

PROJECT - PHASE 1

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Introduction

Report for Numerical Methods **Project - Phase I**, which requires implementation of a program that calculates the possible solutions of a system of linear equations using:

- **Gauss Elimination**
- **Gauss-Jordan Elimination**
- **LU Decomposition**
- **Gauss-Seidel Iterative Method**
- **Jacobi Iteration Method**

Assumptions

- In order to run the program correctly through **Angular**, please start the terminal and type: **"npm install"** to install node modules. You must have node.js installed on your personal computer. **Next in the terminal**, type **"ng serve"** to run the program.
- Please notice that the program runs in → localhost:4200 ← click here to go directly to the program page.
- The user has to enter the number of equations in order to proceed.
- The coefficients have to be numbers and not letters.
- The equations are processed and evaluated using functions implemented within each component.
- **TWO** functions are responsible for handling user input and the pseudocode for these functions is:
 - Function to get number of equations.
 - Let n be the number of equations and get it from the user.
 - If (**n is null or n == 0**) → **return**
 - Else: Create an array of zeros with size n
 - End function.
 - Function to get the equations.

-
- Recreate the array.
 - **For** $i = 0 : n$
 - Push the equation into the array.
 - End **for**.
 - Let precision be a number and get it from the user.
 - If (**precision is null or == 0**) it is set to 7 and the program continues.
 - Now it is the **hash class** turn to evaluate the equations and split them into a matrix of coefficients and a matrix of results and sometimes an augmented matrix.
 - End function.
 - The program handles many exceptional cases and never happens to crash or do *unlogical* operations at any time.
 - Implemented pivoting whenever we could.
 - Implemented scaling and provided steps in some methods.

Gauss Elimination Documentation

Pseudocode

Defining “gaussSolver(arr[][] of number)” function performs forward elimination and then calls back the substitution method.

- array=clone(arr)
- For i:n
 - pivotAndscale(arr,i) - We will explain this function later as it is an implementation of pivoting logic in iteration #i.
 - For k:n
 - Factor = array[k][i] / array[i][i]
 - if(factor=0)

hasSolution(array)

End function

End if
 - For j : array[k].length - As we process on an augmented matrix
 - array[k][j] = array[k][j] - factor * array[i][j]
 - End for
 - array[k][i] =0
 - End for
- End for

backSub(arr)

Defining “hasSolution(arr[][] of number)” function check if function has no solution or infinite number of solution or have unique number of solution

Inf = false

- For i:n

```
rawZero=true
```

```
For j:n
```

```
    if(arr[i][j]!=0)
```

```
        rawZero=false
```

```
        Break For
```

```
End for
```

```
if( number of variable >number of equation or rank matrix!= rankAgumented)
```

```
    Return "has no solution"
```

```
    End function
```

```
Else if (number of variable <number of equation or rawZero=0)
```

```
    Inf = true
```

```
End Else if
```

```
End for
```

```
if(Inf = true)
```

```
    Return "Has infinite number of solution"
```

```
End if
```

```
Else
```

```
    Return "Has unique solution"
```

```
End else
```

Defining "pivotAndscale(arr[][] of number,pivot index)" array need to pivot or not.

```
Temp =clone(| arr |)
```

```
BigestInRow = getBiggestInrow(temp)
```

- For i:n

```
    For j:n
```

```
        Temp[i][j]=Temp[i][j]/BigestInRow
```

```
    End for
```

```
End for
```

```
BigetInCoulmnAfterScale = getBiggestIncoulmn(Temp)
```

```
if (BigetInCoulmnAfterScale=pivot index)
```

```
    Return
```

```
End if
```

```
Else
```

```
    swap(arr[BigetInCoulmnAfterScale],arr[pivot index])
```

```
End else
```

```
Defining "backSubs(arr[][] of number)" get system solutions
```

```
has Solution=hasSolution(arr)
```

```
If( has Solution="unique")
```

```
    For i=n to 0:
```

```
        sum=0
```

```
        For j=0 to i
```

```
            arr[i][arr.length]=arr[i][arr.length]-arr[i][j]
```

```
        End for
```

```
solution[i]=arr[i][arr.length]/arr[i][i]
```

```
End for
```

