Distributed Systems

18. Bigtable

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Bigtable

- Highly available distributed storage
- Built with semi-structured data in mind
 - URLs: content, metadata, links, anchors, page rank
 - User data: preferences, account info, recent queries
 - Geography: roads, satellite images, points of interest, annotations
- Large scale
 - Petabytes of data across thousands of servers
 - Billions of URLs with many versions per page
 - Hundreds of millions of users
 - Thousands of queries per second
 - 100TB+ satellite image data

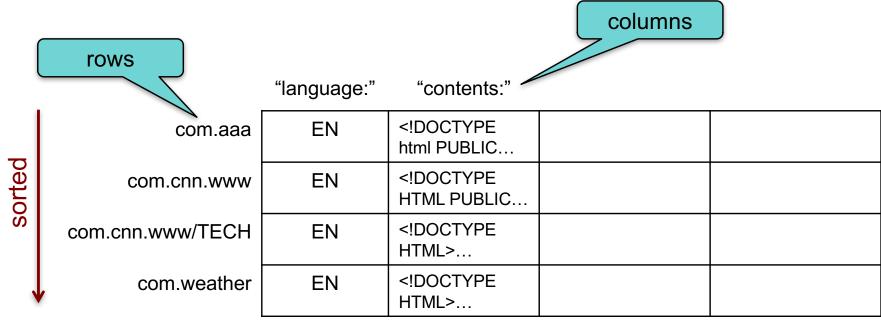
Uses

At Google, used for:

- Google Analytics
- Google Finance
- Personalized search
- Blogger.com
- Google Code hosting
- YouTube
- Gmail
- Google Earth & Google Maps
- Dozens of others...

A big table

- Bigtable is NOT a relational database
- Bigtable appears as a large table
 - "A Bigtable is a sparse, distributed, persistent multidimensional sorted map"*

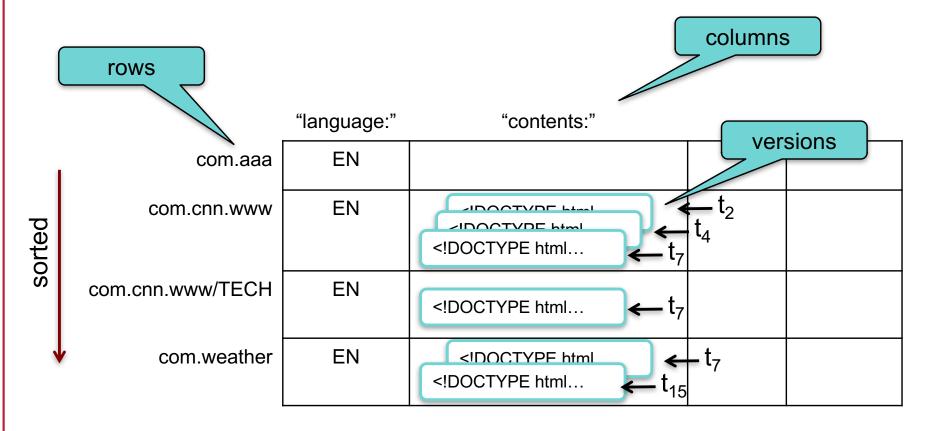


Web table example

*Bigtable: OSDI 2006

Table Model

- (row, column, timestamp) → cell contents
 - Contents are arbitrary strings (arrays of bytes)



Web table example

Tablets: Pieces of a Table

- Row operations are atomic
- Table partitioned dynamically by rows into tablets
- Tablet = range of contiguous rows
 - Unit of distribution and load balancing
 - Nearby rows will usually be served by the same server
 - Accessing nearby rows requires communication with a small # of machines
 - You need to select row keys to ensure good locality
 - E.g., reverse domain names:

com.cnn.www instead of www.cnn.com

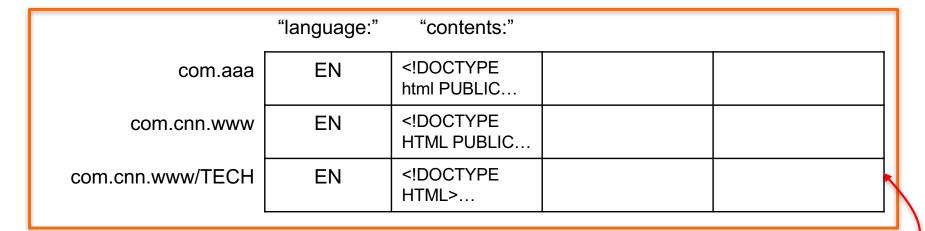
Table splitting

- A table starts as one tablet
- As it grows, it it split into multiple tablets
 - Approximate size: 100-200 MB per tablet by default

