National University of Computer & Emerging Sciences CS 3001 - COMPUTER NETWORKS

Lecture 04
Chapter 1

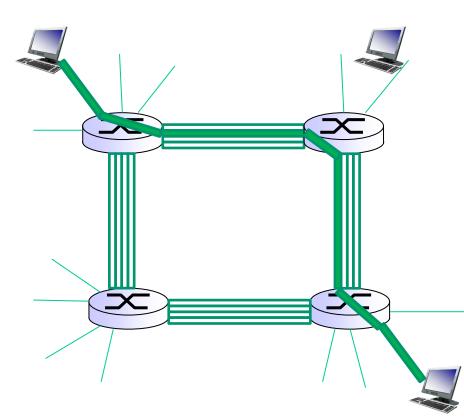
1st September, 2022

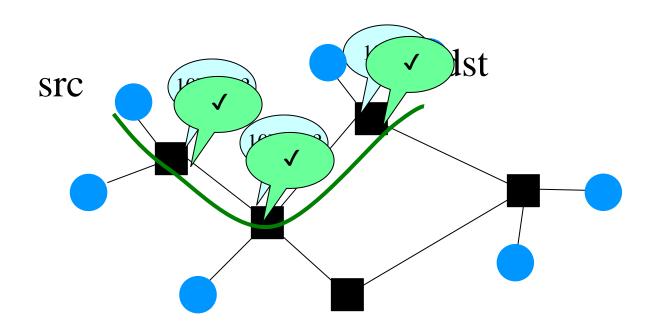
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Office Hours: 02:30 pm till 06:00 pm (Every Tuesday & Thursday)

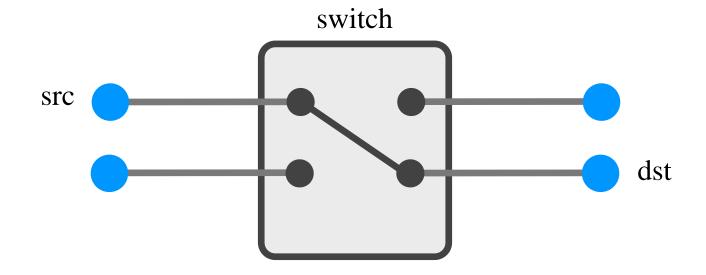
end-end resources allocated to, reserved for "call" between source & dest:

- In diagram, each link has four circuits.
 - call gets 2nd circuit in top link and Ist circuit in right link.
- dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- circuit segment idle if not used by call (no sharing)
- Commonly used in traditional telephone networks





- (I) src sends a reservation request to dst
- (2) Switches "establish a circuit"
- (3) src starts sending data
- (4) src sends a "teardown circuit" message

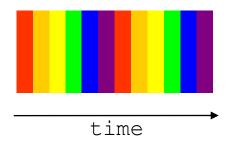


Reservation establishes a "circuit" within a switch

Many kinds of "circuits"

