# Benefits and Drawbacks of SYN Cookies in Linux Kernel

# Ali Ghaffarian Proposal for evaluating SYN cookies in 2024

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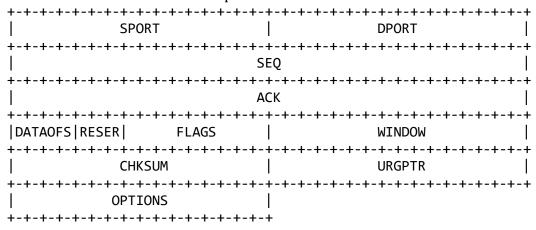
## **Summary**

In an estimated three-week time period, we will examine the costs and benefits of SYN cookies as well as the effects of SYN floods on a modern GNU/Linux host with SYN cookies disabled and the viability of previously known attacks against SYN cookies, such as connection spoofing.

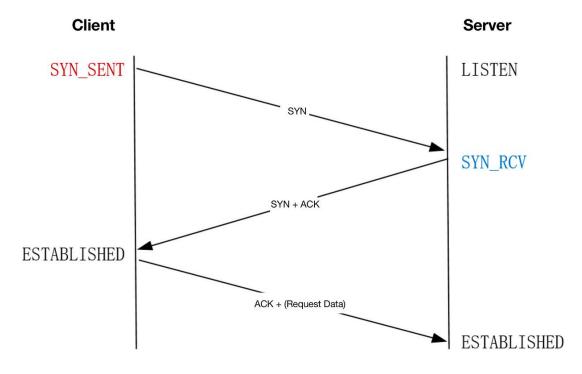
## Introduction

#### **TCP**

Transmission Control Protocol is the most used reliable data transfer protocol in TCP/IP networks and considered a complex one.



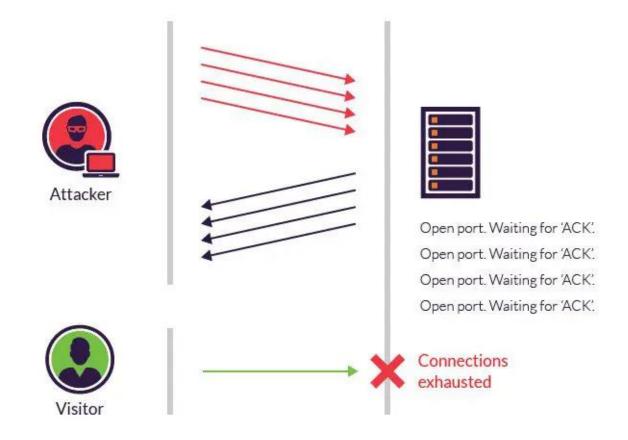
## **Three Way Handshake**



[1]

The two ends of the TCP connection must know each other's initial sequence number before the actual reliable data transfer can begin, one of the responsibilities of TCP three way handshake is exactly this.

#### **SYN Floods**



[2]

TCP SYN flooding is a DoS attack that leverages the reserved resources for pending connections.

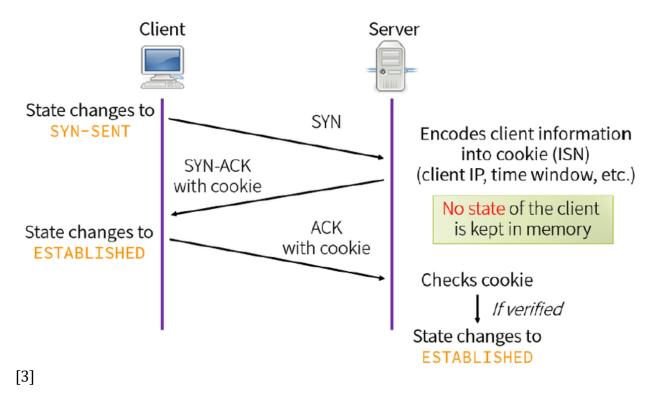
Because TCP servers normally reserve such resources only when a connection is initiated, the attacker can flood the server with SYN requests that include spoofed source IP address.

This way the backlog of the TCP server will eventually fill up, making the server always waiting for the non-existent clients to complete their connections.

When using the server's timeout and backlog as the attack vector, the cost of the attack drops dramatically, as a typical TCP server will wait several seconds until it determines that the client's not going to complete the connection.

In some circumstances, the server's memory will get full if the backlog is big enough.

