



# Welcome!

September 6, 2018



# Github Setup

If you haven't already, set up a github account at: <https://github.com/>

Login, and join our Organization:

<https://github.com/Algorithms-for-CP-and-Interviews>

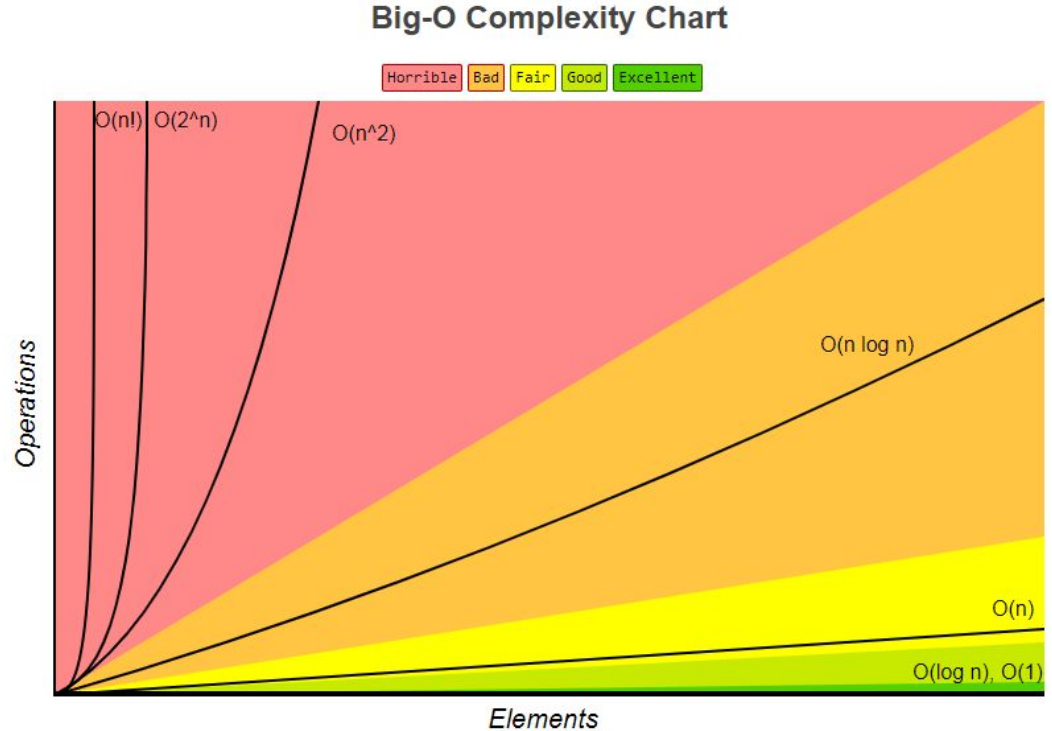
If you've joined the group for extra credit, this is how we will track your progress

# Outline

- Growth of Functions/Asymptotic Concepts
- Insertion Sort & Merge Sort
- Basic ADTs

# Growth of Functions/Asymptotic Concepts

- Use different symbols
  - Quantifies requirements (e.g. space, time, etc)
- Allows for definition of worst, average, best case



# Growth of Functions/Asymptotic Concepts

- $O \approx \leq$

- Upper bound/equal to
- Ex:  $O(n^2) = n^2, n^2 + 4324, n^{1.999}$

- $\Omega \approx \geq$

- Lower bound/equal to
- Ex:  $\Omega(n^2) = n^2, n^2 + 100, n^3$

- $\Theta \approx =$

- Average bound
- Ex:  $\Theta(n^2) \rightarrow n^2, n^2 + 4329482$

- $o \approx <$

- Strictly upper bound
- Ex:  $o(n^2) = n, n^{1.999}$

- $\omega \approx >$

- Strictly lower bound
- Ex:  $\omega(n^2) = n^3, 4n^3 + n \lg n$

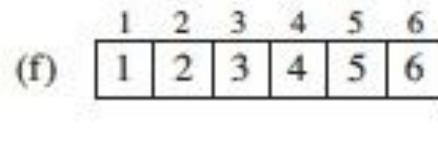
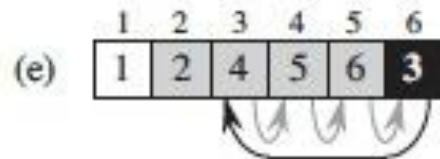
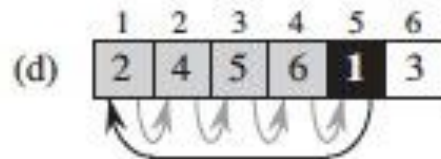
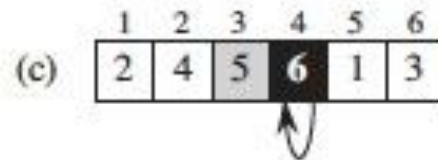
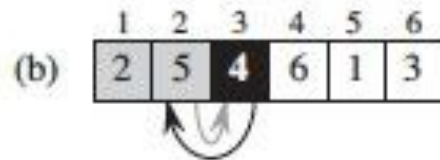
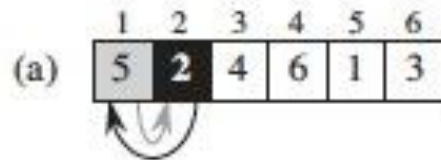
# Insertion Sort

