

## Language Processing Systems

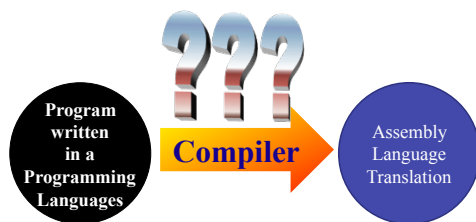
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Software Engineering Lab.  
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## Today's Outline

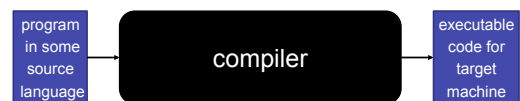
- Anatomy of a compiler
- Compiler front-end and back-end
- Regular expressions

## Anatomy of a Compiler



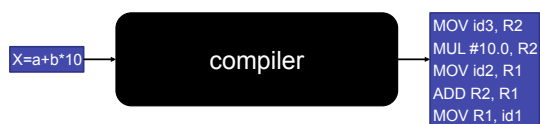
## What is a compiler?

A compiler is a program that reads a program written in one language and translates it into another language.

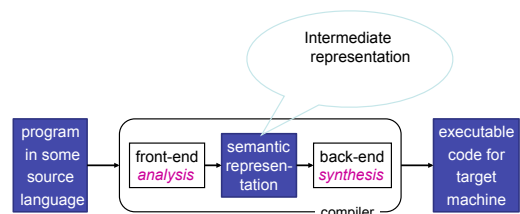


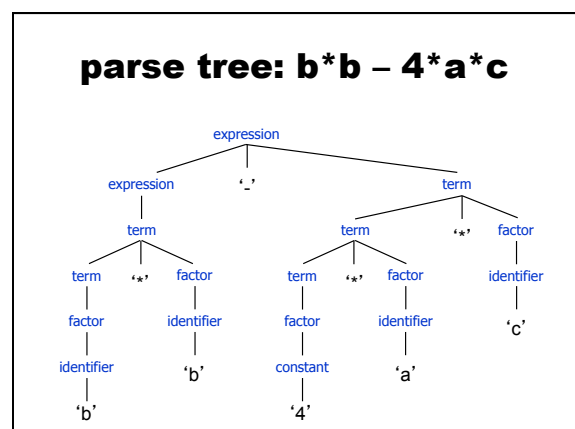
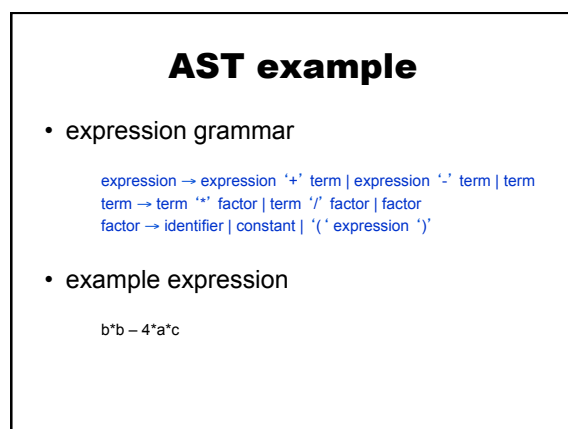
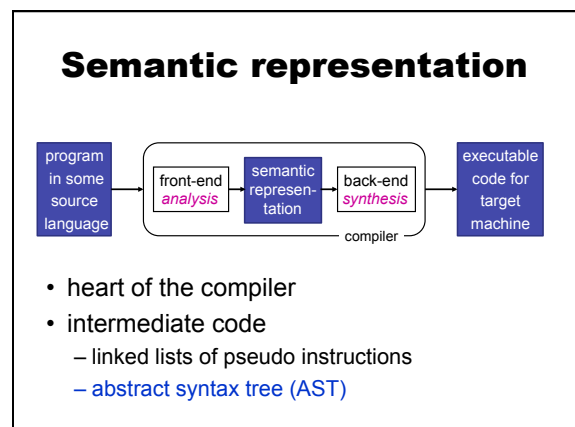
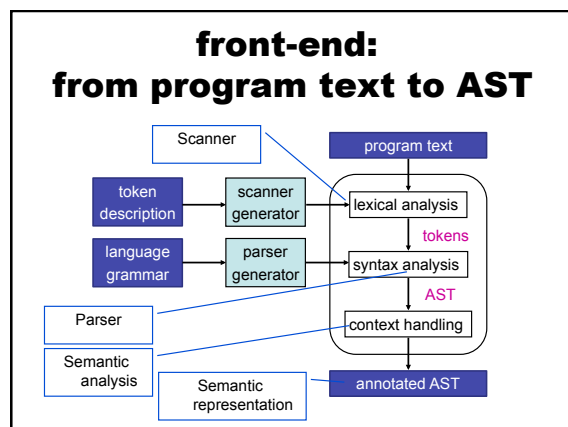
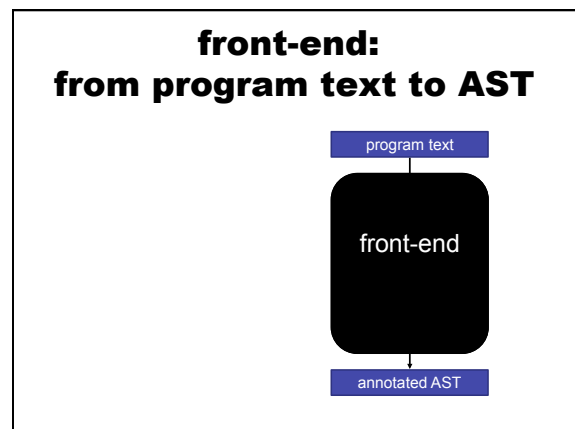
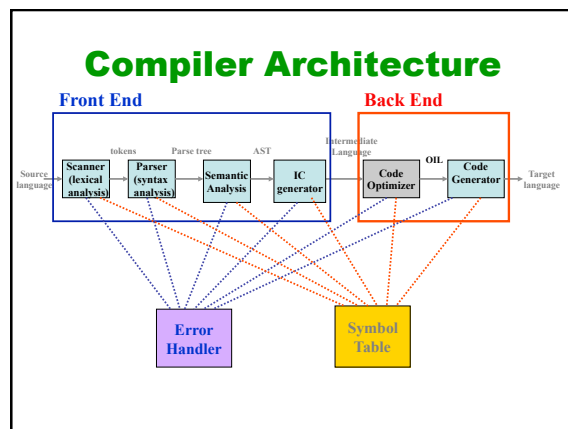
Traditionally, compilers go from high-level languages to low-level languages.

