pre[v]^1].flow -= cap[t];
62
       }
63
    }
65 }
```

# 9. Template

```
1 | #include <bits/stdc++.h>
   using namespace std;
   #define dprint(v) cerr << #v"=" << v << endl //;)
   #define forr(i,a,b) for(int i=(a); i<(b); i++)</pre>
   #define forn(i,n) forr(i,0,n)
   #define dforn(i,n) for(int i=n-1; i>=0; i--)
   #define forall(it,v) for(typeof(v.begin()) it=v.begin();it!=v.end();++it
   #define sz(c) ((int)c.size())
   #define zero(v) memset(v, 0, sizeof(v))
   #define pb push_back
   #define fst first
   #define snd second
   typedef long long 11;
   typedef pair<int,int> ii;
15
   int main() {
     freopen("input.in", "r", stdin);
17
       ios::sync_with_stdio(0);
18
     while(){
19
20
     }
21
```

```
Caloventor en Dos - Universidad Nacional de Rosario
    return 0;
23 }
                             Ayudamemoria
                      10.
                           Cant. decimales
1 #include <iomanip>
cout << setprecision(2) << fixed;</pre>
               Rellenar con espacios(para justificar)
1 #include <iomanip>
cout << setfill(''') << setw(3) << 2 << endl;
                       Leer hasta fin de linea
  #include <sstream>
   //hacer cin.ignore() antes de getline()
   while(getline(cin, line)){
        istringstream is(line);
        while(is >> X)
         cout << X << "";
        cout << endl;</pre>
8
                               Aleatorios
  #define RAND(a, b) (rand()%(b-a+1)+a)
srand(time(NULL));
                            Doubles Comp.
  const double EPS = 1e-9;
  |x == y  <=> fabs(x-y)  < EPS
_3 | x > y <=> x > y + EPS
_4 | x >= y <=> x > y - EPS
                                Limites
1 #include <limits>
2 | numeric_limits<T>
     ::max()
     ::min()
```

::epsilon()

```
Muahaha
```

```
1 #include <signal.h>
  void divzero(int p){
    while(true);}
  |void segm(int p){
    exit(0);}
  //in main
  signal(SIGFPE, divzero);
8 | signal(SIGSEGV, segm);
                          Mejorar velocidad
ios::sync_with_stdio(false);
                         Mejorar velocidad 2
1 //Solo para enteros positivos
 inline void Scanf(int& a){
    char c = 0;
    while(c<33) c = getc(stdin);</pre>
    a = 0;
    while(c>33) a = a*10 + c - '0', c = getc(stdin);
                            Expandir pila
#include <sys/resource.h>
2 rlimit rl;
  getrlimit(RLIMIT_STACK, &rl);
4 | rl.rlim_cur=1024L*1024L*256L;//256mb
5 setrlimit(RLIMIT_STACK, &rl);
                                C++11
1 g++ --std=c++1
                           Leer del teclado
freopen("/dev/tty", "a", stdin);
                         Iterar subconjunto
for(int sbm=bm; sbm; sbm=(sbm-1)&bm)
```

# File setup

```
//tambien se pueden usar comas: {a, x, m, 1} touch {a..l}.in; tee {a..l}.cpp < template.cpp
```