National University of Computer and Emerging Sciences, Lahore Fall Semester 2012

| Course: CS204- <u>DATABASE SYSTEMS</u> | | | Time Allowed: 90 min. | | |
|---|--|--|--|--|--|
| Max Points: 45 | | | | | |
| Midterm 2 | | | | | |
| Section: | Name: | | Roll No: | | |
| a) List the possibb) Show the closc) Suppose that | ition R(V, W, X, Y, Z) with le keys for relation R bas ure for attribute X given t | ed on the FDs above. he FDs above. d into two relations, R1(| $\{X \rightarrow V, VW \rightarrow X\}$. $\{V, W, X\}$ and $\{R2(X, Y, Z)\}$. Is this | | |

| Question 2 (5 points) Given relation R(W, X, Y, Z) and set of FDs F = $\{X \rightarrow W, WZ \rightarrow XY, Y \rightarrow WXZ\}$. Compute the minimal cover for F. |
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| |
| Question 3 (5 points) |
| Given relation R(W, X, Y, Z) with set of FDs F = {Z \rightarrow W, Y \rightarrow XZ, XW \rightarrow Y }. The possible keys are {Y}, {X,Z}, {W,X}. Identify the best normal form that R satisfies (1NF, 2NF, or 3NF). Justify your answer. If the relation is not in 3NF, decompose it until it becomes 3NF. |
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| |

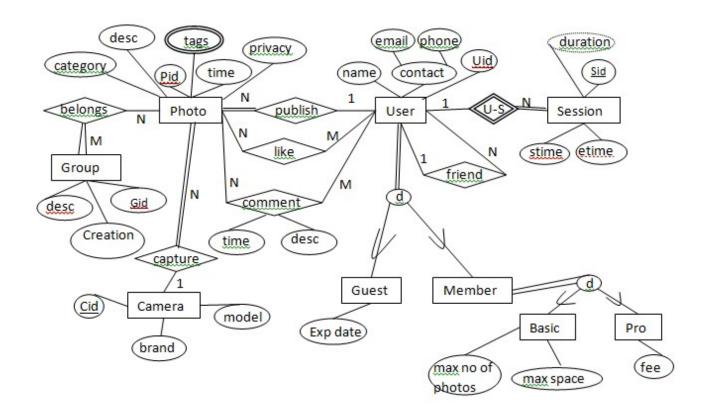
Question 4 (15 points)

You have to design an EER model for an online bookstore. The bookstore consists of books and magazines. Each book is described by a unique ISBN number, a title, its author(s), its publisher, price and year of publication. Each (issue of) magazine is described by a unique name, date of publication, its publisher, price and name of the editor-in-chief. The bookstore stock items in a warehouse and records the warehouse address, phone and unique code. It also records the number of particular item present in a warehouse.

Bookstore keeps track of its customer and for each customer it maintains customer's name, a phone, an address, and unique email. The customer can add bookstore items to a shopping basket and each basket has unique id. Recently, the bookstore adds music cassettes and compact disks to its collection. The same music item may be present in cassette or compact disk format, with differing prices. A shopping basket may contain any combination of books, magazines, music cassettes, or compact disks.

Question 5 (10 points)

Map the EER diagram given below, to relational schema. Clearly specify all the primary keys, foreign keys, and other constraints.



National University of Computer and Emerging Sciences, Lahore Spring Semester 2014

Course: CS204- DATABASE SYSTEMS Time Allowed: 90 min.

Date: 16 April, 2014

Max Points:

SECTI ON(CS-A, B, C, D) Midterm 2

| Section: Name: | Roll No: |
|---|--|
| Question 1(5 points) | |
| Consider the relation R(A, B, C, D, E, F, G, H), with FDs { Boa. Find all the keys for this relation R. (you don't need to lib. Identify the best normal form that R satisfies (1NF, 2 answer. c. If R is not in BCNF, decompose it into a set of BCNF relationship. | ist super keys that are not keys.) 2NF, 3NF, or BCNF). Justify your |
| | |
| | |
| | |
| Question 2 (5 points) | |
| Consider the relation $R(X, Y, Z, W)$, with FDs { $XY \rightarrow Z, XY$ a. Find all the keys for this relation R. (you don't need to lib. Identify the best normal form that R satisfies (1NF, 2 answer. | ist super keys that are not keys.) |
| c. If R is not in BCNF, decompose it into a set of BCNF rela | tions. |
| | |
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| | |

