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# CS 201 Data Structure II (L2 / L5)

## Binary Tree / Treap

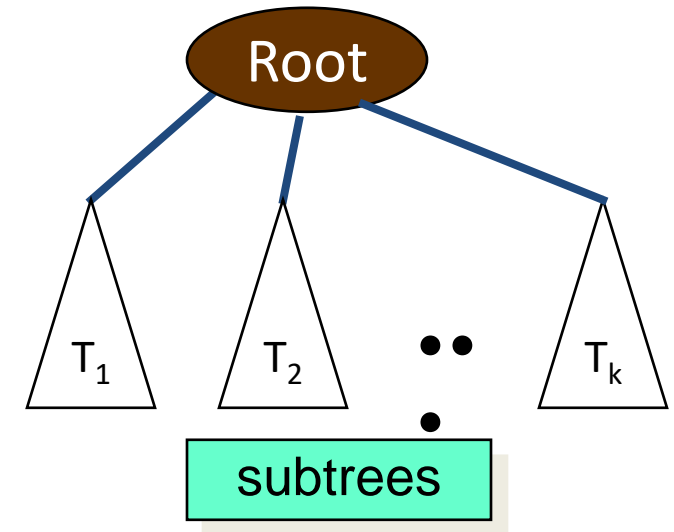
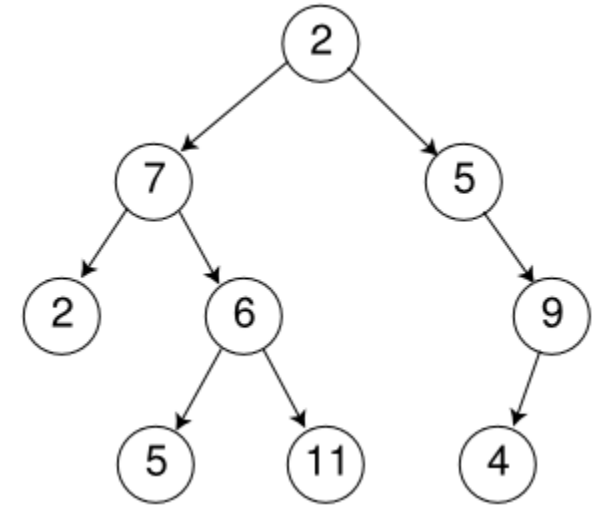
Section 6.1, 6.2, 7.2

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# Tree and Binary Trees

- Hierarchical data structure (non-linear)
- A set nodes and a set of directed edges that connects pair of nodes (non-recursive definition)
- Either a tree is empty or it consists of a root and zero or more nonempty subtrees  $T_1, T_2, \dots, T_k$ , each of whose roots are connected by an edge from the root.(recursive definition)
- Acyclic



# Binary Tree

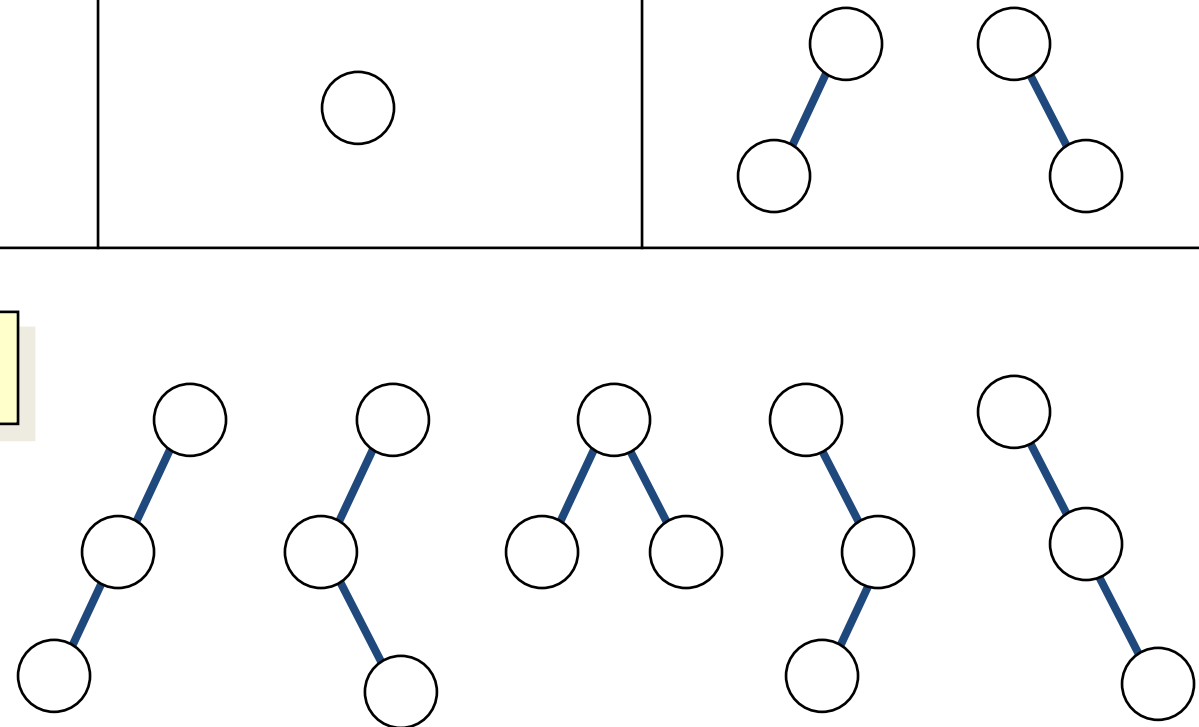
- A tree with no node more than two children
- A binary tree is either **empty**, or it consists of a node called the **root** together with **TWO binary trees** called the left subtree and the right subtree of the root.

