



Topological Sort

CS-6th

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Topological Sort

- A topological sort of a dag (directed acyclic graph) $G=(V,E)$ is a linear ordering of all its vertices such that if G contains an edge (u,v) , then u appears before v in the ordering.
- Only applicable on directed acyclic graphs.
- A directed graph G is acyclic if and only if a depth-first search of G yields no back edges.

Example

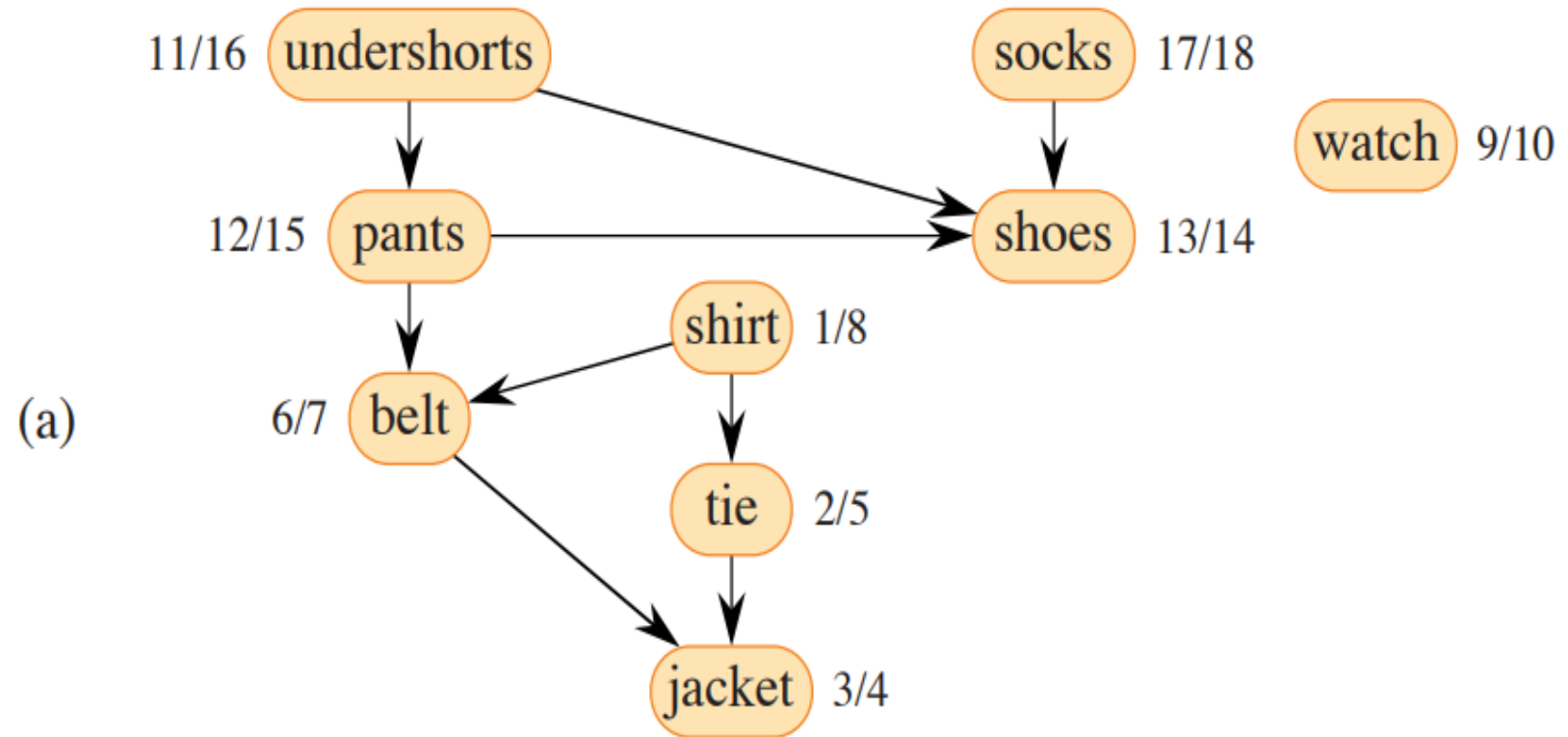
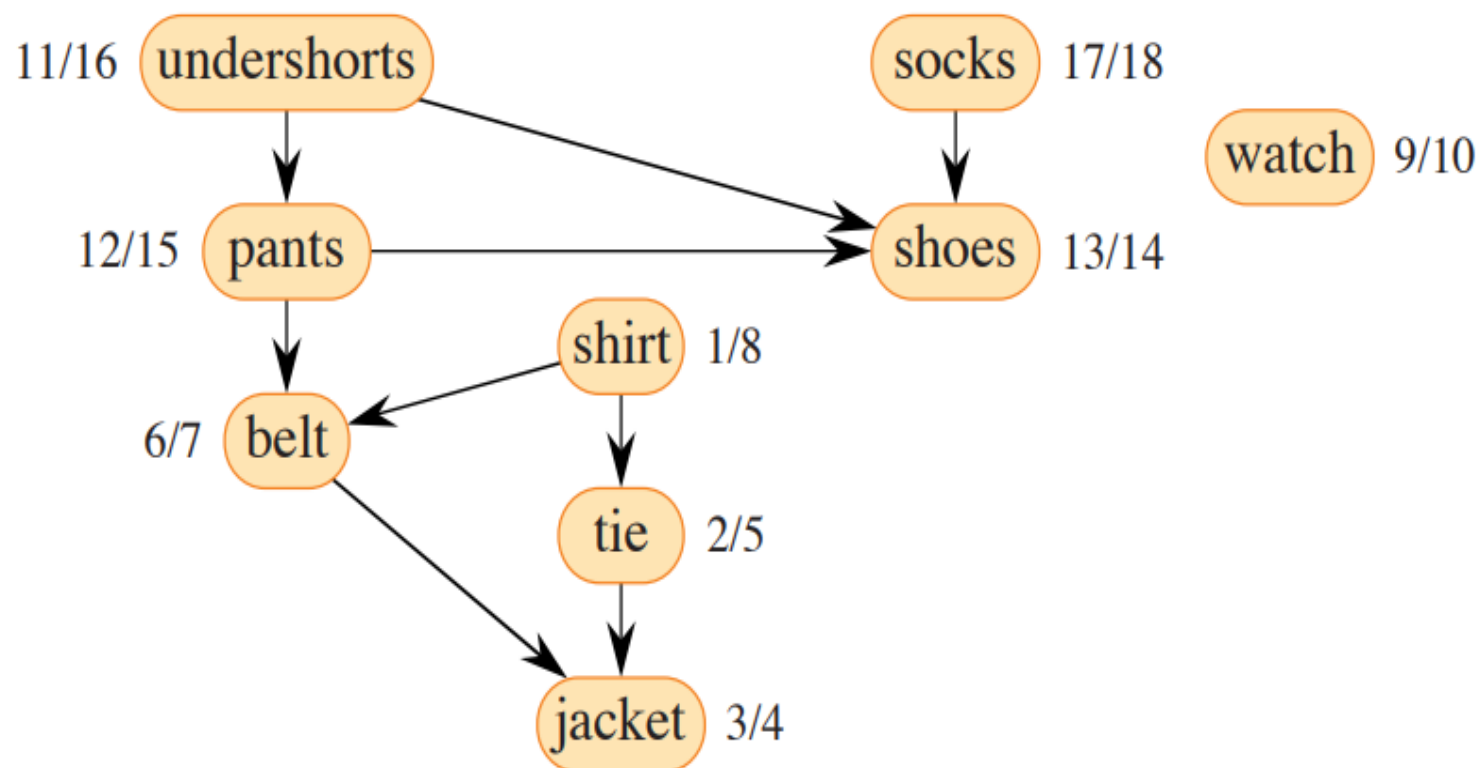


