

Complexity Theory Quiz 05

Q) $A_{DFA} = \{ \langle DFA, x \rangle \mid DFA \text{ accepts string } x \}$.
Show $A_{DFA} \in L$.

* We need to show that there exists a log-space alg that decides A_{DFA} , that is, $DFA \Delta$ accepts x .
 $D = (Q, \Sigma, \delta, q_0, F)$ by definition.

* We can simulate D on x by starting at the start state q_0 & reading each character ~~in~~ x_i in x , and we transition to the next state $q_{i+1} = \delta(q_i, x_i)$.

* Now for each state, we can ~~store~~ ~~label~~ the cells according to the index on the input tape starting from 0.

* We only need to keep track of the current state we are at on, & store it on our work tape. We can store ~~the index of~~ the state, which would take only ~~$O(\log n)$~~ \log space. If we store only the index of the state with ~~res~~ respect to the cell, it would take $(\log n)$ space on the work tape.

* We can reuse that space each time we move onto a new state by ~~erasing~~ overwriting on it as we read each x_i & transition to a new state.

* Once we have completed our simulation, we simply check if the state we are currently on on our work tape exists in the set of Accept states F .

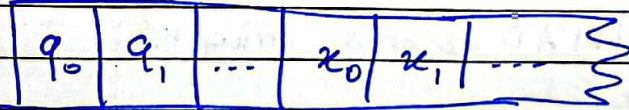
If yes, we accept, else we reject.

Since the whole simulation takes only $\log O(\log)$ space to store the current space, it exists in L .

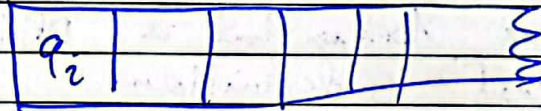
Thus, $A_{DFA} \in L$.

Hence proved!

Input tape



Work Tape



↓ Next State

