Trento University ICPC Team Notebook (2015-16)

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1 Graphs

1.1 Max-flow

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
//
// Running time:
// O(|V|^2 |E|)
//
// INPUT:
// - graph, constructed using AddEdge()
// - source and sink
//
// OUTPUT:
// - maximum flow value
// - To obtain actual flow values, look at edges with capacity > 0
// (zero capacity edges are residual edges).
```

```
#include<cstdio>
#include<vector>
#include<queue>
using namespace std;
typedef long long LL;
struct Edge {
  int u, v;
  LL cap, flow;
  Edge() {}
 Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0) {}
struct Dinic {
  vector<Edge> E;
  vector<vector<int>> g;
  vector<int> d, pt;
  Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
    if (u != v) {
      E.emplace_back(Edge(u, v, cap));
      g[u].emplace_back(E.size() - 1);
      E.emplace_back(Edge(v, u, 0));
      g[v].emplace_back(E.size() - 1);
  bool BFS(int S, int T) {
    queue<int> q({S});
     fill(d.begin(), d.end(), N + 1);
    while(!q.empty()) {
      int u = q.front(); q.pop();
if (u == T) break;
      for (int k: g[u]) {
        Edge &e = E[k];
if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
  d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
    return d[T] != N + 1;
  LL DFS (int u, int T, LL flow = -1) {
    if (u == T || flow == 0) return flow;
    Edge &e = E[g[u][i];

Edge &e = E[g[u][i]^1];

if (d[e.v] == d[e.u] + 1) {

   LL amt = e.cap - e.flow;

   if (flow != -1 && amt > flow) amt = flow;
        if (LL pushed = DFS(e.v, T, amt)) {
          e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
    return 0;
  LL MaxFlow(int S, int T) {
    LL total = 0;
    while (BFS(S, T)) {
     fill(pt.begin(), pt.end(), 0);
while (LL flow = DFS(S, T))
        total += flow;
    return total;
};
// The following code solves SPOJ problem #4110: Fast Maximum Flow (FASTFLOW)
int main()
  scanf("%d%d", &N, &E);
  Dinic dinic(N);
  for (int i = 0; i < E; i++)
    int u, v;
    scanf("%d%d%lld", &u, &v, &cap);
    dinic.AddEdge(u - 1, v - 1, cap);
```

```
dinic.AddEdge(v - 1, u - 1, cap);
}
printf("%lld\n", dinic.MaxFlow(0, N - 1));
return 0;
}
// END CUT
```

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
      max flow:
                            O(|V|^3) augmentations
       min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
      - graph, constructed using AddEdge()
       - source
      - sink
// OUTPUT:
       - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII:
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found;
  VL dist, pi, width;
  VPII dad;
  MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
    this->cap[from][to] = cap;
this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
  L val = dist[s] + pi[s] - pi[k] + cost;
    if (cap && val < dist[k]) {</pre>
      dist[k] = val;
      dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1;
       found[s] = true;
      for (int k = 0; k < N; k++) {
   if (found[k]) continue;</pre>
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
Relax(s, k, flow[k][s], -cost[k][s], -1);
if (best == -1 || dist[k] < dist[best]) best = k;
      s = best;
```

```
for (int k = 0; k < N; k++)
     pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
    L totflow = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
           totcost += amt * cost[dad[x].first][x];
        | else {
          flow[x][dad[x].first] -= amt;
totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
// BEGIN CUT
// The following code solves UVA problem #10594: Data Flow
int main() {
 int N. M:
  while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
      scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {
    mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);</pre>
      mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
    pair<L, L> res = mcmf.GetMaxFlow(0, N);
    if (res.first == D) {
      printf("%Ld\n", res.second);
    else
     printf("Impossible.\n");
  return 0;
// END CUT
```

1.3 Min-cost matching

```
// Min cost bipartite matching via shortest augmenting paths
 // This is an O(n^3) implementation of a shortest augmenting path
 // algorithm for finding min cost perfect matchings in dense
 // graphs. In practice, it solves 1000 \times 1
 // second.
                         cost[i][j] = cost for pairing left node i with right node j
Lmate[i] = index of right node that left node i pairs with
                         Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
#include <algorithm>
 #include <cstdio>
 #include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
 double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
```

```
int n = int(cost.size());
// construct dual feasible solution
VD u(n);
for (int i = 0; i < n; i++) {
  u[i] = cost[i][0];
  for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);
for (int j = 0; j < n; j++) {
  v[j] = cost[0][j] - u[0];</pre>
  for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
// construct primal solution satisfying complementary slackness
Lmate = VI(n, -1);
Rmate = VI(n, -1);
int mated = 0;
for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
   if (Rmate[j] != -1) continue;</pre>
    if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
      Lmate[i] = j;
      Rmate[j] = i;
      mated++;
      break:
VD dist(n);
VI dad(n);
// repeat until primal solution is feasible
while (mated < n) {</pre>
  // find an unmatched left node
  while (Lmate[s] != -1) s++;
  // initialize Dijkstra
 fill(dad.begin(), dad.end(), -1);
 fill(seen.begin(), seen.end(), 0);
for (int k = 0; k < n; k++)</pre>
    dist[k] = cost[s][k] - u[s] - v[k];
  while (true) {
    // find closest
     i = -1:
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      if (j == -1 || dist[k] < dist[j]) j = k;</pre>
    seen[j] = 1;
     // termination condition
    if (Rmate[j] == -1) break;
    // relax neighbors
    const int i = Rmate[j];
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
       \label{eq:const_double} \ \mbox{new\_dist} \ = \ \mbox{dist}[\ j] \ + \ \mbox{cost}[\ i] \ [\ k] \ - \ \mbox{u}[\ i] \ - \ \mbox{v}[\ k];
      if (dist[k] > new_dist) {
        dist[k] = new_dist;
dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++) {
  if (k == j || !seen[k]) continue;
  const int i = Rmate[k];</pre>
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
  // augment along path
  while (dad[j] >= 0) {
  const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
  Lmate[s] = j;
```

```
mated++;
}
double value = 0;
for (int i = 0; i < n; i++)
  value += cost[i][Lmate[i]];
return value;</pre>
```

