Scheme1 Core Typing Relation

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The Scheme1 Core static semantics is given as a three-place relation between a variable typing context Γ, expression e, and type T, written $\Gamma \vdash e : T$, pronounced "under Γ, e has type T". Formally, the static semantics is taken to be the smallest relation closed under the following rules:

Variables and values

$$\frac{}{\Gamma \vdash n : \text{num}} \text{ T-num}$$

$$\frac{}{\Gamma \vdash \mathit{false} : \text{bool}} \text{ T-false}$$

$$\frac{}{\Gamma \vdash \mathit{true} : \text{bool}} \text{ T-true}$$

$$\frac{(x : \text{T}) \in \Gamma}{\Gamma \vdash x : \text{T}} \text{ T-var}$$

Unary operators

$$\frac{\Gamma \vdash e : \text{bool}}{\Gamma \vdash (\text{not } e) : \text{bool}} \text{ T-not}$$

Binary operators

$$\frac{\Gamma \vdash e_1 : \text{num} \qquad \Gamma \vdash e_2 : \text{num} \qquad b \in \{+, *, -, /\}}{\Gamma \vdash (b\ e_1\ e_2) : \text{num}} \text{ T-binop-arith}$$

$$\frac{\Gamma \vdash e_1 : \text{num} \qquad \Gamma \vdash e_2 : \text{num} \qquad b \in \{=, <\}}{\Gamma \vdash (b\ e_1\ e_2) : \text{bool}} \text{ T-binop-comp}$$

Functions

$$\frac{(\Gamma, x : T_1) \vdash e : T_2}{\Gamma \vdash (\text{fun } x\ T_1\ e) : T_2} \text{ T-fun}$$

$$\frac{\Gamma \vdash e_1 : T_1 \to T_2 \qquad \Gamma \vdash e_2 : T_1}{\Gamma \vdash (e_1\ e_2) : T_2} \text{ T-app}$$

Conditionals

$$\frac{\text{T-\textsc{if}}\quad \Gamma \vdash e_{cond} : \text{bool} \qquad \Gamma \vdash e_1 : T \qquad \Gamma \vdash e_2 : T}{\Gamma \vdash (\text{if } e_{cond}\ e_1\ e_2) : T}$$