Figure 20.7 (a) Professor Bumstead **topologically** sorts his clothing when getting dressed. Each directed edge (u, v) means that garment u must be put on before garment v . The discovery and finish times from a depth-first search are shown next to each vertex. (b) The same graph shown **topologically** sorted, with its vertices arranged from left to right in order of decreasing finish time. All directed edges go from left to right.

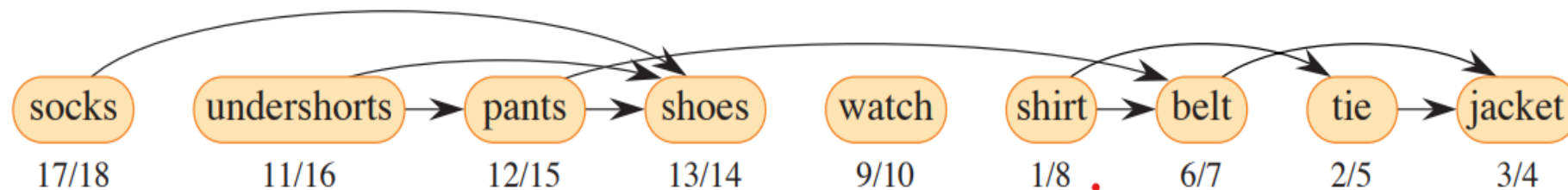
Input

(a)



Output

(b)



Topological Sort Algorithm

TOPOLOGICAL-SORT(G)

- 1 call DFS(G) to compute finish times $v.f$ for each vertex v
- 2 as each vertex is finished, insert it onto the front of a linked list
- 3 **return** the linked list of vertices

Time Complexity

- $\Theta(V+E)$

Depth-First Search Algorithm

DFS(G)

```
1  for each vertex  $u \in G.V$ 
2       $u.color = \text{WHITE}$ 
3       $u.\pi = \text{NIL}$ 
4   $time = 0$ 
5  for each vertex  $u \in G.V$ 
6      if  $u.color == \text{WHITE}$ 
7          DFS-VISIT( $G, u$ )
```

DFS-VISIT(G, u)

```
1   $time = time + 1$                                 // white vertex  $u$  has just been discovered
2   $u.d = time$ 
3   $u.color = \text{GRAY}$ 
4  for each vertex  $v$  in  $G.Adj[u]$  // explore each edge  $(u, v)$ 
5      if  $v.color == \text{WHITE}$ 
6           $v.\pi = u$ 
7          DFS-VISIT( $G, v$ )
8   $time = time + 1$ 
9   $u.f = time$ 
10  $u.color = \text{BLACK}$                                 // blacken  $u$ ; it is finished
```