1.4 Max bipartite matchine

```
// This code performs maximum bipartite matching.
// Running time: O(|E| |V|) -- often much faster in practice
     INPUT: w[i][j] = edge between row node i and column node j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
             mc[j] = assignment for column node j, -1 if unassigned
             function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {</pre>
    if (w[i][j] && !seen[j]) {
      seen[j] = true;
      if (mc[j] < 0 || FindMatch(mc[j], w, mr, mc, seen)) {</pre>
       mr[i] = j;
mc[j] = i;
        return true;
  return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
 mr = VI(w.size(), -1);
 mc = VI(w[0].size(), -1);
 int ct = 0;
for (int i = 0; i < w.size(); i++) {</pre>
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

1.5 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
      0(|V|^3)
// TNPHT.
       - graph, constructed using AddEdge()
// OUTPUT:
      - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
 int N = weights.size();
  VI used(N), cut, best cut;
  int best_weight = -1;
```

```
for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {</pre>
      prev = last;
      last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j]; for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];
        used[last] = true;
        cut.push back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best_cut = cut;
          best_weight = w[last];
        for (int j = 0; j < N; j++)
          w[j] += weights[last][j];
         added[last] = true;
  return make_pair(best_weight, best_cut);
// BEGIN CUT
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
int main() {
  for (int i = 0; i < N; i++) {
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl;
// END CUT
```

1.6 Topological sort (C++)

```
// This function uses performs a non-recursive topological sort.
// Running time: O(|V|^2). If you use adjacency lists (vector<map<int> >),
                 the running time is reduced to O(|E|).
     INPUT: w[i][j] = 1 if i should come before j, 0 otherwise
     OUTPUT: a permutation of 0,...,n-1 (stored in a vector)
              which represents an ordering of the nodes which
              is consistent with w
// If no ordering is possible, false is returned.
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std:
typedef double T:
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool TopologicalSort (const VVI &w, VI &order) {
 int n = w.size();
  VI parents (n);
  queue<int> q;
  order.clear();
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)</pre>
     if (w[i][i]) parents[i]++;
      if (parents[i] == 0) q.push (i);
```

```
while (q.size() > 0){
   int i = q.front();
   q.pop();
   order.push_back (i);
   for (int j = 0; j < n; j++) if (w[i][j]) {
      parents[j]--;
      if (parents[j] == 0) q.push (j);
   }
}
return (order.size() == n);</pre>
```

1.7 Dijkstra and Floyd's algorithm (C++)

```
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
// This function runs Dijkstra's algorithm for single source
// shortest paths. No negative cycles allowed!
// Running time: O(|V|^2)
    INPUT: start, w[i][j] = cost of edge from i to j
    OUTPUT: dist[i] = min weight path from start to i
              prev[i] = previous node on the best path from the
                         start node
void Dijkstra (const VVT &w, VT &dist, VI &prev, int start) {
 int n = w.size();
  VI found (n);
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
  dist[start] = 0;
  while (start != -1) {
    found[start] = true;
    int best = -1;
    for (int k = 0; k < n; k++) if (!found[k]) {
      if (dist[k] > dist[start] + w[start][k]){
        dist[k] = dist[start] + w[start][k];
        prev[k] = start;
      if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    start = best;
// This function runs the Floyd-Warshall algorithm for all-pairs
// shortest paths. Also handles negative edge weights. Returns true
// if a negative weight cycle is found.
// Running time: O(|V|^3)
    INPUT: w[i][j] = weight of edge from i to j
    OUTPUT: w[i][j] = shortest path from i to j
             prev[i][j] = node before j on the best path starting at i
bool FloydWarshall (VVT &w, VVI &prev) {
  int n = w.size();
 prev = VVI (n, VI(n, -1));
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {</pre>
        if (w[i][j] > w[i][k] + w[k][j]) {
         w[i][j] = w[i][k] + w[k][j];
          prev[i][j] = k;
  // check for negative weight cycles
```

```
for(int i=0;i<n;i++)
   if (w[i][i] < 0) return false;
return true;</pre>
```

1.8 Strongly connected components

```
#include <vector>
#include <iostream>
#include <cmath>
using namespace std;
using vi = vector<int>;
using vvi = vector<vector<int>>;
const int UNVISITED = -1;
struct SCC {
       vvi AdjList;
       dfsNumberCounter(0) {}
        void tarjanSCC(int u) {
                dfs_low[u] = dfs_num[u] = dfsNumberCounter++;
               S.push_back(u);
                visited[u] = 1;
               for (int j = 0; j < (int)AdjList[u].size(); j++) {</pre>
                       int v = AdjList[u][j];
                       if (dfs_num[v] == UNVISITED)
     tarjanSCC(v);
                        if (visited[v])
                               dfs_low[u] = min(dfs_low[u], dfs_low[v]);
               if (dfs_low[u] == dfs_num[u]) {
                               int v = S.back(); S.pop_back(); visited[v] = 0;
                               scc[v] = num_scc;
                               if (u == v) {
                                       num scc++;
                                       break;
        void addEdge(int u, int v) {
               AdjList[u].push_back(v);
        void run() {
               for (int i = 0; i < V; i++)</pre>
                       if (dfs_num[i] == UNVISITED)
                               tarjanSCC(i);
};
```

1.9 Eulerian path

1.10 Kruskal's algorithm

```
Uses Kruskal's Algorithm to calculate the weight of the minimum spanning
forest (union of minimum spanning trees of each connected component) of
a possibly disjoint graph, given in the form of a matrix of edge weights
(-1 if no edge exists). Returns the weight of the minimum spanning
forest (also calculates the actual edges - stored in T). Note: uses a disjoint-set data structure with amortized (effectively) constant time per
union/find. Runs in O(E*log(E)) time.
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
typedef int T;
struct edge
  int u. v:
  Td;
};
struct edgeCmp
  int operator()(const edge& a, const edge& b) { return a.d > b.d; }
int find(vector <int>& C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }
T Kruskal (vector <vector <T> >& w)
  int n = w.size():
  T weight = 0;
  vector <int> C(n), R(n);
  for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }</pre>
  vector <edge> T;
  priority_queue <edge, vector <edge>, edgeCmp> E;
  for (int i=0; i<n; i++)</pre>
    for (int j=i+1; j<n; j++)</pre>
      if(w[i][j] >= 0)
        e.u = i; e.v = j; e.d = w[i][j];
        E.push(e);
  while (T.size() < n-1 && !E.empty())
    edge cur = E.top(); E.pop();
    int uc = find(C, cur.u), vc = find(C, cur.v);
    if(uc != vc)
      T.push_back(cur); weight += cur.d;
      if(R[uc] > R[vc]) C[vc] = uc;
      else if(R[vc] > R[uc]) C[uc] = vc;
      else { C[vc] = uc; R[uc]++; }
  return weight;
```

```
}
int main()
{
    int wa[6][6] = {
        { 0, -1, 2, -1, 7, -1 },
        { -1, 0, -1, 2, -1, -1 },
        { 2, -1, 0, -1, 8, 6 },
        { -1, 2, -1, 0, -1, -1 },
        { 7, -1, 8, -1, 0, 4 },
        { -1, -1, 6, -1, 4, 0 } };

vector <vector <int> > w(6, vector <int>(6));

for(int i=0; i<6; i++)
    for(int j=0; j<6; j++)
        w[i][j] = wa[i][j];

cout < Kruskal(w) << endl;
    cin >> wa[0][0];
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise, starting
             with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
// BEGIN CUT
#include <map>
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
  T x, v;
  PT() {}
  PT(T x, T y) : x(x), y(y) {}
bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }
  bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  #endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
```

```
dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn;
#endif
// BEGIN CUT
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
  int t;
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n;
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v);
    map<PT.int> index:
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h):
    double len = 0;
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
  if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y) {}
PT operator + (const PT &p) const { return PT(x+p.x, y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
                                const { return PT(x*c, y*c );
  PT operator * (double c)
  PT operator / (double c)
                                const { return PT(x/c, y/c ); }
double dot(PT p, PT q)
                             { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                               return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
 return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
```

