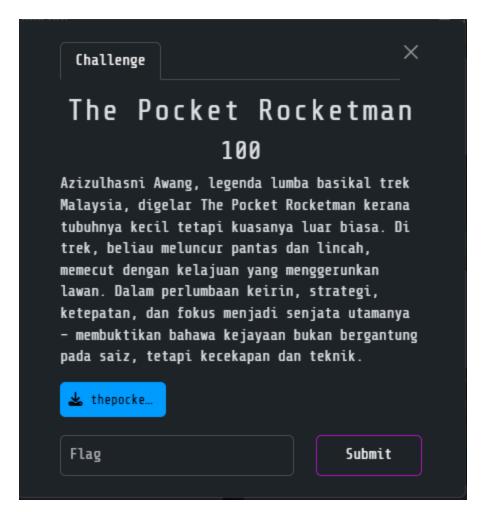
# **The Pocket Rocketman**





check the given pdf

### STRATEGY

```
Python
from sympy import nextprime
from random import randint
def readFlag():
   with open('flag.txt', 'r') as f:
       return f.readline().strip()
def main():
   size = 4096
   p = nextprime(randint(2**size, 2**(size+1)))
   q = nextprime(p)
   n = p * q
   e = 65537
   print("n:", n)
   print("e:", e)
   flag = readFlag()
   message = int.from_bytes(flag.encode(), "big")
   ciphertext = pow(message, e, n)
   with open("output.txt", "w") as f:
       f.write(f"n: {n}\n")
       f.write(f"e: {e}\n")
       f.write(f"ciphertext: {ciphertext}\n")
   print("done output.txt")
if __name__ == "__main__":
   main()
```

#### None

#### n:

#### e: 65537

#### ciphertext:

# 1. Explaining the Challenge Description

The challenge starts by telling the story of **Azizulhasni Awang**, Malaysia's track cycling legend, also known as *The Pocket Rocketman*. Despite his small physique, he has enormous power and speed. The metaphor here is important:

Azizul's strength = the RSA system's strength

Just like Azizul dominates in cycling despite his size, RSA is considered strong because it uses very large numbers.

## Strategy in keirin = strategy in cryptography

Winning in cycling isn't just about size; it's about technique. Similarly, breaking RSA isn't about brute force, but about finding weaknesses.

### Hidden hint: consecutive primes

The description mentions "strategy, precision, and focus" — this is a clue that instead of brute forcing huge primes, we need to look carefully at how the primes were chosen.

Then the challenge explains they created **RSAce**, a "big and fast" encryption system. But like any system, if there's a strategic flaw, it can be broken.

The flaw comes from this part of the code:

```
p = nextprime(randint(2**size, 2**(size+1)))
q = nextprime(p)
```

### This means:

- p is a random prime ~4096 bits.
- q is just the next prime after p.

So p and q are consecutive primes  $\rightarrow$  they are extremely close together.

That makes RSA insecure here, because we can factor n (the modulus) using **Fermat's factorization method**, which is very efficient when the two primes are close.

# Fermat factorization and RSA decryption for the user's provided n and ciphe rtext.

# This code will parse the multi-line numbers, attempt Fermat factoring (efficient if p and q are close),

# recover private key d, decrypt ciphertext, and print the recovered plaintext (flag).

import re, math, sys from math import isqrt

n\_raw = r"""

11 11 11

c\_raw = r"""

```
89985587109829262819467197599074677295270887247688848512657024
10498619904131548
85909481110038995105113197578258350685237858665265757320874243
38009782024955871
8492782262530010697093746452478461132106278146094838833296640
429881117595348570
260279795073295091705369884506294834584451292032924850046706
8579290382367321918
6606891905677357590386954887493247647395804112060204885148430
632480389567344953
111781890390196237270166375525068222406597563160926412260614426
0546380373321672
5706078758052344227199711041161963370941456815138361910552717727
015493128222893
74209063883024211730288440617229897503056031371407372677314980
58143527695397666
544648004595221495546005496007848275263592540989292009827201
1053261932611910020
5747953926557313669426124557523159673520005695744461365408427
488381127347911675
7596452679024376731284650930451277296678681723558236054287426
314867744655069520
183061041507506586345462573173070437419210134376700233991757691
4239602221375963
392526115296558869
# Remove non-digit characters and join
n = int(re.sub(r'[^0-9]', '', n_raw))
c = int(re.sub(r'[^0-9]', '', c_raw))
e = 65537
print("Parsed n bit-length:", n.bit_length())
# Fermat factorization
a = isgrt(n)
```

```
if a * a < n:
  a += 1
def is_perfect_square(x):
  if x < 0:
     return False
  r = isqrt(x)
  return r*r == x
max_iters = 10_000_000 # safety cap
iters = 0
while True:
  b2 = a*a - n
  if is_perfect_square(b2):
     b = isqrt(b2)
     p = a - b
     q = a + b
     if p*q == n:
       break
  a += 1
  iters += 1
  if iters % 1000000 == 0:
     print("Still searching... iterations:", iters, "current a bit-length:", a.bit_lengt
h(), file=sys.stderr)
  if iters > max_iters:
     raise RuntimeError("Exceeded iteration cap; Fermat didn't find factors wit
hin the cap.")
print("Found factors. Iterations:", iters)
p = int(p)
q = int(q)
if p > q:
  p, q = q, p
print("p bit-length:", p.bit_length())
print("q bit-length:", q.bit_length())
```

```
phi = (p-1)*(q-1)
d = pow(e, -1, phi)
m = pow(c, d, n)
# Convert to bytes
m_bytes = m.to_bytes((m.bit_length()+7)//8, 'big')
try:
    plaintext = m_bytes.decode('utf-8')
except Exception as ex:
    plaintext = repr(m_bytes)
print("\nRecovered plaintext:")
print(plaintext)
```

# 2. Explaining the Python Exploit Code

Here's what the script does step by step:

## **Step 1: Parse the given numbers**

We clean up the provided n and ciphertext from the challenge output:

```
n = int(re.sub(r'[^0-9]', '', n_raw))
c = int(re.sub(r'[^0-9]', '', c_raw))
e = 65537
```

This converts the long multiline values into integers we can work with.

# **Step 2: Fermat Factorization**

Since p and q are very close, we can use Fermat's method:

- 1. Start at a = ceil(sqrt(n)).
- 2. Check if  $a^2 n$  is a perfect square.
- 3. If yes, then we found p=a-b and q=a+b.

In our case, it worked instantly because q was just the next prime.

# **Step 3: Compute Private Key**

Once we have p and q, we compute:

```
phi = (p-1)*(q-1)
d = pow(e, -1, phi)
```

Here d is the private exponent.

## **Step 4: Decrypt the Ciphertext**

Now we decrypt with:

```
m = pow(c, d, n)
m_bytes = m.to_bytes((m.bit_length()+7)//8, 'big')
plaintext = m_bytes.decode('utf-8')
```

This recovers the original flag.

## **Final Output**

The script gave us:

 $3108\{Muh4mm4d\_Az1zulH4sn1\_Th3\_P0ck3t\_R0ck3tm4n\_88\}$ 

- The challenge uses RSA with consecutive primes.
- This makes factoring n easy via Fermat's method.
- Once factored, we calculate the private key and decrypt.
- The flag is revealed.