

# CPSC 471: Computer Communications

## Network Performance

Figures from [Computer Networks: A Systems Approach](#), version 6.02dev  
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# Performance:

## Bandwidth/Throughput

- ⦿ Also called data rate
- ⦿ Number of bits that can be transmitted over the network in a certain period of time
  - E.g., 10 Mbps (Megabits per second or million bits per second)
    - Takes  $0.1 \mu\text{s}$  (microsecond) to transmit each bit
- ⦿ Mbits  $\neq$  MB
  - Mbits =  $10^6$  bits per second
  - MB (megabyte) =  $2^{20}$  bits

# Pulse Width of Bits

- ⦿ (a) Bits transmitted at 1 Mbps
  - Each bit is  $1\ \mu\text{s}$  wide
- ⦿ (b) Bits transmitted at 2 Mbps
  - Each bit is  $0.5\ \mu\text{s}$  wide

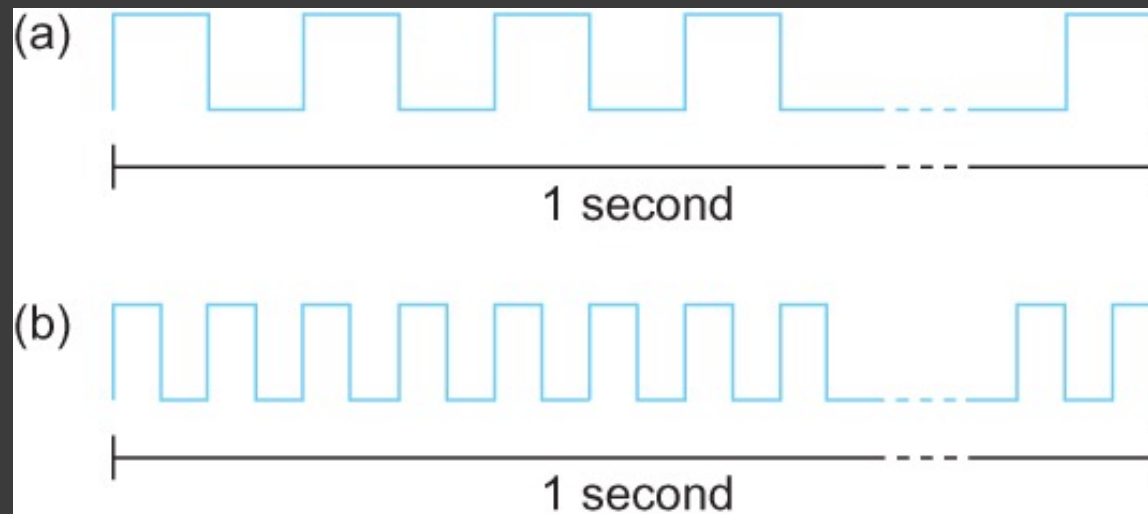


Figure 16

# Bandwidth continued

- ⦿ Analyze bandwidth of
  - Network as a whole
  - Single physical link
  - Logical process-to-process channel
    - Influenced by how often software that handles the channel must handle/transform the data

# Bandwidth Definitions

- ⦿ Bandwidth can also refer to range of signals that can be accommodated
  - Expressed in frequency
- ⦿ **Throughput:** the measured performance of a system
  - E.g., A link with 10 Mbps may have a lower actual throughput
- ⦿ Bandwidth Requirements
  - The number of bits per second an application needs to perform acceptably

# Performance: Latency/Delay

- ⦿ How long it takes a message to travel from source to destination
- ⦿ Latency of
  - Single link
  - End-to-end channel
- ⦿ Measured using time
  - Transcontinental network may have latency of 24 ms (one-way)
- ⦿ Round-Trip Time (RTT) is often more useful

# Components of Latency

- ◎ Speed of light propagation delay
  - Propagation delay: time for signal to propagate from one end of link to other
  - $3 \times 10^8$  m/s in vacuum
  - $2.3 \times 10^8$  m/s in a copper cable
  - $2 \times 10^8$  m/s in an optical fiber
  - $\text{Propagation} = \text{Distance} / \text{Speed\_of\_Light}$

# Components of Latency

## ⦿ Transmit Delay

- Amount of time it takes to transmit a unit of data
- $\text{Transmit} = \text{Packet\_Size} / \text{Bandwidth}$

## ⦿ Queuing Delays

- Switches usually need to store packets for some time before forwarding them on

## ⦿ $\text{Latency} = \text{Propagation} + \text{Transmit} + \text{Queue}$



# Length of Bit Example

- ⦿ How wide is a bit on a 100-Mbps link?
- ⦿ Assume this link is a copper wire with a propagation speed of  $2.3 \times 10^8$  m/s

# Solution to Length of Bit

## Example

- ① 100 Mbps link =  $100 \times 10^6$  bps =  $10^8$  bits/s
- ① Therefore,  $10^{-8}$  s/bit:
  - Each bit is  $10^{-8}$  s or 10 ns wide
- ① If copper wire has a propagation speed of  $2.3 \times 10^8$  m/s:
  - $10^{-8}$  s/bit \*  $2.3 \times 10^8$  m/s = 2.3 m/bit
- ① Each bit is 2.3 m long in this wire

# Example Latency Problem

- Assume that hosts A and B are connected by a switch S:      A-----S-----B
- Assume that the links have a bandwidth of 1 Gbps and a propagation delay of  $15\ \mu\text{s}$
- Assume that the switch is a store-and-forward device and forwards the received packet  $25\ \mu\text{s}$  after receiving it
- Create a timeline of sending two consecutive 100,000-bit packets

# Solution to Example Latency Problem (1/2)

- First, calculate the transmission delay:  
 $100,000 \text{ bits} / 1 \text{ Gbps} = 100 \mu\text{s}$
- At  $t=0 \mu\text{s}$ : A starts transmitting Packet 1 to S
- At  $t=15 \mu\text{s}$ : 1<sup>st</sup> bit of Pkt 1 arrives at S
- At  $t=100 \mu\text{s}$ : A finishes transmitting Pkt 1 to S  
A starts transmitting Packet 2 to S
- At  $t=115 \mu\text{s}$ : Last bit of Pkt 1 arrives at S  
1<sup>st</sup> bit of Pkt 2 arrives at S
- At  $t=140 \mu\text{s}$ : S starts transmitting Pkt 1 to B

# Solution to Example Latency Problem (2/2)

- At  $t=200\ \mu\text{s}$ : A finishes transmitting Pkt2 to S
- At  $t=215\ \mu\text{s}$ : Last bit of Pkt 2 arrives at S
- At  $t=240\ \mu\text{s}$ : 1<sup>st</sup> bit of Pkt 1 arrives at B  
S finishes transmitting Pkt 1 to B  
S starts transmitting Pkt2 to B
- At  $t=255\ \mu\text{s}$ : Last bit of Pkt 1 arrives at B  
1<sup>st</sup> bit of Pkt 2 arrives at B
- At  $t=340\ \mu\text{s}$ : S finishes transmitting Pkt 2 to B
- At  $t=355\ \mu\text{s}$ : Last bit of Pkt 2 arrives at B

# Bandwidth and Latency Impact on Applications

- Latency dominates for small messages
- Bandwidth dominates for large messages
- Example: Consider message size for different applications

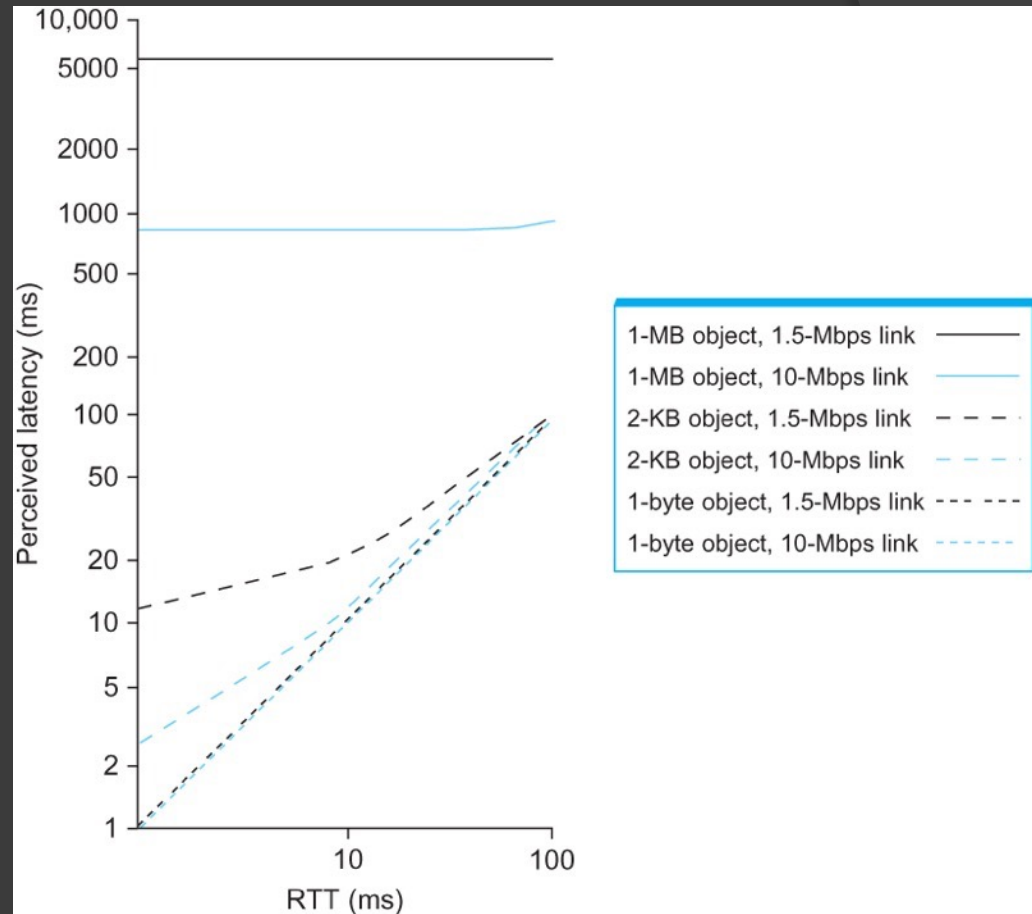


Figure 17

# Delay x Bandwidth Product

- Number of bits that can be in transit (in flight) at a given instant
- How many bits the sender must transmit before the first bit arrives at the receiver
- Usually uses the RTT delay



Figure 18

# Delay x Bandwidth Products

Link Type	Bandwidth	One-Way Distance	RTT	RTT x Bandwidth
Wireless LAN	54 Mbps	50 m	0.33 $\mu$ s	18 bits
Satellite	1 Gbps	35,000 km	230 ms	230 Mb
Cross-country fiber	10 Gbps	4,000 km	40 ms	400 Mb

Table 1



# Example Performance Problem

- Suppose a 128-kbps point-to-point link is set up between the Earth and a rover on Mars. Assume the distance between Earth and Mars is 55 million kilometers and data travels over the link at the speed of light ( $3 \times 10^8$  m/s).
  - a. First, how quickly can a 5 MB image taken by the rover be transmitted to Earth?
  - b. Next, calculate the delay-bandwidth product for this link. Use the RTT of part a's delay.

# Solution to Example

## Performance Problem (part a)

- a. One-way latency = Propagation Delay + Transmit Delay + Queuing Delay
- Assuming Queuing Delay = 0 seconds
  - Prop. Delay =  $55 \times 10^9 \text{ m} / 3 \times 10^8 \text{ m/s}$   
Propagation Delay = 183.33 seconds
  - Image Size = 5 MB =  $5 * 2^{20} \text{ B} = 5 * 1,048,576 \text{ B}$   
Image Size =  $5,242,880 \text{ B} * 8 \text{ bits/Byte}$   
= 41,943,040 b or approx. 41.9 Mb
  - Transmit Delay =  $41,943,040 \text{ b} / 128 \times 10^3 \text{ b/s}$   
= 327.68 s
  - One-way latency =  $183.33 + 327.68 = 511.01 \text{ s}$   
or approx. 8.52 minutes

# Solution to Example Performance Problem (part b)

## b. Delay-Bandwidth Product

- $RTT = 2 * \text{One-way delay} = 2 * 511.01 \text{ s}$   
 $= 1022.02 \text{ s}$  or approx. 17.03 minutes
- $RTT * \text{Bandwidth} = 1022.02 \text{ s} * 128 \times 10^3 \text{ b/s}$   
 $= 130,818,560 \text{ b}$  or approx. 130.8 Mb  
 $= 16,352,320 \text{ B}$  or approx. 15.59 MB

# High-Speed Networks

- Bandwidth can increase for high-speed networks, but usually latency stays constant
  - Latency starts to dominate

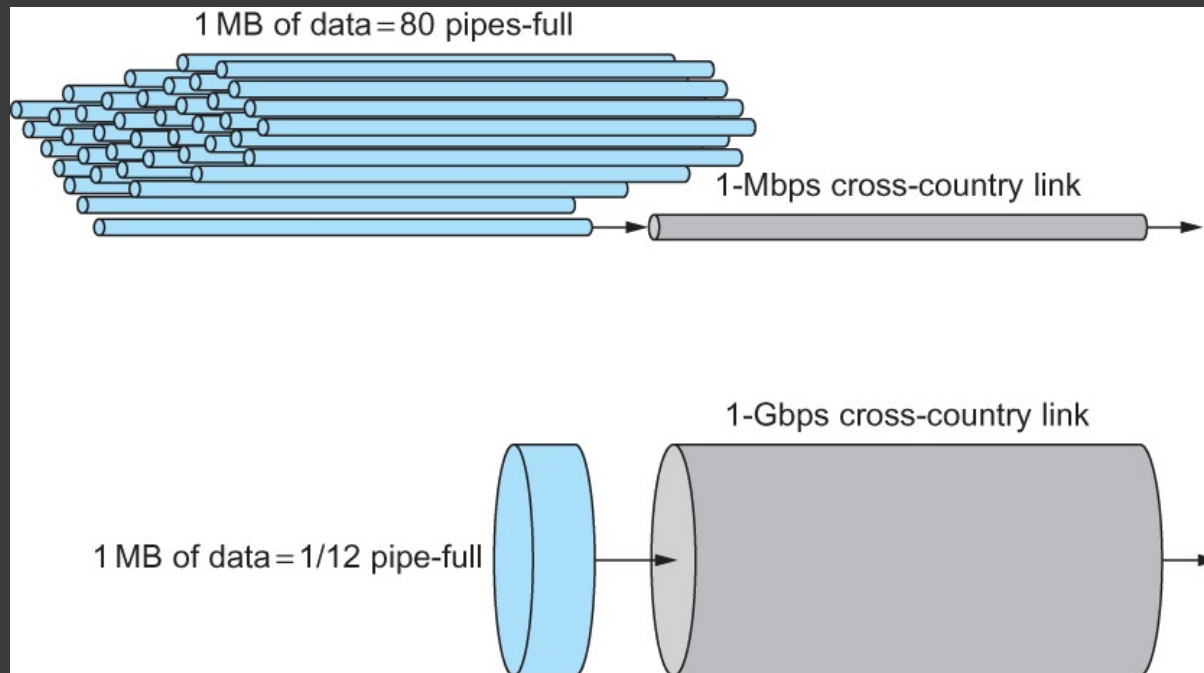


Figure 19

# Effective Throughput

- Effective/Measured Throughput (end-to-end)  
$$= \text{Transfer\_Size} / \text{Transfer\_Time}$$

- Transfer\_Time =  
$$\text{RTT} + 1/\text{Bandwidth} \times \text{Transfer\_Size}$$

# Application Performance Needs

## ⦿ E.g., Video Applications

- Have a resolution
  - E.g.,  $\frac{1}{4}$  of standard def. (352 x 240 pixels)
- Have a color depth
  - E.g., 24-bit color
- Requires 247.5 KB per frame
- Have a frame rate
  - If frame rate is 30 fps, may need throughput rate of 75 Mbps

# Average Bandwidth

- ⦿ Can be misleading
- ⦿ Example: An application has average bandwidth of 2 Mbps and the maximum channel bandwidth is 2 Mbps
  - What happens if application's actual bandwidth varies between 1 and 3 Mbps?

# Peak Rate

- ⦿ Need to know the peak rate or burst
- ⦿ Can allocate storage/buffer capacity to handle a burst
  - Depends how large and how frequent the bursts are
- ⦿ Need to consider queueing theory



# Jitter

- ⦿ Variation in latency
- ⦿ A problem for video applications
- ⦿ If spacing between packets (interpacket gap) varies, jitter is introduced into packet stream
  - Can happen due to different queuing delays



Figure 20