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$\underline{\text{Contest}} \ (1)$	
template.cpp	0.11

using 11 = long long; int main() { cin.tie(0)->sync_with_stdio(0); cin.exceptions(cin.failbit);

#include <bits/stdc++.h>

using namespace std;

Data structures (2)

BIT.h

Description: Query [l, r] sums, and point updates. kth() returns the smallest index i s.t. query(0, i) >= k**Time:** $\mathcal{O}(\log n)$ for all ops.

```
33f78c, 22 lines
template <typename T>
struct BIT {
  vector<T> s;
  int n;
  BIT(int n) : s(n + 1), n(n) {}
  void update(int i, T v) {
    for (i++; i <= n; i += i & -i) s[i] += v;
  T query(int i) {
   T ans = 0:
    for (i++; i; i -= i & -i) ans += s[i];
    return ans;
  T query(int 1, int r) { return query(r) - query(1 - 1); }
  int kth(T k) \{ // returns n if k > sum of tree \}
   if (k <= 0) return -1;
    int i = 0;
    for (int pw = 1 << __lg(n); pw; pw >>= 1)
     if (i + pw <= n && s[i + pw] < k) k -= s[i += pw];</pre>
    return i:
};
```

KDBIT.h

Description: k-dimensional BIT. BIT<int, N, M> gives an $N \times M$ BIT. Query: bit.query(x1, x2, y1, y2) Update: bit.update(x, y, delta) Time: $O(\log^k n)$ Status: Tested 3b9692, 28 lines

```
struct BIT {
 T val = 0;
 void update(T v) { val += v; }
 T query() { return val; }
};
template <class T, int N, int... Ns>
struct BIT<T, N, Ns...> {
 BIT<T, Ns...> bit[N + 1];
 // map<int, BIT<T, Ns...>> bit;
  // if the memory use is too high
 template <class... Args>
 void update(int i, Args... args) {
    for (i++; i <= N; i += i & -i) bit[i].update(args...);</pre>
 template <class... Args>
 T query(int i, Args... args) {
    T ans = 0;
    for (i++; i; i -= i & -i) ans += bit[i].query(args...);
    return ans;
  template <class... Args,
            enable if t<(sizeof...(Args) ==
                         2 * sizeof...(Ns))>* = nullptr>
 T query(int 1, int r, Args... args) {
    return query(r, args...) - query(l - 1, args...);
};
```

9 lines

Description: Maintains union of disjoint sets

Time: $\mathcal{O}(\alpha(N))$

template <class T, int... Ns>

```
struct DSU {
 vector<int> s;
 DSU(int n) : s(n, -1) {}
 int find(int i) { return s[i] < 0 ? i : s[i] = find(s[i]); }</pre>
 bool join(int a, int b) {
   a = find(a), b = find(b);
   if (a == b) return false;
   if (s[a] > s[b]) swap(a, b);
   s[a] += s[b], s[b] = a;
   return true;
 int size(int i) { return -s[find(i)]; }
 bool same(int a, int b) { return find(a) == find(b); }
```

RMQ.h

Description: Constant time subarray min/max queries for a fixed array Time: O(nlogn) initialization and O(1) queries. Status: Tested 536eac, 15 lines

```
template <typename T, class Compare = less<T>>
struct RMQ {
 vector<vector<T>> t;
 Compare cmp;
 RMQ(vector < T > \& a) : t(\underline{lg(a.size())} + 1, a) {
    int n = a.size(), lg = __lg(n);
    for (int k = 1, len = 1; k <= lg; k++, len <<= 1)</pre>
      for (int i = 0; i + 2 * len - 1 < n; i++)</pre>
        t[k][i] = min(t[k-1][i], t[k-1][i+len], cmp);
 T query(int a, int b) {
    int k = __lq(b - a + 1), len = 1 << k;</pre>
    return min(t[k][a], t[k][b - len + 1], cmp);
};
```

```
Splay.h
```

c22586, 14 lines

Description: An implicit balanced BST. You only need to change update () and prop().