| Question 3 (points) |
|---|
| Given the following relation instance. |
| X Y Z |
| 1 4 2 1 5 3 1 6 3 3 2 2 |
| Which of the following functional dependencies are satisfied by the instance? (a) $XY \rightarrow Z$ and $Z \rightarrow Y$ (b) $YZ \rightarrow X$ and $Y \rightarrow Z$ (c) $YZ \rightarrow X$ and $X \rightarrow Z$ (d) $XZ \rightarrow Y$ and $Y \rightarrow X$ |
| |
| |
| |
| |
| Question 4 (points) |
| Consider a relation scheme $R = (A, B, C, D, E, H)$ on which the following functional dependencies hold: $\{A->B, BC->D, E->C, D->A\}$. What are the candidate keys of R? |
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| Oues | tion | 5 (| points) | |
|------|-------|-----|---------|--|
| Ques | LIUII | 9 (| pomaj | |

| | | | _ |
|------------------------------|------------------------|------------------|------------------------------------|
| Consider two sets of functi | innal danandancias for | a diven Relation | R - (A R C D F F) |
| Consider two sets of fulleti | ional dependencies for | a given nelation | $\Gamma = (\Gamma, D, C, D, L, I)$ |

$$F = \{ A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow F \}$$

$$G = \{ A \rightarrow CD, E \rightarrow AF \}$$

Is F equivalent to G?

Question 6 (points)

Find a minimal cover for the following set of functional dependencies: Relation R = (A, B, C, D, E, F)

$$F = \{ A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow F \}$$

| Question 7 (points) |
|---|
| Consider the following (incomplete) schedule S : R1(X), R1(Y), W1(X), R2(Y), W3(Y), W1(X), R2(Y) |
| Can you determine the serializability graph for this schedule? Assuming that all three transactions eventually commit, show the serializability graph. |
| |
| |
| |
| |
| |
| Question 8 (points) |
| Determine whether each schedule is strict, cascadeless, recoverable, or nonrecoverable Determine the strictest recoverability condition that each schedule satisfies and show your working. |
| S1: R1(X), R1(Y), W1(X), R2(Y), W3(Y), a3,W1(X), R2(Y), c1, c2 |

FAST – National University of Computer and Emerging Sciences, Lahore Fundamentals of Database Systems CS213 – Spring 2015 Midterm 2 Solution

Time allowed: 1 hour

Date: 2-May-2015 Total Marks: 22

Question 1: [Marks : 4+5=9]

a) Consider the following database state and discuss all integrity constraints violated by each operation, if any

Department(Name is the key)

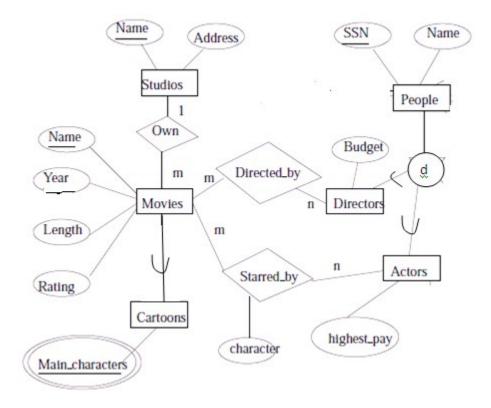
| Name | Block |
|-------------|-------|
| CS | A-1 |
| Management | A-2 |
| Engineering | A-3 |

Student(RollNum is the key) Department Name is foreign key

| FirstName | LastName | RollNum | Department_Name |
|-----------|----------|---------|-----------------|
| Imtiaz | Umar | 20 | CS |
| Aftab | Ahmad | 21 | Engineering |

- I. Insert < Mechnical, A-4> into Department Successful
- II. Insert < Aftab, Ahmad, null, Civil> into Student Entity Integrity Constraint violated
- III. Update RollNum to 21 where FirstName is Imtiaz in Student Key constraint violated
- IV. Delete row where Department is Engineering Refrential integrity constraint violated.

b) Give the relational model corresponding to the following EER diagram



Question 2: [Marks: 2+2+3+3+3=13]

Consider the following relational model and write SQL queries to answer the following questions:

Supplier

| Supplier_id | sname | statu | city |
|-------------|-------|-------|------|
| | | S | |

Part

| <u>Part_id</u> | pname | color | city |
|----------------|-------|-------|------|

Shipment

| Supplier_id | Project_id | <u>Part_i</u> | quantity |
|-------------|------------|---------------|----------|
| | | <u>d</u> | |

Project

| Project id | proiname | city |
|------------|----------|------|
| Project id | promame | CILY |

a) Give a unique and ordered list of names of suppliers who supply some red colored part to project "projX'.