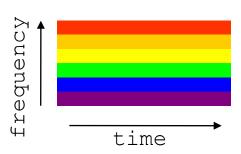
Time division multiplexing

- divide time in time slots
- separate time slot per circuit

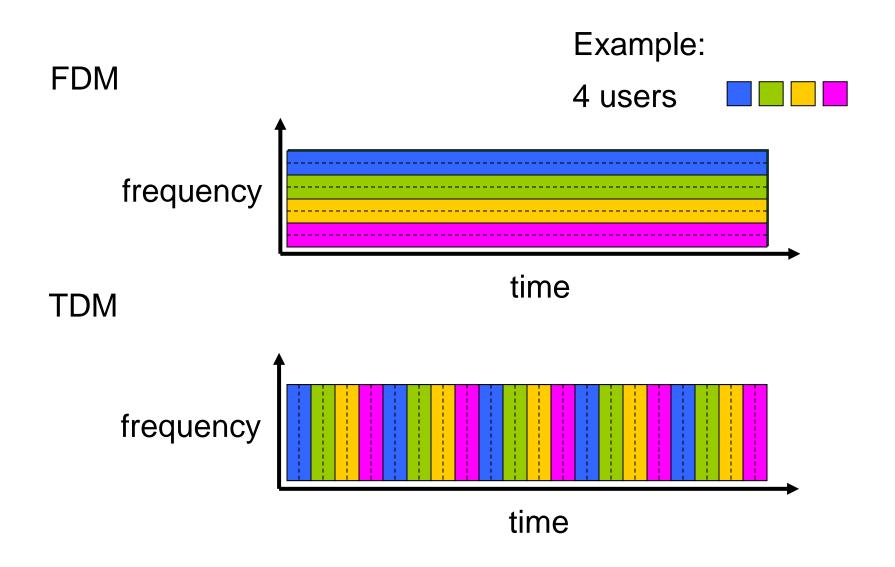


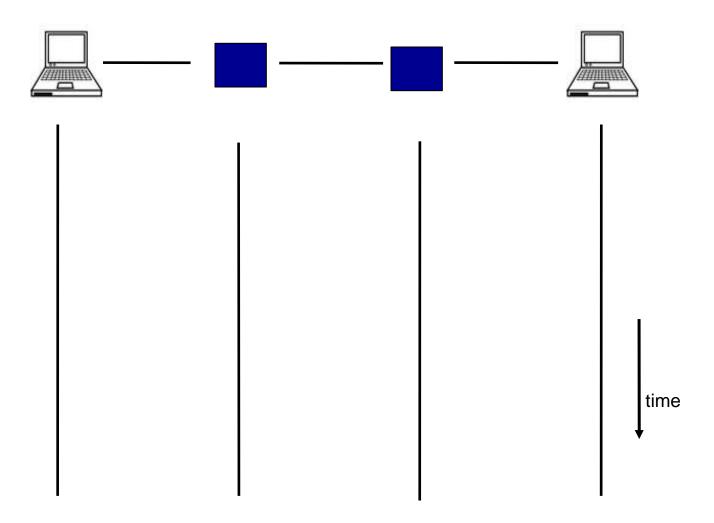
Frequency division multiplexing

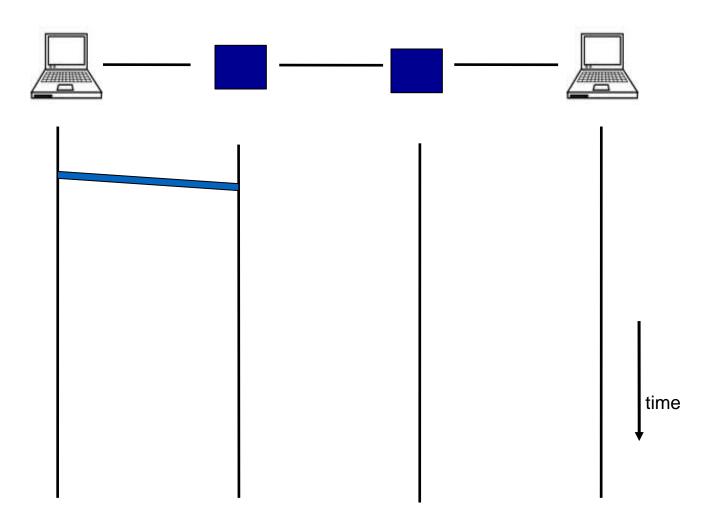
- divide frequency spectrum in frequency bands
- separate frequency band per circuit

