

Real-Time Systems

Week 2
Schedulers

Outline

- ▶ Scheduling approaches
 - ▶ clock driven
 - ▶ priority driven
- ▶ Preemptive vs. Non-preemptive scheduling
- ▶ Offline scheduling
- ▶ Online scheduling

Clock Driven Approach

- ▶ Decisions about which job executes and at what time are made at predefined times.
- ▶ Typically, decision points are scheduled at regular intervals.
- ▶ When system starts, scheduler chooses and schedules the jobs that will execute until the next decision point.
- ▶ It then sets a hardware timer, and sleeps until timer expires.
- ▶ Scheduler awakes and then repeats the process.

Round Robin

- ▶ Round-robin approach is typically used to schedule time-shared applications.
- ▶ When a job becomes ready to execute, it is added to a first-in-first-out (FIFO) queue.
- ▶ The job at head of queue then executes for at most a single time slice.
- ▶ A **time-slice** is basic unit of time allocated to a job, typically on order of tens of milliseconds
- ▶ if job doesn't finish by end of time-slice, it is preempted and put at end of queue.
- ▶ If n jobs waiting, each job gets $1/n^{\text{th}}$ of the processor.

Round Robin Example

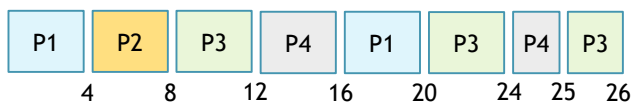
Parameters

- ▶ quantum
 - ▶ 4
- ▶ no priority based preemption
- ▶ what is the **average wait time** ?

Workload Model

Process #	Release Time	Execution Time
1	0	8
2	1	4
3	2	9
4	3	5

Round Robin Example - Solution



Process #	Release Time	Completion Time	Response Time
1	0	20	20
2	1	8	7
3	2	26	24
4	3	25	22
Avg.	-	-	18.5

Weighted Round Robin

- ▶ Weighted round-robin approach is similar but each job gets a weighted share of the processor.
- ▶ If a job has weight wt , then it is allocated wt time slots each round.
- ▶ Length of a round is equal to sum of weights of all waiting jobs.
- ▶ By changing weights, we can speed up or slow down a job.
- ▶ Round-robin methods delays completion of all job, as each is only given a fraction of the processor.
- ▶ Method thus **not suitable** for scheduling precedence constrained jobs as response time of job chain could be quite large.

Priority Driven Approach

- ▶ Priority-driven approach maintains one or more queues of jobs ready to execute, sorted by priority.
- ▶ Method never leaves a resource idle on purpose.
- ▶ If resource becomes available, then highest priority waiting job is given the resource.
- ▶ **Scheduling decisions are made when a job becomes ready or a job completes.**
- ▶ How algorithm decides priority defines algorithm.
 - ▶ Could assign priority based on release time such as FIFO, and LIFO (last-in-first-out) algorithms.
 - ▶ Could also assign priority based on job execution time such as SETF (shortest-execution-time-first), and LETF (longest-execution-time-first) algorithms.

Priority Driven Approach

- ▶ Most scheduling algorithms used in **non real-time systems** are priority-driven
 - ▶ Release time
 - ▶ First-In-First-Out
 - ▶ Last-In-First-Out
 - ▶ Execution time
 - ▶ Shortest-Execution-Time-First
 - ▶ Longest-Execution-Time-First
- ▶ Real-time priority scheduling assigns priorities based on deadline or some other timing constraint
 - ▶ Earliest deadline first
 - ▶ Least slack time first

Preemptive vs. Non-preemptive

- ▶ Not known in general when preemptive is better than non-preemptive scheduling.
- ▶ Usually though, preemptive gives better results.
- ▶ In special case when all jobs have zero release time, preemptive is better if we ignore cost of preemption.

Fixed and Dynamic Priority Algorithms

- ▶ Two classifications for algorithms that schedule periodic tasks:
 - ▶ Fixed-priority algorithms give same priority to each job in a given task.
 - ▶ Means priority of a given task is fixed relative to other tasks.
 - ▶ Dynamic-priority algorithms give different priorities to individual jobs in same task.
- ▶ In both cases, most algorithms give a fixed priority to a given job. i.e. once a job is given a priority, it is not changed.

