

Figure 1: First Circuit

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Quantum Computing  
Problem set 7

## Problem 1

### 0.1 Part a

We have the following circuit shown in Figure 1

We start with states  $|0\rangle|0\rangle|0\rangle$ . When acting on state  $|0\rangle|0\rangle|0\rangle$  the circuit shown first acts with  $H$  on each of the  $|0\rangle$  states. Therefore at position (1) indicated by Figure 1 and reading the qubits from top to bottom, this gives

$$H^2 \otimes H^1 \otimes H^0 |0\rangle|0\rangle|0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \quad (1)$$

Where here my notation is  $H^n$  means  $H$  being applied to the  $n$ th qubit. Now we are at position (2), applying the phase gates we get

$$\left( P_{l\pi} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \right) \otimes \left( P_{l\pi/2} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \right) \otimes \left( P_{l\pi/4} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \right) \quad (2)$$

The controlled two-qubit  $Z_k$  gate is defined as  $Z_k = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/2^k} \end{pmatrix}$ , or equivalently

$$Z_k = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\pi/2^k} \end{pmatrix}$$

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acting on the two-qubit standard basis states  $\{|00\rangle, |01\rangle, |10\rangle, |11\rangle\}$  (it only applied  $e^{i\pi/2^k}$  to the target state if the control state is 1.)

Similar to  $Z_k$ ,  $P_\phi$  is defined as  $P_\phi = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix}$ , it does nothing to the first ( $|0\rangle$ ) state and applies  $e^{i\phi}$  to the second. Therefore (2) becomes

$$\left(\frac{1}{\sqrt{2}}(|0\rangle + e^{il\pi}|1\rangle)\right) \otimes \left(\frac{1}{\sqrt{2}}(|0\rangle + e^{il\pi/2}|1\rangle)\right) \otimes \left(\frac{1}{\sqrt{2}}(|0\rangle + e^{il\pi/4}|1\rangle)\right) \quad (3)$$

Now consider  $|\psi_l\rangle$ ; we use the definition

$$|\psi_l\rangle = \frac{1}{2^{3/2}} \sum_{x=0}^7 e^{2\pi i x l / 8} |x_2 x_1 x_0\rangle \quad (4)$$

Now using

$$x = x_{number} = x_0^{bin} + 2x_1^{bin} + 4x_2^{bin} \quad (5)$$

Where  $x_{bin}$  a binary number (0 or 1), therefore (4) becomes

$$|\psi_l\rangle = \frac{1}{2^{3/2}} \sum_{x=0}^7 e^{2\pi i (x_0 + 2x_1 + 4x_2) l / 8} |x_2 x_1 x_0\rangle \quad (6)$$

We can break this up as a sum on each individual qubit, i.e.

$$\begin{aligned} |\psi_l\rangle &= \left(\frac{1}{\sqrt{2}} \sum_{x_0=0}^1 e^{2\pi i x_0 l / 8} |x_0\rangle\right) \otimes \left(\frac{1}{\sqrt{2}} \sum_{x_1=0}^1 e^{2\pi i (2x_1) l / 8} |x_1\rangle\right) \otimes \left(\frac{1}{\sqrt{2}} \sum_{x_2=0}^1 e^{2\pi i (4x_2) l / 8} |x_2\rangle\right) \\ &= \left(\frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi l/4}|1\rangle)\right) \otimes \left(\frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi l/2}|1\rangle)\right) \otimes \left(\frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi l}|1\rangle)\right) \end{aligned} \quad (7)$$

This is precicely what we got in equation (3)!!! Hence we have verified the requested relation!

### 0.2 Part b

Starting with the definition of  $|\psi_l\rangle$  in (4), we have

$$|\psi_{-l}\rangle = \frac{1}{2^{3/2}} \sum_{x=0}^7 e^{-2\pi i x l / 8} |x_2 x_1 x_0\rangle \quad (8)$$

With the definition of the quantum Fourier transform,

$$U_{FT}|x\rangle_n = \frac{1}{2^{n/2}} \sum_{y=0}^{2^n-1} e^{2\pi i x y / 2^n} |y\rangle_n \quad (9)$$

