

SOLUTION ASSIGNMENT-03 6C & 6A

PART-01

REVIEW QUESTIONS

R3.

Source port number y and destination port number x.

R11.

A timer would still be necessary in the protocol rdt 3.0. If the round trip time is known then the only advantage will be that, the sender knows for sure that either the packet or the ACK (or NACK) for the packet has been lost, as compared to the real scenario, where the ACK (or NACK) might still be on the way to the sender, after the timer expires. However, to detect the loss, for each packet, a timer of constant duration will still be necessary at the sender..

PROBLEMS

Problem 2.

As the starting byte number=500

From byte 500, we need 46 more bytes. So, the next free byte is 546. Hence, the acknowledgement that will be sent back is 546. It indicates that all bytes till 546 have been received and the next expected number is 546.

Acknowledgement number=546

Problem 3 and 4

Q3:

Adding the two bytes gives $11111010 + 01100101 = 01011111$; $00000001 + 01011111$ (wrap around) = 01100000 ; The one's complement gives 10011111 Hence the check sum is 10011111 . At the receiver, all 3 bytes including the 2 bytes and checksum are added. If no errors are introduced into the packet, then clearly the sum at the receiver will be 11111111 . If one of the bits is a 0, then we know that errors have been introduced into the packet.

Q4:

a) 20 bytes

b) 90

Problem 40

- a) TCP slowstart is operating in the intervals [1,6] and [23,26]
- b) TCP congestion avoidance is operating in the intervals [6,16] and [17,22]
- c) After the 16th transmission round, packet loss is recognized by a triple duplicate ACK. If there was a timeout, the congestion window size would have dropped to 1.
- d) After the 22nd transmission round, segment loss is detected due to timeout, and hence the congestion window size is set to 1.
- e) The threshold is initially 32, since it is at this window size that slow start stops and congestion avoidance begins.
- f) The threshold is set to half the value of the congestion window when packet loss is detected. When loss is detected during transmission round 16, the congestion window size is 42. Hence the threshold is 21 during the 18th transmission round.
- g) The threshold is set to half the value of the congestion window when packet loss is detected. When loss is detected during transmission round 22, the congestion window size is 29. Hence the threshold is 14 (taking lower floor of 14.5) during the 24th transmission round.
- h) During the 1st transmission round, packet 1 is sent; packet 2-3 are sent in the 2nd transmission round; packets 4-7 are sent in the 3rd transmission round; packets 8-15 are sent in the 4th transmission round; packets 16-31 are sent in the 5th transmission round; packets 32-63 are sent in the 6th transmission round; packets 64 – 96 are sent in the 7th transmission round. Thus packet 70 is sent in the 7th transmission round.
- i) The threshold will be set to half the current value of the congestion window (8) when the loss occurred and congestion window will be set to the new threshold value + 3 MSS . Thus the new values of the threshold and window will be 4 and 7 respectively.
- j) threshold is 21, and congestion window size is 1.
- k) round 17, 1 packet; round 18, 2 packets; round 19, 4 packets; round 20, 8 packets; round 21, 16 packets; round 22, 21 packets. So, the total number is 52.

Q1

- For each segment, the transport layer examines both the source and destination port numbers (along with the IP addresses). Here's how the segments would be demultiplexed:
1. Segment 1: Source IP: 10.1.1.1, Source Port: 55000, Destination IP: 192.168.1.2, Destination Port: 80.
 - This segment is directed to the application running on port 80 (typically a web server) on host 192.168.1.2.
 2. Segment 2: Source IP: 10.1.1.2, Source Port: 55001, Destination IP: 192.168.1.2, Destination Port: 80.
 - This segment is also directed to the web server on port 80, but the source IP and port distinguish it from Segment 1.
 3. Segment 3: Source IP: 10.1.1.1, Source Port: 55000, Destination IP: 192.168.1.2, Destination Port: 8080.
 - This segment goes to the application listening on port 8080 (commonly used for alternative web services or proxy servers).
 4. Segment 4: Source IP: 10.1.1.2, Source Port: 55001, Destination IP: 192.168.1.2, Destination Port: 443.
 - This segment is sent to the application using HTTPS (port 443)

Q2

If the probability of packet loss is 0.2, the probability that a packet will be successfully transmitted and acknowledged without retransmission is: $P(\text{success}) = (1 - 0.2) \times (1 - 0.2) = 0.64$

Thus, there is a 64% chance that a single packet will be successfully transmitted and acknowledged without retransmission.

Q3

- (a) The congestion window size determines the throughput:
 $\text{Throughput} = \text{Window Size} / \text{RTT} = (10 \times 1500 \text{ bytes}) / (100 \text{ ms}) = 150000 \text{ bytes / sec}$
- (b) When a packet loss occurs, TCP reduces the congestion window size to half. So, the throughput will approximately halve in the next RTT. Fast retransmit and fast recovery will allow the connection to recover faster than with slow start, but there will still be a noticeable decrease.
- (c) As RTT increases, the throughput decreases since the sender must wait longer for ACKs.