	"language:"	"contents:"	
com.aaa	EN	br html PUBLIC	
com.cnn.www	EN	br HTML PUBLIC	
com.cnn.www/TECH	EN	br HTML>	
com.weather	EN	br HTML>	

tablet

Splitting a tablet



EN	br HTML>			K
EN	br HTML>			
EN	br HTML>			
EN	br HTML>			
	EN EN	HTML> EN HTML EN HTML EN HTML	HTML> EN HTML EN HTML EN /td <td>HTML> EN <!DOCTYPE HTML> EN <!DOCTYPE HTML> EN <!DOCTYPE /td> </td>	HTML> EN HTML EN HTML EN /td

Split

Columns and Column Families

Column Family

- Group of column keys
- Column family is the basic unit of data access
- Data in a column family is typically of the same type
- Implementation compresses data in the same column family

Operations

- (1) Create column family
- (2) Store data in any key within the family
- Column families will typically be small
 - ≤ hundreds of keys; a table may have an unlimited # of column families
- Identified by family:qualifier

Column Families: example

Three column families

- "language:" language for the web page
- "contents:" contents of the web page
- "anchor:" contains text of anchors that reference this page.
 - www.cnn.com is referenced by Sports Illustrated (cnnsi.com) and My-Look (mlook.ca)
 - The value of ("com.cnn.www", "anchor:cnnsi.com") is "CNN", the reference text from cnnsi.com.

				Column family <i>anchor</i>		
		"language:"	"contents:"	anchor:cnnsi.com	anchor:mylook.ca	
sorted	com.aaa	EN	br html PUBLIC			
	com.cnn.www	EN	br HTML PUBLIC	"CNN"	"CNN.com"	
	com.cnn.www/TECH	EN	br HTML>			
	com.weather	EN	br HTML>			

Timestamps

- Each column family may contain multiple versions
- Version indexed by a 64-bit timestamp
 - Real time or assigned by client
- Per-column-family settings for garbage collection
 - Keep only latest n versions
 - Or keep only versions written since time t
- Retrieve most recent version if no version specified
 - If specified, return version where timestamp ≤ requested time

API: Operations on Bigtable

- Create/delete tables & column families
- Change cluster, table, and column family metadata (e.g., access control rights)
- Write or delete values in cells
- Read values from specific rows
- Iterate over a subset of data in a table
 - All members of a column family
 - Multiple column families
 - E.g., regular expressions, such as anchor: *.cnn.com
 - Multiple timestamps
 - Multiple rows
- Atomic read-modify-write row operations
- Allow clients to execute scripts (written in Sawzall) for processing data on the servers

Implementation: Supporting Services

GFS

For storing log and data files

Cluster management system

- For scheduling jobs, monitoring health, dealing with failures
- Google SSTable (Sorted String Table)
 - Internal file format optimized for streaming I/O and storing <key,value> data
 - Provides a persistent, ordered, *immutable* map from keys to values
 - Append-only
 - Memory or disk based; indexes are cached in memory
 - If there are additions/deletions/changes to rows
 - New SSTables are written out with the deleted data removed
 - Periodic compaction merges SSTables and removes old retired ones

See http://goo.gl/McD6ex for a description of SSTable

Implementation: Supporting Services

Chubby

- Highly-available & persistent distributed lock (lease) service & file system
- Five active replicas; one elected as master to serve requests
- Majority must be running
- Paxos algorithm used to elect master & keep replicas consistent
- Provides namespace of files & directories
 - Each file or directory can be used as a lock

In Bigtable, Chubby is used to:

- Ensure there is only one active master
- Store bootstrap location of Bigtable data
- Discover tablet servers
- Store Bigtable schema information
- Store access control lists

Implementation

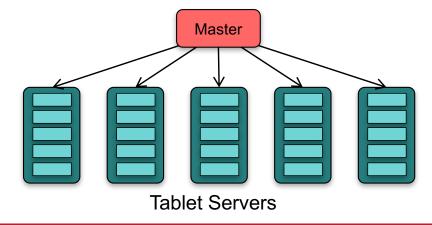
1. Many tablet servers – coordinate requests to tablets

- Can be added or removed dynamically
- Each manages a set of tablets (typically 10-1,000 tablets/server)
- Handles read/write requests to tablets
- Splits tablets when too large

2. One master server

- Assigns tablets to tablet server
- Balances tablet server load
- Garbage collection of unneeded files in GFS
- Schema changes (table & column family creation)

