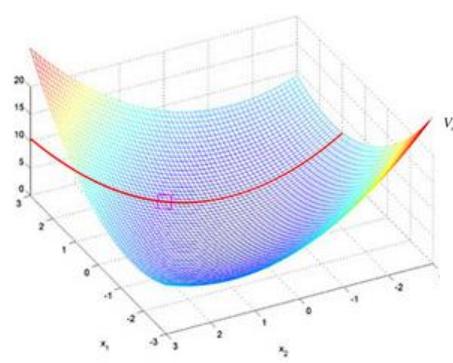
MEGR 7090/8090: Advanced Optimal Control

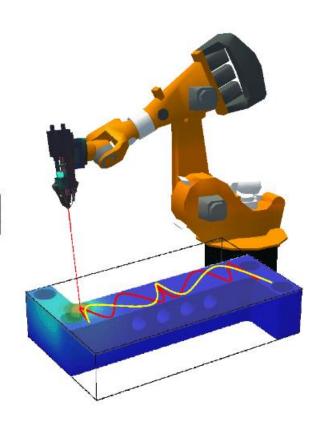




$$V_{n}\left(\mathbf{x}_{n}\right) = \min_{\left\{\mathbf{u}_{n}, \mathbf{u}_{n+1}, \cdots, \mathbf{u}_{N-1}\right\}} \left[\frac{1}{2} \sum_{k=n}^{N-1} \left(\mathbf{x}_{k}^{T} \mathbf{Q}_{k} \mathbf{x}_{k} + \mathbf{u}_{k}^{T} \mathbf{R} \mathbf{u}_{k}\right) + \frac{1}{2} \mathbf{x}_{N}^{T} \mathbf{Q}_{N} \mathbf{x}_{N} \right]$$

$$\begin{aligned} V_{n}\left(\mathbf{x}_{n}\right) &= \min_{\left[\mathbf{u}_{n}, \mathbf{u}_{n-1}, \cdots, \mathbf{u}_{N-1}\right]} \left[\frac{1}{2} \sum_{k=n}^{N-1} \left(\mathbf{x}_{k}^{T} \mathbf{Q}_{k} \mathbf{x}_{k} + \mathbf{u}_{k}^{T} \mathbf{R} \mathbf{u}_{k}\right) + \frac{1}{2} \mathbf{x}_{N}^{T} \mathbf{Q}_{N} \mathbf{x}_{N}\right] \\ &= \min_{\mathbf{u}_{n}} \left[\frac{1}{2} \left(\mathbf{x}_{n}^{T} \mathbf{Q}_{n} \mathbf{x}_{n} + \mathbf{u}_{n}^{T} \mathbf{R} \mathbf{u}_{n}\right) + \min_{\left[\mathbf{u}_{n-1}, \cdots, \mathbf{u}_{N-1}\right]} \left[\frac{1}{2} \sum_{k=n+1}^{N-1} \left(\mathbf{x}_{k}^{T} \mathbf{Q}_{k} \mathbf{x}_{k} + \mathbf{u}_{k}^{T} \mathbf{R} \mathbf{u}_{k}\right) + \frac{1}{2} \mathbf{x}_{N}^{T} \mathbf{Q}_{N} \mathbf{x}_{N}\right] \right] \\ &= \min_{\mathbf{u}_{n}} \left[\frac{1}{2} \left(\mathbf{x}_{n}^{T} \mathbf{Q}_{n} \mathbf{x}_{n} + \mathbf{u}_{n}^{T} \mathbf{R} \mathbf{u}_{n}\right) + V_{n+1} \left(\mathbf{x}_{n+1}\right)\right] \end{aligned}$$

$$V_{n}\left(\mathbf{x}_{n}\right) = \min_{\mathbf{u}_{n}} \left[\frac{1}{2} \left(\mathbf{x}_{n}^{T} \mathbf{Q}_{n} \mathbf{x}_{n} + \mathbf{u}_{n}^{T} \mathbf{R} \mathbf{u}_{n} \right) + V_{n+1} \left(\mathbf{x}_{n+1} \right) \right]$$



Lecture 16 October 17, 2017

Dynamic Programming – Backward Recursion – Reminder



Getting set up:

- ullet Quantize control variables into p discrete values
- Quantize state variables into q discrete values

Solution algorithm:

- Start at step N-1. For each of the q allowable state variables, calculate the stage cost $(g(\mathbf{x}(N-1), u(N-1)))$ for each of the p allowable control variables that lead to constraint satisfaction. Control variables that do not satisfy constraints are termed *inadmissible*. Record the optimal control signals and corresponding stage costs for each originating state.
- Move to step N-2. For each of the q allowable state variables, calculate the stage cost $(g(\mathbf{x}(N-1), u(N-1)))$ and associated intermediate state, $\mathbf{x}_i(N-1)$, for each of the p allowable control variables that lead to constraint satisfaction. To determine which control variable is optimal, compute the total cost to go as $J^*_{\mathbf{x}_o(N-2)\to\mathbf{x}_i(N-1)\to\mathbf{x}_f(N)} = J^*_{\mathbf{x}_o(N-2)\to\mathbf{x}_i(N-1)} + J^*_{\mathbf{x}_i(N-1)\to\mathbf{x}_f(N)}$
- Move to step N-3 and repeat the process (total cost to go is now $J_{\mathbf{x}_o(N-3)\to\mathbf{x}_i(N-2)\to\mathbf{x}_f(N)}^* = J_{\mathbf{x}_o(N-3)\to\mathbf{x}_i(N-2)}^* + J_{\mathbf{x}_i(N-2)\to\mathbf{x}_f(N)}^*$). Keep stepping backward in time until step 0.

Dynamic Programming – Forward Recursion – Reminder



Getting set up:

- Quantize control variables into p discrete values
- Quantize state variables into q discrete values
- Usually the initial state, $\mathbf{x}(0)$, is specified (you don't get to choose your initial state)

Solution algorithm:

- Start at step 1. For each of the q allowable values of $\mathbf{x}(1)$, back-compute the value of $\mathbf{x}(0)$ and stage cost that results from each of the p allowable control values. For each admissible control value (whose associated $\mathbf{x}(0)$ satisfies initial condition requirements), determine the optimal stage cost—this cost, denoted by $J^*_{\mathbf{x}_0(0)\to\mathbf{x}_t(1)}$, is the **optimal cost to arrive** at state $\mathbf{x}_t(1)$.
- Move to step 2. For each of the q allowable values of $\mathbf{x}(2)$, back-compute $\mathbf{x}(1)$ and the associated stage cost for each of the p allowable control variables that lead to constraint satisfaction. To determine which control sequence is optimal, compute the total cost to arrive as $J^*_{\mathbf{x}_o(0) \to \mathbf{x}_i(1) \to \mathbf{x}_f(2)} = J^*_{\mathbf{x}_o(0) \to \mathbf{x}_i(1)} + J^*_{\mathbf{x}_i(1) \to \mathbf{x}_f(2)}$
- Move to step 2 and repeat the process. Keep stepping forward in time until step N.

Dynamic Programming – Some Comments



Backward vs. forward recursion:

Why backward recursion is usually preferred? 1) Result of DP: u* for every (quantited) Xo. - DP is really slow. - Can implement a lookup table based on DP result. - Can use result for benchmarking. 2) At step 1 of fund recursion, for every istate value, we Compute Xo, for control value.

DP vs. SQP:

SQP: Local, continuous (u can take on any value)
DP: @Global, gridded

Dynamic Programming – More Realistic Example UNC CHARLOTTE



Suppose a vehicle's dynamics are given by: $m\dot{v} = u - C_{rr}mg - 0.5\rho v^2C_dA_{ref}$ $\dot{x} = v$

Parameter values:
$$m=1000kg$$
, $g=9.8\frac{m}{s^2}$, $C_{rr}=0.01$, $\rho=1.2\frac{kg}{m^3}$, $C_d=0.4$, $A_{ref}=5m^2$

Suppose that our goals are:

- Given an initial position of x=0, achieve a final position of x=1600 (approximately one "metric" mile of travel) after 60 seconds
- Minimize total energy expended over 60 seconds

Set up a nonlinear optimization problem and solve using dynamic programming

- Identify the states and control signal
- For different quantization levels, assess computational complexity

Dynamic Programming – More Realistic Example



States (2): Position (x) and velocity (v)

• Taking q_x and q_v as the number of quantized positions and velocities, respectively, the number of state variable combinations is given by $q=q_xq_v$

Control variable (1): Applied force (u)... Number of quantized forces denoted p

Horizon length: $N = \frac{60}{\Delta T}$ (where ΔT is the discrete time step)

Computational complexity: Number of cost function evaluations (E) given by:

$$E = Npq = Npq_xq_v$$

Example code available on Canvas.

