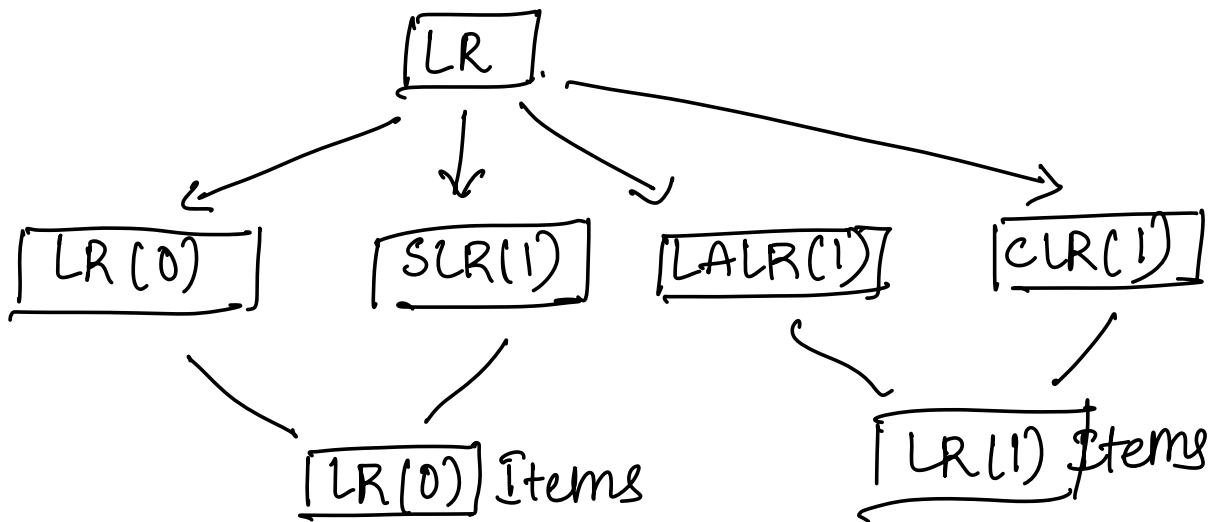
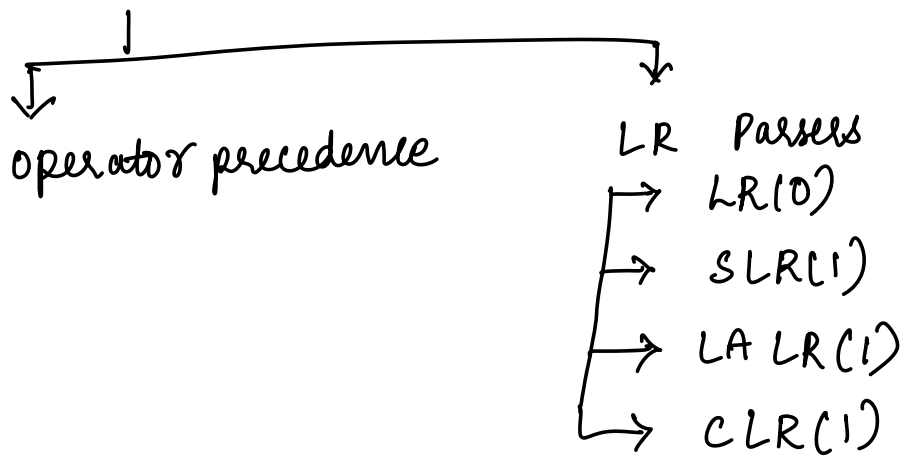
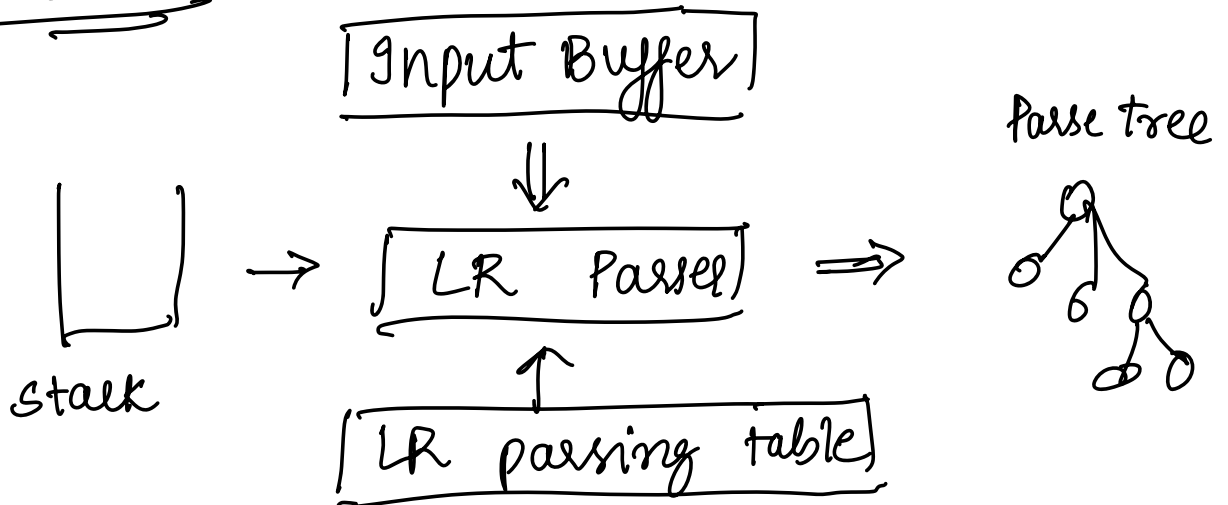


Bottom-up parsers:



Architecture:



SLR(1) : Simple LR

LALR(1) : Look Ahead LR

CLR(1) : Canonical LR

① LR(0) Parser:

Steps:

1. Augment the given grammar
2. Draw the canonical collection of LR(0) items.
3. Number the productions.
4. Create the parsing Table
5. Stack Implementation
6. Build the parse tree.

Example:

Given grammar is

$E \rightarrow BB$

$B \rightarrow cB/d$

Input string is

"ccdd"

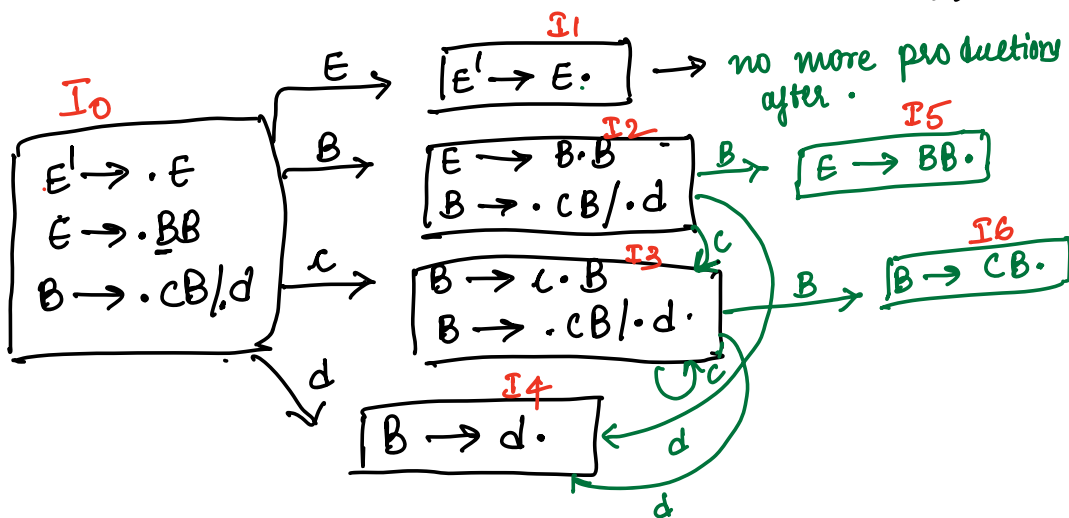
1. Augment the given grammar

Augmented Grammar  $E' \rightarrow E$  (Add a new production by calling start symbol)  
 $E \rightarrow BB$   
 $B \rightarrow cB/d$