If used for link-cut tree, code everything up to splay(). Time: amortized $O(\log n)$ for all operations

```
struct node {
 node *ch[2] = \{0\}, *p = 0;
 int cnt = 1, val;
 node (int val, node * l = 0, node * r = 0)
    : ch{l, r}, val(val) {}
int cnt(node* x) { return x ? x->cnt : 0; }
int dir(node* p, node* x) { return p && p->ch[0] != x; }
void setLink(node* p, node* x, int d) {
 if (p) p\rightarrow ch[d] = x;
 if (x) x->p = p;
node* update(node* x) {
 if (!x) return 0;
 x - cnt = 1 + cnt(x - ch[0]) + cnt(x - ch[1]);
 setLink(x, x->ch[0], 0);
 setLink(x, x->ch[1], 1);
 return x;
void prop(node* x) {
 if (!x) return;
 // update(x); // needed if prop() can change subtree sizes
void rotate(node* x, int d) {
 if (!x || !x->ch[d]) return;
 node *y = x - > ch[d], *z = x - > p;
  setLink(x, y->ch[d ^ 1], d);
  setLink(y, x, d^1);
 setLink(z, y, dir(z, x));
 update(x);
 update(y);
node* splay(node* x) {
 while (x && x->p) {
    node *y = x->p, *z = y->p;
    // prop(z), prop(y), prop(x); // needed for LCT
    int dy = dir(y, x), dz = dir(z, y);
    if (!z)
      rotate(y, dy);
    else if (dy == dz)
      rotate(z, dz), rotate(y, dy);
      rotate(y, dy), rotate(z, dz);
 return x;
// the returned node becomes the new root, update the root
// pointer!
node* nodeAt (node* x, int pos) {
 if (!x) return 0;
 while (prop(x), cnt(x->ch[0]) != pos)
    if (pos < cnt(x->ch[0]))
      x = x->ch[0];
      pos -= cnt(x->ch[0]) + 1, x = x->ch[1];
  return splay(x);
```

```
node* merge(node* 1, node* r) {
 if (!1 || !r) return 1 ?: r;
 1 = nodeAt(1, cnt(1) - 1);
  setLink(l, r, 1);
 return update(1);
// first is everything < pos, second is >= pos
pair<node*, node*> split(node* t, int pos) {
  if (pos <= 0 || !t) return {0, t};</pre>
  node *1 = nodeAt(t, pos - 1), *r = 1->ch[1];
  if (r) 1 \rightarrow ch[1] = r \rightarrow p = 0;
  return {update(1), update(r)};
// insert a new node between pos-1 and pos
node* insert(node* t, int pos, int val) {
  auto [1, r] = split(t, pos);
  return update(new node(val, 1, r));
// apply lambda to all nodes in an inorder traversal
template <class F>
void each(node* x, F f) {
 if (x) \text{ prop}(x), \text{ each}(x->\text{ch}[0], f), f(x), \text{ each}(x->\text{ch}[1], f);
```

Geometry (3)

Graphs (4)

Time: $\mathcal{O}(|V| + |E|)$

SCCTarjan.h

Description: Finds strongly connected components of a directed graph. Visits/indexes SCCs in reverse topological order.

Usage: scc(graph) returns an array that has the ID of each node's SCC. scc(graph, [&](vector<int>& v) { ... }) calls the lambda on each SCC, and returns the same array.

```
namespace SCCTarjan {
  vector<int> val, comp, z, cont;
  int Time, ncomps;
  template <class G, class F>
```

```
int dfs(int j, G& g, F& f) {
 int low = val[j] = ++Time, x;
 z.push_back(j);
 for (auto e : g[j])
   if (comp[e] < 0) low = min(low, val[e] ?: dfs(e, q, f));</pre>
 if (low == val[j]) {
   do {
     x = z.back();
     z.pop_back();
     comp[x] = ncomps;
     cont.push_back(x);
    } while (x != j);
   f(cont);
   cont.clear();
   ncomps++;
 return val[j] = low;
template <class G, class F>
vector<int> scc(G& g, F f) {
 int n = q.size();
 val.assign(n, 0);
```

```
comp.assign(n, -1);
   Time = ncomps = 0;
   for (int i = 0; i < n; i++)</pre>
     if (comp[i] < 0) dfs(i, g, f);</pre>
   return comp;
 template <class G> // convenience function w/o lambda
 vector<int> scc(G& g) {
   return scc(g, [](auto& v) {});
} // namespace SCCTarjan
```

