

## SP After Mids:

27/11/2023

Background Process Skipped.

Select :-

```
int select ( int maxfd, fd_set *readset,  
            fd_set *writeset, fd_set *errorset,  
            struct timeval *t );
```

```
int main( ) { // infinite delay
    printf("Hello");
    select(1, NULL, NULL, NULL, NULL); // blocks
}

```

Now we create delay using select.

// 5 sec & 10usec delay

Now

```
int main() {
    struct timeval mytime; printf("Hello\n");
    mytime.tv_sec = 5; mytime.tv_usec = 10;
    int f = 0 Select(1, [NULL, NULL, NULL, &mytime]
    printf("wired\n");
```

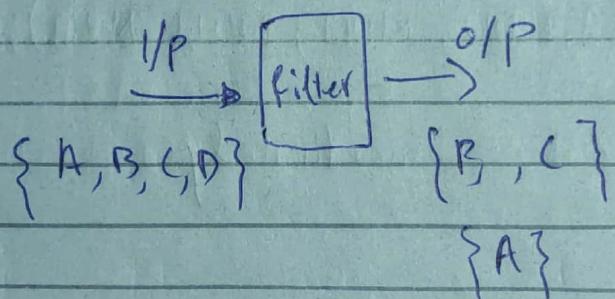
```

int main() {
    ; // monitor 2 files
    ;
    int s = select(maxfd + 1, &readset, NULL,
                   NULL, &mytime);

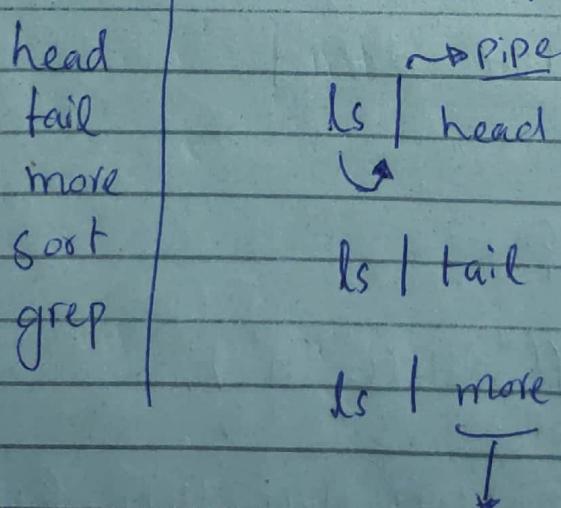
```

printf("Remaining time = %d lt %d usec  
 mytime.tv\_sec, mytime.tv\_nsec);

## Filters & Redirections -



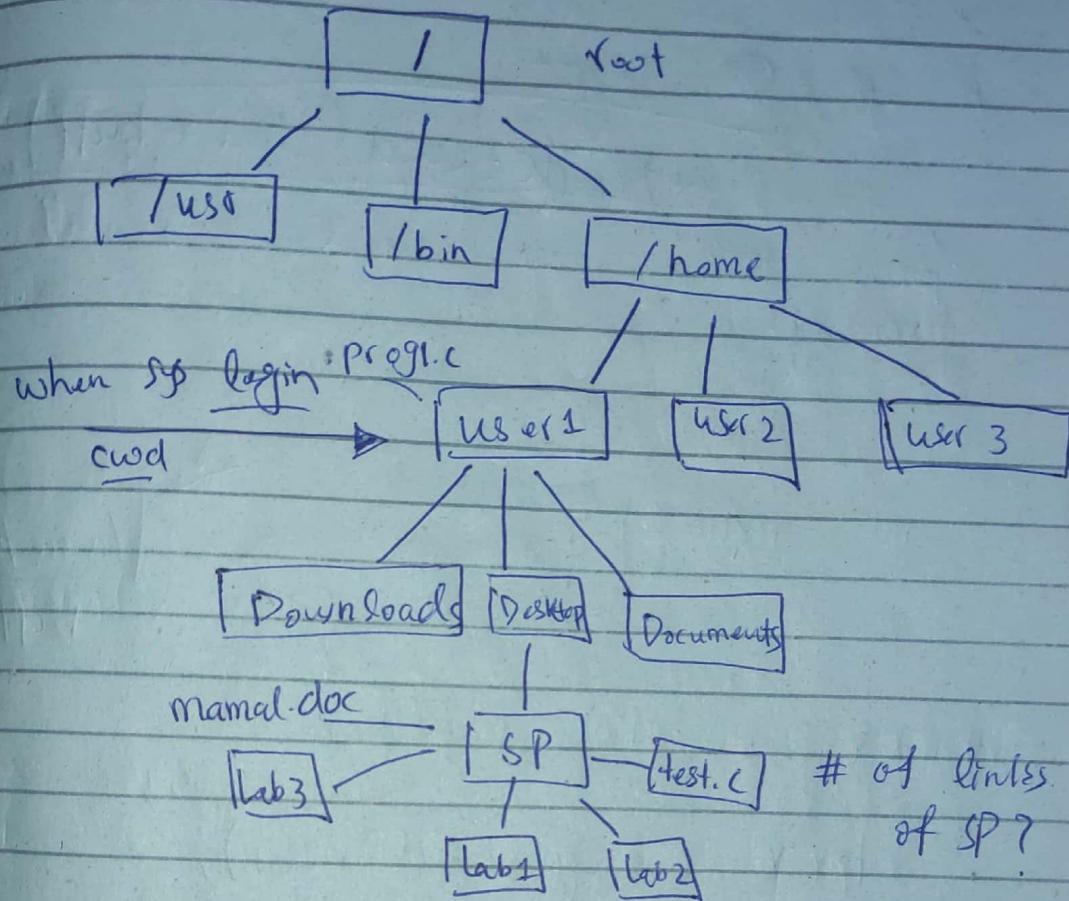
O/P is a subset of I/P



Space bar for more

## Chapter 5: Files & Directories

Directories → hierarchical



cd /home/usr1/Desktop/SP  
cd ./Desktop/SP  
cd .. /usr1/Desktop/SP ] relative  
cd Desktop/SP

cd /bin/.. /home/usr1/Desktop/ ] absolute

Prog1.c

main() {

```
char * val = getenv("PWD");  
printf("my CWD = %s\n", val);
```

const char \*

No need to allocate space for val.

A sys. call exist.

`char * getcwd (char * buff);`

for this we need to allocate memory.

int main() {

    ↓  
    [ /home/ksr ]  
    buff

~~char \*~~ = char buff [255];

    char \* x = getcwd (buff);

} Now we want to change our current dir.

error -1 ← int chdir ( const char \* path );

success 0 →

int main() {

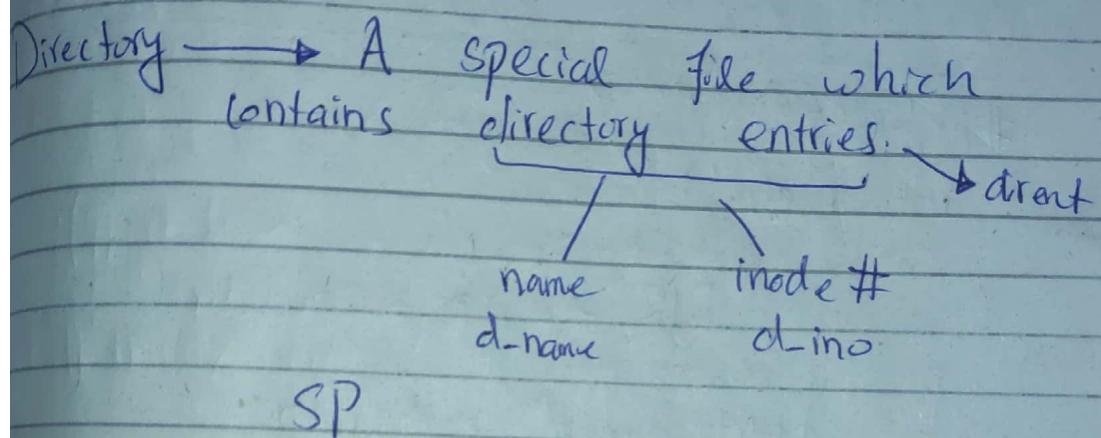
    char buff [255];

    char \* x = getcwd ( buff );  
    printf ("My CWD = %s \n", buff );

Can be absolute/relative Path.

int xv = chdir("Desktop/SP");

④ x = getcwd(buff);  
printf("Path after chdir is: %s\n", buff);



name	inode #
test.l	87
mammal.doc	126
Lab1	500
Lab2	1023
Lab3	1110
..	1234
..	1552
NULL	

We want to open SP directory

ys (all).

DIR \* ~~char~~ opendir (const char \* dirname)

Now we want to read from  
SP

another Sys Call,

~~read~~  
struct dirent \* readdir ( DIR \* dirp );

#include <dirent.h>

int main () {

DIR \* dirp = opendir ("SP");

struct dirent \* direntp;

while (1) {

    direntp = readdir (dirp);

    if (direntp == NULL)

        break;

    printf ("Dirent name = %s it's, dirent  
inode = %d\n", direntp->d\_name,  
direntp->d\_ino);

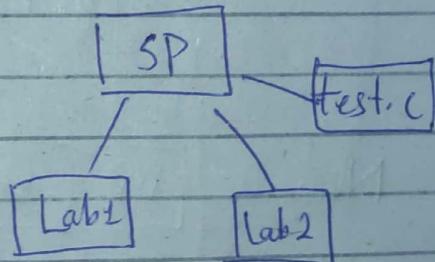
}

|

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Now to get other info,

a sys call exist.



{Syst start.h}

```
int stat( const char* name, struct stat *statbuff);
```

Struct Stat

1) int st_dev	4) int st_Mode	8) time st_ctime
2) int st_ino	5) int st_link	9) time st_mtime
3) int st_nlink	6) int st_uid	10) time st_atime

main() {

DIR \*mydir = opendir (" . " );

struct dirent \*mydirent;

struct stat statbuff;

while ((mydirent = readdir (mydir)) != NULL)

printf (" Name = %s \t ", mydirent->d\_name)

stat (mydirent->d\_name, & statbuff);

printf (" Time = %s \t ", ctime (& statbuff.st\_atime))

printf (" Size = %d \t , gid = %d \t uid = %d  
Links = %d \t ", statbuff.st\_size, statbuff.  
st\_gid, statbuff.st\_uid, statbuff.st\_nlink)

Now

[ int st-mode ]

MACROS

Type of file

S\_ISDIR(m)

S\_ISREG(m)

S\_ISFIFO(m)

S\_ISOCKET(m)