Select distinct s.sname

From supplier s inner join shipment sh on sh.supplier_id=s.supplier_id inner join part p on sh.part_id=p.part_id inner join project pr on pr.project_id=sh.project_id
Where p.color='red' and pr.prohname='projX'
Order by s.sname

b) Get the ids of suppliers who supply both part with part_id 6 and part with part_id 11. Select supplier id

From shipment where part_id=6

Intersect

Select supplier_id

From shipment where part_id=11

c) Get names of suppliers who have supplied the same total quantity of part_id 7 as the supplier with supplier id 36 but have a status lower than that of supplier with id 36.

Select s.sname

From supplier s, shipment sh

Where s.supplier id=sh.supplier id and sh.part id=7

and s.status<(select status from supplier where supplier_id=36)

Group by s.sname

Having sum(sh.quantity) =(select sum(quantity) from shipment where supplier_id=36 and part_id=7)

```
Select s.sname, s.status
From supplier s, shipment sh
Where s.supplier_id=sh.supplier_id and sh.part_id=7
Group by s.sname, s.status
Having sum(sh.quantity) =(select sum(shi.quantity) from shipment shi, supplier si where shi.supplier_id=si.supplier_id and shi.supplier_id=36 and shi.part_id=7 and si.status>s.status)
```

d) Get project_ids of projects supplied entirely by supplier with id 88.

```
Select project_id
From shipment where supplier_id=88
minus
Select project_id
From shipment where supplier_id<>88
```

e) The following query gets the part_id and part name of parts supplied to all projects in London. Please fill in the blanks:

```
select p.part_id, p.pname
from part p
where _not exists
(
  select * from project_
  where city='London' and project_id not in
    ( select sh.project_id_ from shipment sh
      where p.part_id=sh.part_id_
    )
)
```

Roll No. Name Section

National University of Computer and Emerging Sciences, Lahore Campus



Course: **Database Systems** Program: **BS(Computer Science)**

Duration: 60 Minutes Paper Date: 12-Apr-18 Section: ALL

Exam: Midterm-2 Course Code: CS203 Semester: Spring 2018

Total Marks: 40 Weight 15% Page(s): 5

Instruction/Notes:

Scratch sheet can be used for rough work however, all the questions and steps are to be shown on question paper. No extra/rough sheets should be submitted with question paper. You will not get any credit if you do not show proper working, reasoning and steps as asked in question statements.

Q1. (10 points) Consider a relation with schema R(A, B, C, D), with FDs $F = \{BC \rightarrow A, AD \rightarrow B, CD \rightarrow B, CD$ $AC \rightarrow D$.

Assume possible keys of this relation are {BC}, {CD}, and {AC}. Identify the best normal form that R satisfies (1NF, 2NF, 3NF, or BCNF). Justify your answer. If R is not in BCNF, decompose it into a set of BCNF relations and show your steps. Indicate which dependencies if any are not preserved by the BCNF decomposition.

ANSWER:

Highest NF= 3NF; due to violation of FD2: AD → B BCNF relations schemas are R1(A, C, D) & R2(A, D, B); but FD1: $BC \rightarrow A$ & FD3: $CD \rightarrow B$ are lost.

| Roll No. | Name | Section | |
|----------|------|---------|--|
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Q2. (4+6=10 points) Consider the relation schema R(A, B, C, D), with FDs $F = \{AC \rightarrow B, B \rightarrow A, BD \rightarrow C, D \rightarrow A\}$.

a) Which of the following FDs may or may not hold over schema R? Justify your answer.

i)
$$CD \rightarrow B$$

ii)
$$AC \rightarrow D$$

iv)
$$A \rightarrow B$$

b) Find all the candidate keys for this relation R (You do not need to list superkeys that are not keys). Provide proper reason.

