

DATE: / /

SUBJECT:

19016024 على مصطفى السيد الديك لادوى

$$\boxed{1} \quad 32 - (15 - 5) = 22 \text{ elements}$$

$$\boxed{2} \text{ @ Void enqueue (object data) \{$$

$$Q[r] \leftarrow \text{data};$$

$$r++; \}$$

$$\text{Void dequeue \{ object front} \leftarrow Q[0];$$

$$\text{for (j = 0; j < r-1; j++) \{}$$

$$Q[j] \leftarrow Q[j+1]; \}$$

$$r--;$$

$$\text{return front} \}$$

* The dequeue operation has $O(n)$ Time

Complexity. The size of the Q is fixed

\textcircled{b} i. Q is empty: $f = r = 0$ | * Fixed size of the array
 ii. Q is Full: $r = N$ | * waste of memory

\textcircled{c} i. Q is empty: $f = r$ | * When $f = r$, we can't tell
 ii. Q is full: $f = r$ | whether the Q is full or empty

\textcircled{d} i. Q is empty: $f = r$
 Q is full: $f = r + 1$

$$\text{Void enqueue (object data) \{}$$

$$\text{if (f == r + 1) \{ raise FullQ \}}$$

$$Q[r] \leftarrow \text{data};$$

$$r \leftarrow (r + 1) \% N; \}$$


```

[3] void find(Q, S, X) {
    found ← false;
    for (j = 0; j ≤ n-1; j++) {
        curr ← S.Pop();
        if curr == X {
            found ← true;
        }
        Q.enqueue(curr);
    }
    for (j = 0; j ≤ n-1; j++) {
        S.push(Q.dequeue());
    }
    for (j = 0; j ≤ n-1; j++) {
        Q.enqueue(S.Pop());
    }
    for (j = 0; j ≤ n-1; j++) {
        S.push(Q.dequeue());
    }
    return found;
}

```

```

[4] void Push(object data) {
    Q1.enqueue(data);
}

object Pop() {
    if is Empty() { raise EmptyQ Exception; }
    for (j = 0; j < Q1.size(); j++) {
        Q2.enqueue(Q1.dequeue());
    }
    object curr ← Q1.dequeue();
    for (j = 0; j < Q2.size(); j++) {
        Q1.enqueue(Q2.dequeue());
    }
    return curr;
}

```



```

[5] void Push(item i) { int Priority ← 0;
    Q.insert(i, Priority);
    Priority ++; }
item Pop() { if Q.size == 0 { raise EmptyStack }
    Priority --; return Q.DeleteMax(); }

```

[6] A Priority Queue

[7] A Priority Queue will be useful in this Problem because the events will be automatically sorted if the time stamp is provided as a key which means that the smallest time stamp will be always found at the front.

```

[8] int cntLeafNodes(Node root) {
    if root is Null { return 0 }
    if root.left() is Null and root.right() is Null {
        return 1 }
    return cntLeafNodes(root.right())
        + cntLeafNodes(root.left()) + 1
}

```

```

[9] bool isEq(Node root1, Node root2) {
    if root1 is Null and root2 is Null {
        return True }
    if root1.data() != root2.data() {
        return False }
    return isEq(root1.right(), root2.right())
        and isEq(root1.left(), root2.left())
}

```


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```
10. void swapTree(Node T) {  
    if T is Null { return }  
    temp = T.left().data()  
    T.left().data() = T.right().data()  
    T.right().data() = temp  
    swapTree(T.left())  
    swapTree(T.Right())  
}
```