

Exercises

1. Suppose that $S = \{1, 2, \dots, n\}$. Write a backtracking algorithm to generate all permutations of S .
2. Suppose that $S = \{1, 2, \dots, n\}$. Write a backtracking algorithm to generate all k -subsets of S .
3. (Set Cover Problem) Suppose that $S = \{1, 2, \dots, n\}$ and $C \subseteq \text{Powerset}(S)$ is a collection of subsets of S . Write a backtracking algorithm to find a $C' \subseteq C$ such that:

$$\bigcup_{c \in C'} c = S$$

, where $|C'|$ is minimum.

4. Devise a backtracking algorithm to solve the SUDOKU puzzle from an initial state.
5. A derangement is a permutation p of $\{1, 2, \dots, n\}$ such that no item is in its proper position, i.e. $p_i \neq i$ for all $1 \leq i \leq n$. Write a backtracking program that constructs all the derangements of n items.
6. For a given number n , write a backtracking algorithm to generate all its partitions, i.e.

$$\begin{aligned} 4 &= 1 + 1 + 1 + 1, \\ 4 &= 2 + 1 + 1, \\ 4 &= 2 + 2, \\ 4 &= 3 + 1, \\ 4 &= 4. \end{aligned}$$

Exercises

7. (Set Cover Problem) Suppose that $S = \{1, 2, \dots, n\}$ and $C \subseteq \text{Powerset}(S)$ is a collection of subsets of S . Write a Branch and Bound algorithm to find a $C' \subseteq C$ such that:

$$\bigcup_{c \in C'} c = S$$

, where $|C'|$ is minimum.

8. Suppose that $G = (V, E)$ is a graph. A Clique is a complete subgraph of G . A Max-Clique is a clique containing maximum vertices.
 - a. Write a Backtracking algorithm to find a Max-Clique of G .
 - b. Change the Backtracking algorithm to Branch and Bound algorithm to find a Max-Clique of G .

Exercises

9. Describe the **Adjacency Multi-list** representation of a graph.
10. The **square** of a directed graph $G = (V, E)$ is the graph $G^2 = (V, E^2)$ such that $(u, w) \in E^2$ if and only if for some $v \in V$, both $(u, v) \in E$ and $(v, w) \in E$. That is, G^2 contains an edge between u and w whenever G contains a path with exactly two edges between u and w . Describe efficient algorithms for computing G^2 from G for both the adjacency-list and adjacency-matrix representations of G . Analyze the running times of your algorithms.
11. The **incidence matrix** of a directed graph $G = (V, E)$ is a $|V| \times |E|$ matrix $B = (b_{ij})$ such that

$$b_{ij} = \begin{cases} -1 & \text{if edge } j \text{ leaves vertex } i, \\ 1 & \text{if edge } j \text{ enters vertex } i, \\ 0 & \text{otherwise.} \end{cases}$$

Describe what the entries of the matrix product $B \times B^T$ represent, where B^T is the transpose of B .

12. The **diameter** of a tree $T = (V, E)$ is given by

$$\max \delta(u, v) ; u, v \in V$$

that is, the diameter is the largest of all shortest-path distances in the tree. Give an efficient algorithm to compute the diameter of a tree, and analyze the running time of your algorithm.