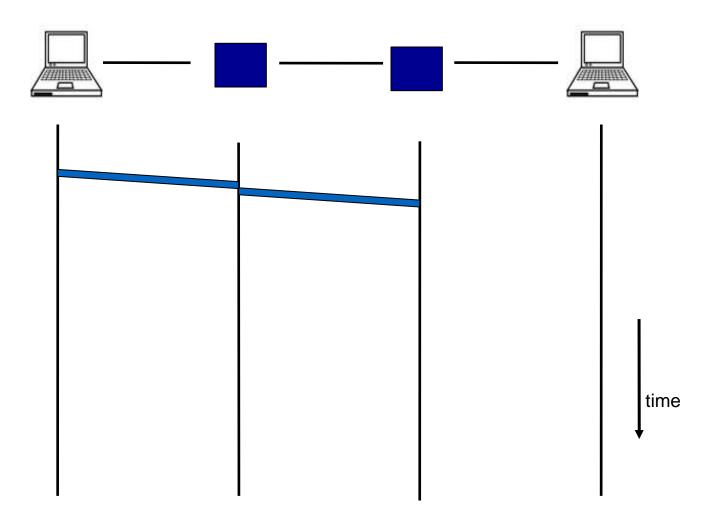


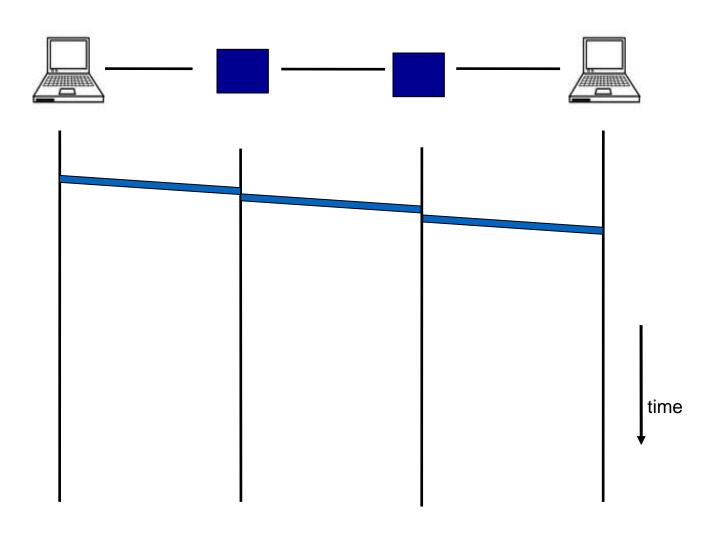
Circuit switching: FDM versus TDM

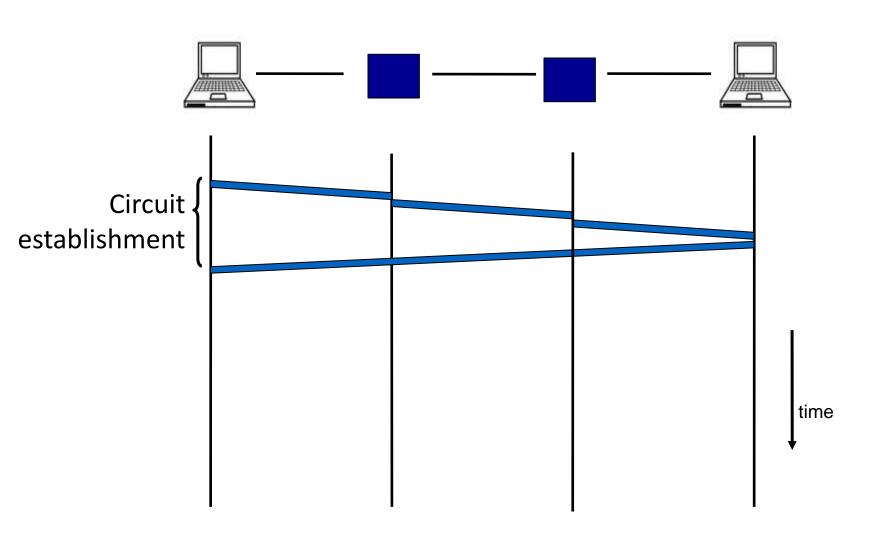


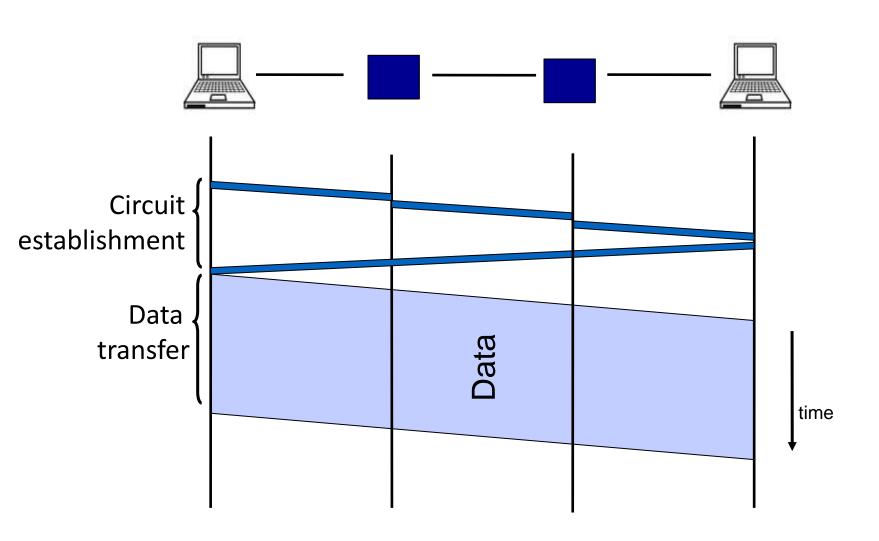


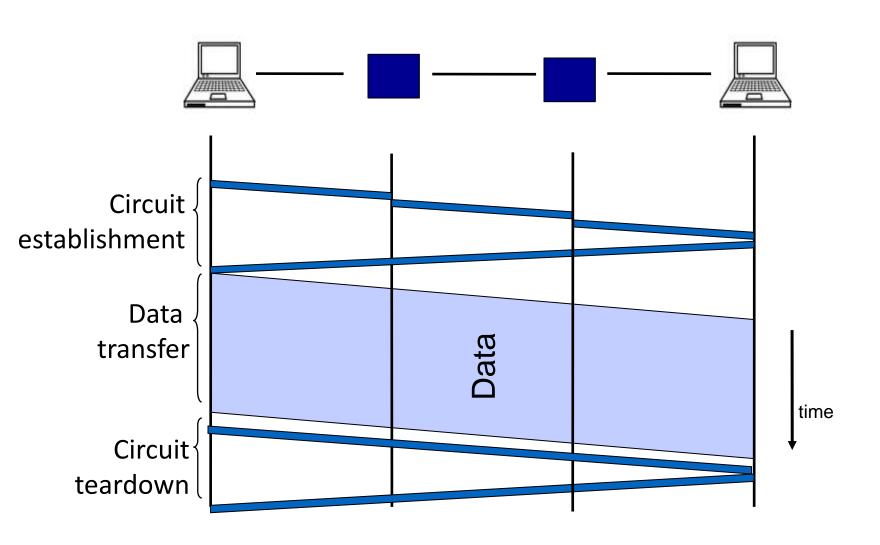












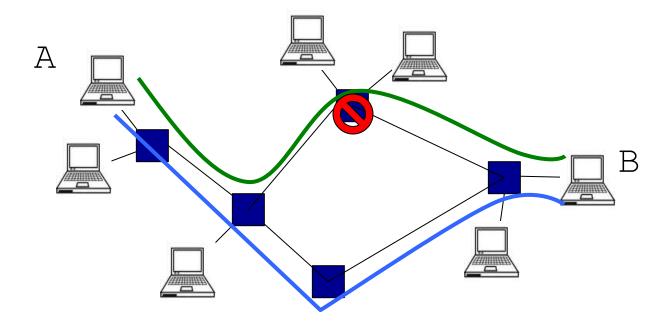
Numerical Example

- How long does it take to send a file of 80 Kbytes from host A to host B over a circuit-switched network?
- All links are 1.536 Mbps
- Each link uses TDM with 24 slots/sec
- Time to establish end-to-end circuit is 500 msec

Let's work it out!

Numerical Example: Solution

- 80 Kbytes is 640,000 bits
- NOTE: networks in bits, end systems in bytes
- NOTE: 8 bits to a byte
- Each circuit has a rate of 1.536Mbps / 24=> 1536000/24 = 64000bps
- So, it takes 640000 bits / 64000 bps = 10 seconds to transmit the file
- Need to add the circuit establishment time $(\frac{1}{2}$ second)
- So, 10.5 seconds



Circuit switching doesn't "route around trouble"

Pros

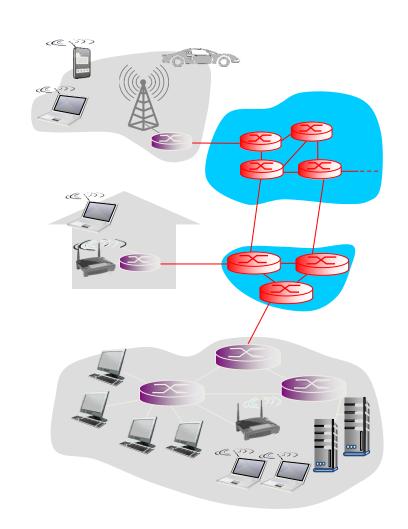
- predictable performance
- simple/fast switching (once circuit established)

Cons

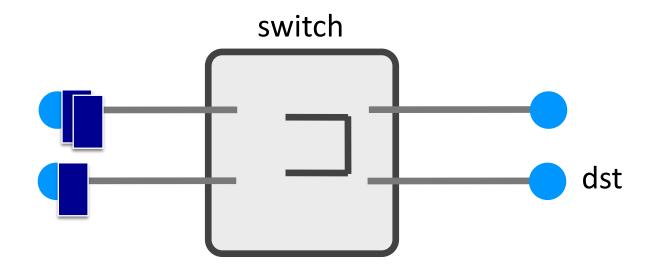
- complexity of circuit setup/teardown
- inefficient when traffic is bursty
- circuit setup adds delay
- switch fails → its circuit(s) fails

Packet Switching

- mesh of interconnected routers
- packet-switching: hosts break application-layer messages into packets
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



Packet switching



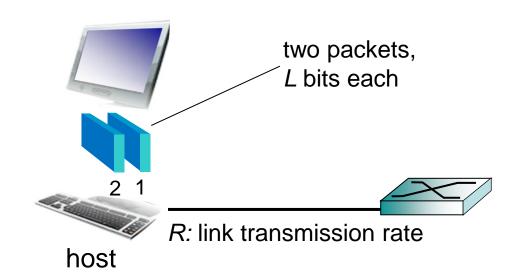
Each packet contains destination (dst) Each packet treated independently

With buffers to absorb transient overloads

Host: sends packets of data

host sending function:

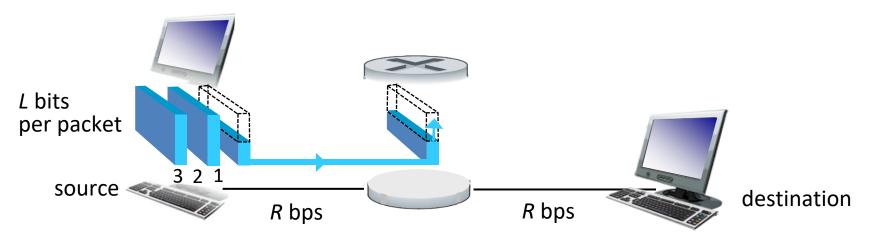
- takes application message
- breaks into smaller chunks, known as packets, of length L bits
- transmits packet into access network at transmission rate R
 - link transmission rate, aka



Link capacity, aka link bandwidth

transmission delay time needed to transmit
$$L$$
-bit packet into link $= \frac{L \text{ (bits)}}{R \text{ (bits/sec)}}$

Packet-switching: store-and-forward

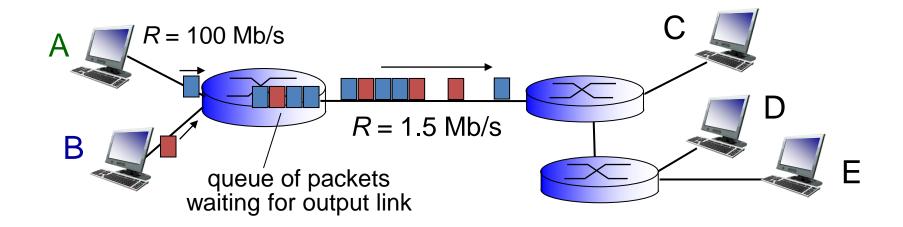


- takes L/R seconds to transmit (push out) L-bit packet into link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link
- end-end delay = 2L/R (assuming zero propagation delay)

one-hop numerical example:

- L = 7.5 Mbits
- R = 1.5 Mbps
- one-hop transmission delay = 5 sec
- How much delay till dest.?
- How much delay till dest. and back to source?

Packet Switching: queueing delay, loss



queuing and loss:

- If arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
 - packets will queue, wait to be transmitted on link
 - packets can be dropped (lost) if memory (buffer) fills up

Packet Switching

Pros

- -efficient use of network resources
- -simpler to implement
- -robust: can "route around trouble"

Cons

- unpredictable performance
- requires buffer management and congestion control

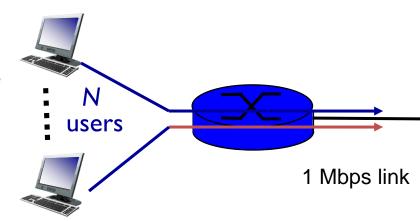
On-demand or reserve?

Packet switching versus circuit switching

packet switching allows more users to use network!

example:

- 1 Mbps link
- 10 total users
- 9 users silent
- Only one active user generates:
 - One thousand 1000-bit packets (1 Million bits)
- circuit-switching:
 - Will take one user 10 seconds to transmit complete
 1 Million bits
 - Only 10 users can be serviced via this model simultaneously
 - Thus if only 1 active, remaining 9 slots wasted
 - But guaranteed service and quality to 10 users
- packet switching:
 - Only 1 active user will transmit complete 1 Million bits in 1 sec (@ the full capacity of the link i.e. 1 Mbps)
 - Depending on probability of simultaneously active users (if they are less or equal to 10), and arrival rate of data <= 1 Mbps (Capacity of the link), packets flow without delay. When users > 10, arrival rate of data increases > 1 Mbps and packets queue, causing delays
 - If probability of simultaneously active users > 10 is miniscule, packet switching provides the same performance (quality) as circuit switching but serving more users



Packet switching versus circuit switching

is packet switching a "slam dunk winner?"

- great for bursty data
 - resource sharing
 - simpler, no call setup
- excessive congestion possible: packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 7)
- Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?