Permissions

S\_IRUSR

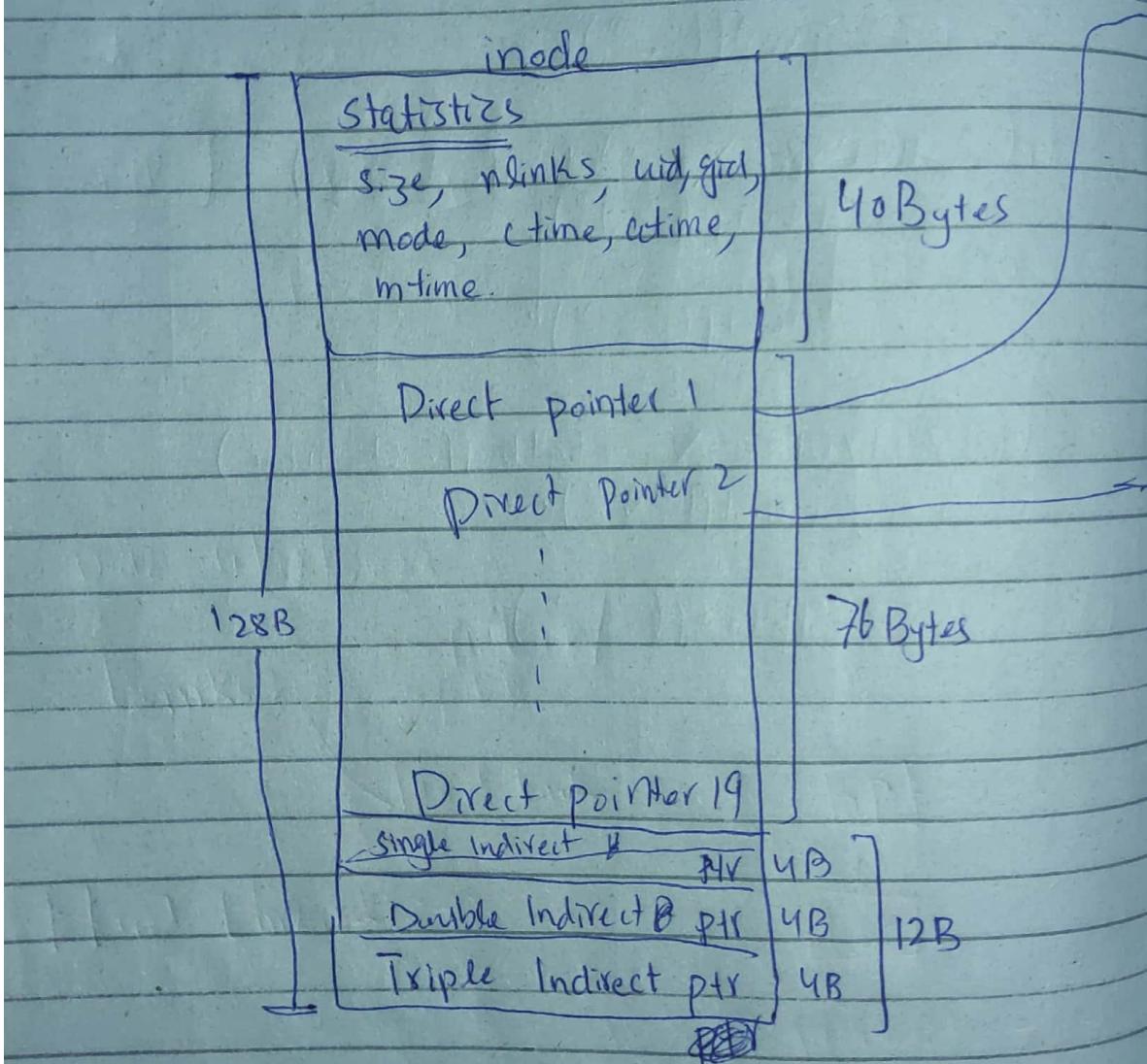
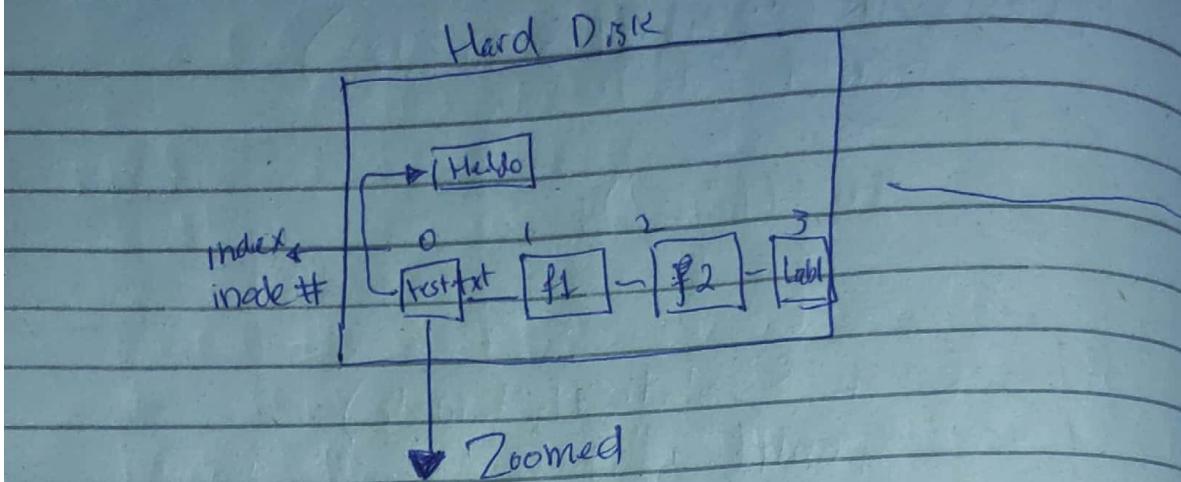
```
→ if (S_ISDIR (statbuff.st_mode))  
    printf ("d");  
else  
    printf ("-");
```

```
if (S_IRUSR & statbuff.st_mode)  
    printf ("r");  
else  
    printf ("-");
```

```
if (S_IWUSR & statbuff.st_mode)  
    printf ("w");  
else  
    printf ("-");
```