Empty tree

Tree of size 1

Tree of size 2

Tree of size 3



# Binary Tree Representation



[how to represent binary tree using node class]



# Terminologies

- Root node
- Leaf / parent / sibling
- Sub-tree
- Path/Path-length
- Height of a node / tree
- Size of a node / tree
- Depth of a node / tree



# Size / Height of a binary tree

[recursive code for size/height]



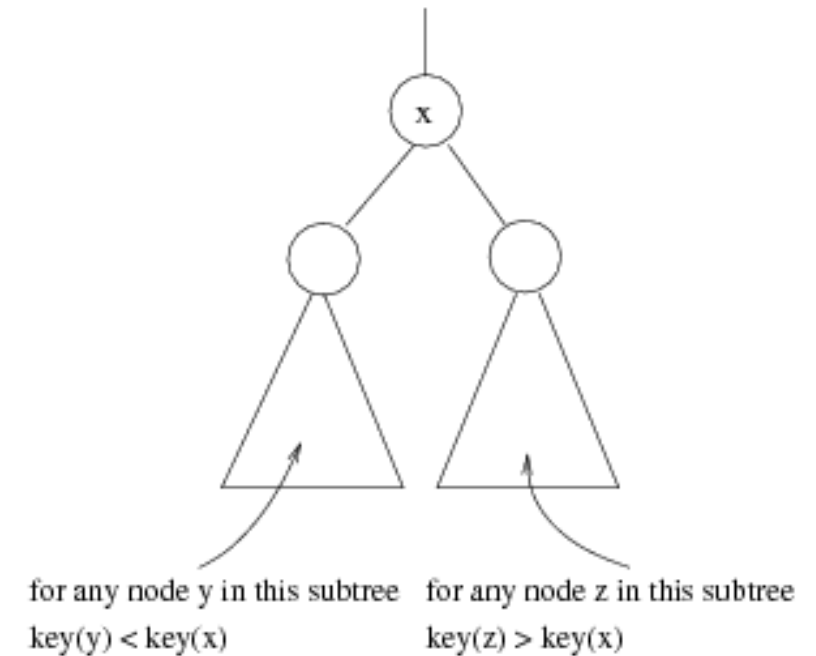
# Traversal

- Pre/Post/In-order

[traversal through activity]