~1

```
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                            double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
  if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
    dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
      return false;
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
 // compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2:
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
 // tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
     \begin{array}{l} \textbf{if } (\{p[i],y \leq q,y \ \& \ q,y < p[j],y \ || \\ p[j],y < q,y \ \& \& \ q,y < p[i],y) \ \& \& \\ q,x < p[i],x + (p[j],x - p[i],x) + (q,y - p[i],y) \ / \ (p[j],y - p[i],y)) \end{array} 
      c = !c:
  return c;
 // determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
```

```
for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true:
// compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
 b = b-a;
  a = a-c:
  double A = dot(b, b);
double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret .push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius \boldsymbol{r}
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret:
  double d = sqrt(dist2(a, b));
  if (d > r+R || d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
 PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
 for (int i = 0; i < p.size(); i++) {
  for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 || j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
int main() {
  // expected: (-5,2)
cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5.2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
```

```
// expected: (5,2) (7.5,3) (2.5,1)
cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
     << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
     << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
// expected: 6.78903
cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
    << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
     << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 0 0 1
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1.2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
// expected: (1.1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
    << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "
     << PointInPolygon(v, PT(5,2)) << " "
     << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
            (5,4) (4,5)
            hlank line
            (4.5) (5.4)
            blank line
            (4.5) (5.4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
```

2.3 3D geometry

```
public class Geom3D {
   // distance from point (x, y, z) to plane aX + bY + cZ + d = 0
   public static double ptPlaneDist(double x, double y, double z,
        double a, double b, double c, double d) {
```

```
return Math.abs(a*x + b*y + c*z + d) / Math.sqrt(a*a + b*b + c*c);
// distance between parallel planes aX + bY + cZ + d1 = 0 and
// aX + bY + cZ + d2 = 0
public static double planePlaneDist(double a, double b, double c,
   double d1, double d2) {
 return Math.abs(d1 - d2) / Math.sqrt(a*a + b*b + c*c);
// distance from point (px, py, pz) to line (x1, y1, z1)-(x2, y2, z2)
// (or ray, or segment; in the case of the ray, the endpoint is the
// first point)
public static final int LINE = 0:
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1, double z1,
   double x2, double y2, double z2, double px, double py, double pz,
    int type) {
 double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1-z2)*(z1-z2);
 if (pd2 == 0) {
   x = x1:
   y = y1;
z = z1;
 } else {
    double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (pz-z1)*(z2-z1)) / pd2;
    x = x1 + u * (x2 - x1);
    y = y1 + u * (y2 - y1);
    z = z1 + u * (z2 - z1);
    if (type != LINE && u < 0) {
     x = x1;
      y = y1;
    if (type == SEGMENT && u > 1.0) {
      x = x2;
      v = v2;
      z = z2;
 \textbf{return} \  \  (x-px) \star (x-px) \  \  + \  \  (y-py) \star (y-py) \  \  + \  \  (z-pz) \star (z-pz) \, ;
public static double ptLineDist(double x1, double y1, double z1,
    double x2, double y2, double z2, double px, double py, double pz,
  return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, z2, px, py, pz, type));
```

3 Math

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int ged(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
```

```
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret:
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a:
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret;
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)
                         ret.push_back(mod(x + i*(n / g), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (g > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
//\ z\ \$\ m1\ =\ r1,\ z\ \$\ m2\ =\ r2.\quad Here,\ z\ is\ unique\ modulo\ M\ =\ lcm\,(m1,\ m2)\ .
// Return (z, M). On failure, M = -1.
PII chinese remainder theorem(int ml, int rl, int m2, int r2) {
       int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / q, m1*m2 / q);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int = 1; i < m.size(); i++) {
    ret = chinese remainder theorem (ret.second, ret.first, m[i], r[i]);
    if (ret.second = -1) break;</pre>
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                if (c) return false;
                return true;
        if (!a)
                if (c % b) return false;
                x = 0; y = c / b;
                return true;
        if (!b)
```

```
if (c % a) return false;
                   x = c / a; y = 0;
                   return true;
         int g = gcd(a, b);
         if (c % g) return false;
         x = c / g * mod_inverse(a / g, b / g);
         y = (c - a*x) / b;
         return true;
int main() {
         // expected: 2
         cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
         int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
         // expected: 95 451
         VI sols = modular_linear_equation_solver(14, 30, 100);
         for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
         cout << endl:
         // expected: 8
         cout << mod_inverse(8, 9) << endl;</pre>
         // expected: 23 105
                     11 12
         PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
         cout << ret.first << " " << ret.second << endl;</pre>
         ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;</pre>
          // expected: 5 -15
         if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;</pre>
         return 0;
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
    (1) solving systems of linear equations (AX=B)
    (2) inverting matrices (AX=I)
    (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT: a[][] = an nxn matrix
             b[][] = an nxm matrix
// OUTPUT: X
                     = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std:
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1:
  for (int i = 0; i < n; i++) {
   int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])
  for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
        if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
```