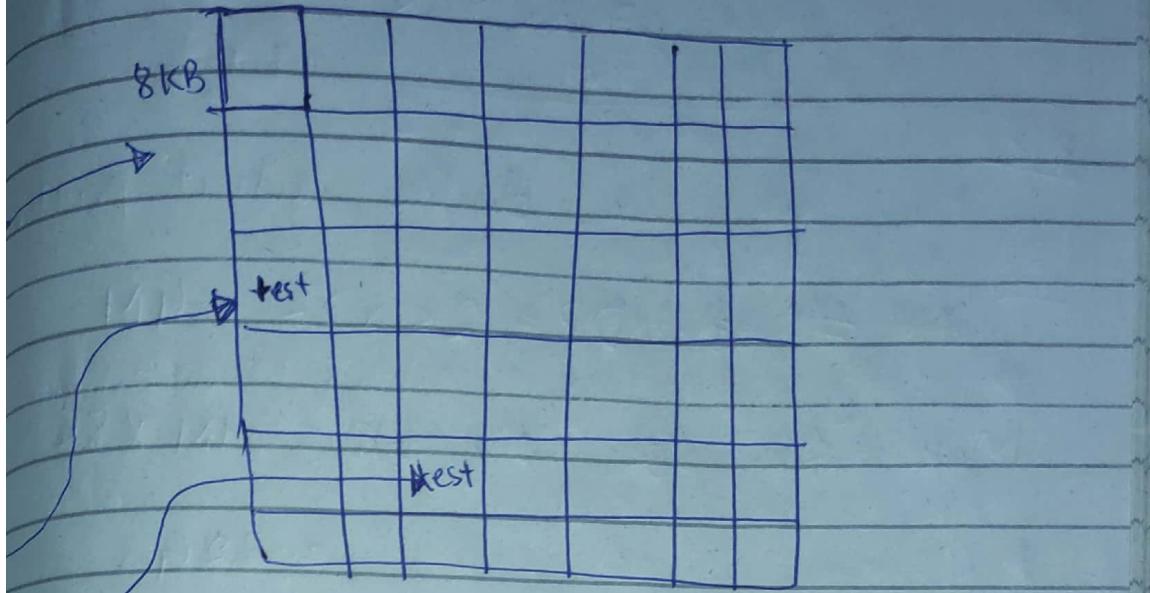
```
if (S_IXUSR & statbuff.st_mode)  
    printf ("x");  
else  
    printf ("-");
```

```
 }  
 }  
 }
```



$$\text{Total memory for DP} = 128 - 40 - 12 = 76$$

Total # of direct pointers =  $\frac{\text{T. available memory}}{\text{size of single}}$



$$= \frac{76B}{4B} = \underline{19 \text{ pointers}}$$

Max file using DPs = # of ptrs & size of  
1 block

$$= 19 \times 8KB = 152KB$$

SID

Total Memory available for DP = 8KB

$$\text{Total \# of DP} = \frac{\text{TAM}}{\text{Size of P}} = \frac{8KB - 2K}{4B} = 2048$$

Max file using SID = # of DPs \* Size of  
block

$$= 2K \times 8K = 16M$$

DIP]

How we use double indirect pointer

$$\text{Total \# of DP} = 2K \times 2K = 4M$$

$$\text{Max file size using SID} = 4M \times 8K$$

$$/\!/ \quad / \quad / = 32G$$

TIP]

$$\text{Total \# of DP} = 2K \times 2K \times 2K = 8G$$

$$\text{Max file size using TIP} = 64TB$$

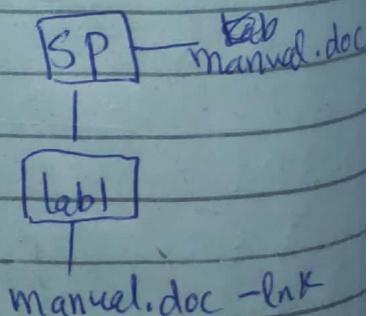
Date 4/12/2023

Links:

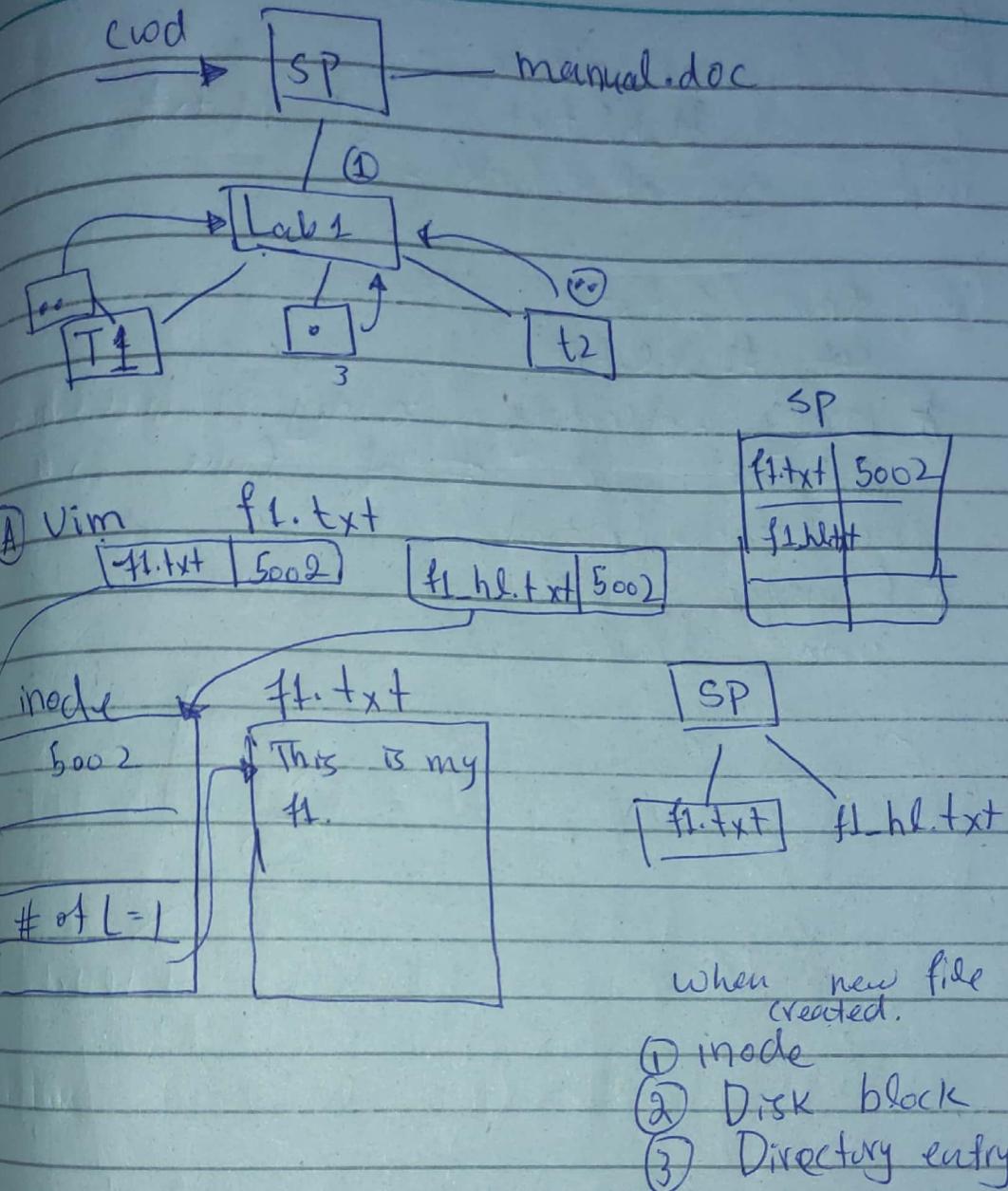
Soft Link /  
Symbolic

Hard Link

Hard Links:-



ls -l  
drwxr-- # links



## B) Link

Command: `ln f1.txt f1-hl.txt`

Sys call: `int link ( const char *src,  
const char *dst);`

Now `ln f1.txt f1-hl.txt`

When we create a hard link only directory entry is created in backend. and the links incremented.

(C)

Now if we write

cat f1\_hl.txt | cat f1.txt

and no. of links = 2

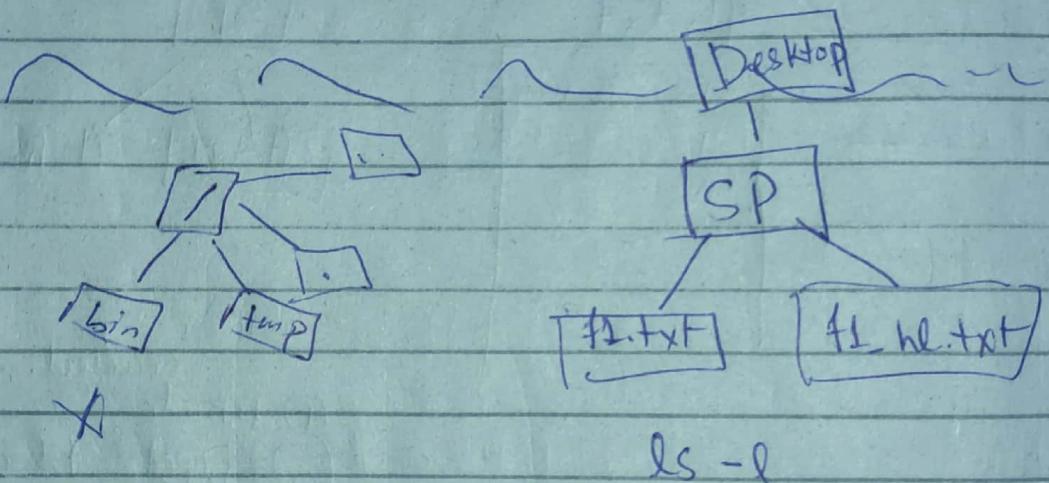
(D)

rm f1.txt

cat f1\_hl.txt

\* only directory entry will be deleted and # of links will be decremented.

\* Changes will be same in both links.



# of Link	Name
2	f1.txt
2	f1_hl.txt

Soft Link :-

Command: ln -s f1.txt f1\_sl.txt

Sys Call : int SymLink (const char \*src, const char \*dst);

① vim f1.txt

② ln -s f1.txt  
f1\_sl.txt

③ cat f1\_sl.txt  
f1.txt

Both show  
Same context.

Now if we update  
the link by

④ vim f1\_sl.txt

⑤ vim f1.txt

⑥ cat f1\_sl.txt

file doesn't  
exist error.

