CS303 Data Structures Assignment #1

attachments and source available at
https://github.com/alexskc/cs303

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1

Choose an arbitrary n_0, eg. $n_0 = 1$.

$Cn_0^3 > n_0^3 - 5n_0^2 + 20n_0 - 10$

$C * 1^3 > 1^3 - 5 * 1^2 + 20 * 1 - 10$

$C > 1 - 5 + 20 - 10$

$C = 6, n_0 = 1$

For all n, where $n > 1$, $6n^3 > n^3 - 5n^2 + 20n - 10$.

2

See attached `comparegrowth.cpp`. Output:

```
y1: 10
y2: 2
y1: 1010
y2: 502
y1: 2010
y2: 2002
y1: 3010
y2: 4502
y1: 4010
y2: 8002
y1: 5010
```

```
y2:  12502
y1:  6010
y2:  18002
y1:  7010
y2:  24502
y1:  8010
y2:  32002
y1:  9010
y2:  40502
y1:  10010
y2:  50002
```

The results here are to be expected. y1 is initially larger because of the constant 100 rather than 5, but y2 quickly overtakes it because of how much faster n^2 grows than n. This would be true regardless of what constant y1 used. If $y1 = 10000n + 20$, y2 would still outgrow it.

3

3.1

The inner loop is run i^2 times, with i being every number from 0 to n-1. Therefore, we get the sum: $T(n) = 1^2 + 2^2 + 3^2 \cdots + n^2$ or
$T(n) = \sum_{i=1}^{n} i^2$, which is simply a geometric series, so $T(n) = \frac{1}{6}n(n+1)(2n+1)$, or $O(n) = n^3$.

3.2

This is the same as a simple "i to n" loop, but backwards, which makes no difference, and skipping every other number, which halves the iterations. For loops with odd numbers, we get to 1, and then decrement again, so we take the ceiling of that fraction.
$T(n) = \lceil \frac{1}{2}n \rceil$, and $O(n) = n$.

3.3

This is similar to **??**, except instead of subcracting, we're dividing. So, instead of dividing i, we use the logarithm. We have to add one, because we continue to divide after getting to 2, and we take the floor because numbers that aren't powers of two will fall just faster than the power of two above them.
$T(n) = \lfloor \log 1 + 1 \rfloor + \lfloor \log 2 + 1 \rfloor + \lfloor \log 3 + 1 \rfloor \dots \lfloor \log n + 1 \rfloor.$
$T(n) = \sum_{i=1}^{n} \lfloor \log i + 1 \rfloor$
This can be imagined as $1+2+2+3+3+3+3+4+4+4+4+4+4+4+4+5\dots$, where the number jumps up every time $\log i$ returns a whole number. Every power of two. So $T(n) = \sum_{i=1}^{n} i*2^{i-1}$ and $T(n) = n2^n - 2^n + 1$. $O(n) = n2^n$.

4

pop_back() removes the last element in the vector. So we'd remove the element `5` at location `4` in that vector, resulting in an array like this:

0	1	2	3	4	5	6	7	8	9
1	2	3	4	null	null	null	null	null	null

with the_data pointing to that array for both v1 and v2.