Then

$$\begin{aligned}
 U_{FT}|\psi_{-l}\rangle_3 &= U_{FT} \frac{1}{2^{3/2}} \sum_{x=0}^7 e^{-2\pi i x l / 8} |x\rangle_3 \\
 &= \frac{1}{2^{3/2}} \sum_{x=0}^7 e^{-2\pi i x l / 8} U_{FT} |x\rangle_3 \\
 &= \frac{1}{8} \sum_{x=0}^7 e^{-2\pi i x l / 8} \sum_{y=0}^7 e^{2\pi i x y / 8} |y\rangle_3 \\
 &= \frac{1}{8} \sum_{x=0}^7 \sum_{y=0}^7 e^{2\pi i x (y-l) / 8} |y\rangle_3
 \end{aligned} \tag{10}$$

Aside (not completely required): Now let's review a bit on Fourier transforms to complete the calculation above. Recall that if  $f(x)$  is periodic such that  $f(x+L) = f(x)$ , then the Fourier expansion of it is

$$f(x) = \sum_{n=-\infty}^{\infty} c_n e^{2\pi i n x / L} \tag{11}$$

If we define  $I(l) = \int_{-L/2}^{L/2} e^{2\pi i x l / L} dx$  then  $I(l) = \frac{L}{2\pi i l} (e^{2\pi i l} - e^{-2\pi i l})$  if  $l \neq 0$ .

If on the other hand  $l = 0$ , then  $e^{2\pi i x l} = e^{-2\pi i x l} (= 1)$ <sup>a</sup> and  $I(l) = 0$ . If  $l = 0$  then  $I(0) = \int_{-L/2}^{L/2} 1 dx = L$ , therefore it follows that

$$\int_{-L/2}^{L/2} e^{2\pi i (n-m)x / L} dx = L \delta_{nm} \tag{12}$$

Then the coefficients of (11) can be found by multiplying (11) by  $e^{-2\pi i m x} / L$  and integrating from  $-L/2$  to  $L/2$ .

<sup>a</sup>Because  $e^{i\pi} = -1$ , therefore  $e^{i\pi l} = (e^{i\pi})^l = (-1)^l$ . Therefore,  $e^{2\pi i l} = (-1)^{2l} = 1$  for all any integer  $l$ .

Now that this Fourier transform aside is complete, we can continue computing (10) by treating it just as a math problem without having to worry about the bits into their binary representations. Doing the  $x$  sum first,

$$U_{FT}|\psi_{-l}\rangle_3 = \sum_{y=0}^7 \left( \frac{1}{8} \sum_{x=0}^7 e^{2\pi i x (y-l) / 8} \right) |y\rangle_3 \tag{13}$$

It's clear to see that if  $y = l$  the sum in (13) is a sum of 1's (8 1's times 1/8) which gives 1. Further, this is a unitary transformation so we know that all the other amplitudes (for  $y \neq l$ ) give zero. Hence clearly this sum is *(something)*  $\times \delta_{yl}$ .

In fact, (13) can just be written as

$$U_{FT}|\psi_{-l}\rangle_3 = \sum_{y=0}^7 \tilde{h}(y) |y\rangle_3 \tag{14}$$

Where  $\tilde{h}(y)$  are the phases in the expansion above ( $\tilde{h}(y) = \frac{1}{8} \sum_{x=0}^7 e^{2\pi i x (y-l) / 8} = 1$  if  $y = l$ ). The amplitude square of these factors give the corresponding probability that is required. For example the probability that the state is found in  $y = l$  is

