NUMERICAL METHODS USING ANGULAR

PROJECT - PHASE 1

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Introduction

Report for Numerical Methods **Project - Phase I**, which requires implementation of a program that calculates the possible solutions of a system of linear equations using:

- Gauss Elimination
- Gauss-Jordan Elimination
- LU Decomposition
- Gauss-Seidel Iterative Method
- Jacobi Iteration Method

Assumptions

- In order to run the program correctly through **Angular**, please start the terminal and type: "*npm install*" to install node modules. You must have node.js installed on your personal computer. **Next in the terminal**, type "*ng serve*" to run the program.
- Please notice that the program runs in → <u>localhost:4200</u> ← click here to go directly to the program page.
- The user has to enter the number of equations in order to proceed.
- The coefficients have to be numbers and not letters.
- The equations are processed and evaluated using functions implemented within each component.
- **TWO** functions are responsible for handling user input and the pseudocode for these functions is:
 - Function to get number of equations.
 - Let n be the number of equations and get it from the user.
 - If (n is null or n == 0) \rightarrow return
 - Else: Create an array of zeros with size n
 - End function.
 - Function to get the equations.

- Recreate the array.
- For i = 0:n
 - Push the equation into the array.
- End for.
- Let precision be a number and get it from the user.
- If (**precision is null or == 0**) it is set to 7 and the program continues.
- Now it is the **hash class** turn to evaluate the equations and split them into a matrix of coefficients and a matrix of results and sometimes an augmented matrix.
- End function.
- The program handles many exceptional cases and never happens to crash or do *unlogical* operations at any time.
- Implemented pivoting whenever we could.
- Implemented scaling and provided steps in some methods.

Gauss Elimination Documentation

Pseudocode

Defining "gaussSolver(arr[][] of number)" function performs forward elimination and then calls back the substitution method.

- array=clone(arr)
- For i:n
 - o pivotAndscale(arr,i) We will explain this function later as it is an implementation of pivoting logic in iteration #i.
 - o For k:n
 - Factor = array[k][i] / / array[i][i]
 - if(factor=0)

hasSolution(array)

End function

End if

- For j : array[k].length As we process on an augmented matrix
 - array[k][j] = array[k][j] factor * array[i][j]
- End for
 - array[k][i] =0
- End for
- End for

backSub(arr)

Defining "hasSolution(arr[][] of number)" function check if function has no solution or infinite number of solution or have unique number of solution

Inf = false

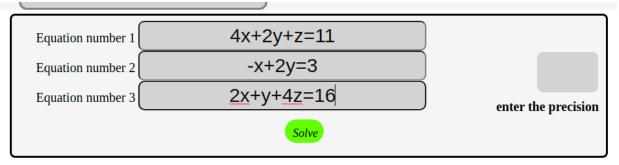
For i:n

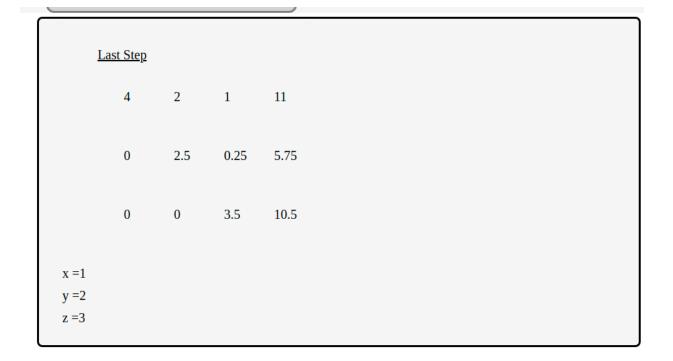
```
rawZero=true
       For j:n
              if(arr[i][j]!=0)
                     rawZero=false
                     Break For
       End for
       if( number of variable >number of equation or rank matrix!= rankAgumented)
              Return "has no solution"
              End function
       Else if (number of variable <number of equation or rawZero=0)
              Inf = true
       End Else if
       End for
       if(inf = true)
              Return "Has infinite number of solution"
       End if
       Else
              Return "Has unique solution"
       End else
Defining "pivotAndscale(arr[][] of number, pivot index)" array need to pivot or not.
Temp =clone(|arr|)
```