```
swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c; for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
      a[p][pk] = 0;
      for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det:
int main() {
  const int n = 4:
  const int m = 2:
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
    a[i] = VT(A[i], A[i] + n);
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;
  // expected: -0.233333 0.166667 0.133333 0.0666667
                 0.166667 0.166667 0.333333 -0.333333
                  0.233333 0.833333 -0.133333 -0.0666667
                 0.05 -0.75 -0.1 0.2
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
    cout << a[i][j] << ' ';
    cout << endl:
  // expected: 1.63333 1.3
                 -0.166667 0.5
                 2.36667 1.7
                 -1.85 -1.35
  cout << "Solution: " << endl;
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++)
cout << b[i][j] << ' ';
    cout << endl;
```

3.3 Fast Fourier transform

```
#include <cassert>
#include <catdio>
#include <cmath>

struct cpx
{
    cpx(){}
    cpx(double aa):a(aa),b(0){}
    cpx(double aa, double bb):a(aa),b(bb){}
    double a;
    double b;
    double bmodsq(void) const
    {
        return a * a + b * b;
    }
    cpx bar(void) const
    {
        return cpx(a, -b);
    }
};
```

```
cpx operator + (cpx a, cpx b)
  return cpx(a.a + b.a, a.b + b.b);
cpx operator *(cpx a, cpx b)
  return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
  cpx r = a * b.bar():
  return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP (double theta)
  return cpx(cos(theta), sin(theta));
const double two_pi = 4 * acos(0);
// in:
            input array
// out: output array
// step: {SET TO 1} (used internally)
// size: length of the input/output {MUST BE A POWER OF 2}
// dir: either plus or minus one (direction of the FFT)
// RESULT: out \{k\} = \sum_{j=0}^{\infty} \{size - 1\} in[j] * exp(dir * 2pi * i * j * k / size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
  if(size < 1) return;</pre>
  if(size == 1)
    out[0] = in[0];
    return;
  FFT(in, out, step \star 2, size / 2, dir);
  FFT(in + step, out + size / 2, step * 2, size / 2, dir);
  for(int i = 0; i < size / 2; i++)
    cpx even = out[i];
    cpx odd = out[i + size / 2];
    out[i] = even + EXP(dir * two_pi * i / size) * odd;
    out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
// f[0...N-1] and g[0..N-1] are numbers
// Want to compute the convolution h, defined by
// h[n] = \operatorname{sum} \text{ of } f[k]g[n-k] \ (k=0,\ldots,N-1).

// Here, the index is cyclic; f[-1] = f[N-1], f[-2] = f[N-2], etc.

// Let F[0\ldots N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N log N) time, do the following:
    1. Compute F and G (pass dir = 1 as the argument).
// 2. Get H by element-wise multiplying F and G.
     3. Get h by taking the inverse FFT (use dir = -1 as the argument)
        and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
  printf("If rows come in identical pairs, then everything works.\n");
  cpx \ a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\};
  cpx b[8] = \{1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2\};
  cpx A[8];
  cpx B[8];
  FFT(a, A, 1, 8, 1);
FFT(b, B, 1, 8, 1);
  for(int i = 0; i < 8; i++)
    printf("%7.21f%7.21f", A[i].a, A[i].b);
   printf("\n");
   for(int i = 0 ; i < 8 ; i++)
     cpx Ai(0,0);
    for (int j = 0; j < 8; j++)
      Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
    printf("%7.21f%7.21f", Ai.a, Ai.b);
  printf("\n");
   cpx AB[8];
   for (int i = 0; i < 8; i++)
    AB[i] = A[i] * B[i];
   cpx aconvb[8];
```

3.4 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
        maximize
        subject to Ax <= b
                       x >= 0
// INPUT: A -- an m x n matrix
           b -- an m-dimensional vector
            c -- an n-dimensional vector
            x -- a vector where the optimal solution will be stored
 // OUTPUT: value of the optimal solution (infinity if unbounded
            above, nan if infeasible)
/\!/ To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
  VI B, N;
  VVD D
  LPSolver(const VVD &A, const VD &b, const VD &c) :
     m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
    for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j]; for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i]; } for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
    N[n] = -1; D[m + 1][n] = 1;
  void Pivot(int r, int s) {
     double inv = 1.0 / D[r][s];
     for (int i = 0; i < m + 2; i++) if (i != r)
       for (int j = 0; j < n + 2; j++) if (j != s)
         D[i][j] -= D[r][j] * D[i][s] * inv;
     for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv; for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
     D[r][s] = inv;
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
     while (true) {
       int s = -1;
       for (int j = 0; j <= n; j++) {
  if (phase == 2 && N[j] == -1) continue;</pre>
         if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] < N[s]) s = j;
```

```
if (D[x][s] > -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {</pre>
       if (D[i][s] < EPS) continue;</pre>
       if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||</pre>
          if (r == -1) return false;
     Pivot(r, s);
  DOUBLE Solve(VD &x) {
   int r = 0;
for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
       int s = -1;
       for (int j = 0; j <= n; j++)
         if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] < N[s]) s = j;
       Pivot(i, s);
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];</pre>
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE _A[m][n] = {
    { 6, -1, 0 },
    \{ -1, -5, 0 \},
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };
  DOUBLE _c[n] = { 1, -1, 0 };
  VVD A(m);
  VD b(\underline{b}, \underline{b} + m);
  for (int i = 0; i < m; i++) A[i] = VD(\_A[i], \_A[i] + n);
  LPSolver solver(A, b, c);
  VD x;
 DOUBLE value = solver.Solve(x);
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl;
 return 0:
```

3.5 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL:
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL) (sqrt ((double)(x))+EPS);
  for(LL i=5;i<=s;i+=6)
    if (!(x%i) || !(x%(i+2))) return false;
  return true;
// Primes less than 1000:
                                11
                                                   19
                                                        73
137
                                                                79
                                     61
113
                                                   71
             43
                         53
                                59
                                            67
                                                                     83
       97
                        107
                               109
                                           127
                                                              139
                                                                    149
                                                                           151
                               179
                                                  193
                                                        197
      157
            163
                  167
                                     181
                                           191
                                                              199
                                                                           223
                                                                    211
                        239
                               241
                                     251
317
                                           257
                                                  263
                                                                    277
                                                                           281
      227
            229
                  233
                                                        269
                                                              271
                   307
                                                        347
                                                                     353
                                                                           359
      283
            293
                                                              349
```