$$\text{Prob}(y = l) = |\tilde{h}(y)|^2 = 1 \tag{15}$$

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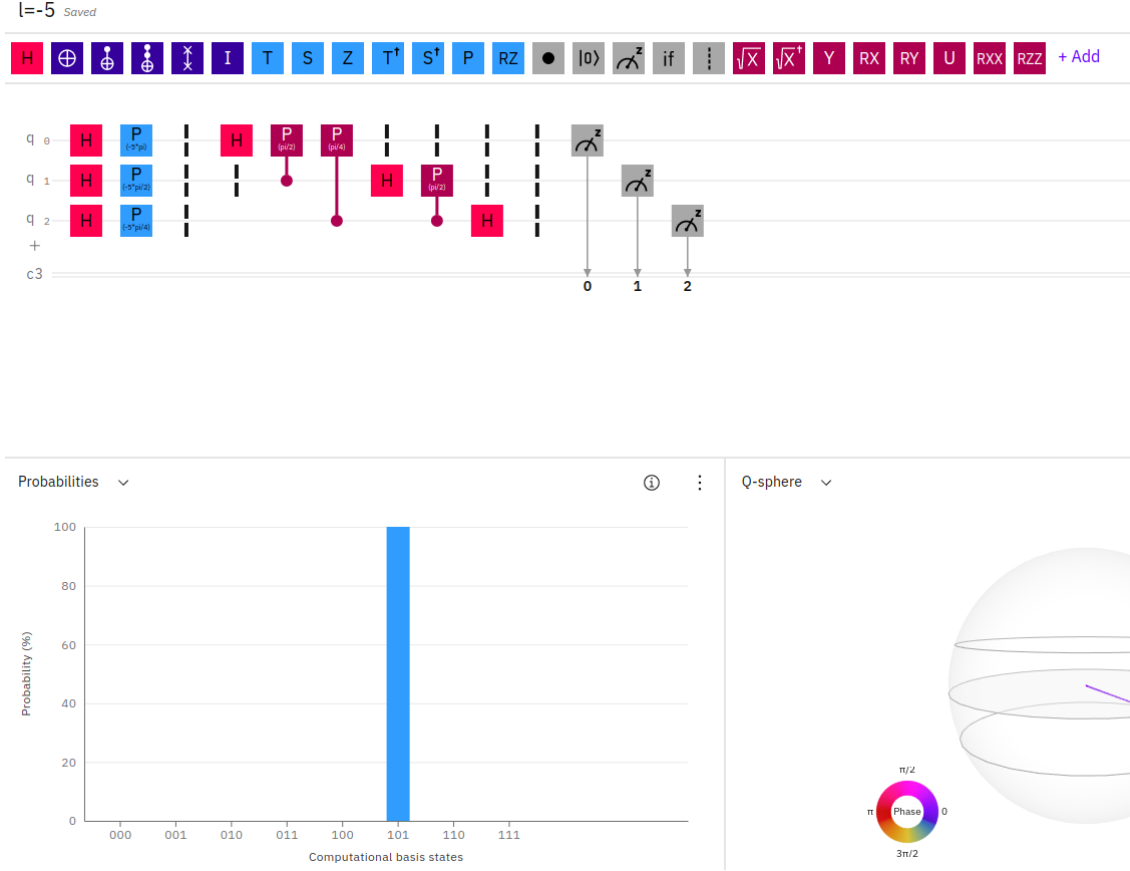


Figure 2: Circuit for  $U_{FT}|\psi_{-l}\rangle$  where  $l = -5$ . The resulting state is  $|l\rangle$  so that  $l = 5$  in binary ( $=101$ ) as expected. Note that technically  $5_{10} = 0101_2$  but we're using only 3 qubits and hence we have  $5 = 101$  in binary

Where the product 1 is calculated and discussed above. This tells us that the state above in which  $y = l$  is by definition the state  $|l\rangle$ , i.e.

$$\begin{aligned} \text{Prob}(y = l) = 1 &\implies \sum_{y=0}^7 \tilde{h}(y)|y\rangle_3 = |l\rangle \\ &\implies U_{FT}|\psi_{-l}\rangle_3 = |l\rangle \end{aligned} \quad (16)$$

## 0.3 Part c

Figure 2 shows the circuit for  $U_{FT}|\psi_{-l}\rangle$  where  $-l = -5$ . The resulting state is  $|l\rangle$  so that  $l = 5$  in binary ( $=101$ ) as expected.

Figure 3 shows the circuit for  $U_{FT}|\psi_{-l}\rangle$  where  $-l = -3$ . The resulting state is  $|l\rangle$  so that  $l = 3$  in binary ( $=011$ ) as expected.