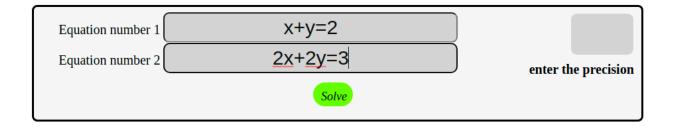
```
BigestInRaw = getBiggestInrow(temp)
  • For i:n
             For j:n
                    Temp[i][j]=Temp[i][j]/BigestInRaw
             End for
   End for
BigetInCoulmnAfterScale = getBiggestIncoulmn(Temp)
if (BigetInCoulmnAfterScale=pivot index)
      Return
End if
Else
      swap (arr[BigetInCoulmnAfterScale], arr[pivot index])
End else
Defining "backSubs(arr[][] of number)" get system solutions
has Solution=hasSolution(arr)
If( has Solution="unique")
      For i=n to 0:
             sum=0
             For j=0 to i
                    arr[i][arr.length] = arr[i][arr.length] - arr[i][j]
             End for
```

End for

Sample Runs







Has No solution

Time Complexity

- Elimination Steps: 2(n³/3).
- Time Complexity: **O(n³)**.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using array with one dimensional and multidimensional it was helpful in computational so all have same type and sure that program do operation to suitable variable
- And using a lot of lists if we didn't know the number of inputs was helpful.

Gauss-Jordan Elimination Documentation

Pseudocode

Assuming we got the user-input matrix we start operating on it as follows.

Defining "gElimination()" function to perform Forward Gauss Elimination, which we would use later in further logic.

- For i = 0:n
 - pivot(i) We will explain this function later as it is an implementation of pivoting logic in iteration #i.
 - o For k = i + 1: n
 - Factor = matrix[k][i] / / matrix[i][i]
 - For **j** = **i** : matrix[k].length As we process on an augmented matrix
 - Matrix[k][j] = matrix[k][j] factor * matrix[i][j]
 - End for.
 - End for.
- End for.

End function.

Defining "pivot(iteration)" function takes an iteration parameter as it pivots per iteration.

- Let pivot = iteration
- Let max = matrix[iteration][iteration]
- For index = iteration + 1 : n
 - Let dummy = matrix[index][iteration]
 - o If dummy > max → Set max to dummy and pivot to i
- End for.
- If pivot != iteration (Checking whether the pivot changed or not)
 - → Swap rows using temporary variables
 - \circ If not \rightarrow do nothing.

End function.

Defining "gjElimination()" function which calls gElimination() then continues backward elimination on the resulting matrix.

- **gElimination()** Apply forward elimination using previously structured Gauss Elimination function.
- If (isSingular()) → return Do not continue as if the matrix is singular, we have a row of zeros and continuing does not make sense.
- For i = n:1
 - pivot(i) Applying pivoting as previously done in gElimination (We did not really have to apply pivoting here, but just in case).
 - o For k = i 1:0
 - Factor = matrix[k][i] / matrix[i][i]
 - For j = n:0
 - Matrix[k][j] = matrix[k][j] factor * matrix[i][j]
 - End for.
 - End for.
- End For.

End function.

Defining "isSingular()" function to check whether the matrix is singular or not.

- Let Determinant = 1
- For index : n, where n is number of equations
 - Determinant = Determinant * matrix[index][index]
- The loop gets the determinant value of the matrix after applying Gauss Elimination
- if(!Determinant) → Singular (boolean attribute) = true

End function.

Defining "**solve()**" function to get matrix solution.

- If (isSingular()) → check if the matrix has no solution or infinite number of solutions and print according to the result.
- If not singular → For i = 0: n

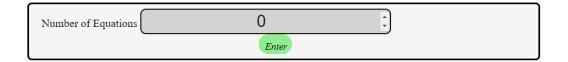
- Let coefficient = matrix[i][i]
- o Let result = matrix[i][n]
- Let solution = result / coefficient
- Push **solution** into the solution array named under **solution**.
- End for.

End Function.

Sample Runs

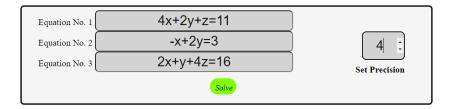
Entering 0 or not entering a number of equations at all does not allow the user to proceed within the program.