## Example

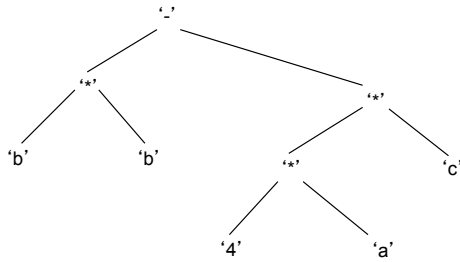


## What is a compiler?

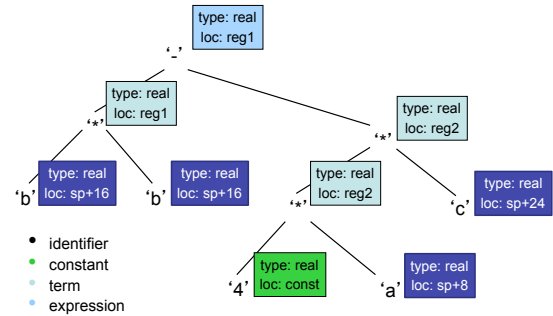




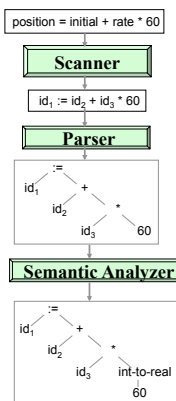
## AST: $b * b - 4 * a * c$



## annotated AST: $b * b - 4 * a * c$



### Example



## AST exercise

- expression grammar

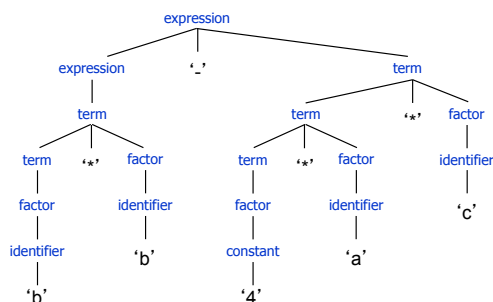
$\text{expression} \rightarrow \text{expression} '+' \text{term} \mid \text{expression} '-' \text{term} \mid \text{term}$   
 $\text{term} \rightarrow \text{term} '*' \text{factor} \mid \text{term} '/' \text{factor} \mid \text{factor}$   
 $\text{factor} \rightarrow \text{identifier} \mid \text{constant} \mid '(' \text{expression} ')'$

- example expression

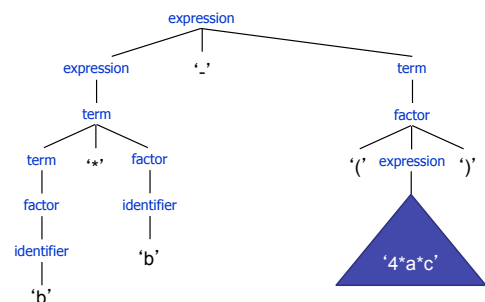
$b * b - (4 * a * c)$

- draw parse tree and AST

## answer parse tree: $b * b - 4 * a * c$



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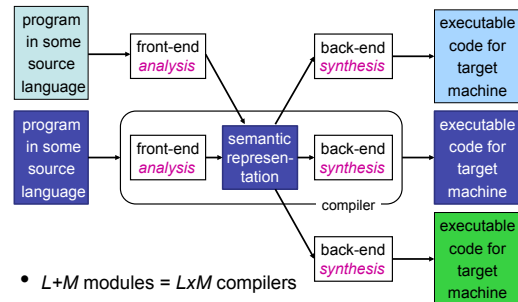


## Advantages of Using Front-end and Back-end

1. **Retargeting** - Build a compiler for a new machine by attaching a new code generator to an existing front-end.
2. **Optimization** - reuse intermediate code optimizers in compilers for different languages and different machines.

**Note:** the terms "intermediate code", "intermediate language", and "intermediate representation" are all used interchangeably.

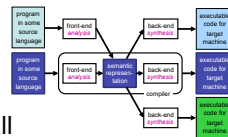
## Compiler structure



- $L+M$  modules =  $LxM$  compilers

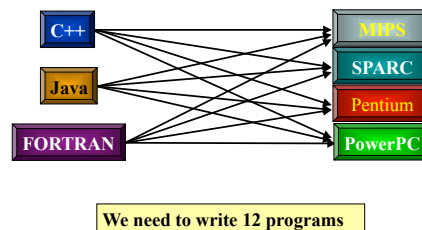
## Limitations of modular approach

- performance
  - generic vs specific
  - loss of information
- variations must be small
  - same programming paradigm
  - similar processor architecture



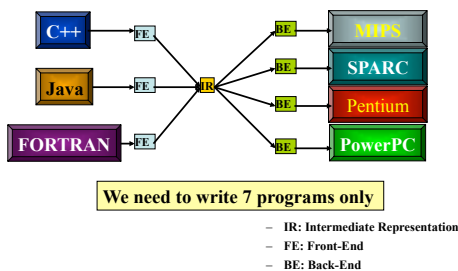
## Front-end and Back-end

- Suppose you want to write 3 compilers to 4 computer platforms:



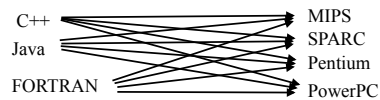
## Front-end and Back-end

- But we can do it better

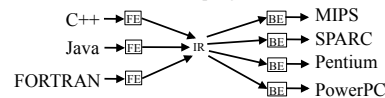


## Front-end and Back-end

- Suppose you want to write compilers from  $m$  source languages to  $n$  computer platforms. A naïve solution requires  $n*m$  programs:

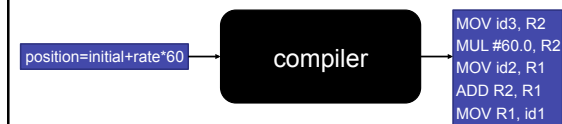


- but we can do it with  $n+m$  programs:

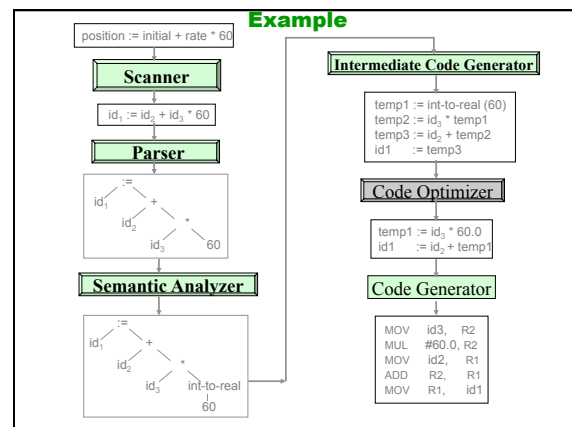


- IR: Intermediate Representation
- FE: Front-End
- BE: Back-End

## Compiler Example



## Example



## Regular Expressions

A **regular expression** is built up out of simpler regular expressions using a set of defining rules.

**Empty Set:**

$\emptyset$  A regular expression formed by **Empty set**.

**Lambda:**

$\lambda$  A regular expression formed by **Empty string**.

**Symbol:**

**a** A regular expression formed by **a**.

**Alternation:**

$M \mid N$  A regular expression formed by **M or N**.

**Concatenation:**

$(M \cdot N)$  A regular expression formed by **M followed by N**.

**Repetition:**

$(M^*)$  A regular expression formed by zero or more repetitions of **M**.

## Regular Expressions

**Operators Precedence:**

$() > * > \cdot > |$

This can simplify regular expressions.

**Example:**

$(a) \mid ((b) * (c))$  can be written as:  $a \mid b * c$ .

**Language:**

The language denoted by a regular expression **r** will be expressed as **L(r)**

**Regular expressions** allows us to define tokens of programming Languages such as identifiers and numbers.

## Regular Expressions

**Examples:**

- $a^*$  is a regular expression denotes the set  $\{\lambda, a, aa, \dots\}$
- $a|b$  is a regular expression denotes the set  $\{a\} \cup \{b\}$
- $a^*|b$  is a regular expression denotes the set  $\{\lambda, a, aa, \dots\} \cup \{b\}$
- $a^*b$  is a regular expression denotes the set  $\{b, ab, aab, \dots\}$

## Match and Create the Regular Expressions

- |                         |           |
|-------------------------|-----------|
| 1. $0(0 1)^*0$          | a. 000000 |
| 2. $((\lambda 0)1^*)^*$ | b. 01010  |
| 3. $((0 1)0(0 1))^*$    | c. 010101 |
|                         | d. 101010 |
|                         | e. 001100 |

- All strings of 0's and 1's that does not contain the substring 011

### Match and Create the Regular Expressions

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3.  $((0|1)0(0|1))^*$  c. 010101
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- e. 001100

- All strings of 0's and 1's that does not contain the substring 011
- $1^*((010)^*0^*)(\lambda|1)$

END