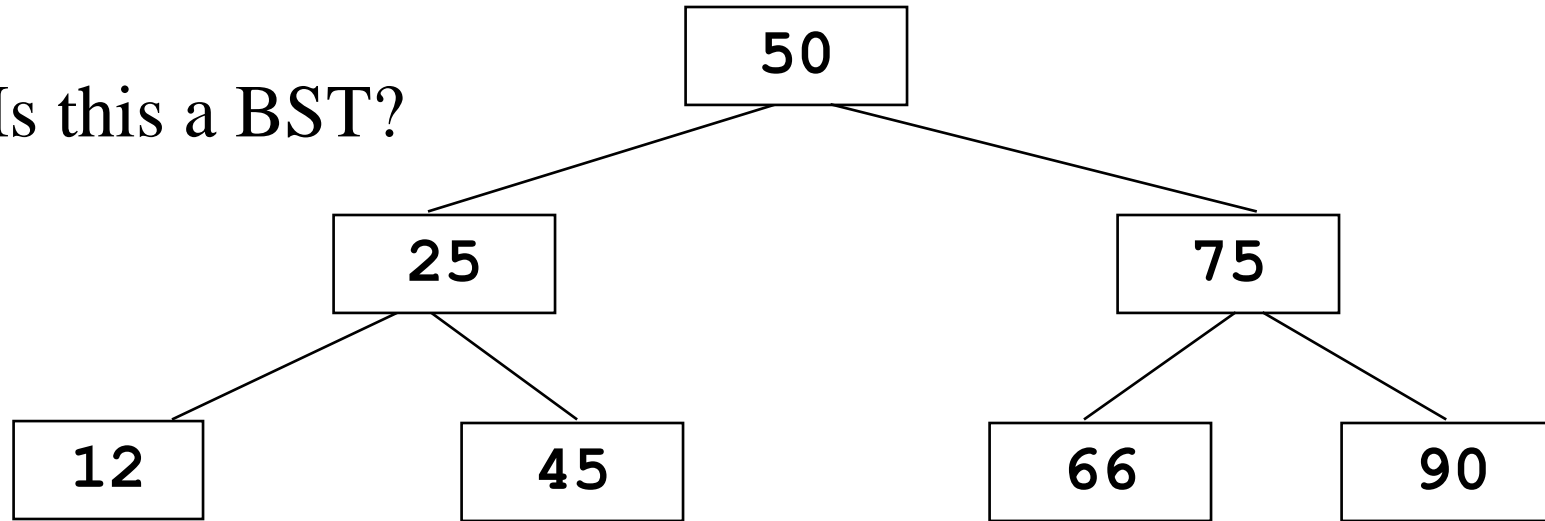
# Binary Search Tree

- A Binary Tree with following properties
- For every node X
  - All the keys in its left sub-tree are smaller than the key value in X
  - All the keys in its right sub-tree are larger than the key value in X

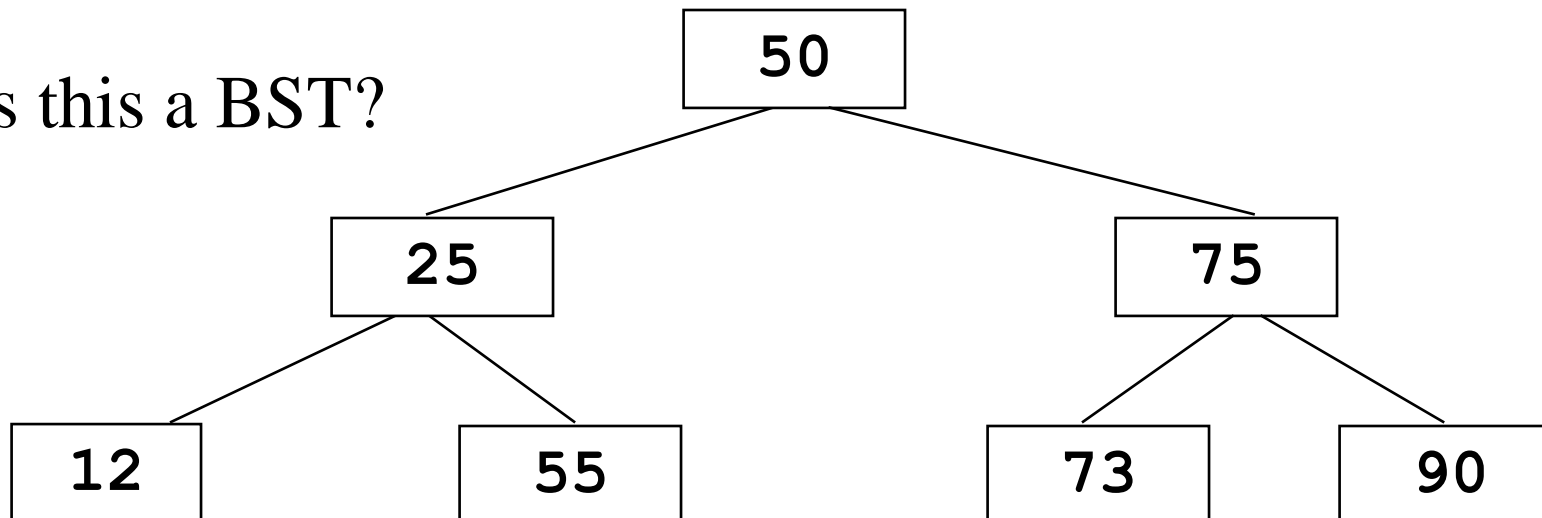




Is this a BST?



Is this a BST?





# Operations on BST

- Search
- Insert
- Delete
  - Case 1 – leave
  - Case 2 – one child
  - Case 3 – two children
- SSET as BST??

# Randomizes Binary Search Tree

- Search –  $O(h)$
- Insert –  $O(h)$
- Delete –  $O(h)$
- The height of the tree depends on the sequence of numbers
  - 0,1,2,3,4,5,6,7,8,9,10,11,12,13,14
    - $h = n$
  - 7,3,11,1,5,9,13,0,2,4,6,8,10,12,14
    - $h = \log n$
- How to prevent  $h = n$ ?

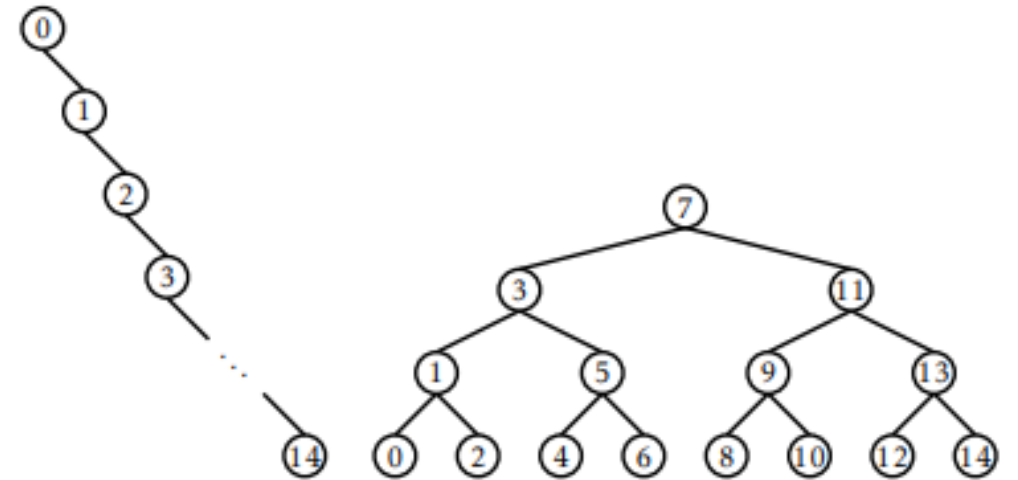
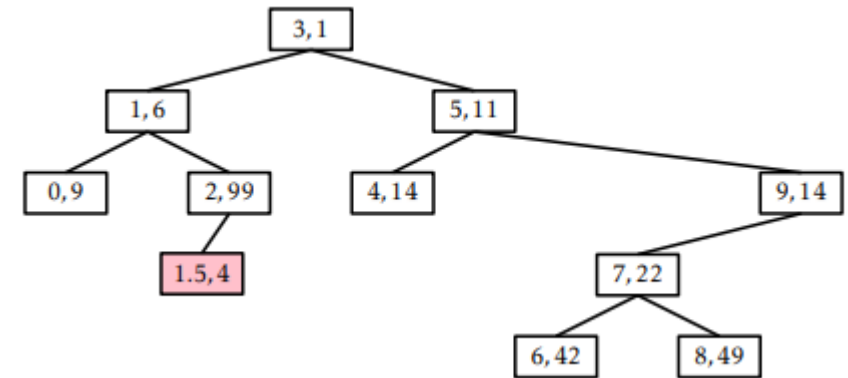


Figure 7.1: Two binary search trees containing the integers 0,..., 14.

# Treap: Tree + Heap

- A Binary Search Tree with an additional value 'p' for each node and one additional property
  - At every node u, except the root,  $u.parent.p < u.p$  (Heap Property)
  - No parent should have a higher 'p' value than its children
  - The value of p – assigned randomly and unique
- add(x):
  - Create a new node 'u' with value x
  - Assign a random value for p
  - Insert using add(x) of BST – 'u' should be a leaf
    - BST property maintains – heap property might not
  - Fix by Bubble Up



# Rotations

- We can fix heap property by performing rotations
- Rotate Right = make the left child as a parent
- Rotate Left = make the right child as a parent
- Decrease (/increase) the depth by one

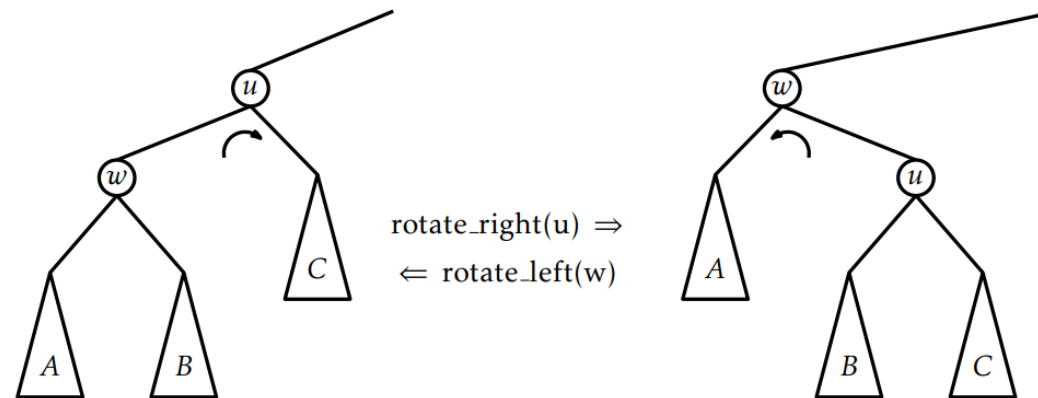
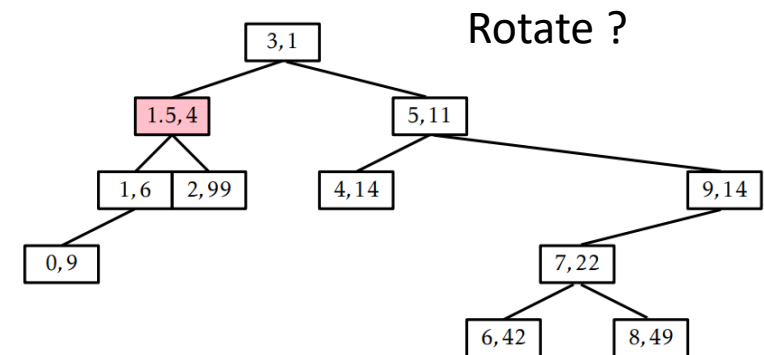
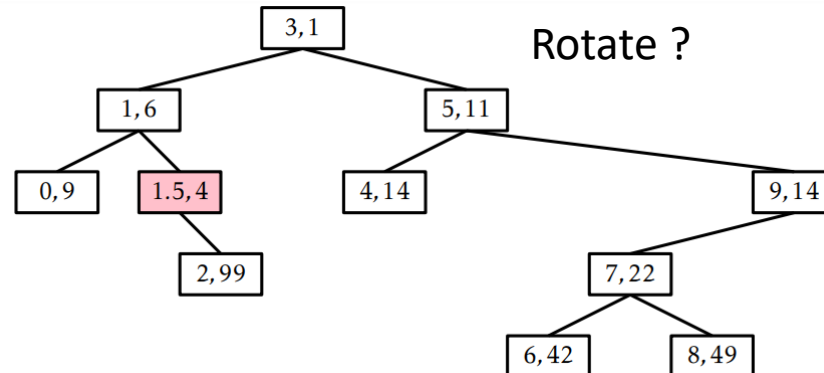
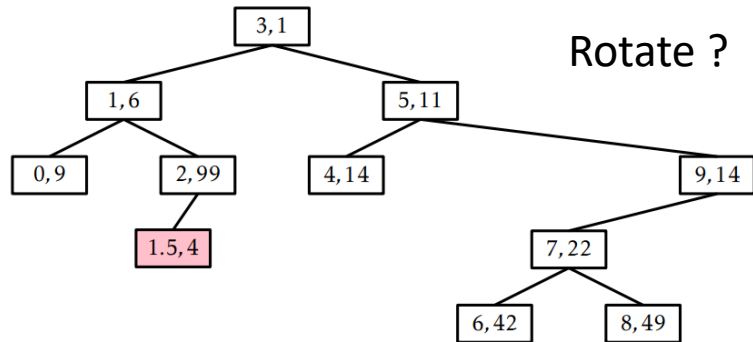


Figure 7.6: Left and right rotations in a binary search tree.

