Concordia University Department of Computer Science and Software Engineering

SOEN 331 Section S: Formal Methods for Software Engineering Assignment 4

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1 Our assignment

1. (10 pts) Find a logically equivalent formula for ϕ W ψ and provide a short reasoning to support your answer. Represent this equivalence between the two expressions with the appropriate logical connective, and support your reasoning.

Solution:

The logically equivalent formula for this would be:

$$(\phi \mathbf{W} \psi) \equiv (\phi \mathbf{U} \psi) \vee \Box(\psi)$$

These are equivalent, because the principle of the strong until operator **U**. Since the ψ is never guaranteed to be true, we would need to add an extra or statement because if the statement ψ becomes true it would mean that ϕ can't be true. This means that ϕ would be true until a certain condition (ψ is true) is met.

2. (10 pts) Find a logically equivalent formula for ϕ U ψ , and provide a short reasoning to support your answer. Represent this equivalence between the two expressions with the appropriate logical connective, and support your reasoning.

Solution:

The logically equivalent formula for this would be:

$$(\phi \mathbf{U} \psi) \equiv (\phi \mathbf{W} \psi) \wedge \Diamond(\psi)$$

These are equivalent, because the principle of the strong until operator U. We know that this ψ will eventually become true. We then now that we can use the weak until clause with an eventually operator, because the only way that ϕ is not true is when the ψ is not true. We add the and operator, because we want to insure that we actually get ψ as being true, because this may never happen.

3. (10 pts) Find a logically equivalent formula for $\phi R \psi$ in terms of W, and provide a short reasoning to support your answer. Represent this equivalence between the two expressions with the appropriate logical connective, and support your reasoning.

Solution:

This paragraph refers to Questions 4 - 5: Consider a railroad with a single rail and a road level-crossing. We introduce the following propositions that represent events:

a: A train is approaching.

b: The barrier is down

c: A train is crossing

l: A light is blinking

- 4. (15 pts) Express each of the following requirements formally. For each one, proceed to find a logically equivalent formula that captures the safety property of the system (i.e. in terms of "something bad never happens"):
 - (a) (5 pts) When a train is crossing, the barrier must be down.

Solution:

 $\Box(c \to \bigcirc b)$

(b) (5 pts) If a train is approaching or crossing, then the light must be blinking. Solution:

 \Box (a $\lor c \to \bigcirc l$)

(c) (5 pts) If the barrier is up and the light is off, then no train is coming or crossing.

<u>Solution</u>:

 $\Box(\neg\ (b\land l)\to\neg\ \bigcirc\ (a\lor c))$

- 5. (10 pts) Express each of the following requirements formally in terms of the liveness property (i.e. in terms of "something good eventually happens"):
 - (a) (5 pts) When a train is approaching, it will eventually cross..

Solution:

 $\Box(\mathbf{a} \to \Diamond c)$

(b) (5 pts) When a train is approaching and no train is crossing, then the barrier will eventually go down before the train crosses.

Solution:

$$\Box((\mathbf{a} \land \neg c) \to \Diamond(b \mathbf{U} a))$$

6. (45 pts) The behavior of a program is expressed by the following temporal formula:

$$\begin{bmatrix} \operatorname{start} \to (\phi \oplus \psi) \\ \operatorname{start} \to \tau \\ \phi \to \bigcirc (\psi \ \mathcal{U} \ \chi) \\ \psi \land \tau \to \bigcirc (\psi \ \mathcal{W} \ \chi) \\ \tau \land \bigcirc \psi \to \bigcirc \omega \\ \\ \psi \land \omega \to \bigcirc^2 \chi \\ \\ \omega \land \bigcirc^2 \chi \to \bigcirc^2 \sigma \\ \\ \psi \land \bigcirc \sigma \to \bigcirc^2 \pi \\ \\ \psi \land \tau \to \sigma \ \mathcal{R} \ \pi \\ \\ \phi \land \bigcirc \psi \to \bigcirc^2 \chi \end{bmatrix}$$

- (a) (20 pts) Visualize all models of behavior.

 <u>Solution</u>:
- (b) (10 pts) Is the set of requirements satisfiable in all models of behavior? Explain why or why not.

Solution:

(c) (10 pts) In the case where the set of requirements is not satisfiable, what modification(s) to the requirements would you make (you may temporarily assume the role of a stakeholder) in order to achieve satisfiability.
Solution:

(d) (5 pts) Having resolved any possible conflicts in requirements, specify conditions (models of behavior), if any exist, under which the program can terminate. If none exist, please indicate so.

Solution: