# Radix Sort

### Radix Sort

- Counting sort is a linear time sorting algorithm that sort in O(n+k) time when elements are in the range from 1 to k.
- What if the elements are in the range from 1 to n<sup>2</sup>?
  We can't use counting sort because counting sort will take O(n<sup>2</sup>) which is worse than comparison-based sorting algorithms. Can we sort such an array in linear time?
- Radix Sort is the answer. The idea of Radix Sort is to do digit by digit sort starting from least significant digit to most significant digit. Radix sort uses counting sort as a subroutine to sort.

### Radix Sort Algorithm

```
RADIX-SORT(A, d)

1 for i = 1 to d

2 use a stable sort to sort array A on digit i
```

## Radix Sort Dry Run

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329		720		720		329
457		355		329		355
657		436		436		436
839	·····j)Dr-	457	·····j)D·	839	····ij)p-	457
436		657		355		657
720		329		457		720
355		839		657		839

### Radix Sort Dry Run

170, 45, 75, 90, 802, 24, 2, 66

Running Time of Radix Sort 
$$2(n+n)$$
  $1-n$ 

$$0 (\log_{10} k (n+1))$$

$$k = \frac{1}{1-n}$$

$$(n(1) + n(1))$$

$$2(n)$$

$$1 - n$$

$$(n(1) + n(1))$$

$$2(n)$$

$$(n(1) + n(1))$$

$$2(n)$$

$$(n(1) + n(1))$$

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$$\frac{b}{n} = n$$

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- Running Time of Radix Sort

   Let there be d digits in input integers.

   Radix Sort takes  $O(d^*(n+b))$  time where b is the base for representing  $= \frac{1}{2} \sqrt{n}$
- numbers
- For example, for the decimal system, b is 10. What is the value of d?
  - If k is the maximum possible value, then d would be  $O(log_h(k))$
  - So overall time complexity is  $O((n+b) * log_h(k))$ .
  - Let us first limit k. Let k <= n<sup>c</sup> where c is a constant. In that case, the complexity becomes  $O(nLog_h(n))$ .

### Running Time of Radix Sort

- So overall time complexity is  $O((n+b) * log_b(k))$ .
- What if we make the value of b larger?
- What should be the value of b to make the time complexity linear?
- If we set b as n, we get the time complexity as O(n).
- In other words, we can sort an array of integers with a range from 1 to n<sup>c</sup> if the numbers are represented in base n (or every digit takes log<sub>2</sub>(n) bits).

# Comparison of Quick Sort and Radix Sort

### **Pros of Radix Sort**

 If we have log<sub>2</sub>n bits for every digit, the running time of Radix appears to be better than Quick Sort for a wide range of input numbers

### Cons of Radix Sort

- The constant factors hidden in asymptotic notation are higher for Radix Sort
- Quick-Sort uses hardware caches more effectively
- Radix sort uses counting sort as a subroutine and counting sort takes extra space to sort numbers.