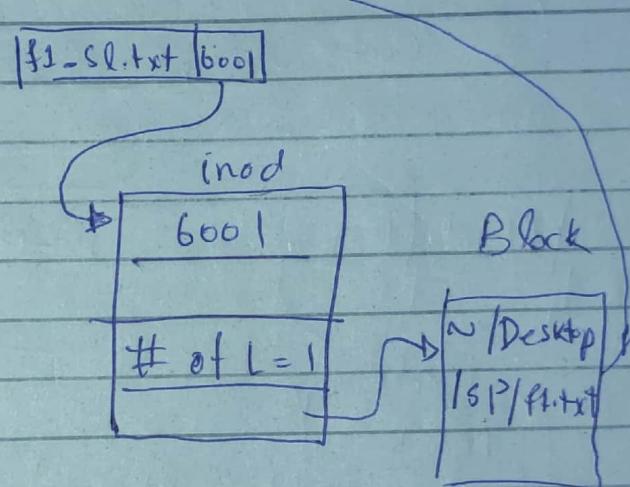
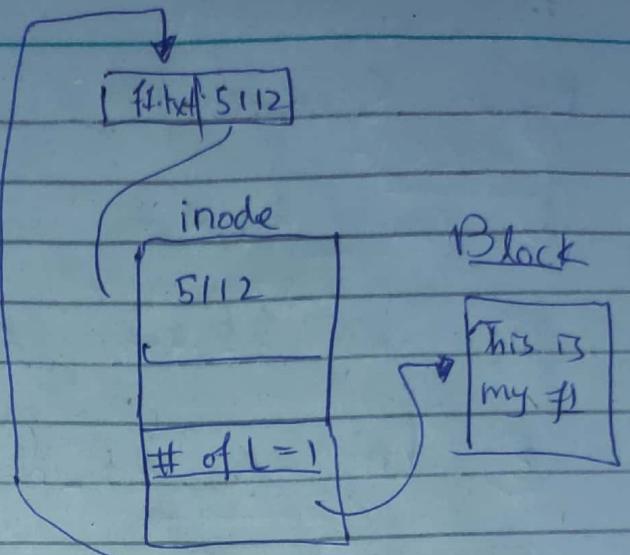
⑦ vim f1.txt

Hello  
World

cat f1\_sl.txt.

will print Hello World.

It means if we create another file  
with same name then we  
can still access it through  
soft link.



Now Scenario:

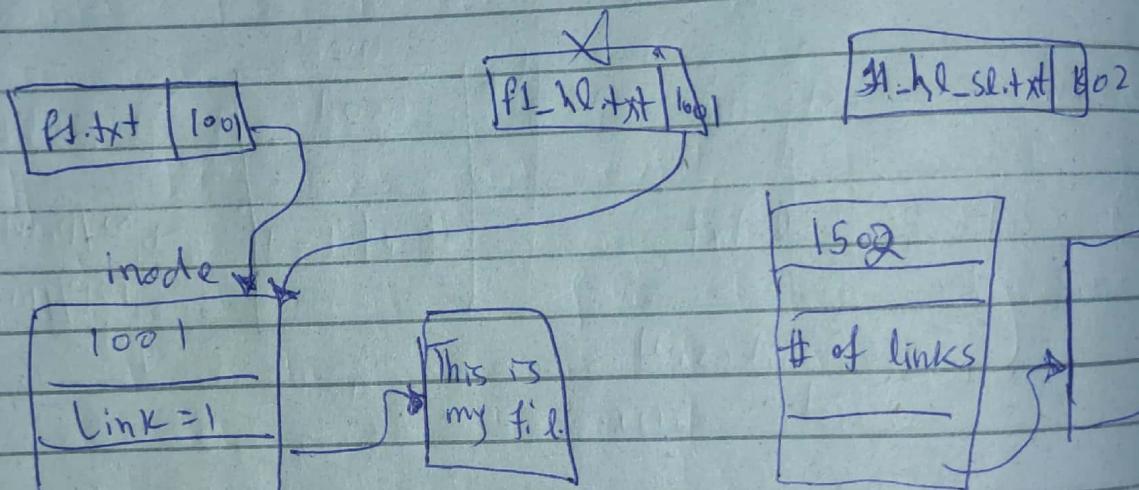
① vim f1.txt

② ln f1.txt f1-hl.txt

③ ln -s f1-hl.txt f1-hl-sl.txt

④ rm f1-hl.txt

⑤ cat f1-hl-sl.txt



⑥ vim f1.txt ("Hello World")

⑦ ln f1.txt f1-hl.txt

⑧ prog.c

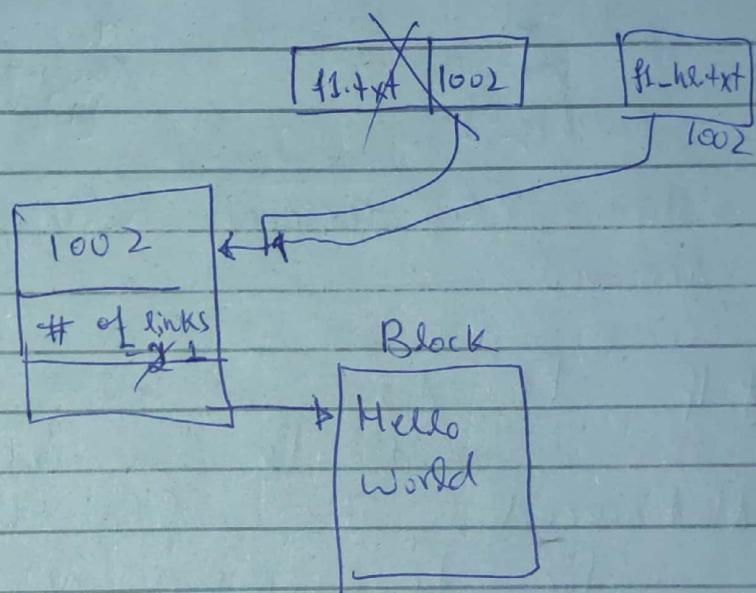
open f1.txt

read f1.txt into buff.