```
509
     521
           523
                541
                     547
                           557
                                563
                                      569
                                           571
599
                613
                                 631
                                      641
                                           643
                                                 647
                           701
                                 709
                                      719
                                                 733
                                                      739
                                                           743
           761
                769
                      773
                           787
                                 797
                                      809
                                           811
                                                 821
                                                      823
                                                           827
     839
           853
                857
                      859
                           863
                                877
                                      881
                                           883
                                                887
                                                      907
                                                           911
     929
           937
                941
                      947
                           953
                                 967
                                      971
The largest prime smaller than 10 is 7.
The largest prime smaller than 100 is 97.
The largest prime smaller than 1000 is 997.
The largest prime smaller than 10000 is 9973. The largest prime smaller than 100000 is 99991.
The largest prime smaller than 1000000 is 999983
The largest prime smaller than 10000000 is 9999991.
The largest prime smaller than 100000000 is 99999989.
The largest prime smaller than 1000000000 is 999999937
The largest prime smaller than 10000000000 is 9999999967.
The largest prime smaller than 10000000000 is 99999999977.
The largest prime smaller than 100000000000 is 999999999999.
The largest prime smaller than 1000000000000 is 999999999971.
The largest prime smaller than 1000000000000 is 9999999999973.
The largest prime smaller than 100000000000000 is 999999999999997.
```

3.6 Ternary Search

4 Data structures

4.1 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N)
    tree[x] += v;
    x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  int res = 0:
  while(x) {
    res += tree[x]:
   x -= (x & -x);
  return res;
```

```
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
  int idx = 0, mask = N;
  while(mask && idx < N) {
   int t = idx + mask;
   if(x >= tree[t]) {
      idx = t;
      x -= tree[t];
   }
  mask >>= 1;
  }
  return idx;
```

4.2 Union-find set

```
#include <iostream>
#include <vector>
using namespace std;
int find(vector<int> &C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }
void merge(vector<int> &C, int x, int y) { C[find(C, x)] = find(C, y); }
int main()
{
    int n = 5;
        vector<int> &C(n);
        for (int i = 0; i < n; i++) C[i] = i;
        merge(C, 0, 2);
        merge(C, 1, 0);
        merge(C, 3, 4);
        for (int i = 0; i < n; i++) cout << i << " " << find(C, i) << endl;
        return 0;
}</pre>
```

4.3 KD-tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation
// that's probably good enough for most things (current it's a
   - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
   - worst case for nearest-neighbor may be linear in pathological
     case
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
bool operator==(const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
    return a.x < b.x:
// sorts points on v-coordinate
bool on_y (const point &a, const point &b)
    return a.y < b.y;</pre>
```

```
// squared distance between points
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox
    ntvpe x0, x1, v0, v1;
    bbox(): x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
       computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {
            x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0)
                                 return pdist2(point(x0, y0), p);
            else if (p.y > y1) return pdist2(point(x0, y1), p);
                                 return pdist2(point(x0, p.y), p);
            else
        else if (p.x > x1) {
                                 return pdist2(point(x1, y0), p);
return pdist2(point(x1, y1), p);
            if (p.y < y0)
             else if (p.y > y1)
                                 return pdist2(point(x1, p.y), p);
        else {
            if (p.y < y0)
                                 return pdist2(point(p.x, y0), p);
             else if (p.y > y1) return pdist2(point(p.x, y1), p);
            else
                                 return 0;
1:
// stores a single node of the kd-tree, either internal or leaf
struct kdnode
    bool leaf;
                     // true if this is a leaf node (has one point)
    point pt;
                     // the single point of this is a leaf
                     // bounding box for set of points in children
    kdnode *first, *second; // two children of this kd-node
    kdnode() : leaf(false), first(0), second(0) {}
    "kdnode() { if (first) delete first; if (second) delete second; }
    // intersect a point with this node (returns squared distance)
    ntype intersect(const point &p) {
        return bound.distance(p);
    // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
         // compute bounding box for points at this node
        bound.compute(vp);
         // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
            leaf = true:
            pt = vp[0];
        else {
             // split on x if the bbox is wider than high (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
             // otherwise split on y-coordinate
             else
                sort(vp.begin(), vp.end(), on_y);
             // divide by taking half the array for each child
             // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnode(); first->construct(v1);
second = new kdnode(); second->construct(vr);
// simple kd-tree class to hold the tree and handle queries
struct kdtree
```

```
kdnode *root:
    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnode();
        root->construct(v);
    "kdtree() { delete root; }
    // recursive search method returns squared distance to nearest point
    ntype search(kdnode *node, const point &p)
        if (node->leaf) {
            // commented special case tells a point not to find itself
              if (p == node->pt) return sentry;
                return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box to search first
         // (note that the other side is also searched if needed)
        if (bfirst < bsecond) {
            ntype best = search(node->first, p);
            if (bsecond < best)</pre>
                best = min(best, search(node->second, p));
            return best;
            ntype best = search(node->second, p);
            if (bfirst < best)</pre>
                best = min(best, search(node->first, p));
            return best;
    // squared distance to the nearest
    ntype nearest(const point &p) {
        return search (root, p);
};
// some basic test code here
int main()
    // generate some random points for a kd-tree
    vector<point> vp;
for (int i = 0; i < 100000; ++i) {</pre>
        vp.push_back(point(rand()%100000, rand()%100000));
    kdtree tree(vp):
    // query some points
    for (int i = 0; i < 10; ++i) {
        point q(rand()%100000, rand()%100000);
        cout << "Closest squared distance to (" << q.x << ", " << q.y << ")"
             << " is " << tree.nearest(q) << endl;
    return 0;
```

4.4 Lazy segment tree