Gauss-Jordan Elimination



Assume the user entered 3, the program continues and asks the user to enter the equations. Please **note** that if the user entered no equations the program shows a message saying "No Solution". Precision is totally optional and if the user ignores it, the program proceeds with maximum precision as possible

Gauss-Jordan Elimination



Given the assumption that a precision of 4 was entered along the equations, this is what a user would normally get.

Gauss-Jordan Elimination

```
Calculated the factor in step 0, Factor = -0.25
Row calculation in step 1, Matrix =
[ 4,2,1,11 ]
[0,2.5,0.25,5.75]
[2,1,4,16]
Calculated the factor in step 2, Factor = 0.5
Row calculation in step 3, Matrix =
[ 4,2,1,11 ]
[0,2.5,0.25,5.75]
[ 0,0,3.5,10.5 ]
Calculated the factor in step 4, Factor = 0
Row calculation in step 5, Matrix =
[ 4,2,1,11 ]
[ 0,2.5,0.25,5.75 ]
[ 0,0,3.5,10.5 ]
Calculated the factor in step 6, Factor = 0.07143
Row calculation in step 7, Matrix =
[ 4,2,1,11 ]
[0,2.5,-0.000005,5]
[ 0,0,3.5,10.5 ]
Calculated the factor in step 8, Factor = 0.2857
Row calculation in step 9, Matrix =
[4,2,0.00005,8]
[ 0,2.5,-0.000005,5 ]
[ 0,0,3.5,10.5 ]
```

```
Calculated the factor in step 10, Factor = 0.8

Row calculation in step 11, Matrix =

[ 4,0,0.000054,4 ]

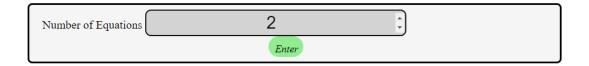
[ 0,2.5,-0.000005,5 ]

[ 0,0,3.5,10.5 ]

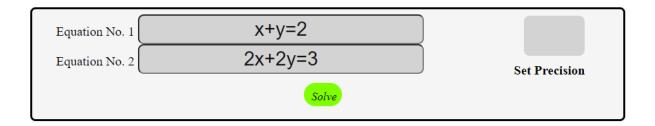
Solution = 1,2,3
```

Let's try another system of linear equations, but this time it consists of 2 equations.

Gauss-Jordan Elimination



Gauss-Jordan Elimination



The program calculates the augmented matrix in steps as follows. This is what the user got for this sample.

Gauss-Jordan Elimination

```
Calculated the factor in step 0, Factor = 0.5
Row calculation in step 1, Matrix =
[ 2,2,3 ]
[ 0,0,0.5 ]
Solution = No solution! Matrix is singular.
```

Time Complexity

• Elimination Steps: 4(n³/3).

• Time Complexity: **O(n³)**.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using an array of one dimension and multidimensions was helpful in computations so all has the same type and making sure that program operates on suitable variables.
- And using a lot of lists if we didn't know the number of inputs was helpful.

LU Decomposition Documentation

Doolittle Decomposition

Pseudocode

Else

1. <u>Defining "Decompose()" method</u>

```
After receiving the coefficients, free terms, and unknowns
                      For i = 0: eqnNo - 1
                               Call partial pivoting method on row i
For j = i + 1: eqnNo
    For k = i : eqnNo
if(k == i) then
MatrixLU [j][i] = matrixLU[j][k] / matrixLU[i][k]
```

2. <u>Defining "partialPivoting(i)" method</u>

matrixLU[j][k] -= matrixLU[j][i] * matrixLU[i][k]

```
For k = i + 1 : eqnNo
         if ( | matrixLU[i][i] | < | matrixLU[k][i] | ) then
```

Swap the 2 rows in the augmented matrix column by column

3. <u>Defining "forward substitution()"method</u>

```
YfreeTerm[0] = stepFreeTerm[0]
For i = 1: eqnNo
        For j = 0:i
            YfreeTerm[i] = stepFreeTerm[i] - matrixLU[i][j] * YfreeTerm[j]
```