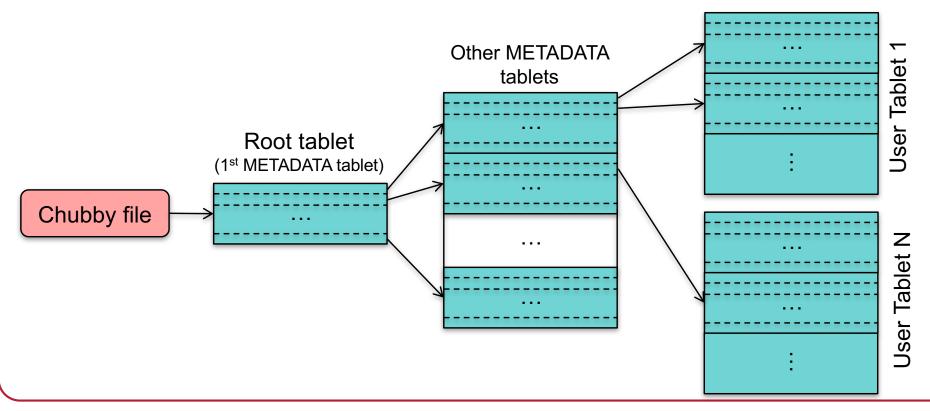
3. Client library



Implementation: METADATA table

Three-level hierarchy

- Balanced structure similar to a B+ tree
- Root tablet contains location of all tablets in a special METADATA table
- Row key of METADATA table contains location of each tablet
 f(table_ID, end_row) ⇒ location of tablet



Implementation

- Tablet assigned to one tablet server at a time
- Chubby keeps track of tablet servers
 - When tablet server starts:
 - It creates & acquires an exclusive lock on a uniquely-named file in a Chubby servers directory
 - Master monitors this directory to discover tablet servers
- When master starts:
 - Grabs a unique master lock in Chubby (prevent multiple masters)
 - Scans the servers directory in Chubby to find live tablet servers
 - Contacts each tablet server to discover what tablets are assigned to that server
 - Scans the METADATA table to learn the full set of tablets
 - Build a list of tablets not assigned to servers
 - These will be assigned by choosing a tablet server & sending it a tablet load request

Implementation

- Tablet assigned to one tablet server at a time
- When master starts:
 - Grabs a unique master lock in Chubby (prevent multiple masters)
 - Scans the servers directory in Chubby to find live tablet servers
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Fault Tolerance

- Fault tolerance is provided by GFS & Chubby
- Dead tablet server
 - Master is responsible for detecting when a tablet server is not working
 - Asks tablet server for status of its lock
 - If the tablet server cannot be reached or has lost its lock
 - Master attempts to grab that server's lock
 - If it succeeds, then the tablet server is dead or cannot reach Chubby
 - Master moves tablets that were assigned to that server into an unassigned state
- Dead master
 - Master kills itself when its Chubby lease expires
 - Cluster management system detects a non-responding master
- Chubby: designed for fault tolerance (5-way replication)
- GFS: stores underlying data designed for n-way replication

Bigtable Replication

 Each table can be configured for replication to multiple Bigtable clusters in different data centers

Eventual consistency model

Sample applications

- Google Analytics
 - Raw Click Table (~200 TB)
 - Row for each end-user session
 - Row name: {website name and time of session}
 - Sessions that visit the same web site are sorted & contiguous
 - Summary Table (~20 TB)
 - Contains various summaries for each crawled website
 - Generated from the Raw Click table via periodic MapReduce jobs

Sample applications

- Personalized Search
 - One Bigtable row per user (unique user ID)
 - Column family per type of action
 - E.g., column family for web queries (your entire search history!)
 - Bigtable timestamp for each element identifies when the event occurred
 - Uses MapReduce over Bigtable to personalize live search results

Sample applications

- Google Maps / Google Earth
 - Preprocessing
 - Table for raw imagery (~70 TB)
 - Each row corresponds to a single geographic segment
 - Rows are named to ensure that adjacent segments are near each other
 - Column family: keep track of sources of data per segment
 (this is a large # of columns one for each raw data image but sparse)
 - MapReduce used to preprocess data
 - Serving
 - Table to index data stored in GFS
 - Small (~500 GB) but serves tens of thousands of queries with low latency

Bigtable outside of Google

Apache HBase

- Built on the Bigtable design
- Small differences (may disappear)
 - access control not enforced per column family
 - Millisecond vs. microsecond timestamps
 - No client script execution to process stored data
 - Built to use HDFS or any other file system
 - No support for memory mapped tablets
 - Improved fault tolerance with multiple masters on standby



Bigtable vs. Amazon Dynamo

- Dynamo targets apps that only need key/value access with a primary focus on high availability
 - key-value store versus column-store (column families and columns within them)
 - Bigtable: distributed DB built on GFS
 - Dynamo: distributed hash table
 - Updates are not rejected even during network partitions or server failures