```
build(2 * curr, begin, mid, list);
                 build(2 * curr + 1, mid+1, end, list);
                 leaf[curr] = leaf[2*curr] + leaf[2*curr+1];
public void update(int begin, int end, int val) {
         update(1,0,origSize-1,begin,end,val);
public void update(int curr, int tBegin, int tEnd, int begin, int end, int val)
        if(tBegin >= begin && tEnd <= end)</pre>
                 update[curr] += val;
        else
                 leaf[curr] += (Math.min(end,tEnd)-Math.max(begin,tBegin)+1) * val;
                 int mid = (tBegin+tEnd)/2;
                 if(mid >= begin && tBegin <= end)
    update(2*curr, tBegin, mid, begin, end, val);</pre>
                 if(tEnd >= begin && mid+1 <= end)</pre>
                          update(2*curr+1, mid+1, tEnd, begin, end, val);
public long query(int begin, int end) {
        return query(1,0,origSize-1,begin,end);
public long query(int curr, int tBegin, int tEnd, int begin, int end) {
        if(tBegin >= begin && tEnd <= end)</pre>
                 if(update[curr] != 0) {
                          leaf[curr] += (tEnd-tBegin+1) * update[curr];
                          if(2*curr < update.length){</pre>
                                   update[2*curr] += update[curr];
update[2*curr+1] += update[curr];
                          update[curr] = 0;
                 return leaf[curr];
        else
                  leaf[curr] += (tEnd-tBegin+1) * update[curr];
                 if(2*curr < update.length){</pre>
                          update[2*curr] += update[curr];
update[2*curr+1] += update[curr];
                 update[curr] = 0;
                 int mid = (tBegin+tEnd)/2;
                 long ret = 0;
                 if(mid >= begin && tBegin <= end)</pre>
                          ret += query(2*curr, tBegin, mid, begin, end);
                 if(tEnd >= begin && mid+1 <= end)</pre>
                          ret += query(2*curr+1, mid+1, tEnd, begin, end);
```

4.5 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num nodes, log num nodes, root;
vector<int> children[max_nodes];
                                          // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1];
                                          // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
      ancestor does not exist
                                          //\ L[i] is the distance between node i and the root
int L[max_nodes];
// floor of the binary logarithm of n
int lb(unsigned int n)
    if(n==0)
       return -1:
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16; }
    if (n >= 1<< 8) { n >>= 8; p += 8; }
if (n >= 1<< 4) { n >>= 4; p += 4; }
    if (n >= 1 << 2) { n >>= 2; p += 2;
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    for(int j = 0; j < children[i].size(); j++)</pre>
        DFS(children[i][j], l+1);
int LCA(int p, int q)
     // ensure node p is at least as deep as node q
    if(L[p] < L[q])
```

```
swap(p, q);
     // "binary search" for the ancestor of node p situated on the same level as q
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1<<i) >= L[q])
             p = A[p][i];
    if(p == q)
        return p;
     // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if(A[p][i] != -1 && A[p][i] != A[q][i])
             p = A[p][i];
             q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
     // read num_nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for (int i = 0; i < num nodes; i++)
        int p:
        // read p, the parent of node i or -1 if node i is the root
         A[i][0] = p;
         if(p != -1)
             children[p] push_back(i);
        else
             root = i;
    // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)
        if(A[i][j-1] != -1)</pre>
                 A[i][j] = A[A[i][j-1]][j-1];
             else
                 A[i][j] = -1;
     // precompute L
    DFS (root, 0);
    return 0:
```

5 Strings

5.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
             of substring s[i...L-1] in the list of sorted suffixes.
             That is, if we take the inverse of the permutation suffix[],
             we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray
  const int L;
  string s;
  vector<vector<int> > P:
  vector<pair<int,int>,int> > M;
  SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) { for (int i = 0; i < L; i++) P[0][i] = int(s[i]); for (int skip = 1, level = 1; skip < L; skip \star = 2, level++) {
       P.push_back(vector<int>(L, 0));
       for (int i = 0; i < L; i++)
```

```
M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
      sort (M.begin(), M.end());
      for (int i = 0; i < L; i++)
        P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
  if (P[k][i] == P[k][j]) {</pre>
         j += 1 << k;
        len += 1 << k;
    return len;
};
// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T:
  cin >> T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    string s;
    cin >> s;
    SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
      int len = 0, count = 0;
      for (int j = i+1; j < s.length(); j++) {</pre>
        int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len;
        bestcount = count;
        bestpos = i;
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;</pre>
    } else {
      cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
      bocel is the 1'st suffix
        ocel is the 6'th suffix
        cel is the 2'nd suffix
         el is the 3'rd suffix
           1 is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
  cout << endl;</pre>
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

5.2 Knuth-Morris-Pratt

```
Searches for the string w in the string s (of length k). Returns the
```

```
0-based index of the first match (k if no match is found). Algorithm
runs in O(k) time.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildTable(string& w, VI& t)
  t = VI(w.length());
 t[0] = -1; t[1] = 0;
  while(i < w.length())</pre>
   if(w[i-1] == w[j]) { t[i] = j+1; i++; j++; }
    else if(j > 0) \bar{j} = t[j];
    else { t[i] = 0; i++; }
int KMP (string& s. string& w)
  int m = 0, i = 0;
 VI t:
  buildTable(w, t);
  while (m+i < s.length())
    if(w[i] == s[m+i])
     if(i == w.length()) return m;
    else
      m += i-t[i];
     if(i > 0) i = t[i];
  return s.length();
  string \ a = (string) "The example above illustrates the general technique for assembling "+
    "the table with a minimum of fuss. The principle is that of the overall search: "+
    "most of the work was already done in getting to the current position, so very "+
    "little needs to be done in leaving it. The only minor complication is that the "+
    "logic which is correct late in the string erroneously gives non-proper "+
    "substrings at the beginning. This necessitates some initialization code.";
  string b = "table";
  int p = KMP(a, b);
  cout << p << ": " << a.substr(p, b.length()) << " " << b << endl;
```

6 Miscellaneous

6.1 Java BigInteger

```
// Java program to find large factorials using BigInteger
import java.util.Scanner;
public class Example
{
    // Returns Factorial of N
    static BigInteger factorial(int N)
    {
        // Initialize result
        BigInteger f = new BigInteger("1"); // Or BigInteger.ONE
        // Multiply f with 2, 3, ...N
        for (int i = 2; i <= N; i++)
              f = f.multiply(BigInteger.valueOf(i));
    return f;
}</pre>
```

```
// Driver method
public static void main(String args[]) throws Exception
    System.out.println(factorial(N));
    int a, b;
    BigInteger A, B;
    a = 54;
    b = 23;
    A = BigInteger.valueOf(54);
    B = BigInteger.valueOf(37);
    A = new BigInteger( 5 4 );
    B = new BigInteger( 123456789123456789);
    A = BigInteger.ONE:
    // Other than this, available constant are BigInteger.ZERO
    // and BigInteger.TEN
    String str = 123456789;
BigInteger C = A.add(new BigInteger(str));
    int val = 123456789;
    BigInteger C = A.add(BigIntger.valueOf(val));
    int x = A.intValue(); // value should be in limit of int x
    long y = A.longValue(); // value should be in limit of long y
    String z = A.toString();
   if (A.compareTo(B) < 0) {} // For BigInteger
if (A.equals(B)) {} // A is equal to B</pre>
```

