Scheme 2 Core Typing Relation

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The Scheme2 Core static semantics is given as a three-place relation between a variable typing context Γ , expression e, and type T, written $\Gamma \vdash e : T$, pronounced "under Γ , e has type T". Formally, the static semantics is taken to be the smallest relation closed under the following rules:

Variables and values

$$\frac{\text{T-num}}{\Gamma \vdash n : \text{num}} \qquad \frac{\text{T-false}}{\Gamma \vdash false : \text{bool}} \qquad \frac{\text{T-true}}{\Gamma \vdash true : \text{bool}} \qquad \frac{\frac{\text{T-var}}{(x : T) \in \Gamma}}{\Gamma \vdash x : T}$$

Unary operators

$$\frac{\Gamma\text{-NOT}}{\Gamma \vdash e : \text{bool}}$$
$$\frac{\Gamma \vdash (\text{not } e) : \text{bool}}{\Gamma \vdash (\text{not } e) : \text{bool}}$$

Binary operators

$$\frac{\Gamma\text{-BINOP-ARITH}}{\Gamma \vdash e_1 : \text{num}} \quad \Gamma \vdash e_2 : \text{num} \quad b \in \{+, *, -, /\}}{\Gamma \vdash (b e_1 e_2) : \text{num}}$$

$$\frac{\Gamma\text{-BINOP-COMP}}{\Gamma \vdash e_1 : \text{num}} \quad \Gamma \vdash e_2 : \text{num} \quad b \in \{=, <\}}{\Gamma \vdash (b e_1 e_2) : \text{bool}}$$

Conditionals

$$\frac{ \substack{ \text{T-IF} \\ \Gamma \vdash e_{cond} : \text{bool} } \quad \Gamma \vdash e_1 : T \qquad \Gamma \vdash e_2 : T }{ \Gamma \vdash (\text{if } e_{cond} \ e_1 \ e_2) : T }$$

Functions

$$\frac{ \begin{array}{l} \text{T-FUN} \\ \Gamma, x: T_1 \vdash e: T_2 \\ \hline \Gamma \vdash (\text{fun } x \, T_1 \, e): T_2 \end{array}}{ \begin{array}{l} \Gamma \vdash (e_1 \, e_2): T_2 \end{array}} \qquad \frac{ \begin{array}{l} \text{T-APP} \\ \Gamma \vdash e_1: T_1 \rightarrow T_2 \end{array}}{ \begin{array}{l} \Gamma \vdash (e_1 \, e_2): T_2 \end{array}}$$

Recursion

$$\frac{\Gamma\text{-REC}}{\Gamma, x: T \vdash e: T} \\ \frac{\Gamma, x: T \vdash e: T}{\Gamma \vdash (\text{rec } x \ T \ e): T}$$

Products

$$\begin{array}{ll} \text{T-pair} & \text{T-fst} \\ \frac{\Gamma \vdash e_1 : T_1}{\Gamma \vdash (\text{pair } e_1 \: e_2) : T_1 * T_2} & \frac{\Gamma \vdash \text{Fst}}{\Gamma \vdash (\text{fst } e) : T_1} & \frac{\Gamma \vdash \text{snd}}{\Gamma \vdash (\text{snd } e) : T_2} \\ \end{array}$$

Sums

$$\begin{split} & \overset{\text{T-INL}}{\Gamma \vdash e : T_1} & \overset{\text{T-INR}}{\Gamma \vdash (\text{inl } T_2 \, e) : T_1 + T_2} \\ & \frac{\Gamma \vdash e : T_2}{\Gamma \vdash (\text{inl } T_1 : e) : T_1 + T_2} \\ & \frac{\text{T-CASE}}{\Gamma \vdash e_1 : T_1 + T_2} & \Gamma \vdash e_2 : T_1 \to T & \Gamma \vdash e_3 : T_2 \to T \\ \hline & \Gamma \vdash (\text{case } e_1 \, e_2 \, e_3) : T \end{split}$$

Lists

$$\begin{split} \frac{\text{T-NIL}}{\Gamma \vdash (\text{nil } T) : \text{list } T} & \frac{\prod\limits_{\Gamma \vdash e_1 : T}^{\text{T-CONS}} \Gamma \vdash e_2 : \text{list } T}{\Gamma \vdash (\text{cons } e_1 e_2) : \text{list } T} \\ \frac{\text{T-FOLD}}{\Gamma \vdash e_1 : \text{list } T_1} & \Gamma \vdash e_2 : T_1 \to T_2 \to T_2 & \Gamma \vdash e_3 : T_2}{\Gamma \vdash (\text{fold } e_1 e_2 e_3) : T_2} \end{split}$$