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# 1. algorithm

#include <algorithm> #include <numeric>

		nclude <numeric></numeric>
Algo	Params	Funcion
sort, stable_sort	f, 1	ordena el intervalo
nth_element	f, nth, l	void ordena el n-esimo, y
		particiona el resto
fill, fill_n	f, l / n, elem	void llena [f, l) o [f,
		f+n) con elem
lower_bound, upper_bound	f, l, elem	it al primer / ultimo donde se
		puede insertar elem para que
		quede ordenada
binary_search	f, l, elem	bool esta elem en [f, l)
copy	f, l, resul	hace $resul+i=f+i \ \forall i$
find, find_if, find_first_of	f, l, elem	$it$ encuentra i $\in$ [f,l) tq. i $=$ elem,
	/ pred / f2, l2	$pred(i), i \in [f2,l2)$
count, count_if	f, l, elem/pred	cuenta elem, pred(i)
search	f, l, f2, l2	busca $[f2,l2) \in [f,l)$
replace, replace_if	f, l, old	cambia old / pred(i) por new
	/ pred, new	
reverse	f, 1	da vuelta
partition, stable_partition	f, l, pred	pred(i) ad, !pred(i) atras
min_element, max_element	f, l, [comp]	$it \min, \max de [f,l]$
lexicographical_compare	f1,l1,f2,l2	bool con [f1,11];[f2,12]
$next/prev\_permutation$	f,1	deja en [f,l) la perm sig, ant
$set\_intersection,$	f1, l1, f2, l2, res	[res,) la op. de conj
set_difference, set_union,		
set_symmetric_difference,		
push_heap, pop_heap,	f, l, e / e /	mete/saca e en heap [f,l),
make_heap		hace un heap de [f,l)
is_heap	f,l	bool es [f,l) un heap
accumulate	f,l,i,[op]	$T = \sum \text{oper de [f,l)}$
inner_product	f1, l1, f2, i	$T = i + [f1, 11) \cdot [f2, \dots)$
partial_sum	f, l, r, [op]	$r+i = \sum /oper de [f,f+i] \forall i \in [f,l)$
builtin_ffs	unsigned int	Pos. del primer 1 desde la derecha
_builtin_clz	unsigned int	Cant. de ceros desde la izquierda.
_builtin_ctz	unsigned int	Cant. de ceros desde la derecha.
_builtin_popcount	unsigned int	Cant. de 1's en x.
_builtin_parity	unsigned int	1 si x es par, 0 si es impar.
builtin_XXXXXXII	unsigned ll	= pero para long long's.

# 2. Estructuras

# 2.1. RMQ (static)

Dado un arreglo y una operacion asociativa *idempotente*, get(i, j) opera sobre el rango [i, j). Restriccion: LVL ≥ ceil(logn); Usar [] para llenar arreglo y luego build().

```
1 | struct RMQ{
     #define LVL 10
2
     tipo vec[LVL] [1<<(LVL+1)];
     tipo &operator[](int p){return vec[0][p];}
     tipo get(int i, int j) {//intervalo [i,j)
5
       int p = 31-_builtin_clz(j-i);
6
       return min(vec[p][i],vec[p][j-(1<<p)]);</pre>
7
     }
8
     void build(int n) {//O(nlogn)
9
       int mp = 31-__builtin_clz(n);
10
       forn(p, mp) forn(x, n-(1<<p))
11
         vec[p+1][x] = min(vec[p][x], vec[p][x+(1<<p)]);
12
     }};
13
```

# 2.2. RMQ (dynamic)

```
1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera
        sobre el rango [i, j).
   #define MAXN 100000
   #define operacion(x, y) max(x, y)
   const int neutro=0;
   struct RMQ{
     int sz;
6
     tipo t[4*MAXN];
7
     tipo &operator[](int p){return t[sz+p];}
8
     void init(int n){//O(nlgn)
9
       sz = 1 \ll (32-\_builtin\_clz(n));
10
       forn(i, 2*sz) t[i]=neutro;
11
12
     void updall(){\frac{}{0}}
13
       dforn(i, sz) t[i]=operacion(t[2*i], t[2*i+1]);}
14
     tipo get(int i, int j){return get(i,j,1,0,sz);}
15
     tipo get(int i, int j, int n, int a, int b){\frac{1}{0}}
16
       if(j<=a || i>=b) return neutro;
17
       if(i<=a && b<=j) return t[n];
18
       int c=(a+b)/2;
19
```

```
return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
20
21
     void set(int p, tipo val){//O(lgn)
22
       for(p+=sz; p>0 && t[p]!=val;){
23
         t[p]=val;
24
         p/=2;
25
         val=operacion(t[p*2], t[p*2+1]);
26
27
     }
28
   }rma;
   //Usage:
31 | cin >> n; rmg.init(n); forn(i, n) cin >> rmg[i]; rmg.updall();
                            2.3. RMQ (lazy)
1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera
        sobre el rango [i, j).
   typedef int Elem; //Elem de los elementos del arreglo
   typedef int Alt;//Elem de la alteracion
   #define operacion(x,y) x+y
   const Elem neutro=0; const Alt neutro2=0;
   #define MAXN 100000
   struct RMQ{
     int sz:
     Elem t[4*MAXN]:
     Alt dirty[4*MAXN];//las alteraciones pueden ser de distinto Elem
10
     Elem &operator[](int p){return t[sz+p];}
11
     void init(int n){//O(nlgn)
12
       sz = 1 \ll (32-\_builtin\_clz(n));
13
       forn(i, 2*sz) t[i]=neutro;
14
       forn(i, 2*sz) dirty[i]=neutro2;
15
16
     void push(int n, int a, int b){//propaga el dirty a sus hijos
17
       if(dirty[n]!=0){
18
         t[n]+=dirty[n]*(b-a);//altera el nodo
19
         if(n \le z)
20
           dirty[2*n]+=dirty[n];
21
           dirty[2*n+1]+=dirty[n];
22
23
         dirty[n]=0;
24
25
26
     }
     Elem get(int i, int j, int n, int a, int b){\frac{1}{0}}
```

```
if(j<=a || i>=b) return neutro;
28
       push(n, a, b);//corrige el valor antes de usarlo
29
       if(i<=a && b<=j) return t[n];</pre>
30
       int c=(a+b)/2;
31
       return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
32
33
     Elem get(int i, int j){return get(i,j,1,0,sz);}
34
     //altera los valores en [i, j) con una alteración de val
35
     void alterar(Alt val, int i, int j, int n, int a, int b)\frac{1}{0}
36
       push(n, a, b);
37
       if(j<=a || i>=b) return;
38
       if(i<=a && b<=j){
39
         dirty[n]+=val;
40
         push(n, a, b);
41
         return;
42
       }
43
       int c=(a+b)/2:
44
       alterar(val, i, j, 2*n, a, c), alterar(val, i, j, 2*n+1, c, b);
45
       t[n]=operacion(t[2*n], t[2*n+1]);//por esto es el push de arriba
46
47
     void alterar(Alt val, int i, int j){alterar(val,i,j,1,0,sz);}
48
49 |}rmq;
```

### 2.4. RMQ (persistente)

```
typedef int tipo;
   tipo oper(const tipo &a, const tipo &b){
       return a+b;
3
4
   struct node{
5
     tipo v; node *1,*r;
6
     node(tipo v):v(v), 1(NULL), r(NULL) {}
7
       node(node *1, node *r) : 1(1), r(r){
8
           if(!1) v=r->v;
9
           else if(!r) v=l->v;
10
           else v=oper(1->v, r->v);
11
       }
12
   };
13
   node *build (tipo *a, int tl, int tr) {//modificar para que tome tipo a
     if (tl+1==tr) return new node(a[tl]);
15
     int tm=(tl + tr)>>1:
16
     return new node(build(a, tl, tm), build(a, tm, tr));
17
18 | }
```

```
node *update(int pos, int new_val, node *t, int tl, int tr){
     if (tl+1==tr) return new node(new_val);
     int tm=(tl+tr)>>1;
21
     if(pos < tm) return new node(update(pos, new_val, t->1, tl, tm), t->r)
22
     else return new node(t->1, update(pos, new_val, t->r, tm, tr));
23
24
   tipo get(int 1, int r, node *t, int tl, int tr){
25
       if(l==tl && tr==r) return t->v;
26
     int tm=(tl + tr)>>1;
27
       if(r<=tm) return get(1, r, t->1, t1, tm);
28
       else if(l>=tm) return get(l, r, t->r, tm, tr);
29
    return oper(get(1, tm, t->1, tl, tm), get(tm, r, t->r, tm, tr));
31 }
```

#### 2.5. Fenwick Tree

```
1 //For 2D threat each column as a Fenwick tree, by adding a nested for in
        each operation
2 struct Fenwick{
     static const int sz=1000001;
     tipo t[sz];
     void adjust(int p, tipo v){//valid with p in [1, sz), O(lgn)
       for(int i=p; i<sz; i+=(i&-i)) t[i]+=v; }</pre>
6
     tipo sum(int p){//cumulative sum in [1, p], O(lgn)
7
       tipo s=0:
8
       for(int i=p; i; i-=(i&-i)) s+=t[i];
       return s;
10
11
     tipo sum(int a, int b){return sum(b)-sum(a-1);}
12
     //get largest value with cumulative sum less than or equal to x;
13
     //for smallest, pass x-1 and add 1 to result
14
     int getind(tipo x) {//O(lgn)
15
         int idx = 0, mask = N;
16
         while(mask && idx < N) {</pre>
17
           int t = idx + mask;
18
         if(x >= tree[t])
19
             idx = t, x -= tree[t]:
20
           mask >>= 1:
21
22
         return idx:
23
    }};
24
```

}

23

### 2.6. Union Find

```
struct UnionFind{
    vector<int> f;//the array contains the parent of each node
2
    void init(int n){f.clear(); f.insert(f.begin(), n, -1);}
3
    int comp(int x){return (f[x]=-1?x:f[x]=comp(f[x]));}//0(1)
4
    bool join(int i, int j) {
5
      bool con=comp(i)==comp(j);
6
      if(!con) f[comp(i)] = comp(j);
      return con;
8
    }};
9
```

# 2.7. Disjoint Intervals

```
| bool operator< (const ii &a, const ii &b) {return a.fst<b.fst;}
   //Stores intervals as [first, second]
   //in case of a collision it joins them in a single interval
   struct disjoint_intervals {
     set<ii>> segs;
5
     void insert(ii v) {//O(lgn)
6
       if(v.snd-v.fst==0.) return;//0J0
       set<ii>>::iterator it,at;
8
       at = it = segs.lower_bound(v);
9
       if (at!=segs.begin() && (--at)->snd >= v.fst)
10
         v.fst = at->fst, --it;
11
       for(; it!=segs.end() && it->fst <= v.snd; segs.erase(it++))</pre>
12
         v.snd=max(v.snd, it->snd);
13
       segs.insert(v);
14
15
<sub>16</sub> | };
```

# 2.8. RMQ (2D)

```
struct RMQ2D{//n filas x m columnas
     int sz;
2
     RMQ t[4*MAXN];
3
     RMQ &operator[](int p){return t[sz/2+p];}//t[i][j]=i fila, j col
4
     void init(int n, int m){\frac{}{\frac{1}{0}(n*m)}}
5
       sz = 1 \ll (32-\_builtin\_clz(n));
6
       forn(i, 2*sz) t[i].init(m); }
7
     void set(int i, int j, tipo val){//O(lgm.lgn)
       for(i+=sz; i>0;){
9
         t[i].set(j, val);
10
         i/=2;
11
```

```
val=operacion(t[i*2][j], t[i*2+1][j]);
12
       } }
13
     tipo get(int i1, int j1, int i2, int j2){return get(i1, j1, i2, j2, 1, 0,
14
     //O(lgm.lgn), rangos cerrado abierto
15
     int get(int i1, int j1, int i2, int j2, int n, int a, int b){
16
       if(i2<=a || i1>=b) return 0;
17
       if(i1<=a && b<=i2) return t[n].get(j1, j2);</pre>
18
       int c=(a+b)/2;
19
       return operacion(get(i1, j1, i2, j2, 2*n, a, c),
            get(i1, j1, i2, j2, 2*n+1, c, b));
21
     }
22
   } rma:
23
   //Example to initialize a grid of M rows and N columns:
   RMQ2D rmq; rmq.init(n,m);
   forn(i, n) forn(j, m){
    int v; cin >> v; rmq.set(i, j, v);}
                               2.9. Big Int
 1 #define BASEXP 6
   #define BASE 1000000
   #define LMAX 1000
   struct bint{
       int 1;
5
       ll n[LMAX]:
6
       bint(11 x=0){
7
           1=1;
8
           forn(i, LMAX){
9
                if (x) l=i+1;
10
                n[i]=x %BASE;
11
                x/=BASE;
12
13
           }
14
15
       bint(string x){
16
       l=(x.size()-1)/BASEXP+1;
17
           fill(n, n+LMAX, 0);
18
           ll r=1;
19
           forn(i, sz(x)){
20
               n[i / BASEXP] += r * (x[x.size()-1-i]-'0');
21
                r*=10; if(r==BASE)r=1;
22
```

```
}
24
       void out(){
25
       cout << n[1-1];
26
       dforn(i, l-1) printf("%6.6llu", n[i]);//6=BASEXP!
27
28
     void invar(){
29
       fill(n+l, n+LMAX, 0);
30
       while(1>1 && !n[1-1]) 1--;
31
    }
32
33
   bint operator+(const bint&a, const bint&b){
34
     bint c;
35
       c.1 = max(a.1, b.1):
36
       11 q = 0;
37
       forn(i, c.1) q += a.n[i]+b.n[i], c.n[i]=q %BASE, q/=BASE;
38
       if(q) c.n[c.l++] = q;
39
       c.invar();
40
       return c;
41
42
   pair<bint, bool> lresta(const bint& a, const bint& b) // c = a - b
44
     bint c;
45
       c.1 = max(a.1, b.1);
46
       11 q = 0;
47
       forn(i, c.l) q += a.n[i]-b.n[i], c.n[i]=(q+BASE) %BASE, q=(q+BASE)/
48
           BASE-1;
       c.invar();
49
       return make_pair(c, !q);
50
51
   bint& operator-= (bint& a, const bint& b){return a=lresta(a, b).first;}
   bint operator- (const bint&a, const bint&b) {return lresta(a, b).first;}
   bool operator< (const bint&a, const bint&b){return !lresta(a, b).second
       :}
   bool operator<= (const bint&a, const bint&b){return lresta(b, a).second
       :}
   bool operator==(const bint&a, const bint&b){return a <= b && b <= a;}
   bint operator*(const bint&a, 11 b){
57
       bint c:
58
59
       forn(i, a.1) q += a.n[i]*b, c.n[i] = q \text{BASE}, q/=BASE;
60
61
       while(q) c.n[c.l++] = q %BASE, q/=BASE;
62
       c.invar();
63
```

```
return c;
64
   }
65
   bint operator*(const bint&a, const bint&b){
66
        bint c;
        c.1 = a.1+b.1;
       fill(c.n, c.n+b.1, 0);
       forn(i, a.1){
           11 q = 0;
            forn(j, b.1) q += a.n[i]*b.n[j]+c.n[i+j], c.n[i+j] = q BASE, q
                /=BASE;
            c.n[i+b.1] = q;
73
74
       c.invar();
75
       return c;
77
    pair<br/>
\frac{b}{c} = a / b; rm = a % b
     bint c:
     11 \text{ rm} = 0;
     dforn(i, a.1){
                rm = rm * BASE + a.n[i];
                c.n[i] = rm / b;
83
                rm %= b;
84
       }
       c.1 = a.1;
        c.invar();
       return make_pair(c, rm);
89
    bint operator/(const bint&a, ll b){return ldiv(a, b).first;}
    11 operator%(const bint&a, 11 b){return ldiv(a, b).second;}
    pair<bint, bint> ldiv(const bint& a, const bint& b){
     bint c;
93
       bint rm = 0;
94
        dforn(i, a.l){
95
            if (rm.l==1 && !rm.n[0])
96
                rm.n[0] = a.n[i];
97
            elsef
98
                dforn(j, rm.l) rm.n[j+1] = rm.n[j];
99
                rm.n[0] = a.n[i]:
100
                rm.l++;
101
102
            ll q = rm.n[b.1] * BASE + rm.n[b.1-1];
103
            ll u = q / (b.n[b.l-1] + 1);
104
           ll v = q / b.n[b.l-1] + 1;
105
```

```
while (u < v-1)
106
                11 m = (u+v)/2;
107
                if (b*m \le rm) u = m;
108
                else v = m;
109
           }
110
           c.n[i]=u;
111
           rm-=b*u;
112
       }
113
     c.l=a.l;
114
       c.invar();
115
       return make_pair(c, rm);
116
117
   bint operator/(const bint&a, const bint&b) {return ldiv(a, b).first;}
   bint operator %(const bint&a, const bint&b) {return ldiv(a, b).second;}
                            2.10. HashTables
   //Compilar: g++ --std=c++11
   struct Hash{
 2
     size_t operator()(const ii &a)const{
       size_t s=hash<int>()(a.fst);
 4
       return hash<int>()(a.snd)+0x9e3779b9+(s<<6)+(s>>2);
 5
 6
     size_t operator()(const vector<int> &v)const{
       size_t s=0;
 8
       for(auto &e : v)
 9
         s = hash<int>()(e)+0x9e3779b9+(s<<6)+(s>>2);
10
       return s;
11
     }
12
13
    unordered_set<ii, Hash> s;
   unordered_map<ii, int, Hash> m;//map<key, value, hasher>
                              2.11. Modnum
   struct mnum{
     static const tipo mod=12582917;
 2
 3
     mnum(tipo v=0): v(v mod) {}
     mnum operator+(mnum b){return v+b.v;}
 5
     mnum operator-(mnum b){return v>=b.v? v-b.v : mod-b.v+v;}
     mnum operator*(mnum b){return v*b.v;}
     mnum operator^(int n){
```

if(!n) return 1;

