Practice Problems (Randomized Algorithms)

CS 6th

Ayesha Enayet

• In HIRE-ASSISTANT, assuming that the candidates are presented in a random order, what is the probability that you hire exactly one time? What is the probability that you hire exactly n times?

- You will hire exactly one time if the best candidate is at first. There are (n-1)! orderings with the best candidate being at first, so the probability that you hire exactly one time is (n-1)!/n!=1/n.
- You will hire exactly n times if the candidates are presented in increasing order. There is only an ordering for this situation, so the probability that you hire exactly n times is 1/n!.

• Use **indicator random** variables to solve the following problem, which is known as the **hat-check problem**. Each of *n* customers gives a hat to a hat-check person at a restaurant. The hat-check person gives the hats back to the customers in a random order. What is the expected number of customers who get back their hat?

Let X be the number of customers who get back their own hat and X_i be the indicator random variable that customer i gets his hat back. The probability that an individual gets his hat back is $\frac{1}{n}$. Thus we have

$$E[X] = E\left[\sum_{i=1}^{n} X_i\right] = \sum_{i=1}^{n} E[X_i] = \sum_{i=1}^{n} \frac{1}{n} = 1.$$

• Let A[1..n] be an array of n distinct numbers. If i < j and A[i]>A[j], then the pair (i,j) is called an **inversion** of A. Suppose that the elements of A form a uniform random permutation of (1,2,...,n). Use indicator random variables to compute the expected number of inversions.

Let $X_{i,j}$ for i < j be the indicator of A[i] > A[j]. We have that the expected number of inversions

$$\begin{split} \mathbf{E} \bigg[\sum_{i < j} X_{i,j} \bigg] &= \sum_{i < j} E[X_{i,j}] \\ &= \sum_{i = 1}^{n-1} \sum_{j = i+1}^{n} \Pr\{A[i] > A[j]\} \\ &= \frac{1}{2} \sum_{i = 1}^{n-1} n - i \qquad \sum_{j = i+1}^{n-1} \frac{1}{2} = \frac{1}{2}(n-i) \\ &= \frac{n(n-1)}{2} - \frac{n(n-1)}{4} \text{ numbers} \\ &= \frac{n(n-1)}{4}. \end{split}$$

$$egin{aligned} E[X] &= E\left[\sum_{i=1}^{n-1}\sum_{j=i+1}^{n}X_{i,j}
ight] \ &= \sum_{i=1}^{n-1}\sum_{j=i+1}^{n}E[X_{i,j}] \ &= \sum_{i=1}^{n-1}\sum_{j=i+1}^{n}rac{1}{2} \ &= rac{n(n-1)}{2} \cdot rac{1}{2} \ &= rac{n(n-1)}{4} \end{aligned}$$

• Write down an algorithm to find the first two largest numbers.	

8.2 Finding The Two Largest

The max-two problem is to find the two largest elements from a sequence of n (unique) numbers. For inspiration, we'll go back and look at the naïve algorithm for the problem:

```
function max2(S) = let

function replace((m_1, m_2), v) =

if v \leq m_2 then (m_1, m_2)

else if v \leq m_1 then (m_1, v)

else (v, m_1)

val start = if S_1 \geq S_2 then (S_1, S_2) else (S_2, S_1)

in iter \ replace \ start \ S \langle 3, \ldots, n \rangle

end
```

Reference: CMU Randomized Algorithm Notes

function
$$Y(e) = \underbrace{1}_{\text{Line } 6} + \underbrace{n-2}_{\text{Line } 3} + \underbrace{\sum_{i=3}^{n} X_i(e)}_{\text{Line } 4}$$

It is common, however, to use the shorthand notation

$$Y = 1 + (n-2) + \sum_{i=3}^{n} X_i$$

Y is total # of comparisons $Y = 1 + (n-2) + \sum_{i=3}^{n} X_i$ is total # of comparisons and Xi is an indicator random variable.

where the argument e and function definition is implied.

We are interested in computing the expected value of Y, that is $\mathbf{E}[Y] = \sum_{e \in \Omega} \mathbf{Pr}[e] Y(e)$. By linearity of expectation, we have

$$\mathbf{E}[Y] = \mathbf{E}\left[1 + (n-2) + \sum_{i=3}^{n} X_i\right]$$
$$= 1 + (n-2) + \sum_{i=3}^{n} \mathbf{E}[X_i].$$

using a counting argument.) Therefore, the probability that T_i is the largest or the second largest element in $\{T_1, \ldots, T_i\}$ is $\frac{1}{i} + \frac{1}{i} = \frac{2}{i}$, so

$$\mathbf{E}\left[X_i\right] = 1 \cdot \frac{2}{i} = 2/i.$$

Plugging this into the expression for $\mathbf{E}[Y]$, we get

$$\mathbf{E}[Y] = 1 + (n-2) + \sum_{i=3}^{n} \mathbf{E}[X_i]$$

$$= 1 + (n-2) + \sum_{i=3}^{n} \frac{2}{i}$$

$$= 1 + (n-2) + 2\left(\frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}\right)$$

$$= n - 4 + 2\left(1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}\right)$$

$$= n - 4 + 2H_n,$$

where H_n is the *n*-th Harmonic number. But we know that $H_n \leq 1 + \log_2 n$, so we get $\mathbf{E}[Y] \leq n - 2 + 2\log_2 n$. We could also use the following sledgehammer: