Maximum Subarray Sum Problem

Maximum Subarray Problem

- The maximum subarray problem is the task of finding the contiguous subarray within a one-dimensional array of numbers which has the largest sum.
- Example

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

- Maximum SubArray is from index 4 till index 7 with sum 8

Brute Force Solution for Maximum Sub Array Sum

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

All sub arrays ending at index 2

All sub arrays ending at index 3

All sub arrays ending at index 4

Brute Force Solution $O(n^2)$

```
MaxSubArraySum(A, n)
{
globalMax = -infinity
for(i = 1 to n)
{
    subArraySum = 0
     for (j = i to 1)
     {
          subArraySum += A[j]
          globalSum = Max (globalSum, subAayraySum)
     }
}
return globalSum
}
```

Divide and Conquer Solution O(n lg n)

- Divide array in two equal halves
- SubArrays can be divide into 3 categories
  - Left subarray (start and end index in left half of array)
  - Right subarray (start and end index in right half of array)
  - Crossing subarray (start in left and end in right half of array)

  - Left and right subarray sum is calculated using recursion
  - Crossing subarray sum is computed using a linear time function

Brute Force Solution

Dry Run

O(n^2) Solution for Maximum Sub Array Sum

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A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

All sub arrays ending at index 2

All sub arrays ending at index 3

All sub arrays ending at index 4

$O(n^2)$ Solution for Maximum Sub Array Sum

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O(n^2) Solution for Maximum Sub Array Sum

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

All sub arrays ending at index 2

All sub arrays ending at index 3

What is a sub problem here?

All subarrays ending at index i is a sub problem of original problem.

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

MaxSum[1] = A[1] = -6

All sub arrays ending at index 2

$$\text{MaxSum}[2] = \text{Max}\begin{cases} A[2] = 2 \\ A[1] + A[2] = -4 \end{cases}$$

All sub arrays ending at index 3

$$\text{MaxSum}[3] = \text{Max}\begin{cases} A[3] = -4 \\ A[2] + A[3] = 2 \\ A[1] + A[2] + A[3] = -8 \end{cases}$$

All sub arrays ending at Index 4

$$\text{MaxSum}[4] = \text{Max}\begin{cases} A[4] = 1 \\ A[3] + A[4] = -3 \\ A[2] + A[3] + A[4] = -1 \\ A[1]+A[2]+A[3]+A[4] = -7 \end{cases}$$

Dynamic Programming Algorithm (Kadane’s Algorithm)

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

MaxSum[1] = A[1] = -6

All sub arrays ending at index 2

$$\text{MaxSum}[2] = \text{Max}\begin{cases} A[2] = 2 \\ A[1] + A[2] = -4 \end{cases}$$

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Can you divide the problem into subproblems such that solution to a bigger subproblem uses solution from smaller subproblem

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

MaxSum[1] = A[1] = -6

All sub arrays ending at index 2

$$\text{MaxSum}[2] = \text{Max}\begin{cases} A[2] = 2 \\ A[1] + A[2] = -4 \end{cases}$$

All sub arrays ending at index 3

$$\text{MaxSum}[3] = \text{Max}\begin{cases} A[3] \\ A[2] + A[3] \\ A[1] + A[2] + A[3] \end{cases}$$

$$\text{MaxSum}[3] = \text{Max}\begin{cases} A[3] \\ \text{Max}\begin{Bmatrix} A[2] \\ A[1] + A[2] \end{Bmatrix} + A[3] \end{cases}$$

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

MaxSum[1] = A[1] = -6

All sub arrays ending at index 2

$$\text{MaxSum[2]} = \text{Max}\begin{cases} A[2] = 2 \\ A[1] + A[2] = -4 \end{cases}$$

All sub arrays ending at index 3

$$\text{MaxSum[3]} = \text{Max}\begin{cases} A[3] = -4 \\ A[2] + A[3] = 2 \\ A[1] + A[2] + A[3] = -8 \end{cases}$$

All sub arrays ending at Index 4

$$\text{MaxSum[4]} = \text{Max}\begin{cases} A[4] \\ A[3] + A[4] \\ A[2] + A[3] + A[4] \\ A[1]+A[2]+A[3]+A[4] \end{cases}$$

$$\text{MaxSum[4]} = \text{Max}\begin{cases} A[4] \\ \text{Max}\begin{Bmatrix} A[3] \\ A[2] + A[3] \\ A[1]+A[2]+A[3] \end{Bmatrix} + A[4] \end{cases}$$

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

$$MaxSum[1] = A[1] = -6$$

All sub arrays ending at index 2

$$MaxSum[2] = Max \begin{cases} A[2] = 2 \\ A[1] + A[2] = -4 \end{cases}$$

All sub arrays ending at index 3

$$MaxSum[3] = Max \begin{cases} A[3] = -4 \\ A[2] + A[3] = 2 \\ A[1] + A[2] + A[3] = -8 \end{cases}$$

All sub arrays ending at Index 4

$$MaxSum[4] = Max \begin{cases} A[4] \\ A[3] + A[4] \\ A[2] + A[3] + A[4] \\ A[1]+A[2]+A[3]+A[4] \end{cases}$$

$$MaxSum[4] = Max \begin{cases} A[4] \\ MaxSum[3] + A[4] \end{cases}$$

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

MaxSum[1] = A[1] = -6

All sub arrays ending at index 2

$$\text{MaxSum}[2] = \text{Max}\begin{cases} A[2] \\ \text{MaxSum}[1] + A[2] \end{cases}$$

All sub arrays

ending at index 3

$$\text{MaxSum}[3] = \text{Max}\begin{cases} A[3] \\ \text{MaxSum}[2] + A[3] \end{cases}$$

$$\text{MaxSum}[4] = \text{Max}\begin{cases} A[4] \\ \text{MaxSum}[3] + A[4] \end{cases}$$

$$\text{MaxSum}[5] = \text{Max}\begin{cases} A[5] \\ \text{MaxSum}[4] + A[5] \end{cases}$$

i	1	2	3	4	5	6	7	8
A[i]	-6	2	-4	1	3	-1	5	-1

All sub arrays ending at index 1

MaxSum[1] = A[1] = -6

All sub arrays ending at index 2

$$\text{MaxSum}[2] = \text{Max}\begin{cases} A[2] \\ \text{MaxSum}[1] + A[2] \end{cases}$$

All sub arrays ending at index 3

$$\text{MaxSum}[3] = \text{Max}\begin{cases} A[3] \\ \text{MaxSum}[2] + A[3] \end{cases}$$

$$\text{MaxSum}[4] = \text{Max}\begin{cases} A[4] \\ \text{MaxSum}[3] + A[4] \end{cases}$$

Recursion for DP Solution

MaxSum[i] = Max (A[i] + MaxSum[i-1] , A[i])

$$\text{MaxSum}[5] = \text{Max}\begin{cases} A[5] \\ \text{MaxSum}[4] + A[5] \end{cases}$$

Brute Force $O(n^2)$ Solution

```
MaxSubArraySum(A, n)
{
globalMax = -infinity
for(i = 1 to n)
{
    subArraySum = 0
    for (j = i to 1)
    {
        subArraySum += A[j]
        globalSum = Max (globalSum, subAayraySum)
    }
}
return globalSum
}
```

DP O(n) Solution

```
MaxSubArraySum(A,n)
{
   globalSum = A[1]
   MaxSum[1] = A[1]
   for (i = 2 to n)
   {
      if (MaxSum[i-1] + A[i] > A[i] )
         MaxSum[i] = MaxSum[i-1] + A[i]
      else
         MaxSum[i] = A[i]
      globalSum = Max (globalSum, MaxSum[i])
   }
   return globalSum
}
```

O(n) Dynamic Programming Algorithm (Kadane's Algorithm)

```
MaxSubArraySum(A,n)
{   globalSum = A[1]
    MaxSum[1] = A[1]
    for (i = 2 to n)
         if (MaxSum[i-1] + A[i] > A[i] )
                MaxSum[i] = MaxSum[i-1] + A[i]
        else
                MaxSum[i] = A[i]
        If (globalSum < MaxSum[i])
            globalSum = MaxSum[i]
            globalEnd = i
    return globalSum
}
```

Recursion for DP Solution

MaxSum[i] = Max (A[i] + MaxSum[i-1] , A[i])

O(n) Dynamic Programming Algorithm (Kadane's Algorithm)

```
1.  MaxSubArraySum(A,n)
2.  {   globalSum = A[1]
3.      MaxSum[1] = A[1]
4.      for (i = 2 to n)
5.          if (MaxSum[i-1] + A[i] > A[i] )
6.              MaxSum[i] = MaxSum[i-1] + A[i]
7.          else
8.              MaxSum[i] = A[i]
9.          If (globalSum < MaxSum[i])
10.             globalSum = MaxSum[i]
11.             globalEnd = i
12.     return globalSum
13. }
```

This algorithm keeps track of end of Max sub array in line 11. Modify this algorithm to keep track of start of Max sub array