9

```
return n \% ? (*this)^(n/2) * (this) : (*this)^(n/2);}
11 };
                           2.12. Treap para set
typedef int Key;
   typedef struct node *pnode;
   struct node{
       Key key;
       int prior, size;
       pnode l,r;
       node(Key key=0): key(key), prior(rand()), size(1), 1(0), r(0) {}
7
   };
8
   static int size(pnode p) { return p ? p->size : 0; }
   void push(pnode p) {
     // modificar v propagar el dirty a los hijos aca(para lazy)
11
12
   // Update function and size from children's Value
   void pull(pnode p) {//recalcular valor del nodo aca (para rmg)
     p->size = 1 + size(p->1) + size(p->r);
16
   //junta dos arreglos
   pnode merge(pnode 1, pnode r) {
     if (!1 || !r) return 1 ? 1 : r;
     push(1), push(r);
     pnode t:
21
     if (1->prior < r->prior) 1->r=merge(1->r, r), t = 1;
     else r\rightarrow l=merge(1, r\rightarrow 1), t = r;
     pull(t);
24
     return t;
25
26
   //parte el arreglo en dos, l<key<=r
   void split(pnode t, Key key, pnode &1, pnode &r) {
       if (!t) return void(1 = r = 0);
29
       push(t);
30
       if (\text{key} \leftarrow \text{t->key}) split(\text{t->l}, \text{key}, l, \text{t->l}), r = t;
31
       else split(t->r, key, t->r, r), l = t;
32
       pull(t):
33
   }
34
35
   void erase(pnode &t, Key key) {
       if (!t) return;
37
       push(t);
38
```

```
if (key == t->key) t=merge(t->1, t->r);
                                                                                           int prior, size;
39
                                                                                    5
       else if (key < t->key) erase(t->1, key);
                                                                                           pnode 1,r,parent;
                                                                                    6
40
                                                                                          node(Value val): val(val), mini(val), dirty(0), prior(rand()), size
       else erase(t->r, key);
41
       if(t) pull(t);
                                                                                               (1), 1(0), r(0), parent(0) {}
42
                                                                                      };
                                                                                    8
43
                                                                                      static int size(pnode p) { return p ? p->size : 0; }
44
                                                                                      void push(pnode p) {//propagar dirty a los hijos(aca para lazy)
   ostream& operator<<(ostream &out, const pnode &t) {
45
     if(!t) return out;
                                                                                        p->val.fst+=p->dirty;
46
       return out << t->l << t->key << ''_' << t->r;
                                                                                        p->mini.fst+=p->dirty;
47
                                                                                   12
                                                                                        if(p->1) p->1->dirty+=p->dirty;
48
   pnode find(pnode t, Key key) {
                                                                                        if(p->r) p->r->dirty+=p->dirty;
                                                                                   14
       if (!t) return 0;
                                                                                        p->dirty=0;
                                                                                   15
50
       if (key == t->key) return t;
                                                                                   16
                                                                                      }
51
       if (key < t->key) return find(t->1, key);
                                                                                      static Value mini(pnode p) { return p ? push(p), p->mini : ii(1e9, -1);
52
       return find(t->r, key);
53
                                                                                      // Update function and size from children's Value
54
                                                                                      void pull(pnode p) {//recalcular valor del nodo aca (para rmq)
   struct treap {
55
                                                                                        p \rightarrow size = 1 + size(p \rightarrow 1) + size(p \rightarrow r);
       pnode root;
56
       treap(pnode root=0): root(root) {}
                                                                                        p->mini = min(min(p->val, mini(p->l)), mini(p->r));//operacion del rmq
57
       int size() { return ::size(root); }
58
       void insert(Key key) {
                                                                                        p->parent=0;
                                                                                   22
59
                                                                                        if(p->1) p->1->parent=p;
           pnode t1, t2; split(root, key, t1, t2);
60
           t1=::merge(t1,new node(key));
                                                                                        if(p->r) p->r->parent=p;
                                                                                   24
61
           root=::merge(t1,t2);
                                                                                   25
62
       }
                                                                                      //junta dos arreglos
                                                                                   26
63
                                                                                      pnode merge(pnode 1, pnode r) {
       void erase(Key key1, Key key2) {
64
           pnode t1,t2,t3;
                                                                                        if (!1 || !r) return 1 ? 1 : r;
65
                                                                                        push(1), push(r);
           split(root,key1,t1,t2);
                                                                                   29
66
           split(t2,key2, t2, t3);
                                                                                        pnode t;
                                                                                   30
67
           root=merge(t1,t3);
                                                                                        if (1->prior < r->prior) 1->r=merge(1->r, r), t = 1;
68
                                                                                   31
                                                                                        else r\rightarrow l=merge(1, r\rightarrow 1), t = r;
                                                                                   32
69
       void erase(Key key) {::erase(root, key);}
                                                                                        pull(t);
70
       pnode find(Key key) { return ::find(root, key); }
                                                                                        return t:
                                                                                   34
71
       Key &operator[](int pos){return find(pos)->key;}//ojito
                                                                                   35
72
   };
                                                                                       //parte el arreglo en dos, sz(l)==tam
73
                                                                                   36
treap merge(treap a, treap b) {return treap(merge(a.root, b.root));}
                                                                                      void split(pnode t, int tam, pnode &1, pnode &r) {
                                                                                        if (!t) return void(1 = r = 0);
                                                                                   38
                       2.13. Treap para arreglo
                                                                                        push(t):
                                                                                   39
                                                                                        if (tam \le size(t->1)) split(t->1, tam, 1, t->1), r = t;
                                                                                        else split(t->r, tam - 1 - size(t->l), t->r, r), l = t;
  typedef struct node *pnode;
                                                                                        pull(t);
  struct node{
                                                                                   ^{42}
       Value val, mini;
                                                                                   43
3
                                                                                   pnode at(pnode t, int pos) {
       int dirty;
4
```

```
if(!t) exit(1);
45
     push(t);
46
     if(pos == size(t->1)) return t;
47
     if(pos < size(t->1)) return at(t->1, pos);
48
     return at(t->r, pos - 1 - size(t->1));
49
50
   int getpos(pnode t){//inversa de at
     if(!t->parent) return size(t->1);
52
     if(t==t->parent->l) return getpos(t->parent)-size(t->r)-1;
53
     return getpos(t->parent)+size(t->1)+1;
54
55
   void split(pnode t, int i, int j, pnode &1, pnode &m, pnode &r) {
     split(t, i, l, t), split(t, j-i, m, r);}
   Value get(pnode &p, int i, int j){//like rmq
     pnode 1,m,r;
59
       split(p, i, j, l, m, r);
60
       Value ret=mini(m);
61
       p=merge(1, merge(m, r));
62
       return ret:
63
64
   void print(const pnode &t) {//for debugging
65
     if(!t) return;
66
       push(t);
67
       print(t->1);
68
       cout << t->val.fst << '';
69
       print(t->r);
70
71 }
```

#### 2.14. Convex Hull Trick

```
struct Line{tipo m,h;};
  tipo inter(Line a, Line b){
       tipo x=b.h-a.h, y=a.m-b.m;
       return x/y+(x\%?!((x>0)^(y>0)):0);//==ceil(x/y)
4
5
  struct CHT {
6
     vector<Line> c;
     bool mx:
8
     int pos;
9
     CHT(bool mx=0):mx(mx),pos(0){}//mx=1 si las query devuelven el max
     inline Line acc(int i){return c[c[0].m>c.back().m? i : sz(c)-1-i];}
11
     inline bool irre(Line x, Line y, Line z){
12
       return c[0].m>z.m? inter(y, z) <= inter(x, y)
13
```

```
: inter(y, z) >= inter(x, y);
14
     }
15
     void add(tipo m, tipo h) {//O(1), los m tienen que entrar ordenados
16
           if (mx) m*=-1, h*=-1;
17
       Line l=(Line){m, h};
18
           if(sz(c) && m==c.back().m) { 1.h=min(h, c.back().h), c.pop_back
19
                (); if(pos) pos--; }
           while(sz(c) \ge 2 \&\& irre(c[sz(c)-2], c[sz(c)-1], 1)) { c.pop_back
20
                (); if(pos) pos--; }
           c.pb(1);
21
22
     inline bool fbin(tipo x, int m) {return inter(acc(m), acc(m+1))>x;}
23
     tipo eval(tipo x){
24
       int n = sz(c);
25
       //query con x no ordenados O(lgn)
26
       int a=-1, b=n-1;
27
       while(b-a>1) { int m = (a+b)/2;
         if(fbin(x, m)) b=m;
         else a=m:
30
       return (acc(b).m*x+acc(b).h)*(mx?-1:1);
32
           //query 0(1)
       while(pos>0 && fbin(x, pos-1)) pos--;
34
       while(pos<n-1 && !fbin(x, pos)) pos++;
       return (acc(pos).m*x+acc(pos).h)*(mx?-1:1);
36
37
38 } ch;
```

# 2.15. Convex Hull Trick (Dynamic)

```
const ll is_query = -(1LL<<62);</pre>
2 struct Line {
       ll m, b;
       mutable multiset<Line>::iterator it;
4
       const Line *succ(multiset<Line>::iterator it) const;
5
       bool operator<(const Line& rhs) const {</pre>
6
            if (rhs.b != is_query) return m < rhs.m;</pre>
7
            const Line *s=succ(it):
8
           if(!s) return 0;
           11 x = rhs.m:
10
            return b - s -> b < (s -> m - m) * x:
11
       }
12
13 };
```

```
14 struct HullDynamic : public multiset<Line>{ // will maintain upper hull
       for maximum
       bool bad(iterator y) {
15
           iterator z = next(y);
16
           if (v == begin()) {
17
               if (z == end()) return 0;
18
               return y->m == z->m && y->b <= z->b;
19
20
           iterator x = prev(y);
21
           if (z == end()) return y->m == x->m && y->b <= x->b;
22
           return (x->b - y->b)*(z->m - y->m) >= (y->b - z->b)*(y->m - x->m)
23
               );
       }
24
       iterator next(iterator y){return ++y;}
25
       iterator prev(iterator y){return --y;}
26
       void insert_line(ll m, ll b) {
27
           iterator y = insert((Line) { m, b });
28
           y->it=y;
29
           if (bad(y)) { erase(y); return; }
30
           while (next(y) != end() && bad(next(y))) erase(next(y));
31
           while (y != begin() && bad(prev(y))) erase(prev(y));
32
       }
33
       ll eval(ll x) {
34
           Line 1 = *lower_bound((Line) { x, is_query });
35
           return 1.m * x + 1.b;
36
37
   }h;
38
   const Line *Line::succ(multiset<Line>::iterator it) const{
       return (++it==h.end()? NULL : &*it);}
40
```

#### 2.16. Gain-Cost Set

```
//esta estructura mantiene pairs(beneficio, costo)
//de tal manera que en el set quedan ordenados
//por beneficio Y COSTO creciente. (va borrando los que no son optimos)
struct V{
   int gain, cost;
   bool operator<(const V &b)const{return gain<b.gain;}
};
set<V> s;
void add(V x){
   set<V>::iterator p=s.lower_bound(x);//primer elemento mayor o igual
   if(p!=s.end() && p->cost <= x.cost) return;//ya hay uno mejor</pre>
```

```
p=s.upper_bound(x);//primer elemento mayor
12
     if(p!=s.begin()){//borro todos los peores (<=beneficio y >=costo)
13
       --p;//ahora es ultimo elemento menor o igual
14
       while(p->cost >= x.cost){
15
         if(p==s.begin()){s.erase(p); break;}
16
         s.erase(p--);
18
    }
19
     s.insert(x);
20
21
   int get(int gain){//minimo costo de obtener tal ganancia
22
     set<V>::iterator p=s.lower_bound((V){gain, 0});
     return p==s.end()? INF : p->cost;}
24
```

### 2.17. Set con busq binaria

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;

typedef tree<int,null_type,less<int>,//key,mapped type, comparator

rb_tree_tag,tree_order_statistics_node_update> set_t;

//find_by_order(i) devuelve iterador al i-esimo elemento

//order_of_key(k): devuelve la pos del lower bound de k

//Ej: 12, 100, 505, 1000, 10000.

//order_of_key(10) == 0, order_of_key(100) == 1,
//order_of_key(707) == 3, order_of_key(9999999) == 5
```

# 3. Algos

# 3.1. Longest Increasing Subsecuence

```
//Para non-increasing, cambiar comparaciones y revisar busq binaria
//Given an array, paint it in the least number of colors so that each
color turns to a non-increasing subsequence.
//Solution:Min number of colors=Length of the longest increasing
subsequence
int N, a[MAXN];//secuencia y su longitud
ii d[MAXN+1];//d[i]=ultimo valor de la subsecuencia de tamanio i
int p[MAXN];//padres
vector<int> R;//respuesta
void rec(int i){
if(i==-1) return;
R.push_back(a[i]);
```

```
rec(p[i]);
11
   }
^{12}
   int lis(){//O(nlogn)
13
     d[0] = ii(-INF, -1); forn(i, N) d[i+1]=ii(INF, -1);
     forn(i, N){
15
       int j = upper_bound(d, d+N+1, ii(a[i], INF))-d;
16
       if (d[j-1].first < a[i]&&a[i] < d[j].first){</pre>
17
         p[i]=d[j-1].second;
18
         d[j] = ii(a[i], i);
19
20
     }
21
     R.clear();
22
     dforn(i, N+1) if(d[i].first!=INF){
23
       rec(d[i].second)://reconstruir
24
       reverse(R.begin(), R.end());
25
       return i;//longitud
26
     }
27
     return 0;
28
29 }
```

# 3.2. Alpha-Beta prunning

```
1 | 11 alphabeta(State &s, bool player = true, int depth = 1e9, 11 alpha = -
       INF, 11 beta = INF) { //player = true -> Maximiza
       if(s.isFinal()) return s.score;
2
     //~ if (!depth) return s.heuristic();
3
       vector<State> children;
4
       s.expand(player, children);
5
       int n = children.size();
6
       forn(i, n) {
7
           ll v = alphabeta(children[i], !player, depth-1, alpha, beta);
8
           if(!player) alpha = max(alpha, v);
9
           else beta = min(beta, v);
10
           if(beta <= alpha) break;</pre>
11
12
       return !player ? alpha : beta;}
13
```

# 3.3. Mo's algorithm

```
6 int ans[MAXN], curans; //ans[i] = ans to ith query
  bool bymos(const Qu &a, const Qu &b){
       if(a.l/sq!=b.l/sq) return a.l<b.l;</pre>
       return (a.1/sq)&1? a.r<b.r : a.r>b.r;
9
   }
10
   void mos(){
11
       forn(i, t) qs[i].id=i;
12
       sort(qs, qs+t, bymos);
       int cl=0, cr=0;
       sq=sqrt(n);
       curans=0;
16
       forn(i, t){ //intervalos cerrado abiertos !!! importante!!
17
            Qu &q=qs[i];
18
            while(cl>q.1) add(--cl);
19
           while(cr<q.r) add(cr++);</pre>
20
            while(cl<q.1) remove(cl++);</pre>
21
            while(cr>q.r) remove(--cr);
22
            ans[q.id]=curans;
23
       }
24
25 }
```

# 4. Strings

#### 4.1. Manacher

```
int d1[MAXN];//d1[i]=long del maximo palindromo impar con centro en i
  int d2[MAXN];//d2[i]=analogo pero para longitud par
  //0 1 2 3 4
   //a a b c c <--d1[2]=3
   //a a b b <--d2[2]=2 (estan uno antes)
   void manacher(){
     int l=0, r=-1, n=sz(s);
    forn(i, n){
8
       int k=(i>r? 1 : min(d1[l+r-i], r-i));
9
       while(i+k<n && i-k>=0 && s[i+k]==s[i-k]) ++k;
10
       d1[i] = k--;
11
       if(i+k > r) l=i-k, r=i+k;
12
    }
13
     l=0, r=-1;
14
     forn(i, n){
15
       int k=(i>r? 0 : min(d2[1+r-i+1], r-i+1))+1;
16
       while(i+k-1<n && i-k>=0 && s[i+k-1]==s[i-k]) k++;
17
       d2[i] = --k;
18
```

```
if(i+k-1 > r) l=i-k, r=i+k-1;
                                                                                    1 #define MAX_N 1000
19
                                                                                      #define rBOUND(x) (x<n? r[x] : 0)</pre>
    }
20
                                                                                      //sa will hold the suffixes in order.
                                 4.2. KMP
                                                                                      int sa[MAX_N], r[MAX_N], n;
                                                                                      string s; //input string, n=sz(s)
  string T;//cadena donde buscar(where)
  string P;//cadena a buscar(what)
                                                                                      int f[MAX_N], tmpsa[MAX_N];
   int b[MAXLEN];//back table b[i] maximo borde de [0..i)
                                                                                       void countingSort(int k){
   void kmppre(){//by gabina with love
                                                                                        zero(f);
       int i = 0, j=-1; b[0]=-1;
5
                                                                                        forn(i, n) f[rBOUND(i+k)]++;
       while(i<sz(P)){</pre>
6
                                                                                        int sum=0;
                                                                                   11
           while(j>=0 && P[i] != P[j]) j=b[j];
                                                                                        forn(i, max(255, n)){
                                                                                   12
           i++, j++;
                                                                                          int t=f[i]; f[i]=sum; sum+=t;}
           b[i] = j;
9
                                                                                        forn(i, n)
                                                                                   14
       }
10
                                                                                           tmpsa[f[rBOUND(sa[i]+k)]++]=sa[i];
                                                                                   15
11
                                                                                        memcpy(sa, tmpsa, sizeof(sa));
                                                                                   16
12
                                                                                   17
   void kmp(){
13
                                                                                       void constructsa(){\frac{1}{0} \text{ (n log n)}}
       int i=0, j=0;
14
                                                                                        n=sz(s):
                                                                                   19
       while(i<sz(T)){</pre>
15
                                                                                        forn(i, n) sa[i]=i, r[i]=s[i];
           while(j>=0 && T[i]!=P[j]) j=b[j];
16
                                                                                        for(int k=1; k<n; k<<=1){
                                                                                   21
           i++, j++;
17
                                                                                           countingSort(k), countingSort(0);
                                                                                   22
           if(j==sz(P)){
18
                                                                                          int rank, tmpr[MAX_N];
                                                                                   23
               printf("P_is_found_at_index_\( \lambda_i\in_T\n\', i-j);
19
                                                                                          tmpr[sa[0]]=rank=0;
                                                                                   24
               j=b[j];
20
                                                                                          forr(i, 1, n)
                                                                                   25
           }
21
                                                                                             tmpr[sa[i]] = r[sa[i-1]] \&\& r[sa[i]+k] = r[sa[i-1]+k] )?
                                                                                   26
       }
22
                                                                                                 rank: ++rank;
23 }
                                                                                          memcpy(r, tmpr, sizeof(r));
                                                                                   27
                                  4.3.
                                        Trie
                                                                                           if(r[sa[n-1]]==n-1) break;
                                                                                   28
                                                                                        }
                                                                                   29
   struct trie{
                                                                                   30
     map<char, trie> m;
                                                                                      void print(){//for debug
2
     void add(const string &s, int p=0){
                                                                                        forn(i, n)
3
                                                                                   32
       if(s[p]) m[s[p]].add(s, p+1);
                                                                                           cout << i << ''' <<
4
     }
                                                                                          s.substr(sa[i], s.find( '$', sa[i])-sa[i]) << endl;}
5
                                                                                   34
     void dfs(){
6
                                                                                                 4.5. String Matching With Suffix Array
       //Do stuff
       forall(it, m)
8
         it->second.dfs();
9
                                                                                    1 //returns (lowerbound, upperbound) of the search
     }
10
                                                                                    2 | ii stringMatching(string P){ //O(sz(P)lgn)
11 };
                                                                                        int lo=0, hi=n-1, mid=lo;
                                                                                        while(lo<hi){</pre>
                   4.4. Suffix Array (largo, nlogn)
                                                                                    4
                                                                                           mid=(lo+hi)/2;
                                                                                    5
```

map<char, trie> next;

4

trie\* tran[256];//transiciones del automata

int idhoja, szhoja;//id de la hoja o 0 si no lo es

```
int res=s.compare(sa[mid], sz(P), P);
6
       if(res>=0) hi=mid;
7
       else lo=mid+1;
8
9
     if(s.compare(sa[lo], sz(P), P)!=0) return ii(-1, -1);
10
     ii ans; ans.fst=lo;
11
     lo=0, hi=n-1, mid;
12
     while(lo<hi){
13
       mid=(lo+hi)/2;
14
       int res=s.compare(sa[mid], sz(P), P);
15
       if(res>0) hi=mid;
16
       else lo=mid+1;
17
     }
18
     if(s.compare(sa[hi], sz(P), P)!=0) hi--;
19
     ans.snd=hi:
     return ans;
21
22 |}
                4.6. LCP (Longest Common Prefix)
   //Calculates the LCP between consecutives suffixes in the Suffix Array.
   //LCP[i] is the length of the LCP between sa[i] and sa[i-1]
  int LCP[MAX_N], phi[MAX_N], PLCP[MAX_N];
   void computeLCP(){\frac{}{0}}
     phi[sa[0]]=-1;
     forr(i, 1, n) phi[sa[i]]=sa[i-1];
     int L=0;
     forn(i, n){
       if(phi[i]==-1) {PLCP[i]=0; continue;}
       while(s[i+L]==s[phi[i]+L]) L++;
10
       PLCP[i]=L;
11
       L=max(L-1, 0);
12
13
     forn(i, n) LCP[i]=PLCP[sa[i]];
15 |}
                              4.7. Corasick
1
  struct trie{
```

```
//link lleva al sufijo mas largo, nxthoja lleva al mas largo pero que
         es hoja
     trie *padre, *link, *nxthoja;
     char pch;//caracter que conecta con padre
     trie(): tran(), idhoja(), padre(), link() {}
     void insert(const string &s, int id=1, int p=0){//id>0!!!
       if(p \le z(s)){
11
         trie &ch=next[s[p]];
12
         tran[(int)s[p]]=&ch;
13
         ch.padre=this, ch.pch=s[p];
         ch.insert(s, id, p+1);
15
16
       else idhoja=id, szhoja=sz(s);
17
     }
18
     trie* get_link() {
19
       if(!link){
20
         if(!padre) link=this;//es la raiz
21
         else if(!padre->padre) link=padre;//hijo de la raiz
         else link=padre->get_link()->get_tran(pch);
23
       return link;
25
26
     trie* get_tran(int c) {
27
       if(!tran[c])
28
         tran[c] = !padre? this : this->get_link()->get_tran(c);
29
       return tran[c];
30
31
     trie *get_nxthoja(){
32
       if(!nxthoja) nxthoja = get_link()->idhoja? link : link->nxthoja;
33
       return nxthoja;
34
35
     void print(int p){
36
       if(idhoia)
37
         cout << "foundu" << idhoja << "uuatupositionu" << p-szhoja << endl
38
       if(get_nxthoja()) get_nxthoja()->print(p);
39
40
     void matching(const string &s, int p=0){
41
       print(p);
42
       if(p<sz(s)) get_tran(s[p])->matching(s, p+1);
```

#### 4.8. Suffix Automaton

```
struct state {
     int len, link;
2
     map<char,int> next;
3
     state() { }
4
   };
5
   const int MAXLEN = 10010;
   state st[MAXLEN*2];
   int sz, last;
   void sa_init() {
     forn(i,sz) st[i].next.clear();
10
     sz = last = 0;
     st[0].len = 0;
12
     st[0].link = -1;
     ++sz;
14
15
   // Es un DAG de una sola fuente y una sola hoja
   // cantidad de endpos = cantidad de apariciones = cantidad de caminos de
        la clase al nodo terminal
  // cantidad de miembros de la clase = st[v].len-st[st[v].link].len (v>0)
        = caminos del inicio a la clase
  // El arbol de los suffix links es el suffix tree de la cadena invertida
       . La string de la arista link(v)->v son los caracteres que difieren
   void sa_extend (char c) {
     int cur = sz++;
21
     st[cur].len = st[last].len + 1;
22
     // en cur agregamos la posicion que estamos extendiendo
23
     //podria agregar tambien un identificador de las cadenas a las cuales
24
         pertenece (si hay varias)
     int p;
25
     for (p=last; p!=-1 && !st[p].next.count(c); p=st[p].link) // modificar
26
          esta linea para hacer separadores unicos entre varias cadenas (c
         =='$')
       st[p].next[c] = cur;
27
     if (p == -1)
28
       st[cur].link = 0;
29
     else {
30
       int q = st[p].next[c];
31
       if (st[p].len + 1 == st[q].len)
32
         st[cur].link = q;
33
       else {
34
         int clone = sz++;
35
         // no le ponemos la posicion actual a clone sino indirectamente
36
             por el link de cur
```

```
st[clone].len = st[p].len + 1;
37
         st[clone].next = st[q].next;
38
         st[clone].link = st[q].link;
39
         for (; p!=-1 && st[p].next.count(c) && st[p].next[c]==q; p=st[p].
40
             link)
           st[p].next[c] = clone;
41
         st[q].link = st[cur].link = clone;
42
43
    }
44
     last = cur;
46 }
```

#### 4.9. Z Function

```
char s[MAXN];
  int z[MAXN]; // z[i] = i==0 ? 0 : max k tq s[0,k) match with s[i,i+k)
   void z_function(char s[],int z[]) {
       int n = strlen(s);
4
       forn(i, n) z[i]=0;
5
       for (int i = 1, l = 0, r = 0; i < n; ++i) {
6
           if (i <= r)
7
               z[i] = min (r - i + 1, z[i - 1]);
8
           while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
9
               ++z[i]:
10
           if (i + z[i] - 1 > r)
11
               1 = i, r = i + z[i] - 1:
12
       }
13
14 }
```

# 5. Geometria

#### 5.1. Punto

```
struct pto{
    double x, y;
    pto(double x=0, double y=0):x(x),y(y){}

pto operator+(pto a){return pto(x+a.x, y+a.y);}

pto operator-(pto a){return pto(x-a.x, y-a.y);}

pto operator+(double a){return pto(x+a, y+a);}

pto operator*(double a){return pto(x*a, y*a);}

pto operator/(double a){return pto(x/a, y/a);}

//dot product, producto interno:

double operator*(pto a){return x*a.x+y*a.y;}
```

```
//module of the cross product or vectorial product:
11
     //if a is less than 180 clockwise from b, a^b>0
12
     double operator^(pto a){return x*a.y-y*a.x;}
13
     //returns true if this is at the left side of line qr
14
     bool left(pto q, pto r){return ((q-*this)^(r-*this))>0;}
     bool operator<(const pto &a) const{return x<a.x-EPS || (abs(x-a.x)<EPS
16
          && y<a.y-EPS);}
   bool operator == (pto a) {return abs(x-a.x) < EPS && abs(y-a.y) < EPS;}
     double norm(){return sqrt(x*x+y*y);}
18
     double norm_sq(){return x*x+v*v;}
19
20
   double dist(pto a, pto b){return (b-a).norm();}
   typedef pto vec;
23
   double angle(pto a, pto o, pto b){
     pto oa=a-o, ob=b-o;
     return atan2(oa^ob, oa*ob);}
26
27
   //rotate p by theta rads CCW w.r.t. origin (0,0)
   pto rotate(pto p, double theta){
     return pto(p.x*cos(theta)-p.y*sin(theta),
30
        p.x*sin(theta)+p.y*cos(theta));
31
32 }
```

### 5.2. Orden radial de puntos

```
struct Cmp{//orden total de puntos alrededor de un punto r
     pto r;
2
     Cmp(pto r):r(r) {}
3
     int cuad(const pto &a) const{
       if(a.x > 0 \&\& a.y >= 0)return 0;
5
       if(a.x <= 0 && a.y > 0)return 1;
6
       if(a.x < 0 && a.y <= 0)return 2;
7
       if(a.x >= 0 \&\& a.y < 0)return 3;
8
       assert(a.x ==0 && a.v==0);
9
       return -1;
10
11
     bool cmp(const pto&p1, const pto&p2)const{
12
       int c1 = cuad(p1), c2 = cuad(p2);
13
       if(c1==c2) return p1.y*p2.x<p1.x*p2.y;</pre>
14
           else return c1 < c2;</pre>
15
     }
16
       bool operator()(const pto&p1, const pto&p2) const{
17
```

```
return cmp(pto(p1.x-r.x,p1.y-r.y),pto(p2.x-r.x,p2.y-r.y));
19
20 };
                                 5.3. Line
int sgn(ll x){return x<0? -1 : !!x;}</pre>
2 struct line{
     line() {}
     double a,b,c;//Ax+By=C
   //pto MUST store float coordinates!
     line(double a, double b, double c):a(a),b(b),c(c){}
     line(pto p, pto q): a(q.y-p.y), b(p.x-q.x), c(a*p.x+b*p.y) {}
     int side(pto p){return sgn(ll(a) * p.x + ll(b) * p.y - c);}
9
   bool parallels(line 11, line 12){return abs(11.a*12.b-12.a*11.b)<EPS;}
   pto inter(line 11, line 12){//intersection
     double det=11.a*12.b-12.a*11.b:
     if(abs(det) < EPS) return pto(INF, INF); //parallels</pre>
13
     return pto(12.b*11.c-11.b*12.c, 11.a*12.c-12.a*11.c)/det;
14
15 }
                              5.4. Segment
1 struct segm{
     pto s,f;
     segm(pto s, pto f):s(s), f(f) {}
     pto closest(pto p) {//use for dist to point
        double 12 = dist_sq(s, f);
5
        if(12==0.) return s;
6
        double t = ((p-s)*(f-s))/12;
7
        if (t<0.) return s://not write if is a line
        else if(t>1.)return f;//not write if is a line
9
        return s+((f-s)*t);
10
     }
11
       bool inside(pto p){return abs(dist(s, p)+dist(p, f)-dist(s, f))<EPS</pre>
12
           ;}
   };
13
14
   pto inter(segm s1, segm s2){
     pto r=inter(line(s1.s, s1.f), line(s2.s, s2.f));
       if(s1.inside(r) && s2.inside(r)) return r;
17
     return pto(INF, INF);
18
19 }
```

# 5.5. Rectangle