6.2 Java input/output

```
// Java program to read data of various types using Scanner class.
import java.util.Scanner;
public class ScannerDemol
    public static void main(String[] args)
        // Declare the object and initialize with
        // predefined standard input object
        Scanner sc = new Scanner(System.in);
        // String input
        String name = sc.nextLine();
        // Character input
        char gender = sc.next().charAt(0);
        // Numerical data input
        // byte, short and float can be read
        // using similar-named functions.
        int age = sc.nextInt();
        long mobileNo = sc.nextLong();
        double cgpa = sc.nextDouble();
        // Print the values to check if input was correctly obtained.
        System.out.println("Name: "+name);
        System.out.println("Gender: "+gender);
        System.out.println("Age: "+age);
        System.out.println("Mobile Number: "+mobileNo);
        System.out.println("CGPA: "+cqpa);
```

6.3 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
{
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);

    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);</pre>
```

```
// Output a '+' before positive values
cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;
cout.unsetf(ios::showpos);

// Output numerical values in hexadecimal
cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;</pre>
```

6.4 C++ std algorithm

```
int main() {
        //Non-modifying sequence operations:
        std::array<int,8> foo = {3,5,7,11,13,17,19,23};
        // Same as any_of, none_of
        std::all_of(foo.begin(), foo.end(), [](int i){return i%2;}); // Return bool
        std::find (myints, myints+4, 30); // Return iterator
        std::find_if (foo.begin(), foo.end(), [](int i){return i%2;}) // Return iterator to first
        // Searches the range [first1,last1) for the last occurrence of the sequence defined by [
               first2, last2)
        std::find_end (foo.begin(), foo.end(), foo.begin() + 5, foo.end(), [](int i, int j) {return i
              == ;;;;;;
         //Same as prev, but return first ocurence
        std::search (foo.begin(), foo.end(), foo.begin() + 5, foo.end(), [](int i, int j) {return i ==
         //Find first occurence of any ellement from 2-nd seg
        std:: find_first_of (haystack.begin(), haystack.end(), needle, needle+3, comp_case_insensitive)
        //Searches the range [first,last) for the first occurrence of two consecutive elements that
        std::adjacent_find (foo.begin(), foo.end(), myfunction); //return iterator to first
        std::count (myvector.begin(), myvector.end(), 20); //return int
        std::count_if (myvector.begin(), myvector.end(), IsOdd); //return int
        // Compares the elements in the range [first1,last1) with those in the range beginning at
               first2.
         // and returns the first element of both sequences that does not match.
        std::mismatch (myvector.begin(), myvector.end(), myvector2.begin(), mypredicate); //Return
              pair of iterators
         //All elements the same
        std::equal (myvector.begin(), myvector.end(), myvector.begin(), mypredicate); //Return bool
        //All elements the same in different order
        std::is_permutation (foo.begin(), foo.end(), bar.begin());
        //Modifying sequence operations:
         //Copies the elements in the range [first,last] into the range beginning at result.
        std::copy (myvector1.begin(), myvector2.begin()+7, myvector.begin());
        std::copy_n ( myvector1.begin(), 7, myvector.begin() );
        std::copy_if (foo.begin(), foo.end(), bar.begin(), [](int i){return !(i<0);} );</pre>
        //Copies the elements in the range [first,last) starting from the end into the range
               terminating at result.
        std::copy_backward (myvector.begin(), myvector.begin()+5, myvector.end());
        //Exchanges the values of each of the elements in the range [first1,last1] with those of their respective elements in the range beginning at first2.
        std::swap_ranges(foo.begin()+1, foo.end()-1, bar.begin());
         //Applies an operation sequentially to the elements
        bar.resize(foo.size());// allocate space
        std::transform (foo.begin(), foo.end(), bar.begin(), op_increase);
        //Assigns new_value to all the elements in the range [first,last] that compare equal to
              old_value.
        std::replace (myvector.begin(), myvector.end(), 20, 99);
        std::replace_if (myvector.begin(), myvector.end(), IsOdd, 0);
         //Assigns val to all the elements in the range [first,last).
        std::fill (myvector.begin()+3, myvector.end()-2,8);
        std::fill n (myvector.begin()+3,3,33);
        //Assigns the value returned by successive calls to gen to the elements in the range [first,
              last).
        std::generate (myvector.begin(), myvector.end(), RandomNumber);
        std::generate_n (myarray, 9, UniqueNumber);
//Transforms the range [first,last) into a range with all the elements that compare equal to
              val removed,
         // and returns an iterator to the new end of that range.
        pend = std::remove (pbegin, pend, 20);
        pend = std::remove_if (pbegin, pend, IsOdd);
         //Removes all but the first element from every consecutive group of equivalent elements in the
                range [first, last).
        std::unique (myvector.begin(), myvector.end(), myfunction); //first sort
        std::reverse(myvector.begin(), myvector.end());
        // Rotates the order of the elements in the range [first,last), 'in such a way that the element pointed by middle becomes the new first element.std::otate(myvector.begin(), myvector.begin()+3, myvector.end());
         // Rearranges the elements in the range [first,last) randomly.
        std::random_shuffle ( myvector.begin(), myvector.end() );
```