- Iterative
- Two parts in thing to be sorted
  - Sorted & unsorted part
- Steps
  - Select first element from unsorted portion
  - Place element in appropriate spot in sorted portion
  - Repeat until all elements are sorted

INSERTION-SORT( $A$ )

```
1  for  $j = 2$  to  $A.length$ 
2       $key = A[j]$ 
3      // Insert  $A[j]$  into the sorted sequence  $A[1..j-1]$ .
4       $i = j - 1$ 
5      while  $i > 0$  and  $A[i] > key$ 
6           $A[i + 1] = A[i]$ 
7           $i = i - 1$ 
8       $A[i + 1] = key$ 
```

# Insertion Sort



# Insertion Sort

- Best case  $\rightarrow O(n)$ 
  - Elements already sorted
  - Process each element once
- Worst case  $\rightarrow O(n^2)$ 
  - Elements in reversed sorted order
- Average case
  - Usually the same as worst case



# Merge Sort

- Divide and Conquer
  - Cut problem into smaller chunks
  - Solve smaller chunks first
  - Combine solved chunks
- Steps
  - Divide → Split array into 2 subarrays
  - Conquer → Recursively sort subarrays
  - Combine → Merge sorted subarrays

```
MERGE-SORT( $A, p, r$ )  
1  if  $p < r$   
2       $q = \lfloor (p + r)/2 \rfloor$   
3      MERGE-SORT( $A, p, q$ )  
4      MERGE-SORT( $A, q + 1, r$ )  
5      MERGE( $A, p, q, r$ )
```

# Merge Sort and Merge Procedure

6 5 3 1 8 7 2 4

Runtime of Merge Procedure:  $O(n)$

MERGE( $A, p, q, r$ )

```
1   $n_1 = q - p + 1$ 
2   $n_2 = r - q$ 
3  let  $L[1..n_1 + 1]$  and  $R[1..n_2 + 1]$  be new arrays
4  for  $i = 1$  to  $n_1$ 
5       $L[i] = A[p + i - 1]$ 
6  for  $j = 1$  to  $n_2$ 
7       $R[j] = A[q + j]$ 
8   $L[n_1 + 1] = \infty$ 
9   $R[n_2 + 1] = \infty$ 
10  $i = 1$ 
11  $j = 1$ 
12 for  $k = p$  to  $r$ 
13     if  $L[i] \leq R[j]$ 
14          $A[k] = L[i]$ 
15          $i = i + 1$ 
16     else  $A[k] = R[j]$ 
17          $j = j + 1$ 
```

# Merge Sort

- Best case time complexity  $\rightarrow \Omega(n \log(n))$
- Average case time complexity  $\rightarrow \Omega(n \log(n))$
- Worst case time complexity  $\rightarrow \Omega(n \log(n))$

# Insertion & Merge Sort

- Insertion Sort

Better on smaller arrays with a best case of  $O(n)$  or better in conditions where the input expected is sorted

- Merge Sort

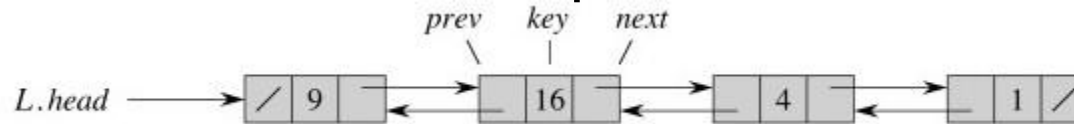
Each case: Best, Average, Worst has the same asymptotic runtime complexity of  $O(n\log(n))$

# Basic ADTs → Linked Lists

- Stores information in a linear order
- Can be sorted or unsorted
- Linked List implementation
  - Singular
    - Nodes have a next pointer pointing to next element
  - Doubly
    - Nodes have a prev (previous) and next pointer
  - Has a head and possibly tail pointer