2. Draw canonical collection of LR(0) items

$E' \rightarrow E$   
 $E \rightarrow BB$   
 $B \rightarrow cB/d$

LR(0) Items:  $E' \rightarrow \cdot E$  (Add a  $\cdot$  to any production at starting of RHS  $\rightarrow$  LR(0) items)  
 $E \rightarrow \cdot BB$   
 $B \rightarrow \cdot cB/d$



3. Number the productions:

$E' \rightarrow E \rightarrow$  Augmented production

$E \rightarrow BB \rightarrow$  ①

$B \rightarrow cB/d \rightarrow$  ③

↓  
②

$B \rightarrow cB$  &  $B \rightarrow d$   
↓                      ↓  
②                      ③

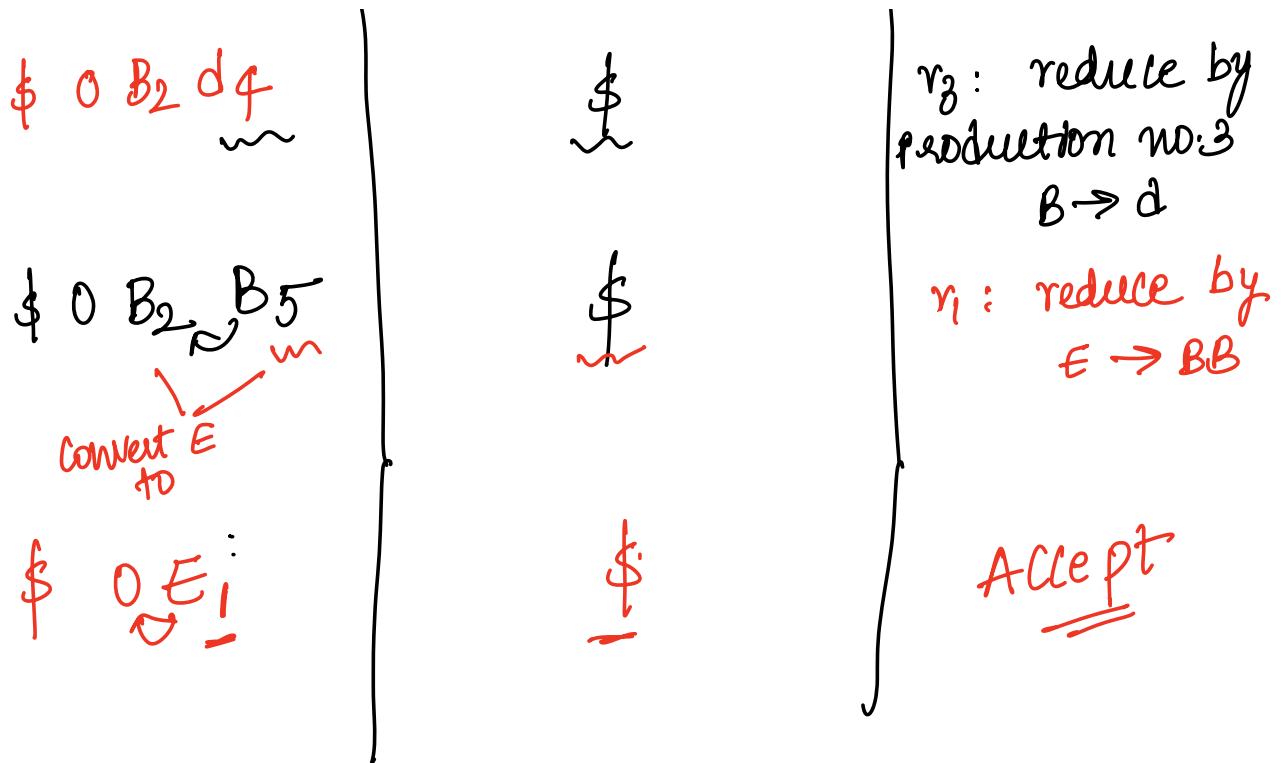
4. Parsing Table:

Action  
↑  
shift (s)  
reduce (r)

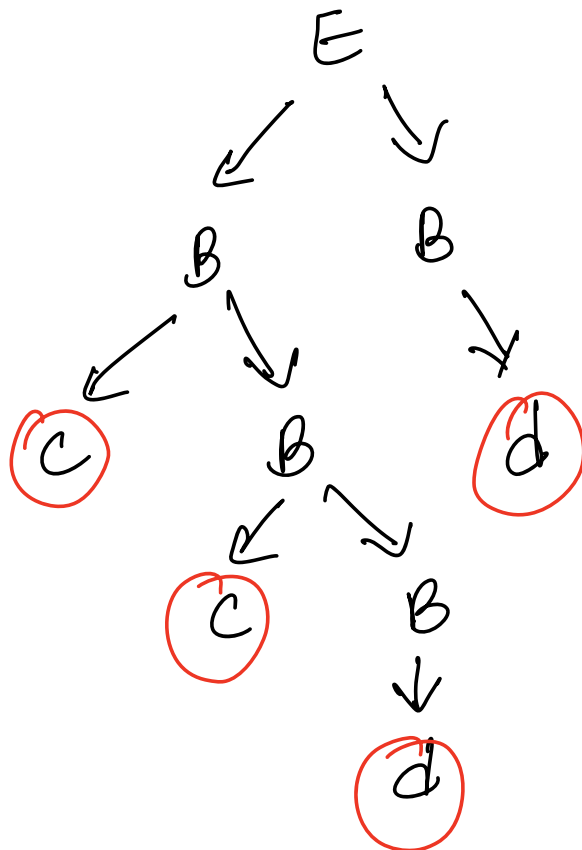
states All I values	Action (Terminals)			Go To (NT)	
	c	d	\$	E	B
0	S <sub>3</sub>	S <sub>4</sub>		1	2
1	Accept				
2	S <sub>3</sub>	S <sub>4</sub>			5
3	S <sub>3</sub>	S <sub>4</sub>			6
4	r <sub>3</sub>	r <sub>3</sub>	r <sub>3</sub>		
5	r <sub>1</sub>	r <sub>1</sub>	r <sub>1</sub>		
6	r <sub>2</sub>	r <sub>2</sub>	r <sub>2</sub>		

### 5. Stack Implementation:

stack	input	Actions
\$ 0 <u>state 0</u>	<u>c</u> c d d \$	<u>s<sub>2</sub></u> : shift c into stack and goto 3
\$ 0 c <u>3</u>	<u>c</u> d d \$	s <sub>3</sub> : shift c into stack & goto 3
\$ 0 c <u>3</u> c <u>3</u>	<u>d</u> d \$	s <sub>4</sub> : shift d & 4
\$ 0 c <u>3</u> c <u>3</u> d <u>4</u>	<u>d</u> \$	<u>r<sub>3</sub></u> : reduce 'd' by B → d
\$ 0 c <u>3</u> c <u>3</u> <u>B</u> <u>6</u> (6 because 3 state and B column value is '6')	<u>d</u> \$	<u>r<sub>2</sub></u> : reduce by B → CB
\$ 0 c <u>3</u> <u>B</u> <u>6</u>	<u>d</u> \$	r <sub>2</sub> : reduce by B → CB
\$ 0 <u>B</u> <u>2</u>	<u>d</u> \$	s <sub>4</sub> : shift 'd' & goto 4



6. Parse tree



ccdd

NOTE:

If the parsing table is having a dual value "rs or rr" (r: reduce, s: shift) together in any of the box, then clearly, you can say it is not a LR(0).

eg:  $\boxed{r_1/r_2 \mid r/s} \rightarrow \text{X } \underline{\text{LR(0)}}$