```
//Returns true if all the elements in the range [first,last) for which pred returns true
      precede those for which it returns false.
std::is_partitioned(foo.begin(),foo.end(),IsOdd)
//Rearranges the elements from the range [first,last), in such a way that all the elements
//for which pred returns true precede all those for which it returns false
bound = std::partition (myvector.begin(), myvector.end(), IsOdd); // Return an iterator that
      points to the first element of the second group of elements
std::stable_partition (myvector.begin(), myvector.end(), IsOdd);
//Returns an iterator to the first element in the partitioned range [first,last)
\ensuremath{//} for which pred is not true, indicating its partition point.
std::partition_point(foo.begin(),foo.end(),IsOdd);
std::sort (myvector.begin()+4, myvector.end(), [](int i, int j) {return i < j;});
std::stable_sort (myvector.begin(), myvector.end());
std::is_sorted(foo.begin(),foo.end());
//Returns an iterator to the first element in the range [first,last) which does not follow an
      ascending order.
std::is_sorted_until(foo.begin(),foo.end());
//Rearranges the elements in the range [first,last), in such a way that the element
//at the nth position is the element that would be in that position in a sorted sequence
std::nth_element (myvector.begin(), myvector.begin()+5, myvector.end(), myfunction);
//Returns an iterator pointing to the first element in the range [first, last) which does not
      compare less than val
// 10 10 10 20 20 20 30 30
low =std::lower_bound (v.begin(), v.end(), 20); // 3
/\!/\!Returns \ an \ iterator \ pointing \ to \ the \ first \ element \ in \ the \ range \ [first,last) \ which \ compares
      greater than val.
up = std::upper_bound (v.begin(), v.end(), 20); // 6
//Returns the bounds of the subrange that includes all the elements of the range [first,last)
      with values equivalent to val.
bounds=std::equal_range (v.begin(), v.end(), 20); // return pair
//Returns true if any element in the range [first,last) is equivalent to val, and false
      otherwise.
std::binary_search (v.begin(), v.end(), 3);
//Rearranges the elements in the range [first,last) in such a way that they form a heap.
std::make_heap (v.begin(),v.end());
std::pop_heap (v.begin(),v.end()); v.pop_back();
v.push_back(99); std::push_heap (v.begin(),v.end());
std::sort heap (v.begin(), v.end());
std::is_heap(foo.begin(),foo.end());
//Returns an iterator pointing to the element with the smallest value in the range [first,last
std::min_element(myints,myints+7,myfn);
std::max_element(myints, myints+7, myfn);
std::minmax_element (foo.begin(), foo.end()); // return pair
//Returns true if the range [first1, last1) compares lexicographically less than the range [
      first2.last2).
std::lexicographical_compare(foo, foo+5, bar, bar+9);
//Rearranges the elements in the range [first, last) into the next lexicographically greater
     permutation
        std::cout << myints[0] << ' ' << myints[1] << ' ' << myints[2] << '\n';
} while ( std::next_permutation(myints, myints+3) );
std::prev_permutation(myints, myints+3)
return -1:
```

6.5 C++ std string

```
// C++ program to demonstrate various function string class
#include <string>
#include <regex>
using namespace std;

int main()
{
    // initialization by part of another string : iteartor version
    string str5(str2.begin(), str2.begin() + 5);
    // clear function deletes all character from string
    str4.clear();
    // both size() and length() return length of string and
    // they work as synonyms
    int len = str6.length(); // Same as "len = str6.size();"
    // a particular character can be accessed using at /
    // [] operator
    char ch = str6.at(2); // Same as "ch = str6[2];"
```

```
char ch_f = str6.front(); // Same as "ch_f = str6[0];"
char ch_b = str6.back(); // Same as below
                            // "ch_b = str6[str6.length() - 1];"
cout << "First char is : " << ch_f << ", Last char is : "
     << ch_b << endl;
// c_str returns null terminated char array version of string
const char* charstr = str6.c_str();
printf("%s\n", charstr);
// append add the argument string at the end
str6.append(" extension");
// same as str6 += " extension'
// another version of appends, which appends part of other
str4.append(str6, 0, 6); // at 0th position 6 character
// find returns index where pattern is found.
// Is pattern is not there it returns predefined constant npos
// whose value is -1
if (str6.find(str4) != string::npos)
    cout << "str4 found in str6 at " << str6.find(str4)</pre>
         << " pos" << endl;
else
    cout << "str4 not found in str6" << endl:
// substr(a, b) function returns a substring of b length
// starting from index a
cout << str6.substr(7, 3) << endl;</pre>
// if second argument is not passed, string till end is
// taken as substring
cout << str6.substr(7) << endl;</pre>
// erase(a, b) deletes b character at index a
str6.erase(7, 4);
cout << str6 << endl;
// iterator version of erase
str6.erase(str6.begin() + 5, str6.end() - 3);
cout << str6 << endl:
str6 = "This is a examples";
// replace(a, b, str) replaces b character from a index by str
str6.replace(2, 7, "ese are test");
if (std::regex_match ("subject", std::regex("(sub)(.*)") ))
std::cout << "string literal matched\n";
std::string s ("subject");</pre>
std::regex e ("(sub)(.*)");
std::smatch sm; // same as std::match_results<const char*> cm;
std::regex_match (s,sm,e);
for (unsigned i=0; i<sm.size(); ++i) {
    std::cout << "[" << sm[i] << "] ";</pre>
//Returns whether some sub-sequence in the target sequence (the subject) matches the regular
      expression rgx (the pattern).
std::regex_search (s,m,e)
//Makes a copy of the target sequence (the subject) with all matches of the regular expression rgx
       (the pattern) replaced by fmt (the replacement).
//The target sequence is either s or the character sequence between first and last, depending on
std::cout << std::regex_replace (s,e,"sub-$2");
```

6.6 Compleate Search

return 0:

```
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```

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```
for(int j = 0; j < vec.size(); j++)</pre>
                  if(bitmask & (1<<j))
                           subset.push_back(vec[j]);
         print_vec(subset);
int main() {
         //All permutations
         vector<int> anagrama {1,2,3};
         do {
         print_vec(anagrama);
} while (next_permutation(anagrama.begin(), anagrama.end()));
         //All subsets
         vector<int> set {20, 30, 40, 50, 60};
         unsigned int pow_set_size = 1 << set.size();</pre>
         for(int mask = 0; mask < pow_set_size; mask++)</pre>
                  print_masked(set, mask);
         cout << endl;</pre>
         // \verb|All subsets of given size k|\\
         int k = 4;
         unsigned int mask = (1 << k) - 1;
unsigned int last = ((1 << k) - 1) << (set.size() - k);
         if (mask == 0)
                  print_masked(set, mask);
         else
                  while (mask <= last) {</pre>
                           print_masked(set, mask);
                           unsigned int t = (mask | (mask - 1)) + 1;
                           mask = t \mid ((((t & -t) / (mask & -mask)) >> 1) - 1);
```

6.7 C++ std set

```
#include <iostream>
#include <regex>
using namespace std;
int main()
        std::set<int> myset;
        myset.insert(20);
        //Searches the container for an element equivalent to val and returns an iterator to it if
             found.
        \ensuremath{//} otherwise it returns an iterator to set::end.
        it=myset.find(20);
        //Searches the container for elements equivalent to val and returns the number of matches.
        myset.count(i)
        //Returns an iterator pointing to the first element in the container which is not considered
              to go before val
        itlow=myset.lower_bound (30);
        //Returns an iterator pointing to the first element in the container which is considered to go
               after val.
        itup=myset.upper_bound (60);
        //Removes from the set container either a single element or a range of elements ([first,last))
        myset.erase (myset.find(40));
       myset.erase(itlow,itup);
       //Returns the bounds of a range that includes all the elements in the container that are equivalent to val.
        auto ret = myset.equal_range(30); // pair of iterators
        //For map multiset and multimap all the same
```