Dynamic Programming – More Realistic Example

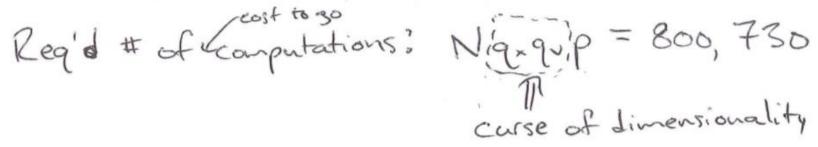


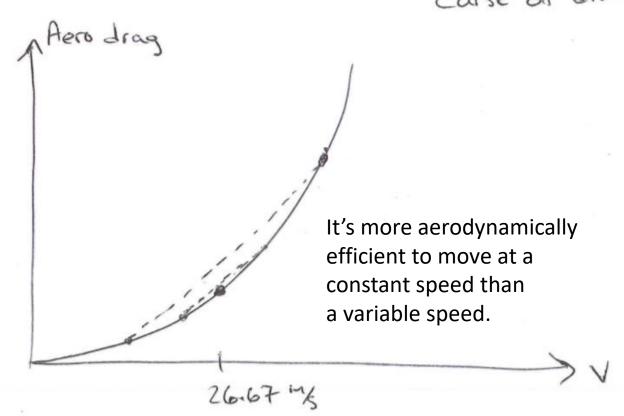
Ex:
$$m\ddot{v} = u - C_{rr}mg - 0.5e N^2C_dA_{ref}$$

 $\ddot{x} = V$
 $V(k+1) = V(k) + \frac{\Delta t}{M}(u(k) - C_{rr}mg - 0.5e C_dA_{ref} V(k)^2)$
 $\chi(k+1) = \chi(k) + \frac{V(k) + V(k+1)}{2} \Delta t$ } trapezoidal approximation
 $J(u; \chi(0)) = \int_0^{60} u(t) \dot{v}(t) dt \approx \int_{c:0}^{2-1} u(c) V(c) \Delta t$
 $Constraints: \chi(N) \ge 1600$
 $V(N) \ge 25$
States: $(\chi_3) \dot{v} \dot{v} = 1000$
 $V(N) \ge 25$
States: $(\chi_3) \dot{v} \dot{v} = 1000$
 $V(N) \ge 25$
 $V(N) \ge 1000$
 $V(N) \ge 1000$

Dynamic Programming – More Realistic Example







KE=12 muz

Work must be done to increase kinetic energy. Therefore, the vehicle should finish at the minimum allowable speed. Note that its *average* speed must be greater than this to achieve x(N) = 1600 meters.

Dynamic Programming and the Principle of Optimality – A Famous Result



Consider the following system dynamics:

$$\mathbf{x}(k+1) = A\mathbf{x}(k) + Bu(k)$$

$$\mathbf{x} \in \mathbb{R}^n$$
, $u \in \mathbb{R}$

...and the following objective function:

$$J(\mathbf{x}(0), \mathbf{u}) = \mathbf{x}^T(N)S\mathbf{x}(N) + \sum_{i=0}^{N-1} (\mathbf{x}^T(i)Q\mathbf{x}(i) + Ru^2(i))$$

where:
$$S > 0$$
, $Q > 0$, $R > 0$

- This is known as the finite-horizon discrete-time linear quadratic regulator problem
- Important notation (seems strange now, will make sense later): $P(0) \triangleq S$

Discrete-Time LQR - Backward Recursion Step N-1



Note:
$$J_{\mathbf{x}(N) \to \mathbf{x}(N)}^* = \mathbf{x}^T(N)S\mathbf{x}(N) = \mathbf{x}^T(N)P(0)\mathbf{x}(N)$$

Beginning backward recursion (step N-1):

$$J_{\mathbf{x}(N-1)\to\mathbf{x}(N)} = \mathbf{x}^{T}(N-1)Q\mathbf{x}(N-1) + Ru^{2}(N-1) + \mathbf{x}^{T}(N)P(0)\mathbf{x}(N)$$

$$= \mathbf{x}^{T}(N-1)Q\mathbf{x}(N-1) + Ru^{2}(N-1) + (A\mathbf{x}(N-1) + Bu(N-1))^{T}P(0)(A\mathbf{x}(N-1) + Bu(N-1))$$