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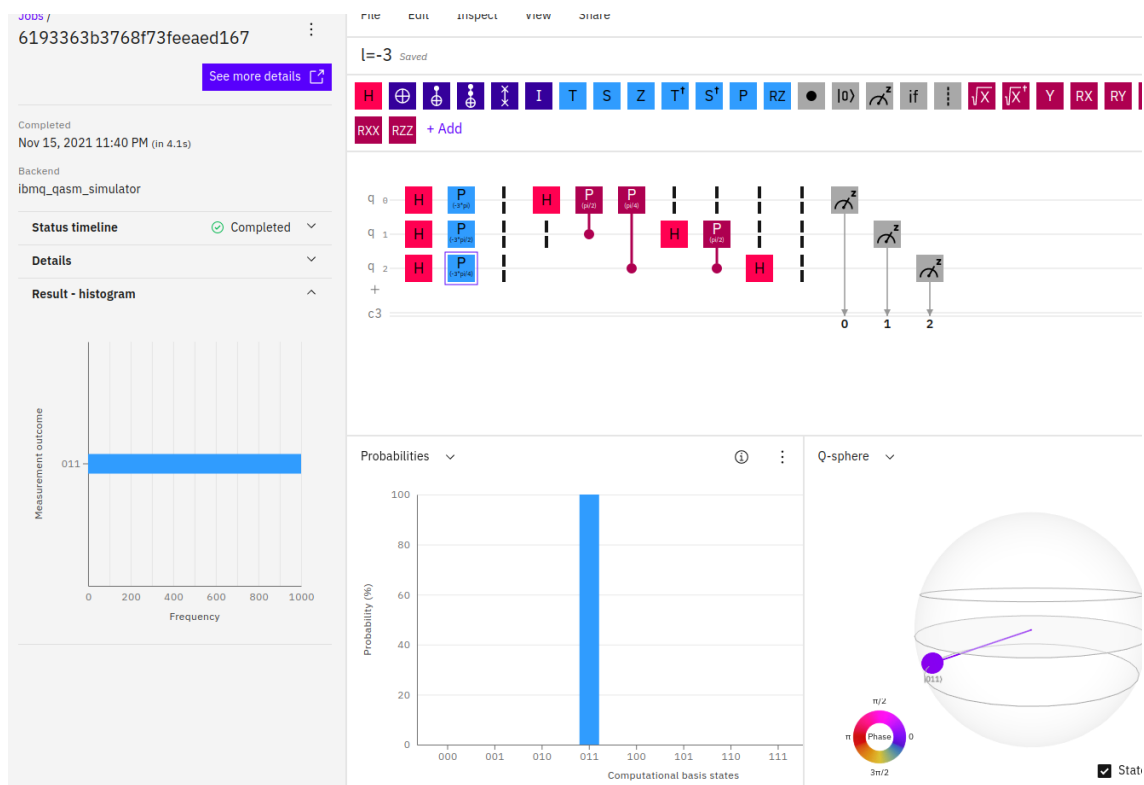


Figure 3: Circuit for  $U_{FT}|\psi_{-l}\rangle$  where  $-l = -3$ . The resulting state is  $|l\rangle$  so that  $l = 3$  in binary ( $=011$ ) as expected.

## 1 Problem 2

All of my code for this problem is provided in an external PDF document that was submitted with this write-up.

### 1.1 Part a

Bob chooses  $p$  and  $q$

$$p = 3, q = 11 \quad (17)$$

Such that

$$N = pq = 3 \times 11 = 33 \quad (18)$$

We now must choose an integer  $c < N \rightarrow c < 33$  that share no common divisors with

$$(p-1)(q-1) = 2 \times 10 = 20 \quad (19)$$

One choice that suffices the above conditions is

$$c = 17 \quad (20)$$

### 1.2 Part b

We determine integer  $d < N$  which satisfies

$$cd \pmod{(p-1)(q-1)} = 1 \quad (21)$$

This can be done quickly with the python code And we attain  $d = 13$ .

### 1.3 Part c

Picking  $a < N$  we compute  $b$  via

$$b = a^c \pmod{N} \quad (22)$$

$$\text{Choosing } a = 5 \text{ gives } b = 14. \quad (23)$$

### 1.4 Part d

Now computing

$$b^d \pmod{N} \quad (24)$$

with our choice of  $b$  from (23) gives 5, i.e. what the value we chose for  $a$ .

### 1.5 Part e

Now we wish to find the period of

$$f(x) = b^x \pmod{N} \quad (25)$$

Since  $N = 33$  is so small, we can find the period by doing a direct check of  $f(x)$ , computationally this means calculating the array of  $f(x)$  where  $x = 1, 2, \dots, 33$  and checking the minimum index in this array in which the output of the function starts repeating. My code attached does this. In this case we get

$$r = 10 \quad (26)$$

**1.6 Part f**

we calculate  $d'$  by

$$cd' \pmod{r} = 1 \tag{27}$$

and attain

$$d' = 3 \tag{28}$$

**1.7 Part h**

The value of  $r$  is even (from above), and

$$b^{r/2} \pmod{N} \neq N - 1 \tag{29}$$

therefore, the answer to both of the questions is yes, and we proceed to part (i).

**1.8 Part i**

$$\text{GCD}(b^{r/2} + 1, N) = 3 \tag{30}$$

$$\text{GCD}(b^{r/2} - 1, N) = 11 \tag{31}$$