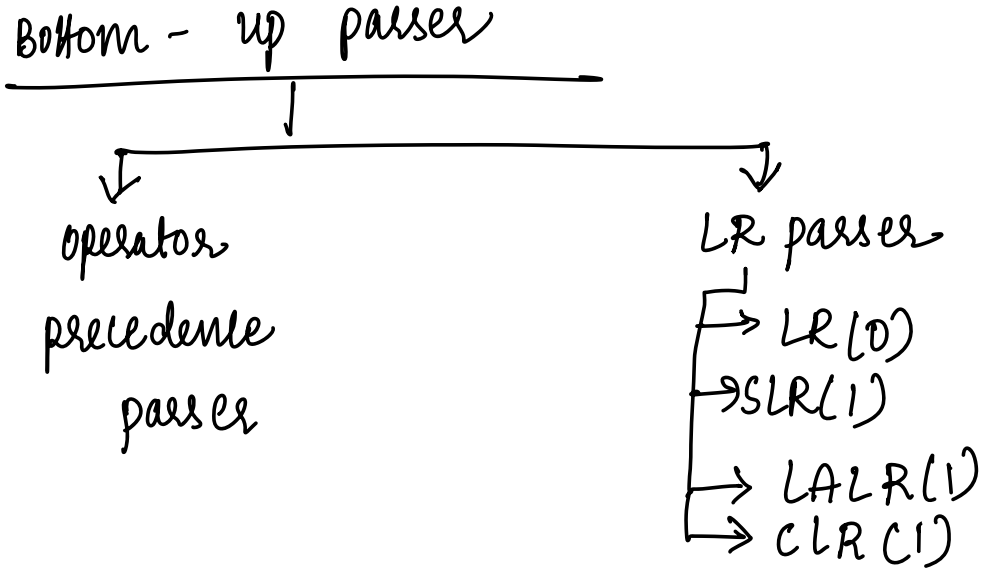
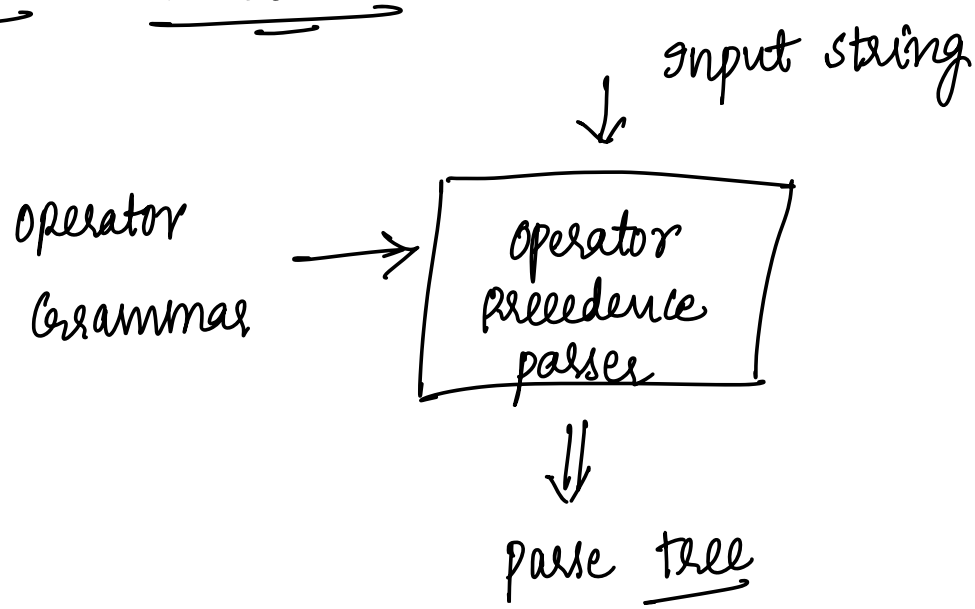
Bottom-up parser:

It constructs the tree from Bottom-up and it considers the Terminals (input) from right-left



* OPERATOR PRECEDENCE



Rules:

operator Grammar:

How to identify? checking the below conditions:

① NO Epsilon (ε) on the right side

② NO adjacent variable. (Having 2 or more N.T together)

Eg:

Adjacent variable yes/NO.

1. AB — YES ✓
2. A+B — NO ✗
3. A-B — NO ✗
4. A/B — ✗ NO

If you find 2, 3, 4 → grammar is OP grammar. (operator precedence)

1. E → E+E / E*E / id ?

Is this OP Grammar YES

2. E → EAE / id NO

3. A → + / * YES

If there are adjacent variables, then you have to convert it into OP Grammar.

Example: $S \rightarrow SAS \mid id$

$A \rightarrow aSa \mid a$

Conversion:

1. $S \rightarrow SAS \mid id$
2. $A \rightarrow aSa \mid a$

In (1) substitute 2.

$S \rightarrow SaSaS \mid S \rightarrow SaS \mid id$

$A \rightarrow aSa \mid a$

Basic knowledge of operators:

id, a, b, c ..., terminals → High ↑

$ → low

+	>	+	(left most has more precedence)
*	>	*	"
id	≠	id	not equal
$	A	$	Accept

① $T \rightarrow T+T \mid T*T \mid id$

with the help of following grammar parse the input string

" id + id * id "

solution

STEPS :

1. check if the grammar is OP or not
2. OP Relation table
3. Parse the given string
4. Generate the parse tree.

1. ✓ grammar is OP.
2. $T \rightarrow T+T \mid T*T \mid id$

High precedence

1. *
2. +
3. —
4. %

	+	*	id	$
+	>	<	<	>
*	>	>	<	>
id	>	>	-	>
$	<	<	<	A

A : Accept

Operator precedence Relation Table

$T \rightarrow T+T$ ✓

$T \rightarrow T*T$ ✓

$T \rightarrow id$

3. Parse the given string. id + id * id $

Stack	Relation	Input	Action (Shift / Reduce)
$	<	id + id * id $	shift id.
compare $ & id.			
$ id	>	+ id * id $	Reduce T → id
if this is greater then reduce.			
$ T	<	+ id * id $	shift +
$ T +	<	id * id $	shift id
$ T + id	>	* id $	Reduce T → id
$ T + T	<	* id $	shift *
$ T + T *	<	id $	shift id
$ T + T * id	>	$	Reduce T → id

$ T + T * T	>	$	Reduce $T \to T * T$
$ T + T	>	$	Reduce $T \to T + T$
$ T	Accept	$	

4. Considering the stack values from Bottom to up we build the parse tree.