```
struct rect{
//lower-left and upper-right corners
pto lw, up;
};

//returns if there's an intersection and stores it in r
bool inter(rect a, rect b, rect &r){
r.lw=pto(max(a.lw.x, b.lw.x), max(a.lw.y, b.lw.y));
r.up=pto(min(a.up.x, b.up.x), min(a.up.y, b.up.y));
//check case when only a edge is common
return r.lw.x<r.up.x && r.lw.y<r.up.y;
}</pre>
```

# 5.6. Polygon Area

```
double area(vector<pto> &p){//O(sz(p))}
double area=0;
forn(i, sz(p)) area+=p[i]^p[(i+1)%sz(p)];
//if points are in clockwise order then area is negative
return abs(area)/2;
}
//Area ellipse = M_PI*a*b where a and b are the semi axis lengths
//Area triangle = sqrt(s*(s-a)(s-b)(s-c)) where s=(a+b+c)/2
```

#### 5.7. Circle

```
|vec perp(vec v){return vec(-v.y, v.x);}
  line bisector(pto x, pto y){
    line l=line(x, y); pto m=(x+y)/2;
     return line(-1.b, 1.a, -1.b*m.x+1.a*m.y);
4
5
   struct Circle{
6
     pto o;
     double r;
8
     Circle(pto x, pto y, pto z){
9
       o=inter(bisector(x, y), bisector(y, z));
10
       r=dist(o, x);
11
12
     pair<pto, pto> ptosTang(pto p){
13
       pto m=(p+o)/2;
14
       tipo d=dist(o, m);
15
       tipo a=r*r/(2*d);
16
       tipo h=sqrt(r*r-a*a);
17
```

```
pto m2=o+(m-o)*a/d;
        vec per=perp(m-o)/d;
19
        return make_pair(m2-per*h, m2+per*h);
20
21
   };
22
   //finds the center of the circle containing p1 and p2 with radius r
   //as there may be two solutions swap p1, p2 to get the other
   bool circle2PtsRad(pto p1, pto p2, double r, pto &c){
            double d2=(p1-p2).norm_sq(), det=r*r/d2-0.25;
26
            if(det<0) return false;</pre>
            c=(p1+p2)/2+perp(p2-p1)*sqrt(det);
28
            return true;
29
30
   #define sqr(a) ((a)*(a))
   #define feq(a,b) (fabs((a)-(b))<EPS)</pre>
   pair<tipo, tipo > ecCuad(tipo a, tipo b, tipo c){//a*x*x+b*x+c=0
     tipo dx = sqrt(b*b-4.0*a*c);
     return make_pair((-b + dx)/(2.0*a), (-b - dx)/(2.0*a));
36
   pair<pto, pto> interCL(Circle c, line 1){
     bool sw=false;
38
     if((sw=feq(0,1.b))){}
39
     swap(1.a, 1.b);
40
     swap(c.o.x, c.o.y);
41
42
     pair<tipo, tipo> rc = ecCuad(
43
     sqr(l.a)+sqr(l.b),
44
     2.0*1.a*1.b*c.o.y-2.0*(sqr(1.b)*c.o.x+1.c*1.a)
45
     sqr(1.b)*(sqr(c.o.x)+sqr(c.o.y)-sqr(c.r))+sqr(1.c)-2.0*1.c*1.b*c.o.y
46
     ):
47
     pair<pto, pto> p( pto(rc.first, (l.c - l.a * rc.first) / l.b),
48
                pto(rc.second, (1.c - 1.a * rc.second) / 1.b) );
49
     if(sw){
50
     swap(p.first.x, p.first.y);
     swap(p.second.x, p.second.y);
52
53
     return p;
54
55
   pair<pto, pto> interCC(Circle c1, Circle c2){
     line 1;
     1.a = c1.o.x-c2.o.x;
     1.b = c1.o.y-c2.o.y;
     1.c = (\operatorname{sqr}(c2.r) - \operatorname{sqr}(c1.r) + \operatorname{sqr}(c1.o.x) - \operatorname{sqr}(c2.o.x) + \operatorname{sqr}(c1.o.y)
```

```
-sqr(c2.o.y))/2.0;
61
     return interCL(c1, 1);
62
63 }
                           5.8. Point in Poly
  //checks if v is inside of P, using ray casting
   //works with convex and concave.
   //excludes boundaries, handle it separately using segment.inside()
   bool inPolygon(pto v, vector<pto>& P) {
     bool c = false;
5
     forn(i, sz(P)){
6
       int j=(i+1) %z(P);
       if((P[j].v>v.v) != (P[i].v > v.v) &&
8
     (v.x < (P[i].x - P[j].x) * (v.y-P[j].y) / (P[i].y - P[j].y) + P[j].x))
9
         c = !c:
10
     }
11
    return c;
12
13 }
                  5.9. Point in Convex Poly log(n)
   void normalize(vector<pto> &pt){//delete collinear points first!
     //this makes it clockwise:
2
       if(pt[2].left(pt[0], pt[1])) reverse(pt.begin(), pt.end());
3
     int n=sz(pt), pi=0;
4
     forn(i, n)
5
       if(pt[i].x<pt[pi].x || (pt[i].x==pt[pi].x && pt[i].y<pt[pi].y))</pre>
6
         pi=i;
7
     vector<pto> shift(n);//puts pi as first point
8
       forn(i, n) shift[i]=pt[(pi+i) %n];
9
       pt.swap(shift);
10
11
   bool inPolygon(pto p, const vector<pto> &pt){
12
     //call normalize first!
13
     if(p.left(pt[0], pt[1]) || p.left(pt[sz(pt)-1], pt[0])) return false;
14
     int a=1, b=sz(pt)-1;
15
     while(b-a>1){
16
       int c=(a+b)/2:
17
       if(!p.left(pt[0], pt[c])) a=c;
18
```

else b=c;

return !p.left(pt[a], pt[a+1]);

19

20

21

22 }

}

#### 5.10. Convex Check CHECK

```
|bool isConvex(vector<int> &p){//O(N), delete collinear points!
     int N=sz(p);
     if(N<3) return false;
     bool isLeft=p[0].left(p[1], p[2]);
     forr(i, 1, N)
       if(p[i].left(p[(i+1) \mathbb{N}], p[(i+2) \mathbb{N}])!=isLeft)
6
         return false:
7
     return true; }
                           5.11. Convex Hull
1 //stores convex hull of P in S, CCW order
   //left must return >=0 to delete collinear points!
   void CH(vector<pto>& P, vector<pto> &S){
     S.clear():
     sort(P.begin(), P.end());//first x, then y
     forn(i, sz(P)){//lower hull
       while(sz(S) \ge 2 \&\& S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back();
       S.pb(P[i]);
8
9
     S.pop_back();
10
     int k=sz(S);
11
     dforn(i, sz(P)){//upper hull
       while(sz(S) \ge k+2 \&\& S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back
13
           ();
       S.pb(P[i]);
14
15
     S.pop_back();
16
                           5.12. Cut Polygon
1 //cuts polygon Q along the line ab
2 //stores the left side (swap a, b for the right one) in P
   void cutPolygon(pto a, pto b, vector<pto> Q, vector<pto> &P){
     P.clear();
     forn(i, sz(Q)){
       double left1=(b-a)^(Q[i]-a), left2=(b-a)^(Q[(i+1) \sz(Q)]-a);
       if(left1>=0) P.pb(Q[i]);
7
```

P.pb(inter(line(Q[i], Q[(i+1) %z(Q)]), line(a, b)));

if(left1\*left2<0)

9

10

}

```
11 }
```

#### 5.13. Bresenham

```
//plot a line approximation in a 2d map
   void bresenham(pto a, pto b){
     pto d=b-a; d.x=abs(d.x), d.y=abs(d.y);
    pto s(a.x<b.x? 1: -1, a.y<b.y? 1: -1);
     int err=d.x-d.y;
     while(1){
6
       m[a.x][a.y]=1;//plot
7
      if(a==b) break:
8
      int e2=err:
9
      if(e2 \ge 0) err=2*d.y, a.x+=s.x;
      if(e2 <= 0) err+= 2*d.x, a.y+= s.y;
    }
12
13 }
```

#### 5.14. Rotate Matrix

```
//rotates matrix t 90 degrees clockwise
//using auxiliary matrix t2(faster)
void rotate(){
forn(x, n) forn(y, n)
t2[n-y-1][x]=t[x][y];
memcpy(t, t2, sizeof(t));
}
```

# 5.15. Interseccion de Circulos en n3log(n)

```
struct event {
       double x; int t;
2
       event(double xx, int tt) : x(xx), t(tt) {}
3
       bool operator <(const event &o) const { return x < o.x; }</pre>
   };
5
   typedef vector<Circle> VC;
   typedef vector<event> VE;
8
   double cuenta(VE &v, double A,double B) {
       sort(v.begin(), v.end());
10
      double res = 0.0, lx = ((v.empty())?0.0:v[0].x);
11
       int contador = 0;
12
       forn(i,sz(v)) {
13
```

```
// interseccion de todos (contador == n), union de todos (
14
               contador > 0),
           // conjunto de puntos cubierto por exacta k Circulos (contador
15
           if (contador == n) res += v[i].x - lx;
           contador += v[i].t;
           lx = v[i].x;
18
       }
19
       return res;
21
   // Primitiva de sqrt(r*r - x*x) como funcion double de una variable x.
   inline double primitiva(double x,double r) {
       if (x \ge r) return r*r*M PI/4.0:
       if (x \le -r) return -r*r*M PI/4.0:
       double raiz = sqrt(r*r-x*x);
       return 0.5 * (x * raiz + r*r*atan(x/raiz));
27
   }
28
   double interCircle(VC &v) {
       vector<double> p; p.reserve(v.size() * (v.size() + 2));
30
       forn(i,sz(v)) {
           p.push_back(v[i].c.x + v[i].r);
32
           p.push_back(v[i].c.x - v[i].r);
33
34
       forn(i,sz(v)) forn(j,i) {
           Circle &a = v[i], b = v[j];
36
           double d = (a.c - b.c).norm();
37
           if (fabs(a.r - b.r) < d \&\& d < a.r + b.r) {
38
               double alfa = acos((sqr(a.r) + sqr(d) - sqr(b.r)) / (2.0 * d)
39
                    * a.r)):
               pto vec = (b.c - a.c) * (a.r / d);
               p.pb((a.c + rotate(vec, alfa)).x);
41
               p.pb((a.c + rotate(vec, -alfa)).x);
43
44
       sort(p.begin(), p.end());
45
       double res = 0.0;
46
       forn(i,sz(p)-1) {
47
           const double A = p[i], B = p[i+1];
           VE ve; ve.reserve(2 * v.size());
49
           forn(j,sz(v)) {
               const Circle &c = v[j];
51
               double arco = primitiva(B-c.c.x,c.r) - primitiva(A-c.c.x,c.r
52
                   );
```

# 6. Math

### 6.1. Identidades

$$\sum_{i=0}^{n} \binom{n}{i} = 2^{n}$$

$$\sum_{i=0}^{n} i \binom{n}{i} = n * 2^{n-1}$$

$$\sum_{i=m}^{n} i = \frac{n(n+1)}{2} - \frac{m(m-1)}{2} = \frac{(n+1-m)(n+m)}{2}$$

$$\sum_{i=0}^{n} i = \sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i=0}^{n} i^{2} = \frac{n(n+1)(2n+1)}{2} = \frac{n^{3}}{3} + \frac{n^{2}}{2} + \frac{n}{6}$$

$$\sum_{i=0}^{n} i(i-1) = \frac{8}{6} (\frac{n}{2})(\frac{n}{2}+1)(n+1) \text{ (doubles)} \rightarrow \text{Sino ver caso impar y par}$$

$$\sum_{i=0}^{n} i^{3} = \left(\frac{n(n+1)}{2}\right)^{2} = \frac{n^{4}}{4} + \frac{n^{3}}{2} + \frac{n^{2}}{4} = \left[\sum_{i=1}^{n} i\right]^{2}$$

$$\sum_{i=0}^{n} i^{4} = \frac{n(n+1)(2n+1)(3n^{2}+3n-1)}{30} = \frac{n^{5}}{5} + \frac{n^{4}}{2} + \frac{n^{3}}{3} - \frac{n}{30}$$

$$\sum_{i=0}^{n} i^{p} = \frac{(n+1)^{p+1}}{p+1} + \sum_{k=1}^{p} \frac{B_{k}}{p-k+1}} \binom{p}{k} (n+1)^{p-k+1}$$

$$r = e - v + k + 1$$

Teorema de Pick: (Area, puntos interiores y puntos en el borde)  $A = I + \tfrac{B}{2} - 1$ 

#### 6.2. Ec. Caracteristica

```
\begin{aligned} a_0T(n) + a_1T(n-1) + \ldots + a_kT(n-k) &= 0 \\ p(x) = a_0x^k + a_1x^{k-1} + \ldots + a_k \end{aligned} Sean r_1, r_2, \ldots, r_q las raíces distintas, de mult. m_1, m_2, \ldots, m_q T(n) = \sum_{i=1}^q \sum_{j=0}^{m_i-1} c_{ij} n^j r_i^n Las constantes c_{ij} se determinan por los casos base.
```

#### 6.3. Combinatorio

```
forn(i, MAXN+1){//comb[i][k]=i tomados de a k
comb[i][0]=comb[i][i]=1;
forr(k, 1, i) comb[i][k]=(comb[i-1][k]+comb[i-1][k-1]) MOD;
}
```

```
6 | 11 lucas (11 n, 11 k, int p){ //Calcula (n,k) %p teniendo comb[p][p]
       precalculado.
     11 \text{ aux} = 1;
     while (n + k){
       aux = (aux * comb[n\%][k\%]) \%;
       n/=p, k/=p;
11
     return aux;
12
13 }
                     6.4. Exp. de Numeros Mod.
1 | 11 expmod (11 b, 11 e, 11 m)\{\frac{1}{0}
     if(!e) return 1;
     11 q= expmod(b,e/2,m); q=(q*q) m;
     return e \%2? (b * q) \%n : q;
5 | }
           6.5. Exp. de Matrices y Fibonacci en log(n)
 1 #define SIZE 350
  int NN:
   void mul(double a[SIZE][SIZE], double b[SIZE][SIZE])
4
       double res[SIZE] [SIZE] = {{0}}:
       forn(i, NN) forn(j, NN) forn(k, NN) res[i][j]+=a[i][k]*b[k][j];
6
       forn(i, NN) forn(j, NN) a[i][j]=res[i][j];
8
   void powmat(double a[SIZE][SIZE], int n, double res[SIZE][SIZE])
10
       forn(i, NN) forn(j, NN) res[i][j]=(i==j);
11
       while(n){
12
           if(n&1) mul(res, a), n--;
13
           else mul(a, a), n/=2;
14
       }
15
   }
16
17
   struct M22{
                    // |a b|
     tipo a,b,c,d;// |c d|
19
     M22 operator*(const M22 &p) const {
20
       return (M22){a*p.a+b*p.c, a*p.b+b*p.d, c*p.a+d*p.c,c*p.b+d*p.d};}
21
   };
22
23
   M22 operator (const M22 &p, int n){
```

if(!n) return (M22){1, 0, 0, 1};//identidad