Earliest-Deadline-First Algorithm

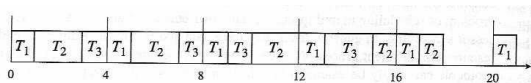
- ▶ Can choose to assign priorities based on deadlines.
- ▶ A common approach is earliest-deadline-first (EDF).
- ▶ For EDF, the earlier the deadline the higher the priority.
- ▶ **Theorem**
 - ▶ When preemption allowed and jobs do not contend for resources, then EDF algorithm is optimal.
 - ▶ Look for the proof at...

Deadline-Monotonic Algorithm

- ▶ For deadline-monotonic (DM) algorithm, tasks with shortest relative deadline are given highest priority.
- ▶ Calculate DM schedule for system with tasks
 - ▶ $T_1 = (\Phi_1, p_1, e_1, D_1) = (50, 50, 25, 100)$,
 - ▶ $T_2 = (0, 62.5, 10, 20)$,
 - ▶ $T_3 = (0, 125, 25, 50)$.
- ▶ What is utilization of the system !?

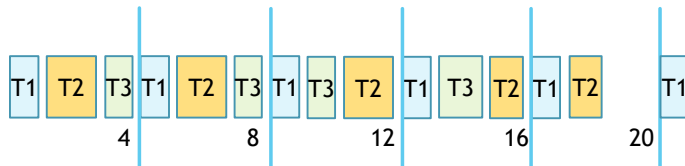
Rate-Monotonic Algorithm

- ▶ For **rate-monotonic (RM)** algorithm, tasks with shortest period are given highest priority.
- ▶ Rate of job releases for a task is inverse of its period.
- ▶ Refer to a schedule produced by the algorithm as an **RM schedule**.
- ▶ Example below shows RM schedule for system with tasks
 - ▶ $T_1 = (p_1, e_1) = (4, 1)$, $T_2 = (5, 2)$, and $T_3 = (20, 5)$.
- ▶ What is the utilization ?



Rate-Monotonic Algorithm

- ▶ Example below shows RM schedule for system with tasks
 - ▶ $T1 = (p_1, e_1) = (4, 1)$, $T2 = (5, 2)$, and $T3 = (20, 5)$.



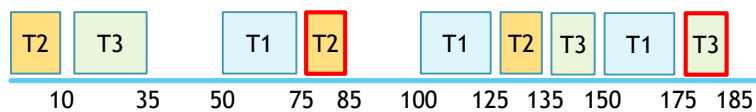
RM versus DM Algorithm

- ▶ When relative deadline of each task is proportional to its period, then RM and DM are equivalent.
- ▶ When relative deadlines arbitrary, DM can sometimes find feasible schedule when RM fails.
- ▶ RM algorithm always fails when DM fails.
- ▶ Consider RM algorithm for tasks
 - ▶ $T1 = (\Phi_1, p_1, e_1, D_1) = (50, 50, 25, 100)$,
 - ▶ $T2 = (0, 62.5, 10, 20)$, and
 - ▶ $T3 = (0, 125, 25, 50)$.

RM versus DM Algorithm

► Consider RM algorithm for tasks

- $T1 = (\Phi_1, p_1, e_1, D_1) = (50, 50, 25, 100)$,
- $T2 = (0, 62.5, 10, 20)$, and
- $T3 = (0, 125, 25, 50)$.



Offline scheduling

- Offline schedules are computed in advance before system begins to execute.
- Requires knowledge of release times and resource requirements for all jobs.
- Disadvantage is that it is inflexible.
- Advantages for deterministic systems is that the schedule is deterministic.
- Also, as scheduling is not done in real-time, complexity of scheduling algorithm not important.

Online scheduling

- ▶ Online schedules are computed on the fly.
- ▶ Distinguished by fact scheduler makes all decisions without knowledge of jobs that will be released in the future.
- ▶ Details of job become known only when job released.
- ▶ Only option when future workload is unpredictable.
- ▶ Advantages are more flexible and adaptable.
- ▶ Disadvantage is it is less likely to make best use of resources.

Next week ...

- ▶ Embedded systems
- ▶ Microprocessor vs. microcontrollers
 - ▶ peripherals
- ▶ **Do NOT forget**
 - ▶ Deadline for project proposals