Collecting like terms, differentiating with respect to u(N-1), and setting the derivative to 0:

$$\frac{\partial J_{N-1 \to N}}{\partial u(N-1)} = 2(Ru(N-1) + B^T P(0)Bu(N-1) + B^T P(0)A\mathbf{x}(N-1))$$

$$u^*(N-1) = -\frac{(R+B^T P(0)B)^{-1}B^T P(0)A\mathbf{x}(N-1)}{\triangleq K(N-1)}$$

Discrete-Time LQR – Backward Recursion Step N-2



Note: $J_{\mathbf{x}(N-1)\to\mathbf{x}(N)}^* = \mathbf{x}^T(N-1)P(1)\mathbf{x}(N-1)$ where:

$$P(1) = (A - BK(N - 1))^{T} P(0) (A - BK(N - 1)) + K^{T}(N - 1)RK(N - 1) + Q$$

Continuing backward recursion (step N-1):

$$J_{\mathbf{x}(N-2)\to\mathbf{x}(N)} = \mathbf{x}^{T}(N-2)Q\mathbf{x}(N-2) + Ru^{2}(N-2) + \mathbf{x}^{T}(N-1)P(1)\mathbf{x}(N-1)$$

$$= \mathbf{x}^{T}(N-2)Q\mathbf{x}(N-2) + Ru^{2}(N-2) + (A\mathbf{x}(N-2) + Bu(N-2))^{T}P(1)(A\mathbf{x}(N-2) + Bu(N-2))$$

Key observation: The cost to go above is **identical** in structure to $J_{\mathbf{x}(N-1)\to\mathbf{x}(N)}$, except that:

- u(N-1) and $\mathbf{x}(N-1)$ have been replaced with u(N-2) and $\mathbf{x}(N-2)$
- P(0) has been replaced with P(1)

Result: The optimal control signal $(u^*(N-2))$ will be the same as before, except that $\mathbf{x}(N-1)$ and P(0) will be replaced with $\mathbf{x}(N-2)$ and P(1), respectively

Discrete-Time LQR – Derivation in Class



Given:
$$\underline{x}(k+1) = \underline{A}\underline{x}(k) + \underline{B}\underline{u}(k)$$

Minimize $J(\underline{u}; \underline{x}(0)) = \underline{x}T(N)S\underline{x}(N) + \sum_{i=0}^{N-1} (\underline{x}T(i)Q\underline{x}(i) + Ruli))$

Equivalent to minimizing

$$J(\underline{u}; \underline{x}(0)) = \sum_{i=0}^{N-1} (\underline{x}T(iH)Q\underline{x}(i+1) + Ru^2(i))$$

when $S = Q$

$$= \underline{x}T(N)Q\underline{x}(N) + \sum_{i=0}^{N-1} (\underline{x}T(i)Q\underline{x}(i) + Ruli)$$

$$= \underline{x}T(N)Q\underline{x}(N) + \sum_{i=0}^{N-1} (\underline{x}T(i)Q\underline{x}(i) + Ruli)$$

$$= \underline{x}T(N)Q\underline{x}(N) + \underline{x}T(i)Q\underline{x}(i) + Ruli)$$

$$= \underline{x}T(N)Q\underline{x}(N) + \underline{x}T(N)Q\underline{x}(N) + \underline{x}T(N)Q\underline{x}(N) + Ruli)$$

$$= \underline{x}T(N)Q\underline{x}(N) + \underline{x}T(N)Q\underline{x$$

$$P(o) \stackrel{\triangleright}{=} 5$$

$$J_{\underline{x}(n) \rightarrow \underline{x}(n)}^{*} = \underline{x}^{T}(n) P(o) \underline{x}(n)$$

Just as the purpose (and Iraq war metaphor) of the strange agrarian community in M. Night Shyamalan's *The Village* is initially hazy but becomes remarkably clear later on, the purpose of the mysterious assignment of P(0) will become clear...with stunning implications for LQR control