```
M22 q=p^(n/2); q=q*q;
25
     return n%2? p * q : q;}
26
27
   11 fibo(ll n){//calcula el fibonacci enesimo
28
     M22 \text{ mat}=(M22)\{0, 1, 1, 1\}^n;
29
     return mat.a*f0+mat.b*f1;//f0 y f1 son los valores iniciales
30
31 |}
                 6.6. Matrices y determinante O(n^3)
  struct Mat {
       vector<vector<double> > vec;
2
       Mat(int n): vec(n, vector<double>(n) ) {}
3
       Mat(int n, int m): vec(n, vector<double>(m) ) {}
4
       vector<double> &operator[](int f){return vec[f];}
5
       const vector<double> &operator[](int f) const {return vec[f];}
6
       int size() const {return sz(vec);}
7
       Mat operator+(Mat &b) { ///this de n x m entonces b de n x m
8
           Mat m(sz(b), sz(b[0]));
9
           forn(i,sz(vec)) forn(j,sz(vec[0])) m[i][j] = vec[i][j] + b[i][j
10
               ];
           return m;
11
       }
12
       Mat operator*(const Mat &b) { ///this de n x m entonces b de m x t
13
           int n = sz(vec), m = sz(vec[0]), t = sz(b[0]);
14
           Mat mat(n,t);
15
           forn(i,n) forn(j,t) {
16
               forn(k,m)
17
                    mat[i][j] += vec[i][k] * b[k][j];
18
           }
19
           return mat;
20
^{21}
       double determinant(){//sacado de e maxx ru
^{22}
           double det = 1;
23
           int n = sz(vec);
24
           Mat m(*this);
25
           forn(i, n){//para cada columna
26
               int k = i:
27
               forr(j, i+1, n)//busco la fila con mayor val abs
28
                    if(abs(m[j][i])>abs(m[k][i]))
29
                        k = j;
30
               if(abs(m[k][i])<1e-9) return 0;
31
               m[i].swap(m[k]);//la swapeo
32
```

```
if(i!=k) det = -det;
33
                det *= m[i][i];
34
                forr(j, i+1, n) m[i][j] /= m[i][i];
35
                //hago 0 todas las otras filas
36
                forn(j, n) if (j!= i && abs(m[j][i])>1e-9)
37
                    forr(k, i+1, n) m[j][k]-=m[i][k]*m[j][i];
38
39
            return det;
40
41
       Mat identidad(int n) {
            Mat m(n);
43
            forn(i,n) m[i][i] = 1;
44
            return m;
45
46
       Mat transpuesta() {
47
            Mat m(sz(vec[0]),sz(vec));
48
            forn(i,sz(vec[0])) forn(j,sz(vec))
49
                m[i][i] = vec[i][i];
50
            return m:
51
       }
52
       void print() {
53
            int n = sz(vec), m = sz(vec[0]);
54
            cout << "********* << endl;
55
            forn(i,n){
56
                forn(j,m) cout << "_"+!j << vec[i][j];
57
                cout << endl;</pre>
58
            }
59
60
61 };
```

#### 6.7. Teorema Chino del Resto

$$y = \sum_{j=1}^{n} (x_j * (\prod_{i=1, i \neq j}^{n} m_i)_{m_j}^{-1} * \prod_{i=1, i \neq j}^{n} m_i)$$

### 6.8. Criba

```
#define MAXP 100000 //no necesariamente primo
int criba[MAXP+1];
void crearcriba(){
int w[] = {4,2,4,2,4,6,2,6};
```

```
for(int p=25;p<=MAXP;p+=10) criba[p]=5;</pre>
5
     for(int p=9;p<=MAXP;p+=6) criba[p]=3;</pre>
6
     for(int p=4;p<=MAXP;p+=2) criba[p]=2;</pre>
     for(int p=7,cur=0;p*p<=MAXP;p+=w[cur++&7]) if (!criba[p])</pre>
       for(int j=p*p;j<=MAXP;j+=(p<<1)) if(!criba[j]) criba[j]=p;</pre>
9
10
11
   vector<int> primos;
   void buscarprimos(){
13
     crearcriba();
14
     forr (i,2,MAXP+1) if (!criba[i]) primos.push_back(i);
16
   //~ Useful for bit trick:
   //~ #define SET(i) ( criba[(i)>>5] |=1<<((i)&31) )
   //~ #define INDEX(i) ( (criba[i>>5]>>((i)&31))&1 )
21 //~ unsigned int criba[MAXP/32+1];
```

# 6.9. Funciones de primos

Sea  $n = \prod p_i^{k_i}$ , fact(n) genera un map donde a cada  $p_i$  le asocia su  $k_i$ 

```
//factoriza bien numeros hasta MAXP^2
   map<ll,ll> fact(ll n){ //O (cant primos)
     map<ll,ll> ret;
     forall(p, primos){
       while(!(n %*p)){
         ret[*p]++;//divisor found
6
         n/=*p;
7
8
9
     if(n>1) ret[n]++;
     return ret;
11
12
13
    //factoriza bien numeros hasta MAXP
   map<11,11> fact2(11 n){ //0 (1g n)}
15
     map<11,11> ret;
16
     while (criba[n]){
17
       ret[criba[n]]++;
18
       n/=criba[n];
19
20
     if(n>1) ret[n]++;
21
     return ret;
22
```

```
23 }
24
   //Usar asi: divisores(fac, divs, fac.begin()); NO ESTA ORDENADO
   void divisores(const map<11,11> &f, vector<11> &divs, map<11,11>::
       iterator it, ll n=1){
       if(it==f.begin()) divs.clear();
       if(it==f.end()) {
           divs.pb(n);
           return;
       ll p=it->fst, k=it->snd; ++it;
32
       forn(_, k+1) divisores(f, divs, it, n), n*=p;
33
   }
34
   11 numPrimeFactors (11 n){
     ll rta = 0;
     map<11,11> f=fact(n);
     forall(it, f) rta += it->second;
     return rta:
40
41
42
   11 numDiffPrimeFactors (11 n){
     ll rta = 0;
44
     map<11,11> f=fact(n);
     forall(it, f) rta += 1;
46
     return rta;
47
48
   11 sumPrimeFactors (11 n){
     11 \text{ rta} = 0;
51
     map<11,11> f=fact(n);
     forall(it, f) rta += it->first;
     return rta:
54
55
   11 numDiv (11 n){
    ll rta = 1;
     map<11,11> f=fact(n);
     forall(it, f) rta *= (it->second + 1);
     return rta;
61
62
64 | 11 sumDiv (11 n) {
```

b /= 2;

return x % c;

8

```
11 }
     ll rta = 1;
65
     map<ll,ll> f=fact(n);
                                                                                      12
66
     forall(it, f) {
67
     11 \text{ pot} = 1, \text{ aux} = 0;
68
     forn(i, it->snd+1) aux += pot, pot *= it->fst;
                                                                                      15
     rta*=aux;
70
                                                                                      16
     }
                                                                                         }
                                                                                      17
71
     return rta;
72
                                                                                      18
73
74
                                                                                      20
   ll eulerPhi (ll n){ // con criba: O(lg n)
75
                                                                                      21
     11 \text{ rta} = n;
76
     map<ll,ll> f=fact(n);
     forall(it, f) rta -= rta / it->first;
                                                                                      24
     return rta;
                                                                                      25
79
                                                                                      26
80
81
   11 eulerPhi2 (11 n){ // 0 (sqrt n)
     11 r = n:
83
     forr (i,2,n+1){
84
       if ((11)i*i > n) break;
85
       if (n \% i == 0){
86
                                                                                      32
         while (n\% == 0) n/=i;
                                                                                      33
87
         r = r/i;
                                                                                      34
88
       }
                                                                                      35
89
90
     if (n != 1) r= r/n;
91
     return r;
92
93 }
                                                                                           forn (j,9)
                                                                                      39
                                                                                      41
                    6.10. Phollard's Rho (rolando)
                                                                                      42
                                                                                      43
   | 11 gcd(11 a, 11 b){return a?gcd(b %a, a):b;}
                                                                                      44
2
   11 mulmod (11 a, 11 b, 11 c) { //returns (a*b) %c, and minimize overfloor
    11 x = 0, y = a\%;
4
                                                                                      48
     while (b > 0){
      if (b \% 2 == 1) x = (x+y) \% c;
       y = (y*2) \% c;
```

```
ll expmod (ll b, ll e, ll m){\frac{1}{0}} \log b
     if(!e) return 1;
    11 q= expmod(b,e/2,m); q=mulmod(q,q,m);
    return e %2? mulmod(b,q,m) : q;
   bool es_primo_prob (ll n, int a)
     if (n == a) return true;
    11 s = 0, d = n-1;
     while (d \% 2 == 0) s++, d/=2:
    ll x = expmod(a,d,n);
     if ((x == 1) \mid | (x+1 == n)) return true;
    forn (i, s-1){
      x = mulmod(x, x, n):
      if (x == 1) return false;
       if (x+1 == n) return true;
     return false;
   bool rabin (ll n){ //devuelve true si n es primo
     if (n == 1) return false;
     const int ar[] = \{2,3,5,7,11,13,17,19,23\};
       if (!es_primo_prob(n,ar[j]))
         return false;
     return true;
   ll rho(ll n){
       if( (n & 1) == 0 ) return 2;
       11 x = 2 , y = 2 , d = 1;
       ll c = rand() % n + 1:
       while (d == 1)
           x = (mulmod(x, x, n) + c) n;
           y = (mulmod(y, y, n) + c) n;
           y = (mulmod( y , y , n ) + c) %n;
           if(x - y \ge 0) d = gcd(x - y, n);
53
```

```
else d = gcd(y - x, n);
54
       }
55
       return d==n? rho(n):d;
56
57
58
   map<ll,ll> prim;
   void factRho (ll n){ //O (lg n)^3. un solo numero
     if (n == 1) return;
     if (rabin(n)){
62
       prim[n]++;
63
       return;
64
65
     11 factor = rho(n):
    factRho(factor):
67
     factRho(n/factor);
69 }
                               6.11. GCD
1 | tipo gcd(tipo a, tipo b){return a?gcd(b %a, a):b;}
                        6.12. Extended Euclid
  void extendedEuclid (ll a, ll b) \{ //a * x + b * y = d \}
     if (!b) { x = 1; y = 0; d = a; return;}
2
     extendedEuclid (b, a%);
3
    11 x1 = y;
4
    11 v1 = x - (a/b) * v;
     x = x1; y = y1;
6
  |}
7
                               6.13. LCM
1 | tipo lcm(tipo a, tipo b){return a / gcd(a,b) * b;}
                             6.14. Inversos
  #define MAXMOD 15485867
  ll inv[MAXMOD];//inv[i]*i=1 mod MOD
   void calc(int p){\frac{1}{0}}
     inv[1]=1:
     forr(i, 2, p) inv[i] = p-((p/i)*inv[p\%i])\%;
5
6
```

int inverso(int x) ${\frac{1}{0}} \log x$ 

```
return expmod(x, eulerphi(MOD)-2);//si mod no es primo(sacar a mano)
     return expmod(x, MOD-2);//si mod es primo
9
10 }
                             6.15. Simpson
   double integral(double a, double b, int n=10000) {//O(n), n=cantdiv
     double area=0, h=(b-a)/n, fa=f(a), fb;
     forn(i, n){
       fb=f(a+h*(i+1)):
       area+=fa+ 4*f(a+h*(i+0.5)) +fb, fa=fb;
5
    }
6
     return area*h/6.;}
                             6.16. Fraction
 tipo mcd(tipo a, tipo b){return a?mcd(b%, a):b;}
   struct frac{
     tipo p,q;
     frac(tipo p=0, tipo q=1):p(p),q(q) {norm();}
     void norm(){
       tipo a = mcd(p,q);
       if(a) p/=a, q/=a;
       else q=1;
       if (q<0) q=-q, p=-p;}
     frac operator+(const frac& o){
       tipo a = mcd(q, o.q);
11
       return frac(p*(o.q/a)+o.p*(q/a), q*(o.q/a));}
     frac operator-(const frac& o){
13
       tipo a = mcd(q, o.q);
       return frac(p*(o.q/a)-o.p*(q/a), q*(o.q/a));}
15
     frac operator*(frac o){
       tipo a = mcd(q,o.p), b = mcd(o.q,p);
       return frac((p/b)*(o.p/a), (q/a)*(o.q/b));}
18
     frac operator/(frac o){
19
       tipo a = mcd(q,o.q), b = mcd(o.p,p);
       return frac((p/b)*(o.q/a),(q/a)*(o.p/b));}
     bool operator<(const frac &o) const{return p*o.q < o.p*q;}</pre>
     bool operator==(frac o){return p==o.p&kq==o.q;}
23
24 };
                            6.17. Polinomio
struct poly {
```

```
vector<tipo> c;//guarda los coeficientes del polinomio
2
       poly(const vector<tipo> &c): c(c) {}
3
       polv() {}
4
     int gr(){//calculates grade of the polynomial
5
       return sz(c); }
     bool isnull() {return c.empty();}
       poly operator+(const poly &o) const {
8
           int m = sz(c), n = sz(o.c);
9
           vector<tipo> res(max(m,n));
10
           forn(i, m) res[i] += c[i];
11
           forn(i, n) res[i] += o.c[i];
12
           return poly(res);
13
       }
14
       poly operator*(const tipo cons) const {
15
       vector<tipo> res(sz(c));
16
           forn(i, sz(c)) res[i]=c[i]*cons;
17
           return poly(res);
18
       }
19
       poly operator*(const poly &o) const {
20
           int m = sz(c), n = sz(o.c);
21
           vector<tipo> res(m+n-1);
22
           forn(i, m) forn(j, n) res[i+j]+=c[i]*o.c[j];
23
           return poly(res);
24
       }
25
     tipo eval(tipo v) {
26
       tipo sum = 0;
27
       dforn(i, sz(c)) sum=sum*v + c[i];
28
       return sum:
29
     }
30
       //poly contains only a vector<int> c (the coeficients)
31
     //the following function generates the roots of the polynomial
32
    //it can be easily modified to return float roots
33
     set<tipo> roots(){
34
       set<tipo> roots;
35
       tipo a0 = abs(c[0]), an = abs(c[sz(c)-1]);
36
       vector<tipo> ps,qs;
37
       forr(p,1,sqrt(a0)+1) if (a0 \% ==0) ps.pb(p),ps.pb(a0/p);
38
       forr(q,1,sqrt(an)+1) if (an)(q=0) qs.pb(q),qs.pb(an/q);
39
       forall(pt,ps)
40
         forall(qt,qs) if ( (*pt) % (*qt)==0 ) {
41
           tipo root = abs((*pt) / (*qt));
^{42}
           if (eval(root)==0) roots.insert(root);
43
         }
44
```

```
return roots;
45
    }
46
   };
47
   pair<poly,tipo> ruffini(const poly p, tipo r) {
     int n = sz(p.c) - 1;
49
     vector<tipo> b(n);
    b[n-1] = p.c[n];
     dforn(k,n-1) b[k] = p.c[k+1] + r*b[k+1];
     tipo resto = p.c[0] + r*b[0];
     polv result(b);
     return make_pair<poly,tipo>(result,resto);
55
56
57
   poly interpolate(const vector<tipo>& x,const vector<tipo>& y) {
       poly A; A.c.pb(1);
59
       forn(i,sz(x)) { poly aux; aux.c.pb(-x[i]); aux.c.pb(1);
60
       A = A * aux:
    }
62
     poly S; S.c.pb(0);
63
    forn(i,sz(x)) { poly Li;
      Li = ruffini(A,x[i]).fst;
       Li = Li * (1.0 / Li.eval(x[i])); // here put a multiple of the
           coefficients instead of 1.0 to avoid using double
       S = S + Li * v[i];
67
    }
68
     return S;
70 }
```

#### 6.18. Ec. Lineales

```
bool resolver_ev(Mat a, Vec v, Vec &x, Mat &ev){
     int n = a.size(), m = n?a[0].size():0, rw = min(n, m);
     vector<int> p; forn(i,m) p.push_back(i);
3
     forn(i, rw) {
4
      int uc=i, uf=i;
5
      forr(f, i, n) forr(c, i, m) if(fabs(a[f][c])>fabs(a[uf][uc])) {uf=f;
6
       if (feq(a[uf][uc], 0)) { rw = i; break; }
7
       forn(j, n) swap(a[j][i], a[j][uc]);
8
       swap(a[i], a[uf]); swap(y[i], y[uf]); swap(p[i], p[uc]);
9
       tipo inv = 1 / a[i][i]; //aca divide
10
       forr(j, i+1, n) {
11
         tipo v = a[j][i] * inv;
12
```

vector<int> rev:

vector<base> wlen\_pw;

inline static void fft(base a[], int n, bool invert) {