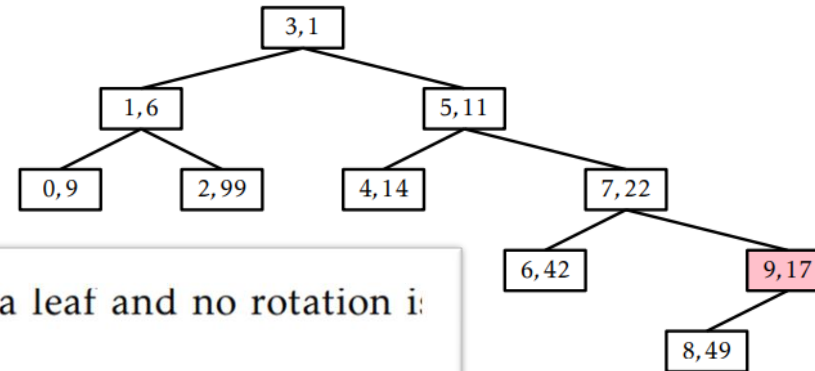
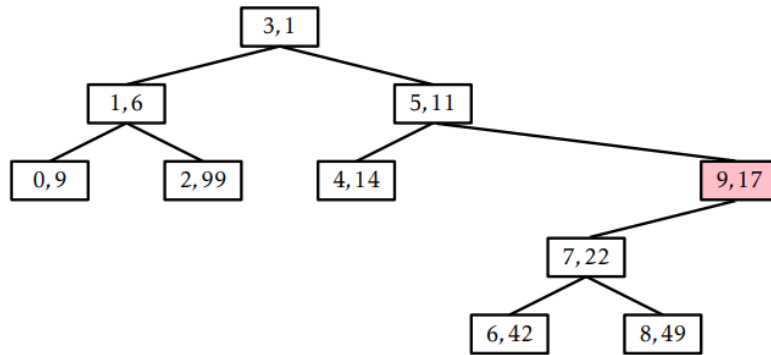
# Bubble Up

- Move element to up by performing rotations

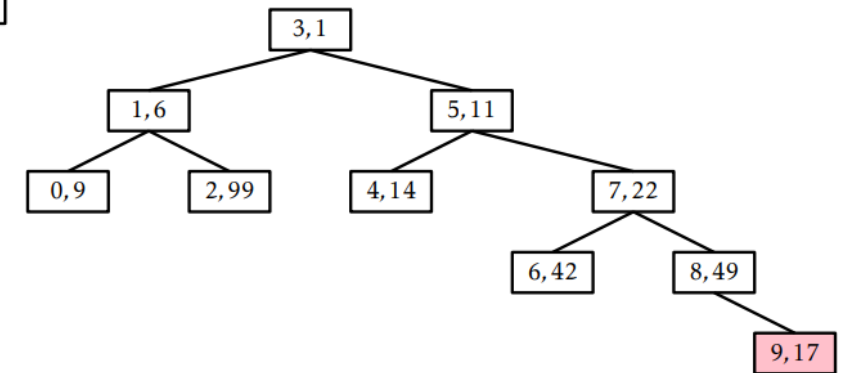


# Remove: Trickle Down (reverse of Bubble Up)

- Move the element as a leaf by performing rotations



1. If  $u.left$  and  $u.right$  are both *nil*, then  $u$  is a leaf and no rotation is performed.
2. If  $u.left$  (or  $u.right$ ) is *nil*, then perform a right (or left, respectively) rotation at  $u$ .
3. If  $u.left.p < u.right.p$  (or  $u.left.p > u.right.p$ ), then perform a right rotation (or left rotation, respectively) at  $u$ .





# Analysis of Operations

- $\text{find}(x) = O(h) \Rightarrow O(\log n)$  [expected]
  - Same as BST
- $\text{add}(x) = O(h) \Rightarrow O(\log n)$  [expected]
  - Insertion –  $O(h)$
  - Bubble up –  $O(h)$
- $\text{remove}(x) = O(h) \Rightarrow O(\log n)$  [expected]
  - Search + trickle down –  $O(h)$





# Construct treap for the following sequences:

1.  $(2,7), (3,6), (5,4), (8,12)$

2. Generate priority randomly.

$(2, \_), (3, \_), (5, \_), (1, \_), (7, \_)$