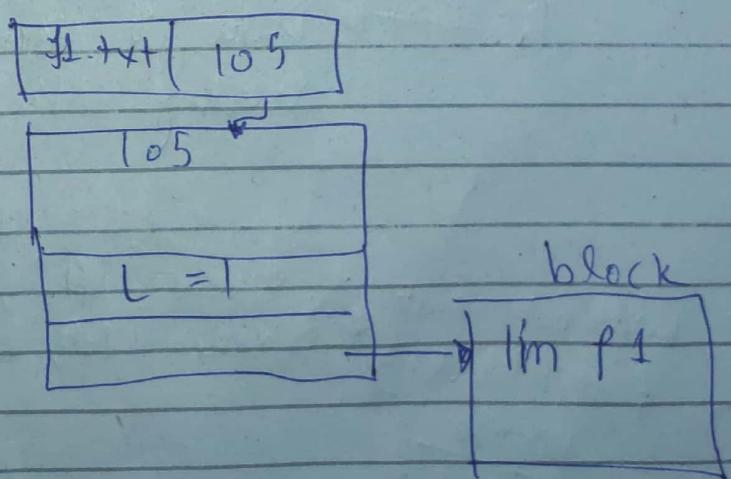
rm f1.txt

```
create f1.txt  
write ("I'm f1");  
close .
```

(4)

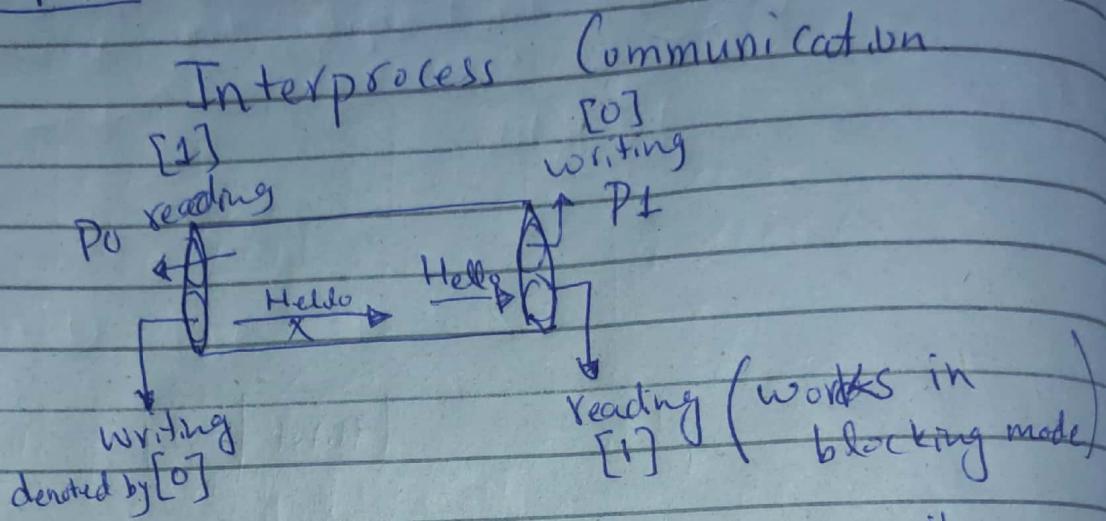


creating f1.txt



Date 7/12/2023

## Pipes:-



When P1 received hello, then it is removed.

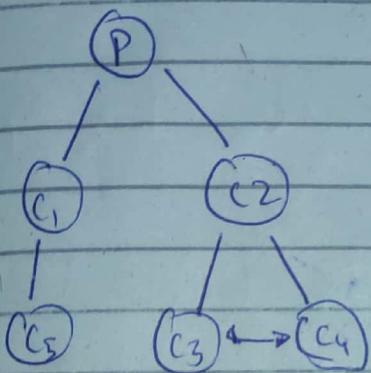
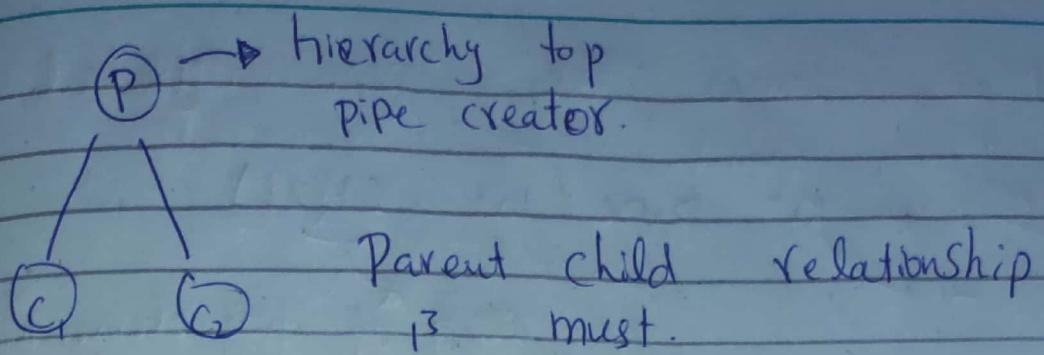
\* If P1 does not exist, then P0 can read its own data.