Sample Runs

Equation number 1

$4x+2y+z=11$

Equation number 2

$-x+2y=3$

Equation number 3

$2x+y+4z=16$

Solve

enter the precision

Last Step

4	2	1	11
0	2.5	0.25	5.75
0	0	3.5	10.5

$x=1$

$y=2$

$z=3$

Equation number 1

$x+y=2$

Equation number 2

$2x+2y=3$

enter the precision

Solve

Has No solution

Time Complexity

- Elimination Steps: $2(n^3/3)$.
- Time Complexity: $O(n^3)$.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using array with one dimensional and multidimensional it was helpful in computational so all have same type and sure that program do operation to suitable variable
- And using a lot of lists if we didn't know the number of inputs was helpful.

Gauss-Jordan Elimination Documentation

Pseudocode

Assuming we got the user-input matrix we start operating on it as follows.

Defining “**gElimination()**” function to perform Forward Gauss Elimination, which we would use later in further logic.

- **For** $i = 0 : n$
 - **pivot(i)** - We will explain this function later as it is an implementation of pivoting logic in iteration #i.
 - **For** $k = i + 1 : n$
 - $\text{Factor} = \text{matrix}[k][i] / \text{matrix}[i][i]$
 - **For** $j = i : \text{matrix}[k].\text{length}$ - As we process on an augmented matrix
 - $\text{Matrix}[k][j] = \text{matrix}[k][j] - \text{factor} * \text{matrix}[i][j]$
 - **End for.**
 - **End for.**
- **End for.**

End function.

Defining “**pivot(iteration)**” function takes an iteration parameter as it pivots per iteration.

- Let $\text{pivot} = \text{iteration}$
- Let $\text{max} = \text{matrix}[\text{iteration}][\text{iteration}]$
- **For** $\text{index} = \text{iteration} + 1 : n$
 - Let $\text{dummy} = \text{matrix}[\text{index}][\text{iteration}]$
 - If $\text{dummy} > \text{max} \rightarrow$ Set max to dummy and pivot to i
- **End for.**
- If $\text{pivot} \neq \text{iteration}$ (Checking whether the pivot changed or not)
 - \rightarrow Swap rows using temporary variables
 - If not \rightarrow do nothing.

End function.

Defining “**gjElimination()**” function which calls **gElimination()** then continues backward elimination on the resulting matrix.

- **gElimination()** Apply forward elimination using previously structured Gauss Elimination function.
- If (**isSingular()**) → **return** Do not continue as if the matrix is singular, we have a row of zeros and continuing does not make sense.
- **For i = n : 1**
 - **pivot(i)** Applying pivoting as previously done in gElimination (We did not really have to apply pivoting here, but just in case).
 - **For k = i - 1 : 0**
 - **Factor** = matrix[k][i] / matrix[i][i]
 - **For j = n : 0**
 - Matrix[k][j] = matrix[k][j] - factor * matrix[i][j]
 - **End for.**
 - **End for.**
- **End For.**

End function.

Defining “**isSingular()**” function to check whether the matrix is singular or not.

- Let Determinant = 1
- For index : n, where n is number of equations
 - Determinant = Determinant * matrix[index][index]
- The loop gets the determinant value of the matrix after applying Gauss Elimination
- if(!Determinant) → Singular (boolean attribute) = true

End function.

Defining “**solve()**” function to get matrix solution.

- If (**isSingular()**) → check if the matrix has no solution or infinite number of solutions and print according to the result.
- If not singular → **For i = 0 : n**

-
- Let **coefficient** = matrix[i][i]
 - Let **result** = matrix[i][n]
 - Let **solution** = **result** / **coefficient**
 - Push **solution** into the solution array named under **solution**.
 - End **for**.

End Function.