#### **SYN Cookies**



SYN floods were so successful because of the stateful nature of the TCP protocol.

One way to stop SYN floods from being so cheap and effective, is to handle the three-way handshake differently, while preserving the rest of the protocol.

In other words, handle the handshake statelessly without breaking the protocol, that is, to forget the client ever initiated a connection, but having a method to recover all connection's crucial data when the client completes the connection.

SYN cookies do this by carefully encoding the connection's data in the server's initial sequence number (ISN) when acknowledging the client's SYN request.

This way when the client sends the final packet of the three-way handshake all of the required data to reconstruct the connection will be recoverable.

#### **TCP Crucial Fields**

The data a TCP server needs to be able to reconstruct the connection includes:

- client's ISN
- server's ISN
- client's IP address
- server's IP address
- client's port number
- server's port number
- max segment size (MSS)

Client's ISN and server's MSS are encoded into server's ISN; The rest can already be recovered from client's ACK packet.

#### **Encoding TCP options**

TCP is an old transport layer protocol, it was designed for networks of its time that tended to lose packets frequently and have lower bandwidth, in which devices reside with small buffers to dedicate to the connection, because of this certain TCP options are needed to be negotiated so that modern devices can use all of their offered load while using TCP as the transport layer protocol.

The most important of which includes:

- Window Scale: each device advertises a window scale value, indicating each declared window size needs to be multiplied by 2^window scale.
- Selective Acknowledgement (SACK): when using SACK as the loss recovery each device can declare the sequence number of the lost data and continue the data transfer without needing to slow down.

Linux encodes these two options in the timestamp field of the SYN-ACK packet only if both devices support timestamps [4].

## **Research Goals**

We try to answer the following questions:

- 1. How does SYN floods effect a SYN cookie disabled 2024 GNU/Linux server?
- 2. Are the known attacks against SYN cookies still viable?
- 3. What are the Drawbacks of SYN cookies in a 2024 GNU/Linux host and how much they impact the performance of the system?

## **Testing Environment**

Rule of the device will probably vary based on the experiment.