\* In Backend, it's file (Special file).  
For two reasons it is special

1 → Unnamed

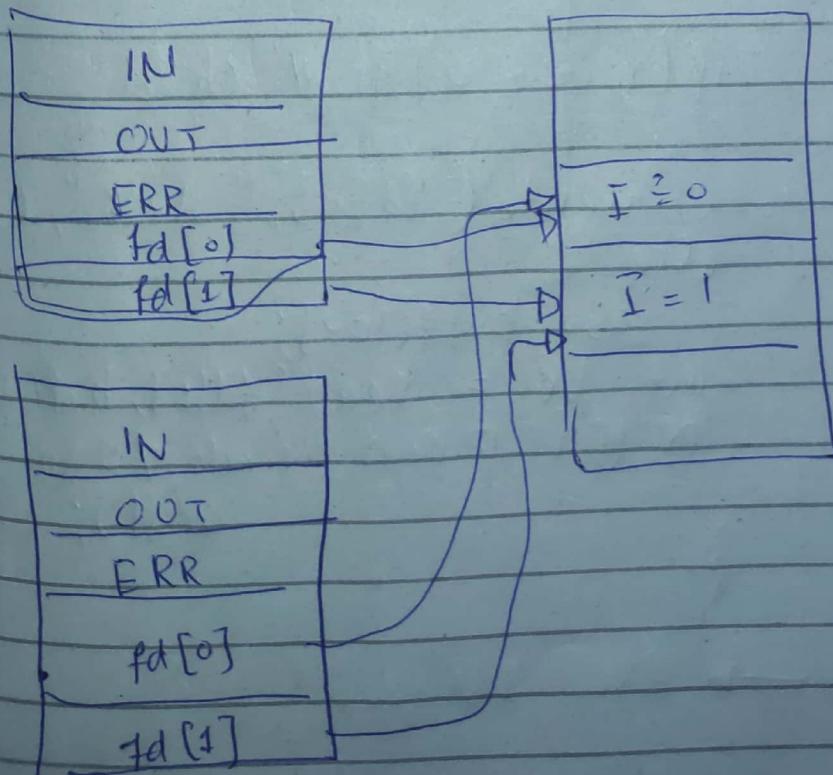
2 → Data is erased after reading.

will not wait if NONBLOCKING mode:  
open ("f1", O\_NONBLOCK)  
read (fd, buff, 100);



Parent

fd table



Now comes the pipe sys call.

success ← int pipe (int fd[2])  
exit -1 ← array of two elements.

Step ① Pipe creation:

```
int main() {
```

```
    int fd[2];
```

```
    int x = pipe(fd);
```

```
    int x = fork();
```

```
    if (x == 0) { // child
```

```
        write(fd[0], "Hello", 6);
```

```
    else {
```

```
        char buff[100];
```

```
        int br = read(fd[1], buff, 100);
        printf("Parent .read: %s\n", buff);
```

```
}
```

Rtp Server → file listing.

Error Case:- File Interrupt  
Resources Exhaust

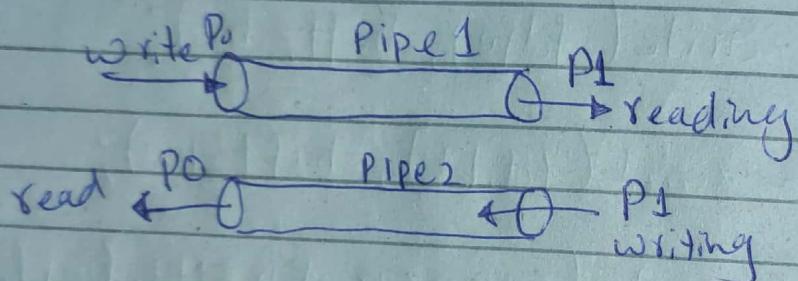
Two way communication:-

```
int main () {  
    int fd[2]; char buff[100];  
    pipe(fd);  
    x = fork();  
    if (x == 0) {  
        br = read (STDIN_FILENO, buff, 100);  
        write (fd[0], buff, br);  
        sleep(1);  
        br = read (fd [1], buff, 100);  
        write (STDOUT_FILENO, buff, br);  
    }  
    else {  
        int br = read (fd[1], buff, 100);  
        write (STDOUT_FILENO, buff, br);  
        br = read (STDIN_FILENO, buff  
        write (fd[0], buff, br);  
    }  
}
```

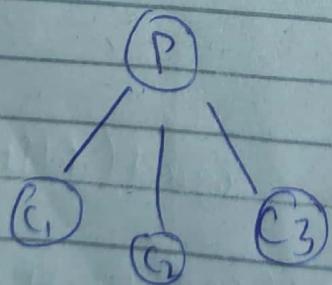
Whenever there is write and then read put sleep(2);

```
write(fd[0])  
sleep(2)  
read(fd[1])
```

Another Solution, make two pipes.



Multiple processes.



```
int main() {
    int fd[2];
    pipe(fd);
    int N = 3;
    for (i = 0; i < 3; i++) {
        x = fork();
        if (x == 0)
            break;
    }
    if (x > 0) {
        for (i = 0; i < N; i++)
            write(fd[i], "Hello", 6);
    } else {
        char buff[100];
        br = read(fd[i], buff, 6);
        printf("child read %s\n", buff);
    }
}
```