SCCKosaraju.h

Description: Finds strongly connected components of a directed graph. Visits/indexes SCCs in topological order.

Usage: scc(graph) returns an array that has the ID of each node's SCC.

Time: $\mathcal{O}(|V| + |E|)$ 9b78e7, 29 lines

```
namespace SCCKosaraju {
 vector<vector<int>> adj, radj;
 vector<int> todo, comp;
 vector<bool> vis;
 void dfs1(int x) {
   vis[x] = 1;
   for (int y : adj[x])
     if (!vis[y]) dfs1(y);
   todo.push_back(x);
 void dfs2(int x, int i) {
   comp[x] = i;
   for (int y : radj[x])
     if (comp[y] == -1) dfs2(y, i);
 vector<int> scc(vector<vector<int>>& _adj) {
   adj = _adj;
   int time = 0, n = adj.size();
    comp.resize(n, -1), radj.resize(n), vis.resize(n);
    for (int x = 0; x < n; x++)
     for (int y : adj[x]) radj[y].push_back(x);
    for (int x = 0; x < n; x++)
     if (!vis[x]) dfs1(x);
    reverse(todo.begin(), todo.end());
    for (int x : todo)
     if (comp[x] == -1) dfs2(x, time++);
    return comp;
}; // namespace SCCKosaraju
```

Mathematics (5)

Fraction.h

358d18, 37 lines

Description: Struct for representing fractions/rationals. All ops are $O(\log N)$ due to GCD in constructor. Uses cross multiplication alde34, 27 lines

```
template <typename T>
struct 0 {
 Ta, b;
 Q(T p, T q = 1) {
   T g = gcd(p, q);
   a = p / q;
   b = q / q;
   if (b < 0) a = -a, b = -b;
 T gcd(T x, T y) const { return __gcd(x, y); }
 O operator+(const O& o) const {
   return {a * o.b + o.a * b, b * o.b};
```

```
O operator-(const O& O) const {
   return *this + O(-o.a, o.b);
 Q operator*(const Q& o) const { return {a * o.a, b * o.b}; }
 Q operator/(const Q& o) const { return *this * Q(o.b, o.a); }
 Q recip() const { return {b, a}; }
 int signum() const { return (a > 0) - (a < 0); }
 bool operator<(const Q& o) const {
   return a * o.b < o.a * b;
 friend ostream& operator<<(ostream& cout, const Q& o) {</pre>
   return cout << o.a << "/" << o.b;
};
```