Discrete-Time LQR – Derivation in Class



Stage N-1:
$$J_{\underline{x}(N-1) \to \underline{x}(N)} = \underline{x}^{T}(N-1)Q_{\underline{x}}(N-1) + Ru^{2}(N-1) + \underline{x}^{T}(N)P(0)\underline{x}(N)$$

$$= \underline{x}^{T}(N-1)Q_{\underline{x}}(N-1) + Ru^{2}(N-1) + [A_{\underline{x}}(N-1)+B_{\underline{u}}(N-1)]^{T}P(0)...$$

$$[A_{\underline{x}}(N-1)+B_{\underline{u}}(N-1)]$$

$$= \underline{\lambda}J_{\underline{x}(N-1)\to \underline{x}(N)} = 2Ru(N-1) + 2B^{T} P(0)B_{\underline{u}}(N-1)$$

$$+ 2B^{T}P(0)A_{\underline{x}}(N-1) = 0 \text{ when}$$

$$u(N-1) = u^{*}(N-1)$$

$$= \underline{\lambda}u^{*}(N-1) = -(R+B^{T}P(0)B)^{-1}(B^{T}P(0)A)\underline{x}(N-1)$$

$$= \underline{\lambda}u^{*}(N-1)$$

Discrete-Time LQR – Derivation in Class



$$J_{\times(N-1)\to\times(N)}^{*} = \chi^{T}(N-1)Q_{\times}(N-1) + \chi^{T}(N-1)K^{T}(N-1)K(N-1)\chi(N-1)$$

$$+ [(A-3K(N-1))\chi(N-1)]^{T}P(0)[(A-3K(N-1))\chi(N-1)]$$

$$= \chi^{T}(N-1)[Q+K^{T}(N-1)RK(N-1)+(A-3K(N-1))^{T}P(0)(A-3K(N-1))$$

$$P(1) \qquad \chi(N-1)$$

Same as at stage N, with P(0) replaced by P(1) (ahh...so that's the reason for P(i)) and $\mathbf{x}(N)$ replaced with $\mathbf{x}(N-1)$

...At stage N-2, we have an identical optimization problem to the one at stage N-1, with $\mathbf{x}(N-1)$ replaced with $\mathbf{x}(N-2)$, u(N-1) replaced with u(N-2), and P(0) replaced with P(1)

...And if the problems are identical (with index substitutions), then the solutions will be identical (with index substitutions)!

Discrete-Time LQR – General Control Law



Remember:
$$u^*(N-1) = -K(N-1)\mathbf{x}(N-1)$$
 and $u^*(N-2) = -K(N-2)\mathbf{x}(N-2)$

where:
$$K(N-1) = -(R + B^T P(0)B)^{-1}B^T P(0)A$$
 and $K(N-2) = -(R + B^T P(1)B)^{-1}B^T P(1)A$

where:
$$P(0) = S$$
 and $P(1) = (A - BK(N - 1))^T P(0)(A - BK(N - 1)) + K^T(N - 1)RK(N - 1) + Q$

Continuing the backward recursion (generic step N-i):

$$u^*(N-i) = -K(N-i)\mathbf{x}(N-i)$$

where:
$$K(N-i) = -(R + B^T P(i-1)B)^{-1} B^T P(i-1)A$$

where:
$$P(i) = (A - BK(N - i))^T P(i - 1)(A - BK(N - i)) + K^T(N - i)RK(N - i) + Q$$

Key takeaway: The resulting controller is a linear, time-varying controller whose gains at each step can be found recursively!

Discrete-Time LQR – Block Diagram and Observations



Key takeaway (reminder): The resulting controller is a linear, time-varying controller whose gains at each step can be found recursively!

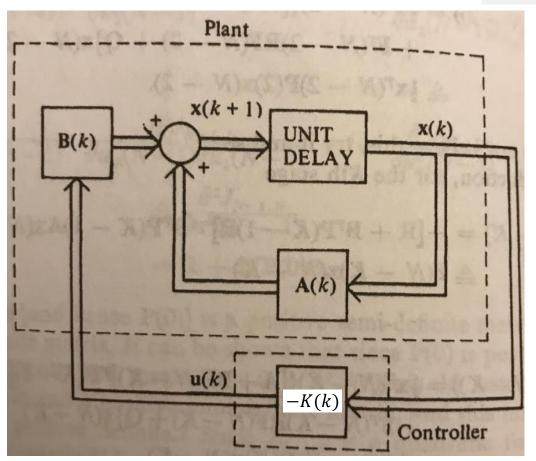


Image source: Kirk, page 82

Key differences from standard DP:

- No quantization of control and states (since cost can be differentiated)
- Unlike DP, where a sequence of control signals is computed all at once, LQR lends itself to a
 feedback control law (which happens to be linear and time-varying)

