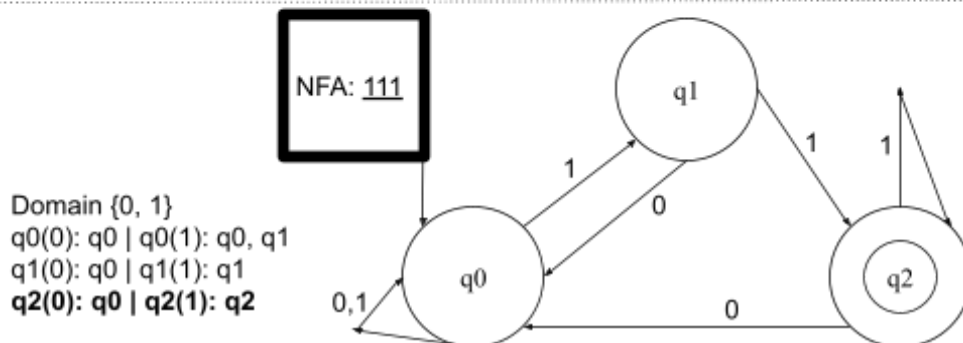
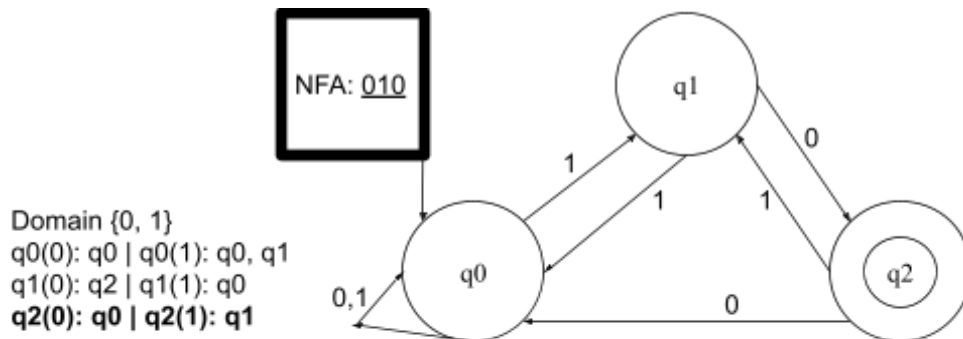
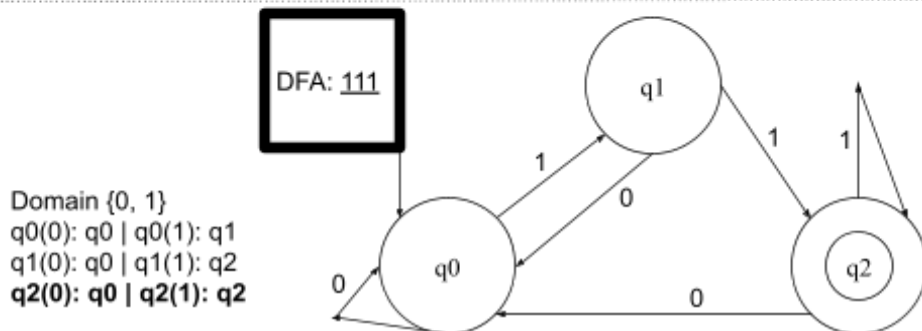
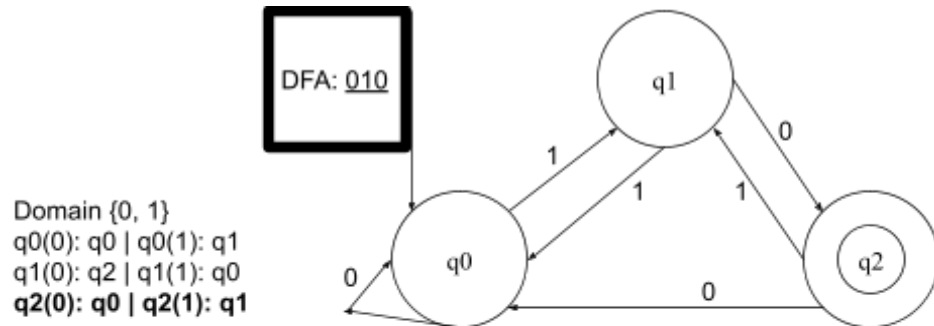


