PART I

1. Exercise 7.2 out of the book, but where k=4 and postings are all 2 bytes

Solution: To prove that the total number of bytes transferred from/to disk is less than 5×s, where s is the list's size in bytes, we can analyze the Inplace strategy with proportional pre-allocation.

In the Inplace strategy with proportional pre-allocation, when a block of postings is fetched from disk, space is pre-allocated for the postings in the in-memory buffer. The pre-allocation factor is denoted by k, and in this case, k=4.

Let's consider the worst-case scenario where every posting is a new term that needs to be loaded into the buffer. In this case, the number of postings loaded into the buffer will be k times the number of postings that can fit in a block.

Let b be the number of postings that fit in one block. Then, the number of postings loaded into the buffer is k×b.

Now, let s be the total size of the list in bytes, and n be the total number of postings in the list. If every posting has a constant size of 2 bytes, then s=2×n.

In the worst case, k×b postings are loaded into the buffer for every disk block. The total number of disk blocks needed is given by the ceiling of n/b, denoted as \(\text{n/b} \) .

Therefore, the total number of postings loaded into the buffer is: k×b×[n/b]

Now, we know that $s=2\times n$, and the worst-case number of postings loaded into the buffer is $k\times b\times \lceil n/b\rceil$. Therefore, the total number of bytes transferred from/to disk is: $2\times k\times b\times \lceil n/b\rceil$

```
Now, let's substitute
b with s/(2\times n) (since s=2\times n):
2 \times k \times (s/(2\times n)) \times \lceil n/(s/(2\times n)) \rceil
```

Simplify the expression: $k \times (s/n) \times \Gamma = 1$

Since $\lceil 2 \rceil = 2$: 2 x k x (s/n)

```
Substitue k = 4:
8 x (s/n)

Now, substitute s = 2 x n:
8 x ((2xn)/n)

Simplify the expression:
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So, the total number of bytes transferred from/to disk is
16×n. Since s=2×n, the total number of bytes transferred is 8×s.
```

Therefore, we have shown that the total number of bytes transferred from/to disk is less than 5 x s.

2. Come up with pseudo-code for maybe also firstDoc, lastDoc, nextDoc, prevDoc which works with an index that is being stored using Logarithmic merging

Solution:

```
firstDoc( term):
  if term in self.inverted index:
     return min(self.inverted index[term])
  else:
     return None
lastDoc(term):
  if term in self.inverted index:
     return max(self.inverted index[term])
  else:
     return None
nextDoc( term, current doc):
  if term in self.inverted index:
     doc ids = self.inverted index[term]
     doc ids.sort()
     for doc id in doc ids:
       if doc id > current doc:
          return doc id
     return None
  else:
```

return None

```
prevDoc(term, current_doc):
    if term in self.inverted_index:
        doc_ids = self.inverted_index[term]
        doc_ids.sort(reverse=True)
        for doc_id in doc_ids:
            if doc_id < current_doc:
                return doc_id
        return None
    else:
        return None</pre>
```

We accumulate postings in an in-memory auxiliary index, and when the limit is reached, we transfer the postings to a new index created on disk. The next time the auxiliary index is full, we merge it with the existing index to create a larger index.

3. In a plain text ASCII document, define a paragraph to be any block of text preceded by a start of document or two blanks line and continuing until an end of document or two blank lines. Consider the punctuation symbols period, comma, and semi-colon. Write a map reduce algorithm to compute the average number of occurrences of these symbols in a paragraph for a corpus of plain text documents.

Solution:

Given a corpus of plain text documents, we want to calculate the average occurrences of the specified punctuation symbols within each paragraph.

Approach:

- 1. Map Phase:
 - For each document in the corpus:
 - Split the document into paragraphs based on the start of document or two blank lines.
 - For each paragraph:
 - Tokenize the paragraph into words.
 - Count the occurrences of the specified punctuation symbols (period, comma, and semi-colon).
 - Emit key-value pairs: (paragraph_id, (punctuation_count, 1)).

- 2. Shuffle and Sort:
 - o Group the key-value pairs by paragraph ID.
- 3. Reduce Phase:
 - For each paragraph ID:
 - Sum up the punctuation counts and the number of paragraphs.
 - Calculate the average punctuation count per paragraph.
 - Emit the result: (paragraph id, average punctuation count).

Pseudo-Code (MapReduce Algorithm):

```
# Mapper
def map(doc id, document):
  paragraphs = document.split("\n\n") # Split by two blank lines
  for para_id, paragraph in enumerate(paragraphs):
    words = tokenize(paragraph) # Tokenize into words
    punctuation_count = count_punctuation(words)
    yield para id, (punctuation count, 1)
# Reducer
def reduce(para id, counts):
  total punctuation, num paragraphs = 0, 0
  for count, num in counts:
    total punctuation += count
    num paragraphs += num
  average punctuation = total punctuation / num paragraphs
  yield para id, average punctuation
# Example usage
corpus = [...] # List of plain text documents
result = {}
for doc id, document in enumerate(corpus):
  for para_id, avg_punctuation in map(doc id, document):
    result[para id] = avg punctuation
# Print the average punctuation counts for each paragraph
for para_id, avg_punctuation in sorted(result.items()):
  print(f"Paragraph {para_id}: Average Punctuation = {avg punctuation:.2f}")
```

4. For a corpus of plain text documents, write a map reduce algorithm which computes for all terms t the conditional probability of t occurring in a paragraph given that the paragraph has a term 'calvacade' in it. i.e., p(t in paragraph |'calvacade' in paragraph).

Solution:

- 1. Map Phase:
 - Read each document (paragraph) from the input corpus.
 - Tokenize the document into words.
 - o For each word **w** in the document:
 - Emit a key-value pair: (w, 1) if w ≠ 'calvacade'.
 - Emit a special key-value pair: ('calvacade', w) if w ≠ 'calvacade'.
 - This step ensures that we capture the co-occurrence of 'calvacade' with other terms.
- 2. Shuffle and Sort Phase:
 - Group the emitted key-value pairs by key (term).
- 3. Reduce Phase:
 - For each term t:
 - Initialize counters: count_t (total occurrences of t) and count_t_given_calvacade (occurrences of t when 'calvacade' is present).
 - For each value **v** associated with **t**:
 - If v = 1, increment count_t.
 - If $v \neq 1$, increment count t given calvacade.
 - Calculate the conditional probability:
 - p(t in paragraph | 'calvacade' in paragraph) = count t given calvacade / count t
- 4. Output Phase:
 - Emit the term t along with its conditional probability.

Here's a simplified Python pseudo-code for the MapReduce algorithm:

Mapper

def mapper(document):
 for word in tokenize(document):
 if word != 'calvacade':
 emit(word, 1)
 else:
 emit('calvacade', word)

Reducer

```
def reducer(term, values):
  count t = 0
  count_t_given_calvacade = 0
  for value in values:
    if value == 1:
       count t += 1
    else:
       count_t_given_calvacade += 1
  if count t > 0:
    conditional_probability = count_t_given_calvacade / count_t
     emit(term, conditional_probability)
# Driver code
for document in input corpus:
  mapper_output = mapper(document)
  shuffled_data = shuffle_and_sort(mapper_output)
  reducer_output = reducer(shuffled_data)
  output(reducer_output)
```