ANSWER:

- a) i. Hold ii. Not Hold iii. Hold iv. Not Hold.
- b) Candidate keys are {BD} and {CD}.

| Roll No | Name | Section |
|---|---|--|
| | sider the relation schema $R(A, B, C, D)$ minimal cover of F (i.e. F_c). |), with FDs $F = \{AB \rightarrow CD, C \rightarrow A, AD \rightarrow C, CD\}$ |
| | | |
| ANSWER: $F_c = \{AB \rightarrow CD, C$ i.e. $F_c = \{AB \rightarrow CD\}$ | →A, AD→C, CD →AB , D →B} , C →A, D →B} | |

| Roll Nol | Name | Section |
|---|--|---|
| Q4. ONLY FOR SECTION (A | A, B, E, F) (3+3+4= 10 point. | s) |
| must be non-serial). | | schedule, if possible (your schedule v3(B), c3. |
| | Hule of three transactions T1, T2 $2(Y)$, r2 (Z) , r1 (Z) , c3, c2, c1. | ?, and T3. |
| serializable or not. If the schedu | | State whether this schedule is (conflict) ne equivalent serial schedule(s) |
| | hy are they important? What a | re transaction commit points, and why |
| | | |
| a) Given these transactions, fin must be non-serial). T1: r1(A), r1(B), w1(B), w1(A), or the following sched S: w3(X), r2(X), w1(Y), r2 Draw the serializability (precede serializable or not. If the schedu otherwise explain why it is not. | and a cascade-free but not strict solution. T2: r2(B), w2(B), c2. T3: was solution. Sulle of three transactions T1, T2 (Y), r2(Z), r1(Z), c3, c2, c1. Sulle is serializable, write down the | schedule, if possible (your schedule v3(B), c3. 2, and T3. State whether this schedule is (conflict) he equivalent serial schedule(s) |

ANSWER:

- a) r1(A), r1(B), r2(B), w3(B), w1(B), w2(B), w1(A), c1, c2, c3.
- b) It is conflict-serializable; Two equivalent serial schedules are T1 \rightarrow T3 \rightarrow T2 and T3 \rightarrow T1 \rightarrow T2.

Edges in graph: T3--^x-->**T2 & T1--**^y-->**T2.**

c) See text book (transaction processing chapter).

| Roll No. | Name | Section |
|----------|------|---------|
| | | |

Q4. ONLY FOR SECTION (C, D) (10 points)

Represent the following requirements as ER model also specify the constraints using **min-max notation**.

Each bank can have multiple branches, and *each branch* have multiple *types of accounts* and offers diverse types of *loans*. A bank is registered with its name and have a unique nine-digit code. The different branches of a bank have branch number that is unique within a bank. Most of the branches are recognized by their location (i.e. town, city, state). To open an account a customer must provide name, CNIC, mobile number and home phone, resident and permanent address and birthdate. A customer must know the unique account number issued by the bank to perform the basic transactions. For each account bank record the account type and balance. The account number consists of 3-digit bank code followed by 7-digit number. Same is the case with loan number. The bank also records the loan type and amount.

Every customer must have at least one account but is restricted to at most two loans at a time, and a bank branch cannot have more than 5000 loans. The customer can access the details of their accounts and loans online. The bank keeps track of the date when the customer last access their account.

National University of Computer and Emerging Sciences, Lahore Campus



| Course Name: | Database Systems | | |
|-----------------|------------------|------------------|-------------|
| Degree Program: | | Course Code: | CS2005 |
| | | Semester: | Spring 2022 |
| Exam Duration: | 60 Minutes | Total Marks: | 25 |
| Paper Date: | Mon 09-May-2022 | Weight | 150/ |
| Section: | ALL | Paga(a) | 15% |
| Exam Type: | Midterm-2 | Total Out of | |
| | | Total Questions: | 5 |

Nabecha Mudasar Name:

Roll No:

BSSE4A Section:

Instruction/Notes:

Scratch sheet can be used for rough work however, all the questions and steps are to be shown on question paper. No extra/rough sheets should be submitted with question paper. You will not get any credit if you do not show proper working, reasoning and steps as asked in question statements.

Q1. (5 points) Consider a relation R (A, B, C, D, E, H, K, L), with the set of FDs F= {A \rightarrow BL, B \rightarrow CE, D \rightarrow BK, K \rightarrow D}. What are the keys of this relation? Prove it.

A > B A > 1 B > C B > E D > B D > K K > D

essential altributes = A, H

maybe = B,D,K

non essential = L,C,E

AH+ = {A,H,B,L,C,E} not key

candidate

AHDT = 2A,H,D,B,L,C,E,K3 -> Key AHK+ = 2 A, H, K, B, L, C, E, D3 -> candidate key

AHB+= 名A,H,B,L,C,E松,等3 > candidate key

candidate keys = 3(A,H,D), (A,H,KZ, (A,H,B)

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= { (AHD), (AHK) = (AHB) 1016

Q2. (5 points) Consider the relation schema R (A, B, C, D, E, H), with FDs F= $\{A \rightarrow BC, B \rightarrow CE, A \rightarrow E, AC \rightarrow H, D \rightarrow B\}$. Find a minimal cover of F (i.e. Fe).