4. Defining "backward substitution()" method

```
For i = eqnNo - 1:0
         For j = i + 1 : eqnNo
```

soln[i] = (Yfreeterm[i] - matrixLU[i][j] * soln[j]) / matrixLU[i][i]

5. <u>Defining "solve()"method</u>

Call decompose method

Call forward substitution method

Call backward substitution method

Sample Runs

Doolittle Decomposition

```
The LU Decomposition: A = L. U

The System

4x+2y+z=11
-x+2y=3
2x+y+4z=16

LMatrix
1
-0.25
1
0.5
0
1
```

<u>U Matrix</u>

4

2

1

2.5

0.25

3.5

Free Term

11

3

16

Solving the Equation

$$L.U.X = A.X = b$$

let
$$U \cdot X = Y$$

By Forward Substitution

$$y_1 = 11$$

$$y_2 = 5.75$$

$$y_3 = 10.5$$

By Backward Substitution

$$x = 1$$

$$y = 2$$

$$z = 3$$

Crout LU Decomposition:

Pseudocode

Assuming we have the augmented matrix we start from SplitMatrices where we split the augmented matrix to coefficient matrix and solution matrix.

Then in LUcroutEvaluate function we evaluate L and U matrices:

• LUcroutEvaluate Function: LUcroutEvaluate()

```
For i = 0 : coefficient row length
              Pivoting Function
              For j = 0 : i+1
                      sumL = 0;
                      For k = 0:j
                             sumL = sumL + lower[i][k] * upper[k][j]
                      End For
                      Lower[i][j] = coefficient[i][j] - sumL
              End For
              For j = i+1: coefficient column length
                      SumJ = 0;
                      For k=0:i
                              sumJ = sumJ + lower[i][k] * upper[k][j]
                      End For
                      upper[i][j] = (coefficient[i][j] - sumJ) / lower[i][i]
              End For
       End For
End Function
```

• Pivotion Function: pivoting(index)

We loop on column [index] and find the row of maximum element in this column and exchange row of index parameter and row of maximum element.

```
For j = index : j : augmented row length
  IF absolute coefficient[j][index] > maxRow
            Then maxRow = j
   End IF;
End For
// replacing 2 row if max is found
IF maxRow != index
           For i = 0:coefficient row length
                   temp = this.coff[index][i]
                   coefficient[index][i] = coefficient[maxRow][i]
                   coefficient[maxRow][i] = temp
           End For
           // replace soln matrix
           temp = soln[maxRow]
           soln[maxRow] = soln[index]
           soln[index] = temp
   Fnd IF
```

• yEvaluate Function: yEvaluate()

For i=0: lower row length

```
sum=0
        For j=0 : i
               sum = sum + lower[i][j]*y[j]
        End For
        Push (soln[i]-sum) to y matrix
 End For
xEvaluate Function: xEvaluate()
 For i=upper row length - 1:-1:i--
        sum=0
        For j=i+1: upper row length
               sum = sum + upper[i][j]*x[j]
        End For
        X[i] = y[i]-sum
 End For
```

Sample Runs

```
LU Crout Decomposition A=LU
Upper Matrix of Crout Decomposition:
    1 2.5 4.5
    0 1 1.9231
    0 0 1
lower matrix of crout decomposition:
    2 0 0
    6 -13 0
    2 -2 -0.1538
free terms
    5
    3
    4
```

```
Solving Equation:
AX=B && A=LU
AX=LUX
Let Y = UX
LY=B
Y matrix:
    2.5
    0.92308
    -5.5017
Y=UX
X matrix:
    -1.5004
    11.503
    -5.5017
```

Cholesky Decomposition

PseudoCode

1. <u>Defining "checkMatrix()" method</u>

```
For i = 0: eqnNo
For j = 0: i
```

```
If coefficients[i][j] != coefficients[j][i] then valid = false
```

Call a part of decompose method in *Doolittle decomposition* to get the upper triangular matrix

```
For i = 0 : eqnNo

If coefficients[i][i] <= 0 then valid = false

////all methods of cholesky will be called if valid == true
```

