

Closure is a property which describes when we combine any two elements of the set; the result is also included in the set.

If we multiply two integer numbers, we will get another integer number. Since this process is always true, it is said that the integer numbers are 'closed under the operation of multiplication'. There is simply no way to escape the set of integer numbers when multiplying.

Let  $S = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, \dots\}$  is a set of integer numbers.

$$12 = 2$$

$$23 = 6$$

$$52 = 10$$

All are included in the set of integer numbers.

We can conclude that integer numbers are closed under the operation of multiplication.

**Theorem 5.3:** Two REs  $L_1$  and  $L_2$  over  $\Sigma$  are closed under union operation.

**Proof:** We have to prove that if  $L_1$  and  $L_2$  are regular over  $\Sigma$ , then their union, i.e.,  $L_1 \cup L_2$  will be also regular.

As  $L_1$  and  $L_2$  are regular over  $\Sigma$ , there must exist FA  $M_1 = (Q_1, \Sigma, \delta_1, q_{01}, F_1)$  and  $M_2 = (Q_2, \Sigma, \delta_2, q_{02}, F_2)$  such that  $L_1 \varepsilon M_1$  and  $L_2 \varepsilon M_2$

Assume that there is no common state between  $Q_1$  and  $Q_2$ , i.e.,  $Q_1 \cap Q_2 = \emptyset$ .

Define another FA,  $M_3 = (Q, \Sigma, \delta, q_0, F)$  where

1.  $Q = Q_1 \cup Q_2 \cup \{q_0\}$ , where  $q_0$  is a new state  $\notin Q_1 \cup Q_2$

2.  $F = F_1 \cup F_2$

3. Transitional function  $\delta$  is defined as  $\delta(q_0, \varepsilon) \rightarrow \{q_{01}, q_{02}\}$

and  $\delta(q, \Sigma) \rightarrow \delta_1(q, \Sigma)$  if  $q \in Q_1$

$$\delta(q, \Sigma) \rightarrow \delta_2(q, \Sigma) \text{ if } q \in Q_2$$

It is clear from the previous discussion that from  $q_0$  we can reach either the initial state  $q_1$  of  $M_1$  or the initial state  $q_2$  of  $M_2$ .

Transitions for the new FA,  $M$ , are similar to the transitions of  $M_1$  and  $M_2$ .

As  $F = F_1 \cup F_2$ , any string accepted by  $M_1$  or  $M_2$  will also be accepted by  $M$ .

Therefore,  $L_1 \cup L_2$  is also regular.

**Example:**

Let  $L_1 = a^*(a + b)b^*$

$$L_2 = ab(a + b)b^*$$

The FA  $M_1$  accepting  $L_1$  is as shown in Fig. 5.38(a).

The FA  $M_2$  accepting  $L_2$  is as shown in Fig. 5.38(b).

The machine  $M$  produced by combining  $M_1$  and  $M_2$  is as shown in Fig. 5.38(c).

It accepts  $L_1 \cup L_2$ .

**Theorem 5.4:** The complement of an RE is also regular.

If  $L$  is regular, we have to prove that  $L^T$  is also regular.

As  $L$  is regular, there must be an FA,  $M = (Q, \Sigma, \delta, q_0, F)$ , accepting  $L$ .

$M$  is an FA, and so  $M$  has a transitional system.

Let us construct another transitional system  $M'$  with the state diagram of  $M$  but reversing the direction of the directed edges.  $M'$  can be designed as follows:

1. The set of states of  $M'$  is the same as  $M$
2. The set of input symbols of  $M'$  is the same as  $M$