#### Device 1:

OS: Arch Linux x86\_64

Host: ASUS TUF Gaming F17 FX706HC\_FX706HC 1.0

Kernel: 6.11.5-arch1-1

CPU: 11th Gen Intel i5-11400H (12) @ 4.500GHz

Memory: 15732MiB

Network controller: MEDIATEK Corp. MT7921 802.11ax PCI Express Wireless

Network Adapter

#### Device 2:

OS: Android 13 aarch64 Host: POCO 2201116PG

Kernel: 5.4.259-qgki-gcb8e1baa6b45

CPU: (8) @ 1.804GHz Memory: 5437MiB

Network controller: Not root on this device, can't check

#### **Device 3:**

OS: Debian GNU/Linux 12 (bookworm) x86\_64

Host: Aspire A315-58G V1.26

Kernel: 6.1.0-22-amd64

CPU: 11th Gen Intel i3-1115G4 (4) @ 4.100GHz

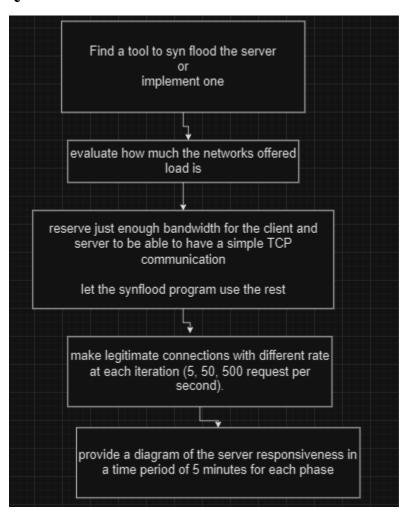
Memory: 11739MiB

Network controller: Intel Corporation Wi-Fi 6 AX201

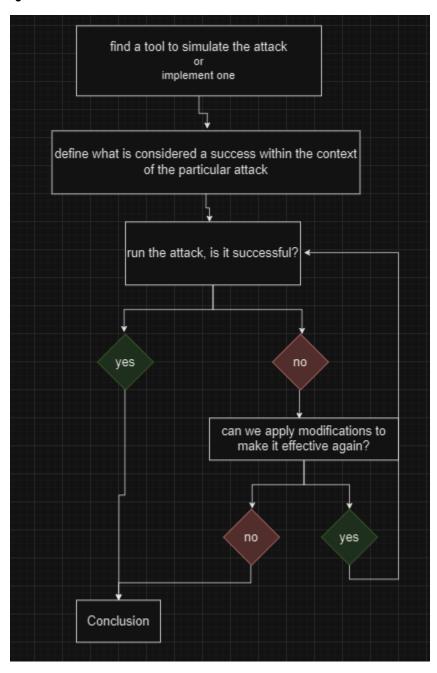
## **Overview of Research Plan**

Questions / Weeks	1	2	3
How does Syn floods effect a syn cookie disabled 2024 gnu/linux server?			
Are the known attacks againts syn cookies still viable?			
What are the Drawbacks of syn cookies in a 2024 linux kernel?			

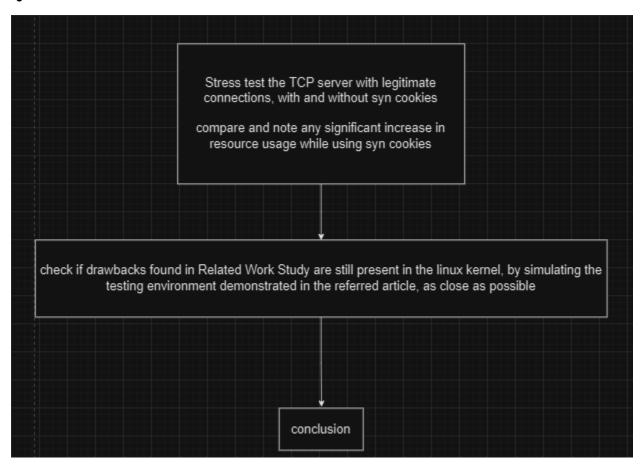
## **Question 1**



## **Question 2**



## **Question 3**



## **Related Work Study**

#### **Linux Kernel**

#### Cookie generation:

The encoded information (client's ISN and MSS) is combined with two cryptographically hashed values (like HASH(secret, saddr, sport, daddr, dport, secret)) secret keys and a count value that is incremented every minute to make the cookie and sent to the client [4][5].

Cookie validation is done via extracting the encoded MSS from client's Acknowledge number and checking if it's out of range [5].

#### **Minisocks**

Linux stores half opened connections via tcp\_request\_sock structure which is 96 bytes compared to a regular tcp\_sock which is a 1616 structure [4].

Considering how smaller the minisock is, the possibility of running out of memory during a SYN flood is extremely low.

#### **Benefits of SYN Cookies**

no need to track any half-opened connections.

#### **Drawbacks**

- Limited support for TCP options.
- SYN cookies are ineffective against SYN flood attacks that target the victim's bandwidth.
- In effective connection spoofing scenario, firewalls focusing on filtering connections by SYN flag will be bypassed.
- Needs more CPU resources to be computed and validated.

#### Packet Loss??

in 3th march 2022, Kevin Graham stated that because of how MSS is encoded into the server's ISN, a TCP segment with 3 bytes of data can be lost once in a connection that is initiated using SYN cookies [6].

Other operating systems might be suffering from the same problem as well, but due to the nature of close source software, we can't state anything unless special experiments are done.

#### **Connection Spoofing Attacks**

As a Linux TCP server in SYN cookie mode accepts connection after receiving a valid SYN cookie from the client, connections can be made by guessing SYN cookies.

In 13th August 2013, Jakob Lell provided a method to guess a valid SYN cookie in around 8 minutes with a gigabit connection to the target [7].

### References

- [1] https://miro.medium.com/v2/resize:fit:1400/1\*BHjxQmlldZv7Qx6TNCG8TQ.png
- [2] https://www.imperva.com/learn/wp-content/uploads/sites/13/2019/01/synflood.jpg.webp
- [3]

 $https://www.researchgate.net/publication/336872866/figure/fig2/AS:95629686606643\\ 3@1605010348044/Mechanism-of-SYN-cookies-4.ppm$ 

- [4] https://lwn.net/Articles/277146/
- [5] https://github.com/torvalds/linux/blob/master/net/ipv4/syncookies.c
- [6] https://wpbolt.com/syn-cookies-ate-my-dog-breaking-tcp-on-linux
- [7] https://www.jakoblell.com/blog/2013/08/13/quick-blind-tcp-connection-spoofing-with-syn-cookie