

# \* Notation $O$ \*

Prouver que:

$$n = O(n)$$

$$\Leftrightarrow \begin{cases} f(n) = O(g(n)) \\ f(n) = n \in \mathbb{N} \\ g(n) = n \in \mathbb{N} \end{cases}$$

$$n = O(n) \Leftrightarrow \exists n_0 \in \mathbb{N}, \exists c > 0, \forall n \geq n_0:$$

on divise  
des deux côtés  
par  $n$ , car  $n \geq 1$ .

$$\Rightarrow$$

$$1 \leq c \cdot n$$

$$1 \leq c$$

Comme  $1 \leq c$  nous assure que,  $n = O(n)$ ,  
il suffit de prendre  $n_0 = 1$  pour que  
la relation reste vraie.

$$2n = O(3n) \Leftrightarrow \begin{cases} f(n) = O(g(n)) \\ f(n) = 2n \in \mathbb{N} \\ g(n) = 3n \in \mathbb{N} \end{cases}$$

$$2n = O(3n) \Leftrightarrow \exists n_0 \in \mathbb{N}, \exists c > 0, \forall n \geq n_0:$$

on divise par  $n$ , car  $n \geq 1$

$$2 \leq c \cdot (3n)$$

$$\Rightarrow 2 \leq 3c$$

$$\Rightarrow 2/3 \leq c$$

Comme  $2/3 \leq c$  nous assure que  $2n = O(3n)$   
il suffit de prendre  $n_0 = 1$ .

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$$\bullet \quad n+2 = O(n) \Leftrightarrow \begin{cases} f(n) = O(g(n)) \\ f(n) = n+2 \in \mathbb{N} \\ g(n) = n \in \mathbb{N} \end{cases}$$

$$n+2 = O(n) \Leftrightarrow \exists n_0 \in \mathbb{N}, \exists c > 0, \forall n > n_0, n+2 \leq c \cdot n$$

~~Il faut trouver un  $c$  et un  $n_0$  tels que  $n+2 \leq c \cdot n$  pour tout  $n > n_0$ .~~

~~On peut prendre  $c=2$  et  $n_0=1$ .~~

$$c=2 \text{ et } n_0=1 \text{ marche, car: } n+2 \leq 2n.$$

$$\bullet \quad \sqrt{n} = O(n) \Leftrightarrow \begin{cases} f(n) = O(g(n)) \\ f(n) = \sqrt{n} \in \mathbb{R}^+ \\ g(n) = n \in \mathbb{N} \end{cases}$$

$$\sqrt{n} = O(n) \Leftrightarrow \exists n_0 \in \mathbb{N}, \exists c > 0, \forall n > n_0, \sqrt{n} \leq c \cdot n$$

$$\Rightarrow \frac{1}{\sqrt{n}} \leq c$$

$$\text{Pour } n \geq 1, \quad n^2 \geq n \\ \Rightarrow \sqrt{n^2} \geq \sqrt{n} \\ \Rightarrow n \geq \sqrt{n}$$

$$c=1 \text{ et } n_0=1.$$



$$\bullet \log(n) = O(n) \quad \left\{ \begin{array}{l} f(n) = O(g(n)) \\ f(n) = \log_2(n) \\ g(n) = n \end{array} \right.$$

$$\log(n) = O(n) \Leftrightarrow \exists n_0 \in \mathbb{N}, \exists c > 0, \forall n > n_0, \log(n) \leq c \cdot n$$

$$\Rightarrow \frac{\log(n)}{n} \leq c$$

Par ailleurs,  $\forall n \geq 1, \log(n) < n$

il suffit de prendre  $c = 1$  et  $n_0 = 1$ .

$$\bullet n = O(n^2) \quad \left\{ \begin{array}{l} f(n) = O(g(n)) \\ f(n) = n \\ g(n) = n^2 \end{array} \right.$$

Ceci est trivialement vrai, donc

$$n = O(n^2) \Leftrightarrow \exists n_0 \in \mathbb{N}, \exists c > 0, \forall n > n_0, n \leq c \cdot n^2$$

$$\Rightarrow \frac{1}{n} \leq c$$

car  $\forall n \geq 1, n^2 \geq n$

il suffit de prendre  $c = n_0 = 1$ .



## \* Notation o \*

$$\bullet \quad \sqrt{n} = o(n) \Leftrightarrow \begin{cases} f(n) = o(g(n)) \\ f(n) = \sqrt{n} \in \mathbb{R}^+ \\ g(n) = n \in \mathbb{N} \end{cases}$$

il suffit de remarquer que :

$$\forall n \geq 1 \quad \frac{\sqrt{n}}{n} = \frac{1}{\sqrt{n}} \xrightarrow{n \rightarrow +\infty} 0$$

ie :  $\lim_{n \rightarrow +\infty} \frac{1}{\sqrt{n}} = 0$

Donc,  $\sqrt{n} = o(n)$ .

$$\bullet \quad \log(n) = o(n) \Leftrightarrow \begin{cases} f(n) = o(g(n)) \\ f(n) = \log(n) \\ g(n) = n \end{cases}$$

il suffit de remarquer que :

$$\forall n \geq 1 \quad \frac{\log(n)}{n} = \frac{\frac{\log(n)}{\log(10)}}{\frac{n}{\log(10)}} = \frac{1}{\log(10)} \left( \frac{\log(n)}{n} \right)$$

Ainsi :  $\lim_{n \rightarrow +\infty} \frac{\log(n)}{n} = \frac{1}{\log(10)} \lim_{n \rightarrow +\infty} \left( \frac{\log(n)}{n} \right)$

situation  $\frac{\infty}{\infty}$   
 $\rightarrow$  Règle d'Hôpital  $= \frac{1}{\log(10)} \lim_{n \rightarrow +\infty} \left( \frac{\log(n)}{n} \right)$

$$= \frac{1}{\log(10)} \lim_{n \rightarrow +\infty} \left( \frac{\log'(n)}{n'} \right)$$

$$= \frac{1}{\log(10)} \lim_{n \rightarrow +\infty} \left( \frac{1/n}{1} \right)$$

$$= 0$$



$$\bullet \quad n = o(n^2) \iff \begin{cases} f(n) = o(g(n)) \\ f(n) = n \\ g(n) = n^2 \end{cases}$$

$$\forall n \geq 1$$

$$\frac{n}{n^2} = \frac{1}{n} \xrightarrow{n \rightarrow \infty} 0$$

$$\bullet \quad \log(n) = o(\sqrt{n}) \iff \begin{cases} f(n) = o(g(n)) \\ f(n) = \log(n) \\ g(n) = \sqrt{n} \end{cases}$$

$$\forall n \geq 1$$

$$\frac{\log(n)}{\sqrt{n}} = \frac{\log(n)}{\sqrt{n}}$$

$$= \left( \frac{1}{\log(40)} \right) \frac{\log(n)}{\sqrt{n}}$$

~~$$\frac{\log(n)}{\sqrt{n}}$$~~

$$= \left( \frac{1}{\log(40)} \right) \cdot 2 \cdot \left( \frac{\frac{1}{2} \log(n)}{\sqrt{n}} \right)$$

$$= \left( \frac{2}{\log(40)} \right) \cdot \frac{\log(n^{1/2})}{n^{1/2}}$$

$$\xrightarrow{n \rightarrow \infty} 0$$