```
forr(k, i, m) a[j][k]-=v * a[i][k];
                                                                                         forn(i, n) if(i<rev[i]) swap(a[i], a[rev[i]]);</pre>
13
                                                                                  18
         y[j] -= v*y[i];
                                                                                       for (int len=2; len<=n; len<<=1) {
                                                                                  19
14
                                                                                         double ang = 2*M_PI/len * (invert?-1:+1);
15
                                                                                  20
     } // rw = rango(a), aca la matriz esta triangulada
                                                                                         int len2 = len >> 1;
                                                                                  21
16
     forr(i, rw, n) if (!feg(y[i],0)) return false; // checkeo de
                                                                                         base wlen (cos(ang), sin(ang));
                                                                                  22
         compatibilidad
                                                                                         wlen_pw[0] = base(1, 0);
                                                                                  23
                                                                                             forr(i, 1, len2) wlen_pw[i] = wlen_pw[i-1] * wlen;
     x = vector < tipo > (m, 0);
                                                                                  24
18
     dforn(i, rw){
                                                                                         for (int i=0; i<n; i+=len) {</pre>
                                                                                  25
19
       tipo s = v[i];
                                                                                           base t,
20
       forr(j, i+1, rw) s -= a[i][j]*x[p[j]];
                                                                                             *pu = a+i,
21
       x[p[i]] = s / a[i][i]; //aca divide
                                                                                             *pv = a+i+len2,
22
                                                                                  28
                                                                                             *pu_end = a+i+len2,
23
                                                                                  29
     ev = Mat(m-rw, Vec(m, 0)); // Esta parte va SOLO si se necesita el ev
                                                                                             *pw = &wlen_pw[0];
     forn(k, m-rw) {
                                                                                           for (; pu!=pu_end; ++pu, ++pv, ++pw) {
25
       ev[k][p[k+rw]] = 1;
                                                                                             t = *pv * *pw;
                                                                                  32
26
       dforn(i, rw){
                                                                                             *pv = *pu - t;
27
         tipo s = -a[i][k+rw];
                                                                                             *pu = *pu + t;
28
         forr(j, i+1, rw) s -= a[i][j]*ev[k][p[j]];
29
         ev[k][p[i]] = s / a[i][i]; //aca divide
                                                                                         }
                                                                                  36
30
31
                                                                                       if (invert) forn(i, n) a[i]/= n;
     }
                                                                                  38
32
     return true;
                                                                                  39
34 }
                                                                                     inline static void calc_rev(int n){//precalculo: llamar antes de fft!!
                                                                                         wlen_pw.resize(n);
                                6.19. FFT
                                                                                         rev.resize(n);
                                                                                  42
                                                                                         int lg=31-__builtin_clz(n);
                                                                                         forn(i, n){
1 //~ typedef complex<double> base; //menos codigo, pero mas lento
                                                                                  44
                                                                                         rev[i] = 0;
  //elegir si usar complejos de c (lento) o estos
                                                                                             forn(k, lg) if(i&(1<<k))
  struct base{
                                                                                                  rev[i] |=1<<(lg-1-k);
       double r,i;
                                                                                  47
4
                                                                                         }
       base(double r=0, double i=0):r(r), i(i){}
                                                                                  48
5
       double real()const{return r;}
6
                                                                                     inline static void multiply(const vector<int> &a, const vector<int> &b,
       void operator/=(const int c){r/=c, i/=c;}
7
                                                                                         vector<int> &res) {
8
                                                                                       vector<base> fa (a.begin(), a.end()), fb (b.begin(), b.end());
   base operator*(const base &a, const base &b){
       return base(a.r*b.r-a.i*b.i, a.r*b.i+a.i*b.r);}
10
                                                                                       while(n < max(sz(a), sz(b))) n <<= 1;
   base operator+(const base &a, const base &b){
11
                                                                                       n <<= 1:
       return base(a.r+b.r, a.i+b.i);}
12
                                                                                        calc_rev(n);
   base operator-(const base &a, const base &b){
                                                                                       fa.resize (n), fb.resize (n);
       return base(a.r-b.r, a.i-b.i);}
```

58

fft (&fa[0], n, false), fft (&fb[0], n, false);

forn(i, n) fa[i] = fa[i] \* fb[i];

fft (&fa[0], n, true);

```
res.resize(n);
forn(i, n) res[i] = int (fa[i].real() + 0.5);

void toPoly(const string &s, vector<int> &P){//convierte un numero a polinomio
    P.clear();
    dforn(i, sz(s)) P.pb(s[i]-'0');
}
```

# 6.20. Tablas y cotas (Primos, Divisores, Factoriales, etc)

Factoriales					
0! = 1	11! = 39.916.800				
1! = 1	$12! = 479.001.600 \ (\in int)$				
2! = 2	13! = 6.227.020.800				
3! = 6	14! = 87.178.291.200				
4! = 24	15! = 1.307.674.368.000				
5! = 120	16! = 20.922.789.888.000				
6! = 720	17! = 355.687.428.096.000				
7! = 5.040	18! = 6.402.373.705.728.000				
8! = 40.320	19! = 121.645.100.408.832.000				
9! = 362.880	$20! = 2.432.902.008.176.640.000 \ (\in \text{tint})$				
10! = 3.628.800	21! = 51.090.942.171.709.400.000				
$\max \text{ signed tint} = 9.223.372.036.854.775.807$					
max unsigned tint = $18.446.744.073.709.551.615$					

#### Primos

 $2\ 3\ 5\ 7\ 11\ 13\ 17\ 19\ 23\ 29\ 31\ 37\ 41\ 43\ 47\ 53\ 59\ 61\ 67\ 71\ 73\ 79\ 83\ 89\ 97\ 101\ 103\ 107\ 109$  $113\ 127\ 131\ 137\ 139\ 149\ 151\ 157\ 163\ 167\ 173\ 179\ 181\ 191\ 193\ 197\ 199\ 211\ 223\ 227$ 229 233 239 241 251 257 263 269 271 277 281 283 293 307 311 313 317 331 337 347 349 353 359 367 373 379 383 389 397 401 409 419 421 431 433 439 443 449 457 461  $463\ 467\ 479\ 487\ 491\ 499\ 503\ 509\ 521\ 523\ 541\ 547\ 557\ 563\ 569\ 571\ 577\ 587\ 593\ 599$  $601\ 607\ 613\ 617\ 619\ 631\ 641\ 643\ 647\ 653\ 659\ 661\ 673\ 677\ 683\ 691\ 701\ 709\ 719\ 727$  $733\ 739\ 743\ 751\ 757\ 761\ 769\ 773\ 787\ 797\ 809\ 811\ 821\ 823\ 827\ 829\ 839\ 853\ 857\ 859$ 863 877 881 883 887 907 911 919 929 937 941 947 953 967 971 977 983 991 997 1009 1013 1019 1021 1031 1033 1039 1049 1051 1061 1063 1069 1087 1091 1093 1097 1103  $1109\ 1117\ 1123\ 1129\ 1151\ 1153\ 1163\ 1171\ 1181\ 1187\ 1193\ 1201\ 1213\ 1217\ 1223\ 1229$ 1231 1237 1249 1259 1277 1279 1283 1289 1291 1297 1301 1303 1307 1319 1321 1327 1361 1367 1373 1381 1399 1409 1423 1427 1429 1433 1439 1447 1451 1453 1459 1471 1481 1483 1487 1489 1493 1499 1511 1523 1531 1543 1549 1553 1559 1567 1571 1579  $1583\ 1597\ 1601\ 1607\ 1609\ 1613\ 1619\ 1621\ 1627\ 1637\ 1657\ 1663\ 1667\ 1669\ 1693\ 1697$  $1699\ 1709\ 1721\ 1723\ 1733\ 1741\ 1747\ 1753\ 1759\ 1777\ 1783\ 1787\ 1789\ 1801\ 1811\ 1823$  $1831\ 1847\ 1861\ 1867\ 1871\ 1873\ 1877\ 1879\ 1889\ 1901\ 1907\ 1913\ 1931\ 1933\ 1949\ 1951$ 1973 1979 1987 1993 1997 1999 2003 2011 2017 2027 2029 2039 2053 2063 2069 2081

#### Primos cercanos a $10^n$

 $\begin{array}{c} 9941\ 9949\ 9967\ 9973\ 10007\ 10009\ 10037\ 10039\ 10061\ 10067\ 10069\ 10079 \\ 99961\ 99971\ 99989\ 99991\ 100003\ 100003\ 1000033\ 1000037\ 1000039 \\ 9999943\ 9999971\ 99999991\ 10000019\ 10000079\ 10000103\ 10000121 \\ 99999941\ 99999959\ 9999971\ 99999989\ 100000007\ 100000037\ 100000039\ 100000049 \\ 99999893\ 99999929\ 99999937\ 1000000007\ 1000000009\ 1000000021\ 1000000033 \end{array}$ 

#### Cantidad de primos menores que $10^n$

```
\pi(10^1) = 4; \pi(10^2) = 25; \pi(10^3) = 168; \pi(10^4) = 1229; \pi(10^5) = 9592

\pi(10^6) = 78.498; \pi(10^7) = 664.579; \pi(10^8) = 5.761.455; \pi(10^9) = 50.847.534

\pi(10^{10}) = 455.052,511; \pi(10^{11}) = 4.118.054.813; \pi(10^{12}) = 37.607.912.018
```

#### Divisores

```
Cantidad de divisores (\sigma_0) para algunos n/\neg \exists n' < n, \sigma_0(n') \geqslant \sigma_0(n)
       \sigma_0(60) = 12; \sigma_0(120) = 16; \sigma_0(180) = 18; \sigma_0(240) = 20; \sigma_0(360) = 24
    \sigma_0(720) = 30; \sigma_0(840) = 32; \sigma_0(1260) = 36; \sigma_0(1680) = 40; \sigma_0(10080) = 72
        \sigma_0(15120) = 80; \sigma_0(50400) = 108; \sigma_0(83160) = 128; \sigma_0(110880) = 144
    \sigma_0(498960) = 200; \sigma_0(554400) = 216; \sigma_0(1081080) = 256; \sigma_0(1441440) = 288
                            \sigma_0(4324320) = 384 : \sigma_0(8648640) = 448
             Suma de divisores (\sigma_1) para algunos n/\neg \exists n' < n, \sigma_1(n') \ge \sigma_1(n)
    \sigma_1(96) = 252; \sigma_1(108) = 280; \sigma_1(120) = 360; \sigma_1(144) = 403; \sigma_1(168) = 480
        \sigma_1(960) = 3048; \sigma_1(1008) = 3224; \sigma_1(1080) = 3600; \sigma_1(1200) = 3844
     \sigma_1(4620) = 16128; \sigma_1(4680) = 16380; \sigma_1(5040) = 19344; \sigma_1(5760) = 19890
   \sigma_1(8820) = 31122; \sigma_1(9240) = 34560; \sigma_1(10080) = 39312; \sigma_1(10920) = 40320
\sigma_1(32760) = 131040; \sigma_1(35280) = 137826; \sigma_1(36960) = 145152; \sigma_1(37800) = 148800
\sigma_1(60480) = 243840; \sigma_1(64680) = 246240; \sigma_1(65520) = 270816; \sigma_1(70560) = 280098
            \sigma_1(95760) = 386880; \sigma_1(98280) = 403200; \sigma_1(100800) = 409448
        \sigma_1(491400) = 2083200; \sigma_1(498960) = 2160576; \sigma_1(514080) = 2177280
        \sigma_1(982800) = 4305280; \sigma_1(997920) = 4390848; \sigma_1(1048320) = 4464096
     \sigma_1(4979520) = 22189440; \sigma_1(4989600) = 22686048; \sigma_1(5045040) = 23154768
     \sigma_1(9896040) = 44323200 : \sigma_1(9959040) = 44553600 : \sigma_1(9979200) = 45732192
```

# 7. Grafos

# 7.1. Dijkstra

```
#define INF 1e9
int N;
#define MAX_V 250001
vector<ii> G[MAX_V];
//To add an edge use
```

```
6 #define add(a, b, w) G[a].pb(make_pair(w, b))
                                                                                    3 int G[MAX_N][MAX_N];
                                                                                      void floyd(){\frac{}{0(N^3)}}
   11 dijkstra(int s, int t){\frac{}{|0(|E| \log |V|)}}
     priority_queue<ii, vector<ii>, greater<ii> > Q;
     vector<ll> dist(N, INF); vector<int> dad(N, -1);
                                                                                      }
                                                                                    7
     Q.push(make_pair(0, s)); dist[s] = 0;
11
     while(sz(Q)){
12
       ii p = Q.top(); Q.pop();
13
       if(p.snd == t) break;
14
       forall(it, G[p.snd])
15
         if(dist[p.snd]+it->first < dist[it->snd]){
16
                                                                                   13
           dist[it->snd] = dist[p.snd] + it->fst;
17
                                                                                   14
           dad[it->snd] = p.snd;
                                                                                   15 }
18
           Q.push(make_pair(dist[it->snd], it->snd));
19
         }
20
     }
21
     return dist[t];
22
     if(dist[t]<INF)//path generator</pre>
23
       for(int i=t: i!=-1: i=dad[i])
24
         printf("%d%c", i, (i==s?'\n':'\_'));
25
26 }
                                                                                    6
                           7.2. Bellman-Ford
                                                                                    8
  vector<ii> G[MAX_N];//ady. list with pairs (weight, dst)
                                                                                    9
   int dist[MAX_N];
                                                                                   10
   void bford(int src){//O(VE)
                                                                                   11
     dist[src]=0;
                                                                                   12
4
                                                                                           }
     forn(i, N-1) forn(j, N) if(dist[j]!=INF) forall(it, G[j])
                                                                                   13
       dist[it->snd]=min(dist[it->snd], dist[j]+it->fst);
6
                                                                                   14
                                                                                   15 }
7
8
   bool hasNegCycle(){
9
     forn(j, N) if(dist[j]!=INF) forall(it, G[j])
       if(dist[it->snd]>dist[j]+it->fst) return true;
     //inside if: all points reachable from it->snd will have -INF distance
         (do bfs)
                                                                                    4
     return false;
                                                                                    5
14 }
                                                                                    6
                          7.3. Floyd-Warshall
                                                                                      }
                                                                                    7
1 |//G[i][j] contains weight of edge (i, j) or INF
2 //G[i][i]=0
```

```
forn(k, N) forn(i, N) if(G[i][k]!=INF) forn(j, N) if(G[k][j]!=INF)
    G[i][j]=min(G[i][j], G[i][k]+G[k][j]);
  bool inNegCycle(int v){
    return G[v][v]<0;}
  //checks if there's a neg. cycle in path from a to b
  bool hasNegCycle(int a, int b){
   forn(i, N) if(G[a][i]!=INF && G[i][i]<0 && G[i][b]!=INF)
      return true;
    return false;
                             7.4. Kruskal
struct Ar{int a,b,w;};
bool operator<(const Ar& a, const Ar &b){return a.w<b.w;}</pre>
  vector<Ar> E:
  ll kruskal(){
      11 cost=0;
      sort(E.begin(), E.end());//ordenar aristas de menor a mayor
      uf.init(n);
      forall(it, E){
          if(uf.comp(it->a)!=uf.comp(it->b)){//si no estan conectados
              uf.unir(it->a, it->b);//conectar
              cost+=it->w;
          }
      return cost;
                               7.5. Prim
 bool taken[MAXN];
 |priority_queue<ii, vector<ii>, greater<ii> > pq;//min heap
  void process(int v){
      taken[v]=true;
      forall(e, G[v])
          if(!taken[e->second]) pq.push(*e);
  | 11 prim(){
      zero(taken);
```

```
process(0);
ll cost=0;
while(sz(pq)){
    ii e=pq.top(); pq.pop();
    if(!taken[e.second]) cost+=e.first, process(e.second);
}
return cost;
}
```

### 7.6. 2-SAT + Tarjan SCC

```
//We have a vertex representing a var and other for his negation.
   //Every edge stored in G represents an implication. To add an equation
       of the form allb, use addor(a, b)
   //MAX=max cant var, n=cant var
   #define addor(a, b) (G[neg(a)].pb(b), G[neg(b)].pb(a))
   vector<int> G[MAX*2];
   //idx[i]=index assigned in the dfs
   //lw[i]=lowest index(closer from the root) reachable from i
   int lw[MAX*2], idx[MAX*2], qidx;
   stack<int> q;
   int qcmp, cmp[MAX*2];
   //verdad[cmp[i]]=valor de la variable i
   bool verdad[MAX*2+1];
13
   int neg(int x) { return x>=n? x-n : x+n;}
   void tin(int v){
15
     lw[v]=idx[v]=++qidx;
16
     q.push(v), cmp[v]=-2;
17
     forall(it, G[v]){
18
       if(!idx[*it] || cmp[*it]==-2){
19
         if(!idx[*it]) tjn(*it);
20
         lw[v]=min(lw[v], lw[*it]);
21
       }
^{22}
     }
23
     if(lw[v]==idx[v]){
24
       qcmp++;
^{25}
       int x:
26
       do{x=q.top(); q.pop(); cmp[x]=qcmp;}while(x!=v);
27
       verdad[qcmp] = (cmp[neg(v)] < 0);</pre>
28
29
30
   //remember to CLEAR G!!!
```

```
bool satisf(){//O(n)
   memset(idx, 0, sizeof(idx)), qidx=0;
   memset(cmp, -1, sizeof(cmp)), qcmp=0;
   forn(i, n){
        if(!idx[i]) tjn(i);
        if(!idx[neg(i)]) tjn(neg(i));
    }
   forn(i, n) if(cmp[i]==cmp[neg(i)]) return false;
   return true;
}
```

#### 7.7. Articulation Points

```
1 int N;
   vector<int> G[1000000];
   //V[i]=node number(if visited), L[i]= lowest V[i] reachable from i
   int qV, V[1000000], L[1000000], P[1000000];
   void dfs(int v, int f){
     L[v]=V[v]=++qV;
     forall(it, G[v])
       if(!V[*it]){
         dfs(*it, v);
9
         L[v] = min(L[v], L[*it]);
         P[v] += L[*it] >= V[v];
11
12
       else if(*it!=f)
13
         L[v]=min(L[v], V[*it]);
14
   }
15
   int cantart(){ //O(n)
     qV=0;
17
     zero(V), zero(P);
     dfs(1, 0); P[1]--;
     int q=0;
     forn(i, N) if(P[i]) q++;
   return q;
23 }
```

# 7.8. Comp. Biconexas y Puentes

```
struct edge {
  int u,v, comp;
  bool bridge;
};
vector<edge> e;
```

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```
6 | void addEdge(int u, int v) {
     G[u].pb(sz(e)), G[v].pb(sz(e));
     e.pb((edge){u,v,-1,false});
9
   //d[i]=id de la dfs
   //b[i]=lowest id reachable from i
   int d[MAXN], b[MAXN], t;
   int nbc;//cant componentes
   int comp[MAXN];//comp[i]=cant comp biconexas a la cual pertenece i
   void initDfs(int n) {
     zero(G), zero(comp);
     e.clear();
     forn(i,n) d[i]=-1;
     nbc = t = 0:
19
20
   stack<int> st;
   void dfs(int u, int pe) \{//0(n + m)\}
     b[u] = d[u] = t++;
23
     comp[u] = (pe != -1);
24
     forall(ne, G[u]) if (*ne != pe){
25
       int v = e[*ne].u ^ e[*ne].v ^ u;
26
       if (d[v] == -1) {
27
         st.push(*ne);
28
         dfs(v,*ne);
29
         if (b[v] > d[u]){
30
            e[*ne].bridge = true; // bridge
31
         }
32
         if (b[v] >= d[u]) \{ // art \}
33
           int last;
34
           do {
35
             last = st.top(); st.pop();
36
              e[last].comp = nbc;
37
           } while (last != *ne);
38
           nbc++:
39
           comp[u]++;
40
41
         b[u] = min(b[u], b[v]);
42
43
       else if (d[v] < d[u]) \{ // back edge
44
         st.push(*ne);
45
         b[u] = min(b[u], d[v]);
46
47
     }
48
```

```
49 }
                          7.9. LCA + Climb
 const int MAXN=100001;
   const int LOGN=20;
   //f[v][k] holds the 2<sup>k</sup> father of v
   //L[v] holds the level of v
   int N, f[MAXN][LOGN], L[MAXN];
   //call before build:
   void dfs(int v, int fa=-1, int lvl=0){//generate required data
    f[v][0]=fa, L[v]=lvl;
    forall(it, G[v])if(*it!=fa)
       dfs(*it, v, lvl+1);
10
11
   void build(){//f[i][0] must be filled previously, O(nlgn)
     forn(k, LOGN-1) forn(i, N) f[i][k+1]=f[f[i][k]][k];}
14
   #define lg(x) (31-_builtin_clz(x))//=floor(log2(x))
16
   int climb(int a, int d){\frac{1}{0}}
     if(!d) return a:
18
     dforn(i, lg(L[a])+1)
19
       if(1<<i<=d)
20
         a=f[a][i], d-=1<<i;
21
       return a;
22
23
   int lca(int a, int b){\frac{1}{0}}
     if(L[a]<L[b]) swap(a, b);</pre>
     a=climb(a, L[a]-L[b]);
     if(a==b) return a;
     dforn(i, lg(L[a])+1)
       if(f[a][i]!=f[b][i])
         a=f[a][i], b=f[b][i];
30
     return f[a][0];
31
32
   int dist(int a, int b) {//returns distance between nodes
    return L[a]+L[b]-2*L[lca(a, b)];}
                 7.10. Heavy Light Decomposition
int treesz[MAXN];//cantidad de nodos en el subarbol del nodo v
int dad[MAXN];//dad[v]=padre del nodo v
void dfs1(int v, int p=-1){//pre-dfs
```

int szt[MAXN];

void calcsz(int v, int p) {

26

reverse(path.begin(), path.end());

```
dad[v]=p;
                                                                                       szt[v] = 1:
4
                                                                                       forall(it,G[v]) if (*it!=p && !taken[*it])
     treesz[v]=1;
5
                                                                                  9
                                                                                         calcsz(*it,v), szt[v]+=szt[*it];
     forall(it, G[v]) if(*it!=p){
6
                                                                                  10
       dfs1(*it, v);
                                                                                  11
       treesz[v]+=treesz[*it];
                                                                                     void centroid(int v=0, int f=-1, int lvl=0, int tam=-1) {//0(nlogn)
                                                                                       if(tam==-1) calcsz(v, -1), tam=szt[v];
9
                                                                                       forall(it, G[v]) if(!taken[*it] && szt[*it]>=tam/2)
10
                                                                                  14
   //PONER Q EN O !!!!!
                                                                                         {szt[v]=0; centroid(*it, f, lvl, tam); return;}
                                                                                  15
   int pos[MAXN], q;//pos[v]=posicion del nodo v en el recorrido de la dfs
                                                                                       taken[v]=true;
                                                                                  16
   //Las cadenas aparecen continuas en el recorrido!
                                                                                       padre[v]=f;
                                                                                  17
                                                                                       forall(it, G[v]) if(!taken[*it])
   int cantcad:
                                                                                  18
   int homecad[MAXN];//dada una cadena devuelve su nodo inicial
                                                                                         centroid(*it, v, lvl+1, -1);
                                                                                  19
   int cad[MAXN];//cad[v]=cadena a la que pertenece el nodo
                                                                                  20 }
   void heavylight(int v, int cur=-1){
                                                                                                             7.12. Euler Cycle
     if(cur==-1) homecad[cur=cantcad++]=v:
18
     pos[v]=q++;
19
     cad[v]=cur;
20
                                                                                  int n,m,ars[MAXE], eq;
     int mx=-1;
                                                                                    vector<int> G[MAXN];//fill G,n,m,ars,eq
21
     forn(i, sz(G[v])) if(G[v][i]!=dad[v])
                                                                                  3 | list<int> path;
22
       if(mx==-1 || treesz[G[v][mx]]<treesz[G[v][i]]) mx=i;</pre>
                                                                                     int used[MAXN]:
23
     if(mx!=-1) heavylight(G[v][mx], cur);
                                                                                     bool usede[MAXE];
24
     forn(i, sz(G[v])) if(i!=mx && G[v][i]!=dad[v])
25
                                                                                     queue<list<int>::iterator> q;
       heavylight(G[v][i], -1);
                                                                                     int get(int v){
26
                                                                                       while(used[v]\leq z(G[v]) && usede[G[v][used[v]]]) used[v]++;
27
   //ejemplo de obtener el maximo numero en el camino entre dos nodos
                                                                                       return used[v]:
28
                                                                                  9
   //RTA: max(query(low, u), query(low, v)), con low=lca(u, v)
                                                                                  10
   //esta funcion va trepando por las cadenas
                                                                                     void explore(int v, int r, list<int>::iterator it){
   int query(int an, int v){//O(logn)}
                                                                                       int ar=G[v][get(v)]; int u=v^ars[ar];
31
                                                                                 12
     //si estan en la misma cadena:
                                                                                       usede[ar]=true;
32
                                                                                 13
     if(cad[an] == cad[v]) return rmq.get(pos[an], pos[v]+1);
                                                                                       list<int>::iterator it2=path.insert(it, u);
33
                                                                                 14
     return max(query(an, dad[homecad[cad[v]]]),
                                                                                       if(u!=r) explore(u, r, it2);
34
                                                                                 15
            rmq.get(pos[homecad[cad[v]]], pos[v]+1));
                                                                                       if(get(v)<sz(G[v])) q.push(it);</pre>
35
                                                                                  16
36 }
                                                                                 17
                                                                                     void euler(){
                                                                                  18
                    7.11. Centroid Decomposition
                                                                                       zero(used), zero(usede);
                                                                                 19
                                                                                       path.clear();
                                                                                 20
                                                                                       q=queue<list<int>::iterator>();
  int n;
                                                                                 21
  vector<int> G[MAXN]:
                                                                                       path.push_back(0); q.push(path.begin());
                                                                                 22
  bool taken[MAXN];//poner todos en FALSE al principio!!
                                                                                       while(sz(q)){
                                                                                 23
   int padre[MAXN];//padre de cada nodo en el centroid tree
                                                                                         list<int>::iterator it=q.front(); q.pop();
                                                                                 24
                                                                                         if(used[*it] < sz(G[*it])) explore(*it, *it, it);</pre>
                                                                                 25
```

```
28 }
   void addEdge(int u, int v){
29
     G[u].pb(eq), G[v].pb(eq);
30
     ars[eq++]=u^v;
31
  |}
32
                         7.13.
                                Diametro árbol
   vector<int> G[MAXN];
   int n,m;
2
   int p[MAXN],d[MAXN],d2[MAXN];
   int bfs(int r, int *d) {
     queue<int> q;
6
     d[r]=0; q.push(r);
7
     int v:
8
     while(sz(q)) { v=q.front(); q.pop();
9
       forall(it,G[v]) if (d[*it]==-1) {
10
         d[*it]=d[v]+1, p[*it]=v;
11
         q.push(*it);
12
       }
13
14
     return v;//ultimo nodo visitado
15
16
17
   vector<int> diams;
   vector<ii> centros;
   void diametros(){
     memset(d,-1,sizeof(d));
21
     memset(d2,-1,sizeof(d2));
     diams.clear(), centros.clear();
23
     forn(i, n) if(d[i]==-1){
24
       int v,c;
25
       c=v=bfs(bfs(i, d2), d);
26
       forn(_,d[v]/2) c=p[c];
27
       diams.pb(d[v]);
28
       if(d[v]&1) centros.pb(ii(c, p[c]));
29
       else centros.pb(ii(c, c));
30
     }
31
32 }
                               7.14. Chu-liu
```

```
void visit(graph &h, int v, int s, int r,
```

```
vector<int> &no, vector< vector<int> > &comp,
     vector<int> &prev, vector< vector<int> > &next, vector<weight> &mcost,
     vector<int> &mark, weight &cost, bool &found) {
     if (mark[v]) {
       vector<int> temp = no;
       found = true;
       do {
         cost += mcost[v];
         v = prev[v];
         if (v != s) {
           while (comp[v].size() > 0) {
12
             no[comp[v].back()] = s;
13
             comp[s].push_back(comp[v].back());
14
             comp[v].pop_back();
           }
16
         }
17
       } while (v != s);
18
       forall(j,comp[s]) if (*j != r) forall(e,h[*j])
         if (no[e->src] != s) e->w -= mcost[ temp[*j] ];
20
     }
21
     mark[v] = true;
22
     forall(i,next[v]) if (no[*i] != no[v] && prev[no[*i]] == v)
       if (!mark[no[*i]] || *i == s)
24
         visit(h, *i, s, r, no, comp, prev, next, mcost, mark, cost, found)
25
26
   weight minimumSpanningArborescence(const graph &g, int r) {
       const int n=sz(g);
     graph h(n);
29
     forn(u,n) forall(e,g[u]) h[e->dst].pb(*e);
     vector<int> no(n);
31
     vector<vector<int> > comp(n);
     forn(u, n) comp[u].pb(no[u] = u);
33
     for (weight cost = 0; ;) {
34
       vector<int> prev(n, -1);
35
       vector<weight> mcost(n, INF);
36
       forn(j,n) if (j != r) forall(e,h[j])
37
         if (no[e->src] != no[i])
           if (e->w < mcost[ no[i] ])</pre>
39
             mcost[no[j]] = e->w, prev[no[j]] = no[e->src];
       vector< vector<int> > next(n);
41
       forn(u,n) if (prev[u] >= 0)
42
         next[ prev[u] ].push_back(u);
43
```

```
bool stop = true;
44
       vector<int> mark(n);
45
       forn(u,n) if (u != r && !mark[u] && !comp[u].empty()) {
46
         bool found = false;
47
         visit(h, u, u, r, no, comp, prev, next, mcost, mark, cost, found);
48
         if (found) stop = false;
49
       }
50
       if (stop) {
51
         forn(u,n) if (prev[u] >= 0) cost += mcost[u];
52
         return cost;
53
       }
54
    }
55
56 }
```

# 7.15. Hungarian

```
1 //Dado un grafo bipartito completo con costos no negativos, encuentra el
        matching perfecto de minimo costo.
2 | tipo cost[N][N], lx[N], ly[N], slack[N]; //llenar: cost=matriz de
       advacencia
  int n, max_match, xy[N], yx[N], slackx[N], prev2[N]; //n=cantidad de nodos
   bool S[N], T[N]; //sets S and T in algorithm
   void add_to_tree(int x, int prevx) {
     S[x] = true, prev2[x] = prevx;
    form(y, n) if (lx[x] + ly[y] - cost[x][y] < slack[y] - EPS)
7
       slack[y] = lx[x] + ly[y] - cost[x][y], slackx[y] = x;
8
9
   void update_labels(){
10
     tipo delta = INF;
11
     forn (y, n) if (!T[y]) delta = min(delta, slack[y]);
12
     form (x, n) if (S[x]) lx[x] -= delta;
13
     forn (y, n) if (T[y]) ly[y] += delta; else slack[y] -= delta;
14
15
   void init_labels(){
16
     zero(lx), zero(ly);
17
     form (x,n) form (y,n) lx[x] = max(lx[x], cost[x][y]);
18
19
   void augment() {
20
     if (max_match == n) return;
21
     int x, y, root, q[N], wr = 0, rd = 0;
     memset(S, false, sizeof(S)), memset(T, false, sizeof(T));
23
     memset(prev2, -1, sizeof(prev2));
24
     forn (x, n) if (xy[x] == -1){
```

```
q[wr++] = root = x, prev2[x] = -2;
26
       S[x] = true; break; }
27
    forn (y, n) slack[y] = lx[root] + ly[y] - cost[root][y], slackx[y] =
28
     while (true){
29
       while (rd < wr){
30
         x = q[rd++];
31
         for (y = 0; y < n; y++) if (cost[x][y] == lx[x] + ly[y] && !T[y]){
           if (yx[y] == -1) break; T[y] = true;
33
           q[wr++] = yx[y], add_to_tree(yx[y], x); }
         if (y < n) break; }
35
       if (y < n) break;
       update_labels(), wr = rd = 0;
37
       for (y = 0; y < n; y++) if (!T[y] && slack[y] == 0){
         if (yx[y] == -1)\{x = slackx[y]; break;\}
         else{
           T[y] = true;
41
           if (!S[yx[y]]) q[wr++] = yx[y], add_to_tree(yx[y], slackx[y]);
         }}
       if (v < n) break; }
     if (v < n){
45
       max_match++;
       for (int cx = x, cy = y, ty; cx != -2; cx = prev2[cx], cy = ty)
47
         tv = xv[cx], vx[cv] = cx, xv[cx] = cv;
       augment(); }
49
50
   tipo hungarian(){
51
     tipo ret = 0; max_match = 0, memset(xy, -1, sizeof(xy));
     memset(yx, -1, sizeof(yx)), init_labels(), augment(); //steps 1-3
     forn (x,n) ret += cost[x][xy[x]]; return ret;
55 }
                     7.16. Dynamic Conectivity
```

```
1 struct UnionFind {
       int n, comp;
2
       vector<int> pre,si,c;
3
       UnionFind(int n=0):n(n), comp(n), pre(n), si(n, 1) {
4
           forn(i,n) pre[i] = i; }
5
      int find(int u){return u==pre[u]?u:find(pre[u]);}
6
       bool merge(int u, int v) {
7
           if((u=find(u))==(v=find(v))) return false;
8
9
           if(si[u]<si[v]) swap(u, v);</pre>
```

```
si[u]+=si[v], pre[v]=u, comp--, c.pb(v);
10
           return true;
11
       }
12
       int snap(){return sz(c);}
13
       void rollback(int snap){
14
           while(sz(c)>snap){
15
               int v = c.back(); c.pop_back();
16
               si[pre[v]] -= si[v], pre[v] = v, comp++;
17
18
19
20
   enum {ADD,DEL,QUERY};
   struct Query {int type,u,v;};
   struct DynCon {
       vector<Query> q;
24
       UnionFind dsu;
25
       vector<int> match,res;
26
       map<ii,int> last;//se puede no usar cuando hay identificador para
27
           cada arista (mejora poco)
       DynCon(int n=0):dsu(n){}
28
       void add(int u, int v) {
29
           if(u>v) swap(u,v);
30
           q.pb((Query){ADD, u, v}), match.pb(-1);
31
           last[ii(u,v)] = sz(q)-1;
32
       }
33
       void remove(int u, int v) {
34
           if(u>v) swap(u,v);
35
           q.pb((Query){DEL, u, v});
36
           int prev = last[ii(u,v)];
37
           match[prev] = sz(q)-1;
38
           match.pb(prev);
39
40
       void query() {//podria pasarle un puntero donde guardar la respuesta
41
           q.pb((Query){QUERY, -1, -1}), match.pb(-1);}
42
       void process() {
43
           forn(i,sz(q)) if (q[i].type == ADD && match[i] == -1) match[i] =
44
                sz(q);
           go(0,sz(q));
45
       }
46
       void go(int 1, int r) {
47
           if(l+1==r){
48
               if (q[1].type == QUERY)//Aqui responder la query usando el
49
                    dsu!
```

```
res.pb(dsu.comp);//aqui query=cantidad de componentes
50
                        conexas
                return;
51
           }
52
           int s=dsu.snap(), m = (1+r) / 2;
53
           forr(i,m,r) if(match[i]!=-1 && match[i]<1) dsu.merge(q[i].u, q[i</pre>
54
                ].v);
           go(1,m);
55
           dsu.rollback(s);
           s = dsu.snap();
           forr(i,1,m) if(match[i]!=-1 && match[i]>=r) dsu.merge(q[i].u, q[
58
                il.v):
           go(m,r):
59
           dsu.rollback(s);
60
62 }dc;
```

### 8. Network Flow

#### 8.1. Dinic

```
const int MAX = 300;
   int nodes, src. dst:
   int dist[MAX], q[MAX], work[MAX];
   struct Edge {
       int to, rev;
       ll f, cap;
       Edge(int to, int rev, 11 f, 11 cap): to(to), rev(rev), f(f), cap(
9
           cap) {}
   };
10
   vector<Edge> G[MAX];
11
12
   // Adds bidirectional edge
   void addEdge(int s, int t, ll cap){
       G[s].pb(Edge(t, sz(G[t]), 0, cap));
15
       G[t].pb(Edge(s, sz(G[s])-1, 0, 0));
16
   }
17
18
   bool dinic_bfs(){
       fill(dist, dist+nodes, -1);
20
       dist[src]=0;
21
```

```
int qt=0;
22
       q[qt++]=src;
23
       for(int qh=0; qh<qt; qh++){</pre>
^{24}
            int u =q[qh];
25
            forall(e, G[u]){
26
                int v=e->to;
27
                if(dist[v]<0 \&\& e->f < e->cap){
28
                     dist[v]=dist[u]+1;
29
                     q[qt++]=v;
30
                }
31
            }
32
       }
33
       return dist[dst]>=0;
34
35
36
   ll dinic_dfs(int u, ll f){
37
       if(u==dst) return f;
38
       for(int &i=work[u]; i<sz(G[u]); i++){</pre>
39
            Edge &e = G[u][i];
40
            if(e.cap<=e.f) continue;</pre>
41
            int v=e.to;
42
            if(dist[v]==dist[u]+1){
43
                     11 df=dinic_dfs(v, min(f, e.cap-e.f));
44
                     if(df>0){
45
                              e.f+=df;
46
                             G[v][e.rev].f-= df;
47
                             return df;
48
                     }
49
            }
50
       }
51
       return 0;
52
53
54
   ll maxFlow(int _src, int _dst){
55
       src=_src;
56
       dst=_dst;
57
       11 result=0;
58
       while(dinic bfs()){
59
            fill(work, work+nodes, 0);
60
            while(ll delta=dinic_dfs(src,INF))
61
                result+=delta;
62
       }
63
       // todos los nodos con dist[v]!=-1 vs los que tienen dist[v]==-1
64
```

```
forman el min-cut
       return result;
65
66 }
                               8.2. Konig
1 // asume que el dinic YA ESTA tirado
   // asume que nodes-1 y nodes-2 son la fuente y destino
   int match[maxnodes]; // match[v]=u si u-v esta en el matching, -1 si v
       no esta matcheado
  int s[maxnodes]; // numero de la bfs del koning
   queue<int> kq;
   // s[e] %2==1 o si e esta en V1 y s[e] ==-1-> lo agarras
   void koning() {//O(n)
     forn(v,nodes-2) s[v] = match[v] = -1;
     forn(v,nodes-2) forall(it,g[v]) if (it->to < nodes-2 && it->f>0)
       { match[v]=it->to; match[it->to]=v;}
10
     forn(v,nodes-2) if (match[v]==-1) {s[v]=0;kq.push(v);}
11
     while(!kq.empty()) {
12
       int e = kq.front(); kq.pop();
13
       if (s[e] %2==1) {
14
         s[match[e]] = s[e]+1;
15
         kq.push(match[e]);
16
       } else {
17
18
         forall(it,g[e]) if (it->to < nodes-2 && s[it->to]==-1) {
19
           s[it->to] = s[e]+1;
20
           kq.push(it->to);
21
22
23
    }
24
25 }
                         8.3. Edmonds Karp's
  #define MAX_V 1000
   #define INF 1e9
   //special nodes
   #define SRC 0
   #define SNK 1
6 map<int, int> G[MAX_V];//limpiar esto
7 //To add an edge use
  #define add(a, b, w) G[a][b]=w
9 int f, p[MAX_V];
```

```
void augment(int v, int minE){
                                                                                          int amt = min(excess[a], ll(G[a][b]));
                                                                                    16
                                                                                          if(height[a] <= height[b] || amt == 0) return;</pre>
     if(v==SRC) f=minE;
                                                                                    17
11
     else if(p[v]!=-1){
12
                                                                                    18
       augment(p[v], min(minE, G[p[v]][v]));
13
                                                                                    19
       G[p[v]][v]-=f, G[v][p[v]]+=f;
                                                                                          enqueue(b);
14
                                                                                    20
15
                                                                                    21
16
                                                                                    22
   11 maxflow(){//0(VE^2)
                                                                                          forn(v, N){
     11 Mf=0;
18
     do{
19
       f=0:
20
                                                                                    26
       char used[MAX_V]; queue<int> q; q.push(SRC);
21
                                                                                    27
       zero(used), memset(p, -1, sizeof(p));
22
                                                                                    28
       while(sz(q)){
                                                                                         }
                                                                                    29
23
         int u=q.front(); q.pop();
                                                                                    30
24
         if(u==SNK) break;
25
         forall(it, G[u])
26
           if(it->snd>0 && !used[it->fst])
27
                                                                                    33
             used[it->fst]=true, q.push(it->fst), p[it->fst]=u;
                                                                                    34
28
       }
29
       augment(SNK, INF);
                                                                                    36
30
       Mf+=f;
31
                                                                                          enqueue(v);
     }while(f);
32
                                                                                    38
     return Mf;
                                                                                    39
33
34 }
                                                                                    40
                       8.4. Push-Relabel O(N3)
                                                                                    42
                                                                                    43
   #define MAX_V 1000
                                                                                    44
   int N;//valid nodes are [0...N-1]
   #define INF 1e9
                                                                                    46
   //special nodes
   #define SRC 0
                                                                                    48
```

#define SNK 1

map<int, int> G[MAX\_V];

#define add(a, b, w) G[a][b]=w

int height[MAX\_V], active[MAX\_V], count[2\*MAX\_V+1];

if (!active[v] && excess[v] > 0) active[v]=true, Q.push(v); }

//To add an edge use

void enqueue(int v) {

void push(int a, int b) {

ll excess[MAX V]:

queue<int> Q;

```
G[a][b]-=amt, G[b][a]+=amt;
     excess[b] += amt, excess[a] -= amt;
   void gap(int k) {
       if (height[v] < k) continue;</pre>
       count[height[v]]--;
       height[v] = max(height[v], N+1);
       count[height[v]]++;
       enqueue(v);
   void relabel(int v) {
     count[height[v]]--;
     height[v] = 2*N;
     forall(it, G[v])
       if(it->snd)
         height[v] = min(height[v], height[it->fst] + 1);
     count[height[v]]++;
   ll maxflow() \{//0(V^3)
     zero(height), zero(active), zero(count), zero(excess);
     count[0] = N-1;
     count[N] = 1;
     height[SRC] = N;
     active[SRC] = active[SNK] = true;
     forall(it, G[SRC]){
       excess[SRC] += it->snd;
       push(SRC, it->fst);
    }
49
     while(sz(Q)) {
50
       int v = Q.front(); Q.pop();
51
       active[v]=false;
52
     forall(it, G[v]) push(v, it->fst);
53
     if(excess[v] > 0)
54
       count[height[v]] == 1? gap(height[v]):relabel(v);
55
     }
56
     ll mf=0;
57
     forall(it, G[SRC]) mf+=G[it->fst][SRC];
```

return mf;

```
60 }
                               Min-cost Max-flow
                        8.5.
   const int MAXN=10000;
   typedef 11 tf;
   typedef 11 tc;
   const tf INFFLUJO = 1e14;
   const tc INFCOSTO = 1e14;
   struct edge {
     int u, v;
     tf cap, flow;
     tc cost;
     tf rem() { return cap - flow; }
10
11
   int nodes; //numero de nodos
12
   vector<int> G[MAXN]; // limpiar!
   vector<edge> e; // limpiar!
   void addEdge(int u, int v, tf cap, tc cost) {
     G[u].pb(sz(e)); e.pb((edge){u,v,cap,0,cost});
     G[v].pb(sz(e)); e.pb((edge){v,u,0,0,-cost});
17
18
   tc dist[MAXN], mnCost;
   int pre[MAXN];
   tf cap[MAXN], mxFlow;
   bool in_queue[MAXN];
   void flow(int s, int t) {
23
     zero(in_queue);
24
     mxFlow=mnCost=0;
25
     while(1){
26
       fill(dist, dist+nodes, INFCOSTO); dist[s] = 0;
27
       memset(pre, -1, sizeof(pre)); pre[s]=0;
28
       zero(cap); cap[s] = INFFLUJO;
29
       queue<int> q; q.push(s); in_queue[s]=1;
30
       while(sz(q)){
31
         int u=q.front(); q.pop(); in_queue[u]=0;
32
         for(auto it:G[u]) {
33
           edge &E = e[it];
34
           if(E.rem() \&\& dist[E.v] > dist[u] + E.cost + 1e-9){ // ojo EPS}
35
             dist[E.v]=dist[u]+E.cost;
36
             pre[E.v] = it;
37
             cap[E.v] = min(cap[u], E.rem());
38
```

```
if(!in_queue[E.v]) q.push(E.v), in_queue[E.v]=1;
39
40
         }
41
       }
42
       if (pre[t] == -1) break;
43
       mxFlow +=cap[t];
44
       mnCost +=cap[t]*dist[t];
45
       for (int v = t; v != s; v = e[pre[v]].u) {
         e[pre[v]].flow += cap[t];
         e[pre[v]^1].flow -= cap[t];
49
    }
50
51 }
```

# 9. Template

```
#include <bits/stdc++.h>
   using namespace std;
   #define dprint(v) cout << #v"=" << v << endl //;)
   #define forr(i,a,b) for(int i=(a); i<(b); i++)</pre>
   #define forn(i,n) forr(i,0,n)
   #define dforn(i,n) for(int i=n-1; i>=0; i--)
   #define forall(it,v) for(auto it=v.begin();it!=v.end();++it)
   #define sz(c) ((int)c.size())
   #define zero(v) memset(v, 0, sizeof(v))
   #define pb push_back
   #define fst first
   #define snd second
   typedef long long 11;
   typedef pair<int,int> ii;
15
   const int MAXN=100100;
   int n;
17
18
   int main() {
19
       freopen("input.in", "r", stdin);
20
       ios::sync_with_stdio(0);
21
       while(cin >> n){
22
23
24
25
       return 0;
26 | }
```

# 10. Ayudamemoria

#### Cant. decimales

```
#include <iomanip>
cout << setprecision(2) << fixed;</pre>
```

# Rellenar con espacios(para justificar)

```
#include <iomanip>
cout << setfill(''') << setw(3) << 2 << endl;</pre>
```

#### Leer hasta fin de linea

```
#include <sstream>
//hacer cin.ignore() antes de getline()

while(getline(cin, line)){
   istringstream is(line);
   while(is >> X)
   cout << X << """;
   cout << endl;
}</pre>
```

#### Aleatorios

```
#define RAND(a, b) (rand() %(b-a+1)+a)
rand(time(NULL));
```

# Doubles Comp.

```
const double EPS = 1e-9;
x == y <=> fabs(x-y) < EPS
x > y <=> x > y + EPS
x > y <=> x > y - EPS
```

### Limites

```
i #include <limits>
  numeric_limits<T>
  i:max()
  i:min()
  ::epsilon()
```

#### Muahaha

```
#include <signal.h>
  void divzero(int p){
    while(true);}
  void segm(int p){
    exit(0);}
  //in main
  signal(SIGFPE, divzero);
8 signal(SIGSEGV, segm);
                         Mejorar velocidad
ios::sync_with_stdio(false);
                        Mejorar velocidad 2
1 //Solo para enteros positivos
  inline void Scanf(int& a){
    char c = 0;
    while(c<33) c = getc(stdin);</pre>
    a = 0:
    while(c>33) a = a*10 + c - '0', c = getc(stdin);
7 }
                            Expandir pila
#include <sys/resource.h>
2 rlimit rl;
  getrlimit(RLIMIT_STACK, &rl);
4 | rl.rlim_cur=1024L*1024L*256L;//256mb
5 setrlimit(RLIMIT_STACK, &rl);
                               C++11
1 g++ --std=c++1
                           Leer del teclado
freopen("/dev/tty", "a", stdin);
                         Iterar subconjunto
for(int sbm=bm; sbm; sbm=(sbm-1)&bm)
                              File setup
```

```
//tambien se pueden usar comas: {a, x, m, 1} touch {a..1}.in; tee {a..1}.cpp < template.cpp
```