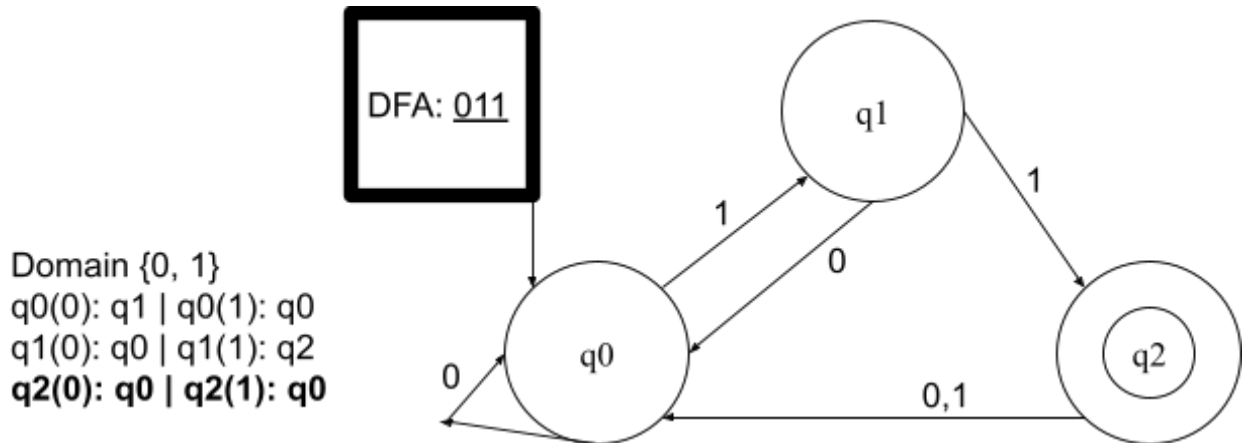
## Assignment 1

1. A compiler is a program dedicated to translating traditional code into machine code. The noun is often used shorthand to describe the holistic process of converting HLLs to absolute code: (preprocessing, compiling, assembling, and loading/linking); however, it more directly refers to the parsing methods of specifically, code compilation: (lexical/syntax/semantic analysis, code generation/optimization, and target code generation).
2. A finite state automata is a set method that allows you to generate code within specific input guidelines. The two large groups of automata are DFAs and NFAs, which are known for having a deterministic and non-deterministic approach- respectively. This difference in approach is represented by the variable output as best shown in DFAs 1-to-1 mapping for edges and individual states; whereas NFAs are known for accepting more than one state per individual edge value.

3. DFA and NFA for { target: 010, 111 }



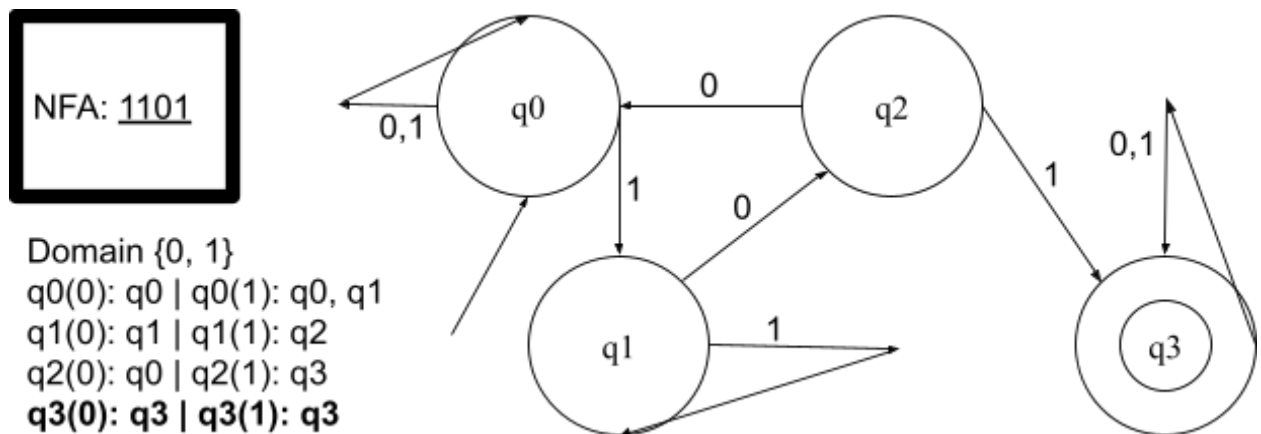
4. DFA; Domain  $\{0,1\}$  {target: 011}



5. Regular Expressions are series' of expressions that describe language acceptance cases.

They first express a domain, and then an expression that uses articles to clarify an acceptance case.

6. NFA {target: 1101}



7. In the shown FSA, there are 3 tuples, each representing a state and their associated transitions (q0, q1, and q2). Input strings with substring '11' will be accepted by the shown FSA. (I thought originally it would be strings that include two 1s but because there's no third loop I chose to write it as shown above).

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8. As a minimal finite automaton, to accept the string 101011, would require 6 unique states.
9. The core difference between a deterministic and non-deterministic approach is that a deterministic has a 1-to-1 functional mapping for each state-input case; whereas, non-deterministic approaches have ambiguous input-output mappings that would require further clarification for implementation.

NFA {domain: 0,1 }

$a_0(0): a_1, a_2 \mid a_0(1): a_2$


$a_1(0): a_2 \mid a_1(1): a_3$

$a_2(0): a_2 \mid a_2(1): a_3$

**$a_3(0): \emptyset \mid a_3(1): \emptyset$**

10. Diagram for the NFA shown:

NFA {domain: 0,1}  
q0(0): q0, q1 | q0(1): q0  
q1(0): q1 | q1(1): q1, q2  
q2(0): q3 | q2(1):  $\emptyset$   
**q3(0): q3 | q3(1): q3**



NFA {domain: 0,1}  
q0 (0): q0q1 | q0(1): q0  
q1 (0): q1 | q1(1): q1q2  
q2 (0): q3 | q2(1):  $\emptyset$   
**q3 (0): q3 | q3(1): q3**  
q0q1(0): q0q1 | q0q1(1): q0q1q2  
q1q2(0): q1q3 | q1q2(1): q1q2  
q1q3(0): q0q1q3 | q1q3(1): q1q2q3  
q0q1q3(0): q0q1q3 | q0q1q3(1): q0q1q2q3  
q1q2q3(0): q1q3 | q1q2q3(1): q1q2q3  
q0q1q2q3(0): q0q1q3 | q0q1q2q3(1): q0q1q2q3