2. <u>Defining "Decompose()" method</u>

3. <u>Defining "forwardSubstitution()"method</u>

4. Defining "backwardSubstitution()" method

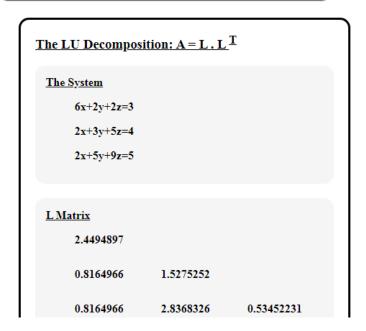
```
For i = eqnNo - 1 : 0

For j = i + 1 : eqnNo

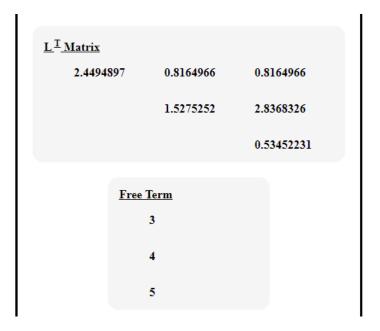
soln[i] = (Yfreeterm[i] - matrixLU[i][j] * soln[j]) / matrixLU[i][i]
```

Sample Runs

Cholesky Decomposition



Note that the matrix is +ve definite symmetric matrix



Solving the Equation

$$L.L^T.X=A.X=b$$

let
$$L^T \cdot X = Y$$

By Forward Substitution

$$y_1 = 1.22$$

$$y_2 = 1.97$$

$$y_3 = -2.96$$

By Backward Substitution

$$x_1 = -1.5140136$$

$$x_2 = 11.573884$$

2.

Cholesky Decomposition

About the Syste	em			
equation 1	2x + 5y = 8		enter the precision	9
equation 2	5x + 5y = 6			
		Solve		

Note that matrix here is symmetric but not positive definite

Cholesky Decomposition

"No Soln"

Time Complexity

- Elimination Steps: **K** (n³/3), where K is the number of equations.
- Time Complexity: **O(n³)**.

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using an array of one dimension and multidimensions was helpful in computations so all has the same type and making sure that program operates on suitable variables.
- And using a lot of lists if we didn't know the number of inputs was helpful.

Iterative Methods Documentation:

Pseudocode → **Jacobi-iterative**

Defining "**implementJacobi()**" function to perform the jacobi iterative method, which we would use later in further logic.

```
gaussSiedelResults = [][]
gaussSiedelResults.push(intialGuess)
   • If (noOflterations != -1)
              For count: noOfIterations
                      For i: intialGuess.length
                             For j: intialGuess.length
                                    tempArray[i] += a[i][j] * intialGuess[j]
                             End for
                             tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]
                      End for
                      intialGuess = tempArray.clone
                     jacobiMethodResults.push(intialGuess)
              End for
      Else
              While (relativeError >= eTolerance)
```

For i: intialGuess.length

For j: intialGuess.length

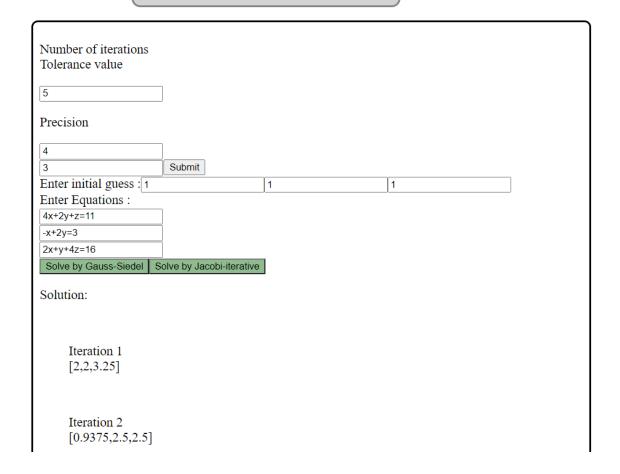
```
tempArray[i] += a[i][j] * intialGuess[j]
                     End for
                     tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]
                     relativeError = Math.max(relativeError, (tempArray[i] - intialGuess[i]) /
       tempArray[i] * 100)
              End for
              intialGuess = tempArray.clone
              jacobiMethodResults.push(intialGuess)
       End for
       Return gaussSiedelResults
End function
Pseudocode → Gauss-Siedel
gaussSiedelResults = [][]
gaussSiedelResults.push(intialGuess)
      If (noOfIterations != -1)
              For count: noOfIterations
                     For i: intialGuess.length
                             For j: intialGuess.length
                                    tempArray[i] += a[i][j] * tempArray[j]
                             End for
                             tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]
```

```
End for
                     intialGuess = tempArray.clone
                     jacobiMethodResults.push(intialGuess)
              End for
      Else
              While (relativeError >= eTolerance)
                     For i: intialGuess.length
                     For j: intialGuess.length
                            tempArray[i] += a[i][j] * tempArray[j]
                     End for
                     tempArray[i] = (a[i][a[0].length - 1] - tempArray[i]) / a[i][i]
                     relativeError = Math.max(relativeError, (tempArray[i] - intialGuess[i]) /
       tempArray[i] * 100)
              End for
              intialGuess = tempArray.clone
              jacobiMethodResults.push(intialGuess)
       End for
       Return gaussSiedelResults
End function.
```

Sample run:

Example 1: Jacobi Iterative Method, Example 2: Gauss-Seidel Method.

Iterative Methods

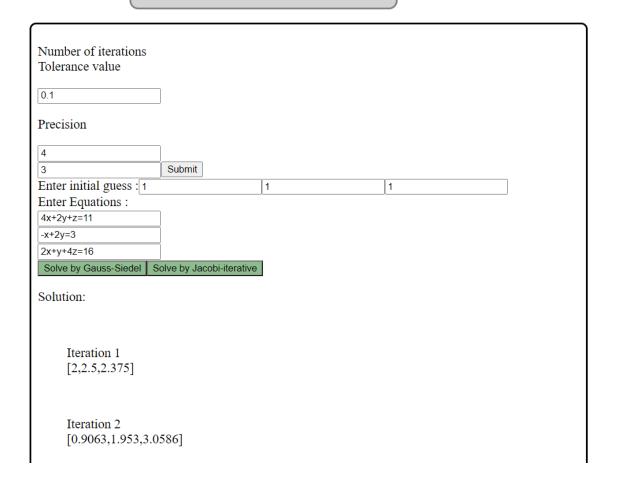


Iteration 3
[0.875,1.96875,2.90625]

Iteration 4
[1.0390625,1.9375,3.0703125]

Iteration 5
[1.014,2.02,2.996]

Iterative Methods



Iteration 4 [0.9992,2,3.0004]

Iteration 5 [0.9999,2,3]

Time Complexity

• Each iteration takes O(n²).

Data Structure Used

- Using Map in evaluating expressions helps in reading expressions and evaluating them by making each key have a coefficient of variable in order.
- Using an array of one dimension and multidimensions was helpful in computations so all has the same type and making sure that program operates on suitable variables.
- And using a lot of lists if we didn't know the number of inputs was helpful.

Comparison

Based on the sample runs, we noticed that **Each Direct Method** outputted the exact solution. **Absolute Error = 0.0**, **Relative Error = 0.0%**.

Exact Solution = [1, 2, 3].

That is not quite the case in the **Iterative Methods** as **Jacobi-Iterative** outputs **[1.014, 2.02, 2.996]**.

• Absolute Error [1]: 0.014 Relative Error [1]: 1.4%

• Absolute Error [2]: 0.02 Relative Error [2]: 1.0%

• Absolute Error [3]: 0.004 Relative Error [3]: 0.1333%

And Gauss-Seidel Method outputs [0.9999, 2, 3]. This is nearly the exact solution.

Absolute Error [1]: 0.0001

 Absolute Error [2]: 0.0
 Absolute Error [2]: 0.0

 Relative Error [2]: 0.0%
 Absolute Error [3]: 0.0
 Relative Error [3]: 0.0%

If we were to trade accuracy for speed and time we would use an Iterative Method as it only costs $O(n^2)$, but when accuracy matters the most we resort to Gauss and Gauss-Jordan Direct Methods but we must know that it costs $O(n^3)$. LU Decomposition can be useful for certain applications with the same time complexity.

Closure

- Most of the code is commented and follows the clean code principles.
- You can also view the video of **GUI Explanation** here.

Thank you.