



1 Hamming distance

In several cases, we are given an input pattern and we need to find similar patterns from a database of patterns. If we are not able to find exact-matching patterns, we need to find approximate-matching patterns. One measure of approximate matching is **hamming distance**: Two patterns have a hamming distance $= d$ if they have equal number of characters and a character-by-character comparison results in d unequal characters. For example, the hamming distance between **AGC** and **ACG** is $d=2$, while the hamming distance between **AGC** and **TGC** is $d=1$. The hamming distance between **AGC** and **AG** is not defined since they have different lengths.

First, we build a trie of the database patterns, and then we use recursion to backtrack over all possibilities as shown in the following example where the input pattern is **AGC** and the database patterns are { **AGA**, **AA**, **AAG**, **GAAG**, **TCG** } and we need to find all patterns with **hamming distance** ≤ 2 :

2 Edit distance

Another measure of approximate matching is **edit distance**: Two patterns have an edit distance $= d$ if the minimum number of insertion, deletion, and/or substitution operations required to convert one pattern to the other one $= d$. For example, the edit distance between **AGC** and **AC** is **d=1** (One deletion operation of **G** from **AGC** will convert it to **AC**), while the edit distance between **AGC** and **GCT** is **d=2** (Deletion of **A** from **AGC** then insertion of **T**) and the edit distance between **AGC** and **GGC** is **d=1** (Substitution of the first character).

We can also use recursion to backtrack over all possibilities as shown in the following example where the input pattern is **AGC** and the database patterns are { **AGA**, **AA**, **AAG**, **GAAG**, **TCG** } and we need to find all patterns with **edit distance** ≤ 1 :