Sample Runs

Entering 0 or not entering a number of equations at all does not allow the user to proceed within the program.

Gauss-Jordan Elimination

Number of Equations

Enter

Assume the user entered 3, the program continues and asks the user to enter the equations. Please **note** that if the user entered no equations the program shows a message saying **"No Solution"**. **Precision** is totally optional and if the user ignores it, the program proceeds with maximum precision as possible

Gauss-Jordan Elimination

Equation No. 1	<input type="text" value="4x+2y+z=11"/>
Equation No. 2	<input type="text" value="-x+2y=3"/>
Equation No. 3	<input type="text" value="2x+y+4z=16"/>

Set Precision

Solve

Given the assumption that a precision of 4 was entered along the equations, this is what a user would normally get.

Gauss-Jordan Elimination

Calculated the factor in step 0, Factor = -0.25

Row calculation in step 1, Matrix =

[4,2,1,11]
[0,2.5,0.25,5.75]
[2,1,4,16]

Calculated the factor in step 2, Factor = 0.5

Row calculation in step 3, Matrix =

[4,2,1,11]
[0,2.5,0.25,5.75]
[0,0,3.5,10.5]

Calculated the factor in step 4, Factor = 0

Row calculation in step 5, Matrix =

[4,2,1,11]
[0,2.5,0.25,5.75]
[0,0,3.5,10.5]

Calculated the factor in step 6, Factor = 0.07143

Row calculation in step 7, Matrix =

[4,2,1,11]
[0,2.5,-0.000005,5]
[0,0,3.5,10.5]

Calculated the factor in step 8, Factor = 0.2857

Row calculation in step 9, Matrix =

[4,2,0.00005,8]
[0,2.5,-0.000005,5]
[0,0,3.5,10.5]

```
Calculated the factor in step 10, Factor = 0.8
Row calculation in step 11, Matrix =
[ 4,0,0.000054,4 ]
[ 0,2.5,-0.000005,5 ]
[ 0,0,3.5,10.5 ]
```

```
Solution = 1,2,3
```

Let's try another system of linear equations, but this time it consists of 2 equations.

Gauss-Jordan Elimination

Number of Equations

2

Enter

Gauss-Jordan Elimination

Equation No. 1

$x+y=2$

Equation No. 2

$2x+2y=3$

Set Precision

Solve

The program calculates the augmented matrix in steps as follows. This is what the user got for this sample.

Gauss-Jordan Elimination

Calculated the factor in step 0, Factor = 0.5
Row calculation in step 1, Matrix =
[2,2,3]
[0,0,0.5]
Solution = No solution! Matrix is singular.

Time Complexity

- Elimination Steps: $4(n^3/3)$.
- Time Complexity: $O(n^3)$.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using an array of one dimension and multidimensions was helpful in computations so all has the same type and making sure that program operates on suitable variables.
- And using a lot of lists if we didn't know the number of inputs was helpful.

LU Decomposition Documentation

Doolittle Decomposition

Pseudocode

1. Defining “Decompose()” method

```
After receiving the coefficients , free terms, and unknowns
For i = 0 : eqnNo - 1
    Call partial pivoting method on row i
For j = i + 1 : eqnNo
    For k = i : eqnNo
        if(k == i) then
            MatrixLU[j][i] = matrixLU[j][k] / matrixLU[i][k]
        Else
            matrixLU[j][k] -= matrixLU[j][i] * matrixLU[i][k]
```

2. Defining “partialPivoting(i)” method

```
For k = i + 1 : eqnNo
    if ( | matrixLU[i][i] | < | matrixLU[k][i] | ) then
Swap the 2 rows in the augmented matrix column by column
```

3. Defining “forward substitution()” method

```
YfreeTerm[0] = stepFreeTerm[0]
For i = 1 : eqnNo
    For j = 0 : i
        YfreeTerm[i] = stepFreeTerm[i] - matrixLU[i][j] * YfreeTerm[j]
```

4. Defining “backward substitution()” method

```
For i = eqnNo - 1 : 0
    For j = i + 1 : eqnNo
```

$$\text{soln}[i] = (\text{Yfreeterm}[i] - \text{matrixLU}[i][j] * \text{soln}[j]) / \text{matrixLU}[i][i]$$

5. Defining "solve()" method

Call decompose method

Call forward substitution method

Call backward substitution method

Sample Runs

Doolittle Decomposition

The LU Decomposition: $A = L \cdot U$

The System

$$4x + 2y + z = 11$$

$$-x + 2y = 3$$

$$2x + y + 4z = 16$$

L Matrix

$$1$$

$$\begin{matrix} -0.25 & 1 \end{matrix}$$

$$\begin{matrix} 0.5 & 0 & 1 \end{matrix}$$

U Matrix

$$\begin{array}{ccc} 4 & 2 & 1 \\ & 2.5 & 0.25 \\ & & 3.5 \end{array}$$

Free Term

$$11$$

$$3$$

$$16$$

Solving the Equation

$$L \cdot U \cdot X = A \cdot X = b$$

$$\text{let } U \cdot X = Y$$

By Forward Substitution

$$y_1 = 11$$

$$y_2 = 5.75$$

$$y_3 = 10.5$$

By Backward Substitution

$$x = 1$$

$$y = 2$$

$$z = 3$$

Crout LU Decomposition:

Pseudocode

Assuming we have the augmented matrix we start from SplitMatrices where we split the augmented matrix to coefficient matrix and solution matrix.

Then in LUcroutEvaluate function we evaluate L and U matrices:

- **LUcroutEvaluate Function: LUcroutEvaluate()**

```
For i = 0 : coefficient row length
```

```
    Pivoting Function
```

```
    For j = 0 : i+1
```

```
        sumL = 0;
```

```
        For k = 0 : j
```

```
            sumL = sumL + lower[i][k] * upper[k][j]
```

```
        End For
```

```
        Lower[i][j] = coefficient[i][j] - sumL
```

```
    End For
```

```
    For j = i+1: coefficient column length
```

```
        SumJ = 0;
```

```
        For k=0:i
```

```
            sumJ = sumJ + lower[i][k] * upper[k][j]
```

```
        End For
```

```
        upper[i][j] = (coefficient[i][j] - sumJ) / lower[i][i]
```

```
    End For
```

```
End For
```

```
End Function
```

- **Pivotion Function: pivoting(index)**

We loop on column [index] and find the row of maximum element in this column and exchange row of index parameter and row of maximum element.

```
For j = index : j : augmented row length
```

```
    IF absolute coefficient[j][index] > maxRow
```

```
        Then maxRow = j
```

```
    End IF;
```

```
End For
```

```
// replacing 2 row if max is found
```

```
IF maxRow != index
```

```
    For i = 0:coefficient row length
```

```
        temp = this.coff[index][i]
```

```
        coefficient[index][i] = coefficient[maxRow][i]
```

```
        coefficient[maxRow][i] = temp
```

```
    End For
```

```
// replace soln matrix
```

```
temp = soln[maxRow]
```

```
soln[maxRow] = soln[index]
```

```
soln[index] = temp
```

```
End IF
```

- **yEvaluate Function: yEvaluate()**

For i=0 : lower row length

sum=0

For j=0 : i

sum = sum + lower[i][j]*y[j]

End For

Push (soln[i]-sum) to y matrix

End For

- **xEvaluate Function: xEvaluate()**

For i=upper row length - 1 : -1 : i - -

sum=0

For j=i+1 : upper row length

sum = sum + upper[i][j]*x[j]

End For

X[i] = y[i]-sum

End For

Sample Runs

LU Crout Decomposition $A=LU$

Upper Matrix of Crout Decomposition:

1 2.5 4.5

0 1 1.9231

0 0 1

lower matrix of crout decomposition:

2 0 0

6 -13 0

2 -2 -0.1538

free terms

5

3

4

Solving Equation:

$AX=B$ && $A=LU$

$AX=LUX$

Let $Y = UX$

$LY=B$

Y matrix:

2.5

0.92308

-5.5017

$Y=UX$

X matrix:

-1.5004

11.503

-5.5017

Cholesky Decomposition

PseudoCode

1. Defining "checkMatrix()" method

For $i = 0 : \text{eqnNo}$

For $j = 0 : i$

If coefficients[i][j] != coefficients[j][i] then valid = false

Call a part of decompose method in *Doolittle decomposition* to get the upper triangular matrix

For i = 0 : eqnNo

If coefficients[i][i] <= 0 then valid = false

/////all methods of cholesky will be called if valid == true

2. Defining “Decompose()” method

For i = 0 : eqnNo

For j = 0 : i

For k = 0 : j

matrixLU[i][j] -= matrixLU[i][k] * matrixLU[j][k]

If i == j then

matrixLU[i][j] = sqrt(matrixLU[i][j])

Else

matrixLU[i][j] /= matrix[j][j]

3. Defining “forwardSubstitution()”method

For i = 0 : eqnNo

For j = 0 : i

YfreeTerm[i] = (stepFreeTerm[i] - matrix[i][j] * stepFreeTerm[j]) / matrix[i][i]

4. Defining “backwardSubstitution()” method

For i = eqnNo - 1 : 0

For j = i + 1 : eqnNo

soln[i] = (Yfreeterm[i] - matrixLU[i][j] * soln[j]) / matrixLU[i][i]

Sample Runs

Cholesky Decomposition

The LU Decomposition: $A = L \cdot L^T$

The System

$$6x + 2y + 2z = 3$$

$$2x + 3y + 5z = 4$$

$$2x + 5y + 9z = 5$$

L Matrix

2.4494897

0.8164966 1.5275252

0.8164966 2.8368326 0.53452231

Note that the matrix is +ve definite symmetric matrix

L^T Matrix

2.4494897 0.8164966 0.8164966

1.5275252 2.8368326

0.53452231

Free Term

3

4

5

Solving the Equation

$$L \cdot L^T \cdot X = A \cdot X = b$$

$$\text{let } L^T \cdot X = Y$$

By Forward Substitution

$$y_1 = 1.22$$

$$y_2 = 1.97$$

$$y_3 = -2.96$$

By Backward Substitution

$$x_1 = -1.5140136$$

$$x_2 = 11.573884$$

$$x_3 = -5.5376547$$

2.

Cholesky Decomposition

About the System

equation 1 $2x + 5y = 8$

enter the
precision

equation 2 $5x + 5y = 6$

Solve

Note that matrix here is symmetric but not positive definite

Cholesky Decomposition

"No Soln"

Time Complexity

- Elimination Steps: $K(n^3/3)$, where K is the number of equations.
- Time Complexity: $O(n^3)$.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using an array of one dimension and multidimensions was helpful in computations so all has the same type and making sure that program operates on suitable variables.
- And using a lot of lists if we didn't know the number of inputs was helpful.

Iterative Methods Documentation :

Pseudocode → Jacobi-iterative

Defining “**implementJacobi()**” function to perform the jacobi iterative method, which we would use later in further logic.

```
gaussSiedelResults = []
```

```
gaussSiedelResults.push(intialGuess)
```

- If (noOfIterations != -1)

```
    For count : noOfIterations
```

```
        For i : intialGuess.length
```

```
            For j : intialGuess.length
```

```
                tempArray[i] += a[i][j] * intialGuess[j]
```

```
            End for
```

```
            tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]
```

```
        End for
```

```
        intialGuess = tempArray.clone
```

```
        jacobiMethodResults.push(intialGuess)
```

```
    End for
```

- Else

```
    While (relativeError >= eTolerance)
```

```
        For i : intialGuess.length
```

```
            For j : intialGuess.length
```

```

        tempArray[i] += a[i][j] * intialGuess[j]

    End for

    tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]

    relativeError = Math.max(relativeError, (tempArray[i] - intialGuess[i]) /
tempArray[i] * 100)

    End for

    intialGuess = tempArray.clone

    jacobiMethodResults.push(intialGuess)

End for

Return gaussSiedelResults

End function