FractionOverflow.h

Description: Safer struct for representing fractions/rationals. Comparison is 100% overflow safe; other ops are safer but can still overflow. All ops are $O(\log N)$. a42e99, 43 lines

```
template <tvpename T>
struct 00 {
 T a, b;
 QO(T p, T q = 1) {
   T g = gcd(p, q);
    a = p / q;
    b = q / g;
   if (b < 0) a = -a, b = -b;
 T gcd(T x, T y) const { return __gcd(x, y); }
  00 operator+(const 00% o) const {
   T q = qcd(b, o.b), bb = b / q, obb = o.b / q;
    return {a * obb + o.a * bb, o.b * obb};
  OO operator-(const OO& O) const {
    return *this + QO(-o.a, o.b);
  OO operator*(const OO& O) const {
    T g1 = gcd(a, o.b), g2 = gcd(o.a, b);
    return { (a / g1) * (o.a / g2), (b / g2) * (o.b / g1) };
  00 operator/(const 00% o) const {
    return *this * 00(o.b, o.a);
  QO recip() const { return {b, a}; }
  int signum() const { return (a > 0) - (a < 0); }
  static bool lessThan(T a, T b, T x, T y) {
    if (a / b != x / y) return a / b < x / y;</pre>
    if (x % y == 0) return false;
    if (a % b == 0) return true;
    return lessThan(y, x % y, b, a % b);
 bool operator<(const QO& o) const {
    if (this->signum() != o.signum() || a == 0) return a < o.a;</pre>
      return lessThan(abs(o.a), o.b, abs(a), b);
      return lessThan(a, b, o.a, o.b);
 friend ostream& operator<<(ostream& cout, const 00& o) {</pre>
    return cout << o.a << "/" << o.b;</pre>
};
```

PrimeSieve.h

Description: Prime sieve for generating all primes up to a certain limit. isprime[i] is true iff i is a prime.

Time: $\lim_{n\to\infty} 100'000'000 \approx 0.8 \text{ s.}$ Runs 30% faster if only odd indices are stored.

5, 14 lines for

```
const int MAX_PR = 5'000'000;
bitset<MAX_PR> isprime;
vector<int> primeSieve(int lim) {
   isprime.set();
   isprime[0] = isprime[1] = 0;
   for (int i = 4; i < lim; i += 2) isprime[i] = 0;
   for (int i = 3; i * i < lim; i += 2)
      if (isprime[i])
        for (int j = i * i; j < lim; j += i * 2) isprime[j] = 0;
   vector<int> pr;
   for (int i = 2; i < lim; i++)
      if (isprime[i]) pr.push_back(i);
   return pr;
}
```

PrimeSieveFast.h

Description: Prime sieve for generating all primes smaller than LIM. Time: LIM= $1e9 \approx 1.5s$ Details: Despite its n log log n complexity, segmented sieve is still faster than other options, including bitset sieves and linear sieves. This is primarily due to its low memory usage, which reduces cache misses. This implementation skips even numbers.

Pulled directly from KACTL, see there for details.

a1933d, 23 lines

```
const int LIM = 1e8;
bitset<LIM> isPrime;
vector<int> primeSieve() {
  const int S = round(sqrt(LIM)), R = LIM / 2;
  vector<int> pr = \{2\}, sieve(S + 1);
  pr.reserve(int(LIM / log(LIM) * 1.1));
  vector<pair<int, int>> cp;
  for (int i = 3; i <= S; i += 2)</pre>
   if (!sieve[i]) {
      cp.push_back(\{i, i * i / 2\});
      for (int j = i * i; j <= S; j += 2 * i) sieve[j] = 1;</pre>
  for (int L = 1; L \le R; L += S) {
    array<bool, S> block{};
    for (auto& [p, idx] : cp)
      for (int i = idx; i < S + L; idx = (i += p))</pre>
       block[i - L] = 1;
    for (int i = 0; i < min(S, R - L); i++)
      if (!block[i]) pr.push_back((L + i) * 2 + 1);
  for (int i : pr) isPrime[i] = 1;
  return pr;
```

Miscellaneous (6)

NDimensionalVector.h

```
3c0f61, 12 lines
```

```
Submasks.h

for (int mask = 0; mask < (1 << n); mask++)

for (int sub = mask; sub; sub = (sub - 1) & mask)

// do thing
```

$\underline{\text{Strings}}$ (7)

ZValues.h

```
Z Values.in
vector<int> zValues(string& s) {
  int n = ( int ) s.length();
  vector<int> z(n);
  for (int i = 1, 1 = 0, r = 0; i < n; ++i) {
    if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
    while (i + z[i] < n && s[z[i]] == s[i + z[i]]) ++z[i];
  if (i + z[i] - 1 > r) 1 = i, r = i + z[i] - 1;
  }
  return z;
```