Lab 03: Processes in UNIX

3.1. UNIX Process Creation and fork

A process can create a new process by calling fork. The calling process becomes the **parent**, and the created process is called the **child**. The fork function copies the parent's memory image so that the new process receives a copy of the address space of the parent. Both processes continue at the instruction after the fork statement (executing in their respective memory images).

```
SYNOPSIS
#include <unistd.h>
pid t fork(void);
```

Creation of two completely identical processes would not be very useful. The fork function **return value** is the critical characteristic that allows the parent and the child to distinguish themselves and to execute different code. The fork function returns 0 to the child and returns the child's process ID to the parent. When fork fails, it returns -1 and sets the errno. If the system does not have the necessary resources to create the child or if limits on the number of processes would be exceeded, fork sets errno to EAGAIN. In case of a failure, the fork does not create a child.

3.2. The wait Function

When a process creates a child, both parent and child proceed with execution from the point of the fork. The parent can execute wait or waitpid to **block** until the child finishes. The wait function causes the caller to suspend execution until a child's status becomes available or until the caller receives a signal. A process status most commonly becomes available after termination, but it can also be available after the process has been stopped.

The waitpid function allows a parent to wait for a particular child. The waitpid function takes three parameters: a pid, a pointer to a location for returning the status and a flag specifying options. If pid is -1, waitpid waits for any child. If pid is greater than 0, waitpid waits for the specific child whose process ID is pid. Two other possibilities are allowed for the pid parameter. If pid is 0, waitpid waits for any child in the same process group as the caller. Finally, if pid is less than -1, waitpid waits for any child in the process group specified by the absolute value of pid.

```
SYNOPSIS
   #include <sys/wait.h>
   pid_t wait(int *stat_loc);
   pid t waitpid(pid t pid, int *stat loc, int options);
```

If wait or waitpid returns because the status of a child is reported, these functions return the process ID of that child. If an error occurs, these functions return -1 and set errno.

4.2.1. Status values

The stat_loc argument of wait or waitpid is a pointer to an integer variable. If it is not NULL, these functions store the return status of the child in this location. The child returns its status by calling exit, _exit, _exit or return from main. A zero return value indicates EXIT_SUCCESS; any other value indicates EXIT_FAILURE. The parent can only access the 8 least significant bits of the child's return status.

POSIX specifies six macros for testing the child's return status. Each takes the status value returned by a child to wait or waitpid as a parameter.

```
SYNOPSIS

#include <sys/wait.h>
WIFEXITED(int stat_val)
WEXITSTATUS(int stat_val)
WIFSIGNALED(int stat_val)
WTERMSIG(int stat_val)
WIFSTOPPED(int stat_val)
WSTOPSIG(int stat_val)
```

The six macros are designed to be used in pairs. The WIFEXITED evaluates to a nonzero value when the child terminates normally. If WIFEXITED evaluates to a nonzero value, then WEXITSTATUS evaluates to the low-order 8 bits returned by the child through _exit(), exit() or return from main.

The WIFSIGNALED evaluates to a nonzero value when the child terminates because of an uncaught signal. If WIFSIGNALED evaluates to a nonzero value, then WTERMSIG evaluates to the number of the signal that caused the termination.

The WIFSTOPPED evaluates to a nonzero value if a child is currently stopped. If WIFSTOPPED evaluates to a nonzero value, then WSTOPSIG evaluates to the number of the signal that caused the child process to stop.

The following function determines the exit status of a child.

Task 1: Create process chain as shown in figure 3.1(b) and fill the figure 3.1 (b) with actual IDs. The program shall take a single command-line argument that specifies the number of processes to be created. Before exiting, each process shall output its i value (loop variable), its process ID (using getpid()), its parent process ID (getppid()) and the process ID of its child (return value of fork). The parent does not execute wait.

Task 2: Create process fan as shown in figure 3.1 (a) and fill the figure 3.1 (a) with actual IDs.

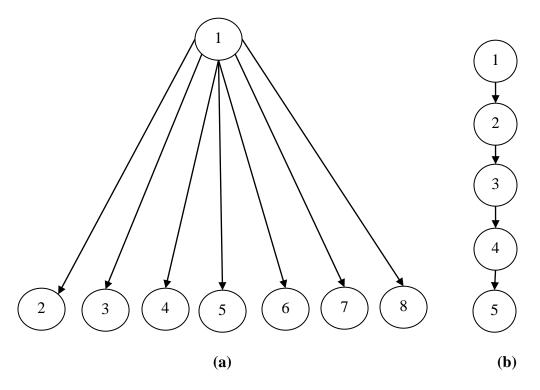


Figure 3.1 Multiple Processes (a) Process Fan (b) Process Chain

Task 3: Create process tree as shown in figure 3.2 and fill figure 3.2 with actual IDs.

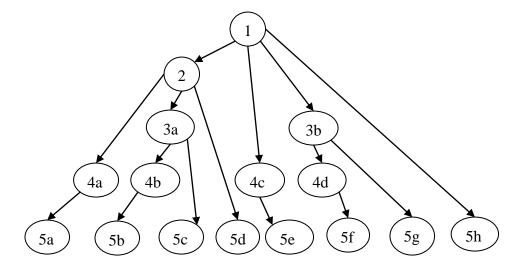


Figure 3.2 Multiple Processes: Process Tree

Task 4: creates a chain of processes. It takes a single command-line argument that specifies the number of processes to create. Before exiting, each process outputs its i value, its process ID, its parent process ID, and the process ID of its child. The parent does not execute wait. If the parent exits before the child, the child becomes an **orphan**. In this case, the child process is adopted by a special system process (which traditionally is a process, init, with process ID of 1). As a result, some of the processes may indicate a parent process ID of 1.

Task 5: Write a program that takes N number of integers as argument and displays the factorials of N integers (print 1 only if integers are not less than zero, 0 or 1). Create separate child process for each integer. Make sure no child is orphan/zombie.

Task 6: Write a program that creates an array of size 100. Initialize the array with random numbers. Create 10 child processes divide the array between them. Each child will add the portion and return their sum to parent process. Parent will add the results and display a final sum.