# Basic ADTs → Linked Lists

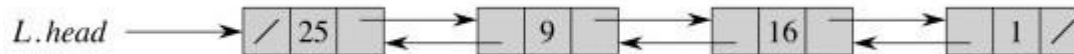
## Linked List Implementation



## Inserting 25



## Deleting 4

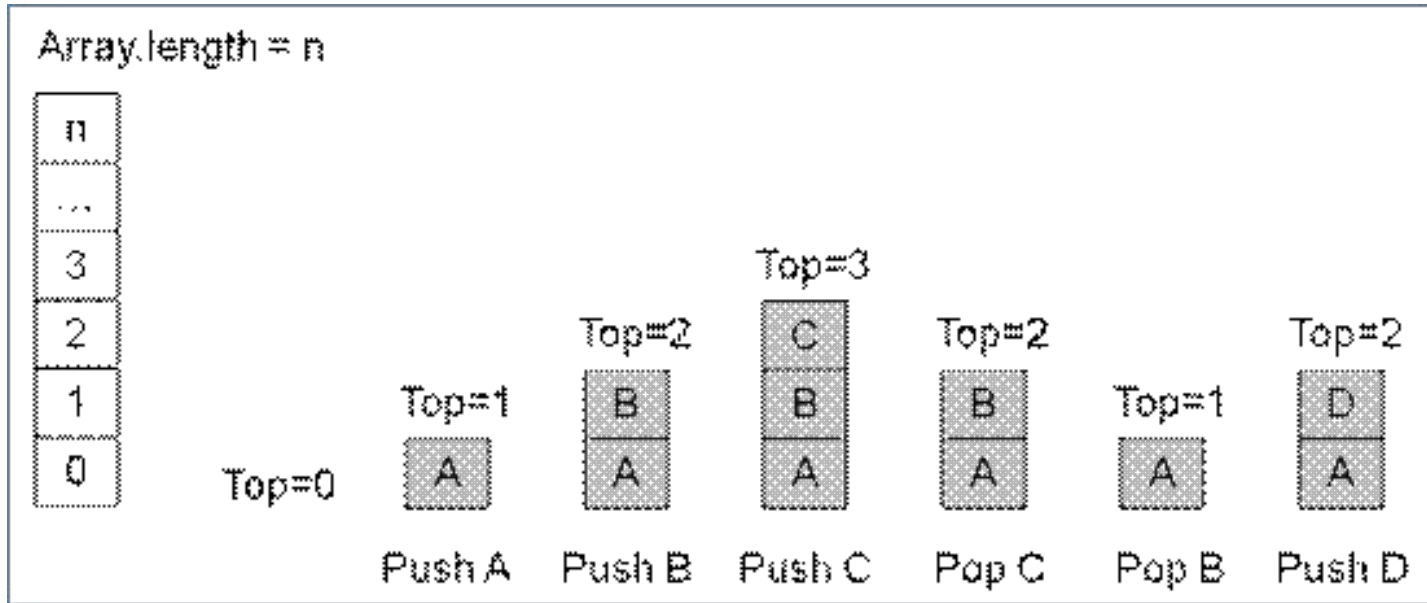


# Basic ADTs → Stacks

- `public Stack();`
  - Creates an instance of a stack and initializes
- `public void push(Object o);`
  - Inserts `o` at the top of the stack
- `public Object pop();`
  - Removes and returns top of the stack
- `public int size();`
  - Returns the number of elements in the stack
- `public boolean isEmpty();`
  - Returns true if stack is empty
- `public Object top();`
  - Returns top of stack without removal
  - Returns error if `isEmpty() == true`

# Basic ADTs → Stacks

- LIFO (last in first out)



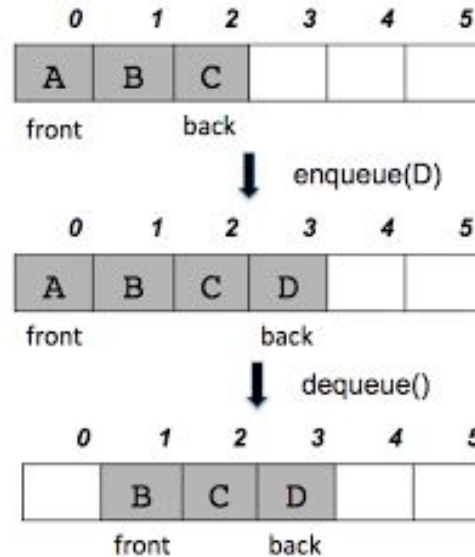


# Basic ADTs → Queues

- `public Queue();`
  - Create an instance of ADT Queue and initialize it to the empty queue.
- `public void enqueue(Object o);`
  - Insert o at the end of the queue
- `public Object dequeue();`
  - Removes & returns object at front of the queue
  - Error if queue is empty
- `public int size();`
  - Returns number of objects in queue
- `public boolean isEmpty();`
  - Returns true if queue is empty
- `public Object front();`
  - Returns front of the queue without removing it

# Basic ADTs → Queues

- FIFO (First in first out)



# Practice

- <https://www.hackerearth.com/practice/algorithms/sorting/merge-sort/practice-problems/algorithm/chandu-and-his-girlfriend/>

# Interview Problem Practice

Question: Given two sorted arrays, find the number of elements in common. The arrays are the same length and each has all distinct elements.

Let's start with a good example. We'll underline the elements in common.

A:	13	27	<u>35</u>	<u>40</u>	49	<u>55</u>	59
B:	17	<u>35</u>	39	<u>40</u>	<u>55</u>	58	60