Infinite Horizon LQR - Setup



Consider the following system dynamics:

$$\mathbf{x}(k+1) = A\mathbf{x}(k) + Bu(k)$$
 $\mathbf{x} \in \mathbb{R}^n, u \in \mathbb{R}$

...and the following objective function:

$$J(\mathbf{x}(0), \mathbf{u}) = \sum_{i=0}^{\infty} (\mathbf{x}^{T}(i)Q\mathbf{x}(i) + Ru^{2}(i))$$

where:
$$Q > 0$$
, $R > 0$

...Same problem as before, except that:

- $N = \infty$
- There is no terminal penalty

Infinite Horizon LQR - Solution



Key facts:

- In the finite-horizon LQR problem, $J_{\mathbf{x}(i)\to\mathbf{x}(N)}^* = \mathbf{x}^T(i)P(i)\mathbf{x}(i)$, regardless of the presence of a terminal penalty
- For a time-invariant system, the infinite horizon cost to go from a given state (x) is the same regardless of the time instant at which the controller is "turned on":

$$J_{\mathbf{x}(i)\to\infty}^* = J_{\mathbf{x}(i+1)\to\infty}^*$$

• Putting these facts together, it must be true that P(i) = P(i+1) = P...P is a constant!

Recall the update laws for **P** and **K**: $K(N-i) = -(R+B^TP(i-1)B)^{-1}B^TP(i-1)A$

$$P(i) = (A - BK(N - i))^{T} P(i - 1)(A - BK(N - i)) + K^{T}(N - i)RK(N - i) + Q$$

The "update" laws now become: $K = -(R + B^T P B)^{-1} B^T P A$

$$P = (A - BK)^T P(A - BK) + K^T RK + Q$$
Riccati equation

Infinite Horizon LQR - Solution



Remember:
$$J_{x(i)}^{*} \rightarrow x(\omega) = J_{x(i+1)}^{*} \rightarrow x(\omega)$$

Same state values

$$J_{x(i)}^{*} \rightarrow x(\omega) = x^{T}(i) P(\omega + x(i))$$

$$J_{x(i+1)}^{*} \rightarrow x(\omega) = x^{T}(i+1) P(\omega - i - 1) x(i+1)$$
If $x(i) = x(i+1)$, $P(\omega - i - 1) = P(\omega - i)$

Infinite Horizon LQR - Solution



Given x: - Ax(1) + Bu, we can design a state feedback based regulator in 2 ways: 1) Pole placement: u=-Kx(k) =) x= (A-BK)x(k) (K+1) Set eigenvalues equal to desired pole locations. 2) LQR design: Specify Q, R (QEIRNXN REIR) where $J(x_0, K) = \sum_{i=0}^{\infty} \{x^T(i)Qx(i) + Ru(i)^2 \}$ =) Determine K based on discrete-time algebraic Riccati

Infinite Horizon LQR Solution in MATLAB



Syntax:

Resulting feedback gain vector

Example to work in class:

$$\mathbf{x}(k+1) = A\mathbf{x}(k) + Bu(k)$$

$$A = \begin{bmatrix} 0 & 1 \\ -2 & -3 \end{bmatrix} \qquad B = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

$$J(\mathbf{x}(0), \mathbf{u}) = \sum_{i=0}^{\infty} (\mathbf{x}^{T}(i)Q\mathbf{x}(i) + Ru^{2}(i))$$

$$Q = \begin{bmatrix} 2 & 0 \\ 0 & 1 \end{bmatrix} \qquad R = 1$$

Observations About the LQR Solution



Good features:

- Closed-form solution available (via differentiation of cost to go), so no need to quantize states and control variables
- Optimal controller is linear state feedback...can be implemented as a feedback controller, rather than just a sequence of control values that begins at time 0
- With an infinite horizon and time-invariant system, the feedback gain is time-invariant

Not so good features:

- Quadratic cost function (in states and control variables) required
- No constraints considered

Question to be addressed in future lectures: Can we obtain the good features above while allowing for more general cost function structures and constraints?

Preview of Upcoming Lectures



Model predictive control – answers the question of "how do we take the trajectory optimization algorithms that we've learned about and use them for real-time feedback control?":

- Basic MPC setup and notation (October 19)
- MPC implementation in MATLAB and Simulink (October 24)
- Stability and robustness of MPC (October 26)