

# Project 2 Report

COT4400

Dalton Splinter

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1.)

Aligning a set of teeth as outlined in the problem is simply the task of minimizing the max height of any single pair of teeth. The minimum height of the two interlocking extending teeth is the min height of the teeth in front of it and the min height of the teeth behind it. For the pair  $a(1)$ ,  $b(1)$  and the pair  $a(n)$ ,  $b(m)$ , their min height is simply the sum of the pair, if you extend one pair to its lower neighbor, you have to extend the opposite pair so they are the same length, the end result is these two pairs of teeth create lower bound of total height.

The inner teeth  $a(2)$  to  $a(n-1)$  and  $b(2)$  to  $b(m-1)$  min height can be found by recursively looking at the min height of the sum of the previous index and the next index.

The min height of the two arrays would be the max(all pairs of teeth) after extending them all.

2.)

The base cases would be when you've recursed through both arrays and found the min height for every position of both upper and lower sets of teeth and so that they are now both equal in length.

3.)

A 2D array to store the sums of each pair with its opposing potential pairs.

4.)

Wrapper(a, b)

    Create Array() with height set to -1

    Return (Min(dynamic(arr[i, j]))

Dynamic(arr[i, j])

    If (arr[i, j] == -1)

        if ((i == 0 && j == 0) || (i == n && j == m))

            arr = max(edge pairs)

        else

            arr[i, j] = min(arr[i, j..m], arr[i...n, j], arr[i...n, j...m])

    Else

        Return arr[i, j]

5.)

The worst case for the algorithm above would be  $O(n * m)$ . You should only need to do a constant amount of operations per index of the matrix.

6.)

NewTopArray

NewBotArray

Iterative()

    Height = 0

    While you haven't reach the end of both arrays

        If top array leads to a smaller equal height

            Increment top array

        If bot array leads to a smaller or equal height

            Increment bot array

        If neither do

            Set new min height

            continue

    Push the current value of the two arrays to the new arrays

7.)

The iterative algorithm can be optimized to only use specific columns or rows that are being looked at and does not need to store the whole matrix in memory at once.

8.)

Algo\_TeethFinding

Input: 2 arrays of upper teeth, lower teeth after extension to get the assumed optimal height of either array but not both

Output: index of individual teeth that need to be extended so both array can have the same size

```
Int Algo_TeethFinding(arr1, arr2){
    Find out which array has larger size
    Find min value in smaller arr
    Return index of that min value
}
Algo TeethExtending(arr1, arr2){
    If(length of array1 doesn't equal length of array2)
    {
        Index = Algo_TeethFinding(arr1, arr2); // get position of teeth to be extend:
        Have all position from index of array with smaller size fill with value of
        position[index] in smaller array until both array are same size
    }
    End when both array same size after extension
}
```