

# Algorithmische Eigenschaften von Wahlsystemen I

Ausgewählte Folien zur Vorlesung

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# Websites

- Vorlesungswebsite:

`http://ccc.cs.uni-duesseldorf.de/~rothe/voting1`

- Anmeldung nicht nur im LSF, sondern auch unter

`http://ccc.cs.uni-duesseldorf.de/verwaltung`

(CCC-System für alle meine Veranstaltungen)

# Literature

- **A Richer Understanding of the Complexity of Election Systems**, P. Faliszewski, E. Hemaspaandra, L. Hemaspaandra, and J. Rothe. Chapter 14 in *Fundamental Problems in Computing: Essays in Honor of Professor Daniel J. Rosenkrantz*, pp. 375–406, S. Ravi and S. Shukla, Editors. Springer, Berlin, Heidelberg, New York, 2009.
- **Computational Aspects of Approval Voting**, D. Baumeister, G. Erdélyi, E. Hemaspaandra, L. Hemaspaandra, and J. Rothe. Chapter 10 in *Handbook on Approval Voting*, pp. 199–251, R. Sanver and J. Laslier, Editors. Springer-Verlag, Berlin, Heidelberg, 2010.

# Literature

- **Voting Procedures**, S. Brams and P. Fishburn. Chapter 4 in Volume 1 of the *Handbook of Social Choice and Welfare*, pp. 173–236, K. Arrow, A. Sen, and K. Suzumura, Editors. North-Holland, 2002.
- **Chaotic Elections! A Mathematician Looks at Voting**, D. Saari. American Mathematical Society, 2001.
- **Original Papers** cited in this book and these book chapters.
- ...

# Elections

The Captain of Starship Enterprise is to be elected:  
Candidates:



Voters:



# Elections

## Definition

- An *election* (or *preference profile*)  $(C, V)$  is specified by a set

$$C = \{c_1, c_2, \dots, c_m\}$$

of candidates and a list

$$V = (v_1, v_2, \dots, v_n)$$

of votes over  $C$ .

- How the voters' preferences are represented depends on the voting system used, e.g., by
  - a *linear order (strict ranking)* or
  - an *approval vector*.

# Elections

## Definition

A *linear order* (or *strict ranking*)  $>$  on  $C$  is a binary relation on  $C$  that is

- *total*: for any two distinct  $c, d \in C$ , either  $c > d$  or  $d > c$ ;
- *transitive*: for all  $c, d, e \in C$ , if  $c > d$  and  $d > e$  then  $c > e$ ;
- *asymmetric*: for all  $c, d \in C$ , if  $c > d$  then  $d > c$  does not hold.

Remark:

- 1 Asymmetry of  $>$  implies irreflexivity of  $>$ .
- 2 We often omit the symbol  $>$  in the linear orders and write, e.g.,

$b \ c \ a \ e \ d$  instead of  $b > c > a > e > d$

to indicate that this voter (strictly) prefers  $b$  to  $c$ ,  $c$  to  $a$ ,  $a$  to  $e$ , and  $e$  to  $d$ . So the leftmost candidate is the most preferred one.

# Elections

Remark:

- 3 Occasionally, by dropping asymmetry voters are allowed to be *indifferent* between candidates, as in:

$$b > c = a > e = d$$

If so, it will be mentioned explicitly.

- 4 One may distinguish between *weighted* and *unweighted* voters.  
Default case: unweighted voters (i.e., each voter has weight one).
- 5 Votes may be represented either *succinctly* or *nonsuccinctly*.  
Default case: nonsuccinct (i.e., one ballot per voter).



# Elections

## Example

Election  $(C, V)$  with  $C = \{a, b, c, d, e\}$  and  $V = \{v_1, \dots, v_7\}$ :

$v_1 : c \quad b \quad a \quad e \quad d$

$v_2 : a \quad e \quad d \quad c \quad b$

$v_3 : b \quad a \quad c \quad e \quad d$

$v_4 : b \quad d \quad e \quad a \quad c$

$v_5 : c \quad b \quad a \quad e \quad d$

$v_6 : c \quad d \quad b \quad e \quad a$

$v_7 : e \quad d \quad a \quad b \quad c$

Who should win this election?

# Election Systems

## Definition

An *election system* is a rule determining the winner(s) of a given election  $(C, V)$ . Formally, letting

- $\mathcal{P}(C)^n$  denote the set of all  $n$ -vote preference profiles (e.g.,  $n$  linear orders or  $n$  approval vectors) over the set  $C$  of candidates and
- $\mathfrak{P}(S)$  the set of all subsets of a set  $S$ ,

an election system defines a *social choice correspondence*

$$f : \mathcal{P}(C)^n \rightarrow \mathfrak{P}(C).$$

Given a preference profile  $P \in \mathcal{P}(C)^n$ ,  $f(P) \subseteq C$  is the *set of winners* (which may be empty and may have more than one winner).

# Election Systems

Remark:

- A *social choice function* is a mapping

$$f : \mathcal{P}(C)^n \rightarrow C$$

that assigns a single winner to each given preference profile.

- Letting  $\mathcal{R}(C)$  denote the set of all transitive, total preference relations over  $C$ , a *social welfare function* is a mapping

$$f : \mathcal{P}(C)^n \rightarrow \mathcal{R}(C)$$

that assigns a complete (possibly nonstrict) ranking to each given preference profile.

# Election Systems: An Incomplete Taxonomy

- Preference-based Systems:
  - Positional scoring protocols (plurality, veto,  $k$ -approval, Borda, ...)
  - Majority-based voting (simple majority, Bucklin voting, ...)
  - Pairwise-comparison-based voting procedures (Condorcet, Dodgson, Young, Kemeny, Copeland, Llull, ...)
  - Point distribution voting procedures (single transferable vote, ...)
- Nonranked Systems:
  - Approval voting
  - Negative voting
  - Plurality voting
  - Multistage voting procedures (plurality with run-off, ...)
- Hybrid Systems:
  - Sincere-strategy preference-based approval voting
  - Fallback voting

# Election Systems: Plurality, Antiplurality, $k$ -Approval

## Definition

- *Plurality-rule elections*: The winners are precisely those candidates who are ranked first by the most voters.
- *Antiplurality-rule (a.k.a. veto) elections*: The winners are precisely those candidates who are ranked last by the fewest voters.
- *$k$ -approval*: Each voter gives one point to each of the  $k$  most preferred candidates. Whoever scores the most points wins.

In our above example,  $c$  is the plurality winner,  $e$  is the antiplurality winner, and both  $a$  and  $b$  are 3-approval winners.

# Election Systems: Borda Count

## Definition



- **Borda Count:** With  $m$  candidates, each voter gives:
  - $m - 1$  points to the candidate ranked at first position,
  - $m - 2$  points to the candidate ranked at second position,
  - $\vdots$
  - 0 points to the candidate ranked at last position.

Whoever scores the most points wins.

In our above example,  $b$  is the Borda winner.

# Election Systems: Borda Count

points :	4	3	2	1	0
$v_1 :$	<i>c</i>	<i>b</i>	<i>a</i>	<i>e</i>	<i>d</i>
$v_2 :$	<i>a</i>	<i>e</i>	<i>d</i>	<i>c</i>	<i>b</i>
$v_3 :$	<i>b</i>	<i>a</i>	<i>c</i>	<i>e</i>	<i>d</i>
$v_4 :$	<i>b</i>	<i>d</i>	<i>e</i>	<i>a</i>	<i>c</i>
$v_5 :$	<i>c</i>	<i>b</i>	<i>a</i>	<i>e</i>	<i>d</i>
$v_6 :$	<i>c</i>	<i>d</i>	<i>b</i>	<i>e</i>	<i>a</i>
$v_7 :$	<i>e</i>	<i>d</i>	<i>a</i>	<i>b</i>	<i>c</i>

Viewed as a social welfare function, the Borda system yields:

ranking	<i>b</i>	>	<i>c</i>	>	<i>a</i>	>	<i>e</i>	>	<i>d</i>
points	17	>	15	>	14	>	13	>	11

# Election Systems: Scoring Protocols

## Definition

A *scoring protocol* for  $m$  candidates is specified by a *scoring vector*,  $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_m)$ , satisfying

$$\alpha_1 \geq \alpha_2 \geq \dots \geq \alpha_m.$$

Votes are linear orders. Each vote contributes

- $\alpha_1$  points to that vote's most preferred candidate,
- $\alpha_2$  points to that vote's second most preferred candidate,
- $\vdots$
- $\alpha_m$  points to that vote's second least preferred candidate.

Whoever scores the most points wins.



# Election Systems: Scoring Protocols for $m$ Candidates

Voting System	Scoring Vector
Plurality	$\alpha = (1, \overbrace{0, \dots, 0}^{m-1})$
Antiplurality (Veto)	$\alpha = (\overbrace{1, \dots, 1}^{m-1}, 0)$
$k$ -Approval ( $(m - k)$ -Veto)	$\alpha = (\overbrace{1, \dots, 1}^k, \overbrace{0, \dots, 0}^{m-k})$
Borda Count	$\alpha = (m - 1, m - 2, \dots, 0)$
$\vdots$	$\vdots$



# Simple Majority and Condorcet Voting

## Definition

A candidate  $c$  wins by *simple majority* if  $c$  is ranked first by more than half of the voters.

In our above example, no candidate wins by simple majority. This obstacle is avoided by, e.g., Bucklin voting.

## Definition

A candidate  $c$  is a *Condorcet winner* if  $c$  defeats every other candidate by a strict majority in pairwise comparisons.

In our above example, there is no Condorcet winner (as we have a top-3-cycle). This obstacle is avoided by, e.g., Dodgson, Young, Copeland, and Kemeny voting.

# Election Systems: Dodgson, Young, and Copeland

Let  $(C, V)$  be a given election where votes are linear orders.

- **Dodgson**: The *Dodgson score of  $c \in C$*  (denoted by  $DScore(c)$ ) is the smallest number of sequential switches needed to make  $c$  a Condorcet winner. Whoever has the smallest Dodgson score wins.
- **Young**: The *Young score of  $c \in C$*  (denoted by  $YScore(c)$ ) is the size of a largest sublist of  $V$  for which  $c$  is a Condorcet winner. Whoever has the maximum Young score wins.
- **Copeland**: For each  $c, d \in C$ ,  $c \neq d$ , let  $N(c, d)$  be the number of voters who prefer  $c$  to  $d$ . Let  $C(c, d) = 1$  if  $N(c, d) > N(d, c)$  and  $C(c, d) = 1/2$  if  $N(c, d) = N(d, c)$ . The *Copeland score of  $c$*  is  $CScore(c) = \sum_{d \neq c} C(c, d)$ . Whoever has the maximum Copeland score wins.

# Election Systems: Bucklin Voting

## Definition

- The *strict majority threshold for a list  $V$  of voters* is

$$\text{maj}(V) = \lfloor \|V\|/2 \rfloor + 1.$$

- Given an election  $(C, V)$  and a candidate  $c \in C$ , define the *level  $i$  score of  $c$  in  $(C, V)$*  (denoted by  $\text{score}_{(C,V)}^i(c)$ ) as the number of votes in  $V$  that rank  $c$  among their top  $i$  positions.
- The *Bucklin score of  $c$  in  $(C, V)$*  is the smallest  $i$  such that

$$\text{score}_{(C,V)}^i(c) \geq \text{maj}(V).$$

- All candidates with a smallest Bucklin score, say  $k$ , and a largest level  $k$  score are the *Bucklin winners (BV winners) in  $(C, V)$* .