A >B&, A >C, B > C, B >E, A >E, AC>H, P>B

Step2 Remove redundances on LHS.

A+= ?A,B,C,E,#3 :: C,A both necessary.

No redundancies on left

<u>Step 3</u> Remove redundancies on right.

 $A \rightarrow B$, A + = A, C, E, H <math> not redundant

 $A \rightarrow C$, $A^{\dagger} = \S A, B, C, E, H \S \rightarrow [redundant]$

 $B \rightarrow C$, $B^{+} = \S B; E, \S$ not redundant

 $B \rightarrow E$, $B^{+} = \frac{9}{2}B, C, \frac{3}{3}$ not redundant $A \rightarrow E$, $A^{\dagger} = {}^{2}A, B, C, E, H 3 \rightarrow [redundant]$

AC>H, AC+= \(\frac{2}{4} \, \chi_{\beta} \, \beta \) not redundant D-B, D+ = & D, & 3 not redundant

redundant depencies = A > C, A > E F=minimal cover: (A -B,B -C,B -E, AC>H, D -B)

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Page 2 of 5

Q3. (5 points) Consider a relation R (A, B, C, D), with the set of FDs F = $\{A \rightarrow B, B \rightarrow C, C \rightarrow D, D \rightarrow A\}$. Show the relation state that must hold all these FDs.

| A | B | C | D |
|---|---|---|---|
| 1 | 4 | 5 | 2 |
| 2 | 3 | 6 | 1 |
| 2 | 3 | 6 | 1 |

5

Q4. (5 points) Consider the relation R (A, B. C, D, E), with FDs $\{AC \rightarrow B, DE \rightarrow B, C \rightarrow E\}$. Key is $\{ACD\}$. State which of the following decompositions of R relation are lossless decomposition. Prove it.

- a. R1(A, C, B), R2(A, C, D), and R3(C, E).
- b. R1(A, C, B), R2(B, D, E), and R3(C, E).

$$\bigcirc$$
 $(R_1 \bowtie R_2) \bowtie R_3$

R, and R2 can be joined using Ac which is the candidate key For R1.

Then this table can be joined with R3 based on C.

So this is lossless decomposition.

This is not lossiess decomposition.
There is no common (candidate) key
there is no common (RINR3) and R2. to join them
between (RINR3) and R2.

Roll No: 201-1080

Q5. (5 points) Consider the relation schema R (A, B, C, D, E, H), with FDs F= {AB \rightarrow C, CD \rightarrow AE, E \rightarrow H}. Keys are {ABD} and {BCD}. Identify the best normal form that R satisfies (1NF, 2NF, 3NF, or BCNF). Justify your answer. If R is not in BCNF, decompose it into a set of BCNF relations and show your steps. Indicate which dependencies if any are not preserved by the BCNF decomposition.

AB -> C CD -> AE E -> H AB > C CD > A CD > E E > H

Prime attributes = A,B,D,C

Non Prime = E, H

Highest NF = 1NF.

-> Check INF = yes (no multivalued attributes)

no.
Partial dependency exists: CD → E -) Check 2NF = no.

(hence cannot be in 2NF, so INF).

· Removing Partial dependencies: RI(ABC), R2(ACD) R3(C,D,E,H)

-> check 3NF = no. Transitive dependency exists in R3

· Removing Transitive dependency:

RI(A,B,C) R2(A,C,D) R3(C,D,E) R4(E,H)

(ABD) (BCD) (ACD) (C,D,E) (E,H)

Joheck BONF = yes. Left sides are super keys.

Final BCNF = R1 (ABC) R2(A,C,D) R3(G,P,E) R4(EH)

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Lost dependency = & None

Page 4 of 5

National University of Computer and Emergin | Course: | Database Systems

Course: Database System
Program: BS (CS, DS, SE)
Go Minutes
Mon 10-Apr-2023
ALL

Exam:

BS (CS, DS, SE)
60 Minutes
Mon 10-Apr-2023
ALL
Midterm-II

Course Code: CS2005
Semester: Spring 2023
Total Marks: 25
Weight 15%
Page(s): 1

Instruction/Notes:

Solve the questions in the given order.

You will not get any credit if you do not show proper working, reasoning, and steps as asked in the question statements.