```

Pseudocode → Gauss-Siedel

```

gaussSiedelResults = []

gaussSiedelResults.push(intialGuess)

• If (noOfIterations != -1)

    For count : noOfIterations

        For i : intialGuess.length

            For j : intialGuess.length

                tempArray[i] += a[i][j] * tempArray[j]

            End for

            tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]

```

```

        End for

        intialGuess = tempArray.clone

        jacobiMethodResults.push(intialGuess)

    End for

    • Else

        While (relativeError >= eTolerance)

            For i : intialGuess.length

                For j : intialGuess.length

                    tempArray[i] += a[i][j] * tempArray[j]

                End for

                tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]

                relativeError = Math.max(relativeError, (tempArray[i] - intialGuess[i] /
tempArray[i] * 100)

            End for

            intialGuess = tempArray.clone

            jacobiMethodResults.push(intialGuess)

        End for

        Return gaussSiedelResults

    End function.

```

Sample run :

Example 1: **Jacobi Iterative Method**, Example 2: **Gauss-Seidel Method**.

Iterative Methods

Number of iterations

Tolerance value

5

Precision

4

3

Submit

Enter initial guess :

1

1

1

Enter Equations :

$4x+2y+z=11$

$-x+2y=3$

$2x+y+4z=16$

Solve by Gauss-Siedel

Solve by Jacobi-iterative

Solution:

Iteration 1

[2,2,3.25]

Iteration 2

[0.9375,2.5,2.5]

Iteration 3
[0.875,1.96875,2.90625]

Iteration 4
[1.0390625,1.9375,3.0703125]

Iteration 5
[1.014,2.02,2.996]

Iterative Methods

Number of iterations

Tolerance value

Precision

Enter initial guess :

Enter Equations :

Solution:

Iteration 1
[2,2.5,2.375]

Iteration 2
[0.9063,1.953,3.0586]

```
Iteration 4
[0.9992,2,3.0004]
```

```
Iteration 5
[0.9999,2,3]
```

Time Complexity

- Each iteration takes $O(n^2)$.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using an array of one dimension and multidimensions was helpful in computations so all has the same type and making sure that program operates on suitable variables.
- And using a lot of lists if we didn't know the number of inputs was helpful.

Comparison

Based on the sample runs, we noticed that **Each Direct Method** outputted the exact solution. **Absolute Error = 0.0, Relative Error = 0.0%.**

Exact Solution = [1, 2, 3].

That is not quite the case in the **Iterative Methods** as **Jacobi-Iterative** outputs **[1.014, 2.02, 2.996]**.

- | | |
|------------------------------------|------------------------------------|
| • Absolute Error [1]: 0.014 | Relative Error [1]: 1.4% |
| • Absolute Error [2]: 0.02 | Relative Error [2]: 1.0% |
| • Absolute Error [3]: 0.004 | Relative Error [3]: 0.1333% |

And **Gauss-Seidel Method** outputs [0.9999, 2, 3]. *This is nearly the exact solution.*

- | | |
|-------------------------------------|----------------------------------|
| • Absolute Error [1]: 0.0001 | Relative Error [1]: 0.01% |
| • Absolute Error [2]: 0.0 | Relative Error [2]: 0.0% |
| • Absolute Error [3]: 0.0 | Relative Error [3]: 0.0% |

If we were to trade **accuracy** for **speed and time** we would use an **Iterative Method** as it only costs $O(n^2)$, but when accuracy matters the most we resort to **Gauss and Gauss-Jordan Direct Methods** but we must know that it costs $O(n^3)$. **LU Decomposition** can be useful for certain applications with the same **time complexity**.

Closure

- Most of the code is commented and follows the clean code principles.
- You can also view the video of **GUI Explanation** [here](#).

Thank you.