- Q1. (5 points) Consider a relation R (A, B, C, D, E, F), with the set of FDs $F = \{A \rightarrow BC, B \rightarrow D, CF \rightarrow E, E \rightarrow F\}$. Find all possible keys of this relation? Prove it.
- Q2. (5 points) Consider two sets of FDs, F and G, $F = \{A \rightarrow BC, B \rightarrow CD, C \rightarrow DE\}$ and $G = \{A \rightarrow B, B \rightarrow C, C \rightarrow D, D \rightarrow E\}$. Are F and G equivalent? Prove it.
- Q3. (5 points) Find a minimal cover of $F = \{A \rightarrow BCD, BC \rightarrow DE, D \rightarrow E\}$. Show all steps.
- Q4. (5 points) Consider the relation schema R (A, B, C, D), with FDs F= {AB → C, BC → D, AD → B}. Keys of this relation R are {AB} and {AD}.Identify the best normal form that R satisfies (1NF, 2NF, 3NF, or BCNF). Justify your answer. If R is not in BCNF, decompose it into a set of BCNF relations and show your steps. Indicate which dependencies if any are not preserved by the BCNF decomposition.
- Q5. (5 points) Consider the relation R (A, B, C, D, E), with FDs F= {A → BC, B → CD, C → DE, D → E}. Identify the best normal form that R satisfies (1NF, 2NF, 3NF, or BCNF). Justify your answer. If R is not in 3NF, decompose it into a set of 3NF relations and show your steps.



National University of Computer and Emerging Sciences Lahore Campus

| Database S | Systems | (CS2005) |
|------------|---------|----------|
|------------|---------|----------|

Date: Fri, 05 April 2024 Course Instructor(s)

Sessional-II Exam

Total Time (Hrs.): **Total Marks: Total Questions:**



| D-UAI- | Section | Student Signature |
|----------------------|------------|-------------------|
| Roll No | Section | |
| Do not write below t | this line. | |

CLO # 3

Noldim

Q. No 1: Consider a relation R (A, B, C, D, E, F), with the set of FDs F= $\{AB \rightarrow C, CD \rightarrow E, EF \rightarrow A, BC \rightarrow D, CD \rightarrow E, EF \rightarrow B, CD \rightarrow E, EF \rightarrow E, \rightarrow E, E$ DE-F). Find all possible keys (i.e. candidate keys) of this relation? Prove it. [5] CK= (AB, BC, BDE, BEF

Q. No 2: Consider the relation schema R (A, B, C, D, E), with FDs $F = \{A \rightarrow BC, BCD \rightarrow E, BC \rightarrow D, A \rightarrow D\}$. Find a minimal cover of F (i.e. Fc). [5]

CLO # 3

Q. No 3: Consider two sets of FDs, F and G, $F = \{A \rightarrow BC, B \rightarrow D, C \rightarrow E, D \rightarrow E\}$ and $G = \{A \rightarrow BC, B \rightarrow D, C \rightarrow E, D \rightarrow E\}$ $C \rightarrow E$, $BD \rightarrow E$, $A \rightarrow D$ }. Are F and G equivalent? Prove it. [5]

010#3 Not Equivalent & covers (7) but Go diesnot cover (F)

Q. No 4: Consider the relation R (A, B. C, D, E), with FDs (AC > B, D > E). State which of the following decompositions of R relation are lossless decomposition. Prove it. [5] Ley = (ACD)

a. R1(A, C, D), R2(A, B, C), and R3(D, E).

b. R1(A, C, D), R2(A, B, C), and R3(D, E).

b. R1(A, B, D), R2 (A, B, C), and R3(D, E).

CLO # 3

Q. No 5: Consider the relation schema R (A, B, C, D, E), with FDs F= $\{AB \rightarrow C, BC \rightarrow D, D \rightarrow E, AE \rightarrow B\}$ Keys of this relation are AB, AD, and AE. Identify the best normal form that R satisfies (1NF, 2NF, 3NF, or BCNF). Justify your answer. If R is not in BCNF, decompose it into a set of BCNF relations and show your steps. Indicate which dependencies if any are not preserved by the BCNF decomposition.

[5] Spring 2024

ANBC, BLANE, BLAD, AND AUDI ANB, ANC, BUDNE, BUND, AND stepas ATB, ATC, BUTE, BUTD, AND At Adix A+= JABD X BE= (BCD) x BC+ = |BCE] X) A >B, A >C, BC-)E, BC-)D A->BC) BC-) DE L