

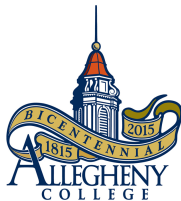
CS201 - PL'S

Names, Scopes, and Bindings

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September 13, 2021



Finishing up on Bye Code

- **Not just used for Java:**

https://en.wikipedia.org/wiki/List_of_JVM_languages

- **JVM instruction list:**

<https://docs.oracle.com/javase/specs/jvms/se8/html/jvms-6.html>

- **Understanding Bytecode makes you a better programmer:**

http://www.ibm.com/developerworks/ibm/library/it-haggar_bytecode/

- **Java Bytecode Fundamentals Blog:**

<http://arhipov.blogspot.com/2011/01/java-bytecode-fundamentals.html>

- A name is exactly what you think it is
 - Most names are identifiers operators
 - symbols (like '+') can also be names
- A **binding** is an association between two things, such as a name and the thing it names.
- The **scope** of a binding is the part of the program (textually) in which the binding is active.

Binding Time

is the point at which a binding is created or, more generally, the point at which any implementation decision is made.

- language design time program structure, possible type
- language implementation time
 - I/O, arithmetic overflow, type equality (if unspecified in manual)

Other Implementation Decisions

- program writing time
 - algorithms, names
- compile time
 - plan for data layout
- link time
 - layout of whole program in memory
- load time

More Implementation Decisions

- run time
 - value/variable bindings, sizes of strings
 - NOTE: run time includes
 - program start-up time
 - module entry time
 - elaboration time (point at which a declaration is first “seen”)
 - procedure entry time

The terms **STATIC** and **DYNAMIC** are generally used to refer to things bound before run time and at run time, respectively.

- In general, early binding times are associated with greater efficiency
- Later binding times are associated with greater flexibility
- Compiled languages tend to have early binding times
- Interpreted languages tend to have later binding times

Scope Rules - control bindings

- Fundamental to all programming languages is the ability to **name data**
 - i.e., to refer to data using symbolic identifiers rather than addresses
- Not all data is named! For example, dynamic storage in C or Pascal is referenced by pointers, not names

```
double *d = (double *)malloc(8);  
*d = 3.14; /* No name is bound to the  
value 3.14 */  
/* The name ``d'' is bound to the  
ADDRESS containing 3.14 */
```

Lifetime and Storage Management

- The period of time from creation to destruction is called the **LIFETIME** of a binding.
- If object outlives binding it's **garbage**.
- If binding outlives object it's a **dangling reference**.
- The textual region of the program in which the binding is active is its **scope**.

Storage Allocation mechanisms

- Static
- Stack
- Heap

Static allocation for

- code
- globals
- static or own variables

Static Example

In C, variables can be global (visible to any function)

```
int i; /* i is global */
int f(int x) {
    return i+x; /* i is visible inside function f */
}
main() {
    int j; /* j is visible only within main */
    i = 10; /* i is visible inside function main */
    j = 99;
    j = f(3); /* sets j to 13 */
}
```

Static Example

When we compile this, i is stored in a fixed location, while j is allocated on the stack

The diagram illustrates the memory layout for variables i and j in ARM assembly. It shows the following assembly code and memory addresses:

```
0000 00000000    i:      ← .space 4
007c 00000000    .L7:      ← .word i
004c 6330A0E3    mov     r3, #99
0050 08300BE5    str     r3, [fp, #-8]
```

Annotations:

- "i" receives an actual memory address (in this case, "0000000")**: Points to the `i:` label.
- "j", however, is just a place on the stack (in this case, 8 bytes below the stack pointer "fp")**: Points to the `[fp, #-8]` offset in the `str` instruction.

(C-to-ARM assembly from <http://assembly.ynh.io/>)

Two Types of Scoping

Static scoping (also called “lexical scoping”)

- most familiar (Java, C)
- scope of variables known at compile time

Dynamic scoping

- scope depends on order of function calls at execution time
- pretty rare nowadays

Static Scope Example (Java)

```
public static int x,y;  
public static void main(String[] args) {  
    x = 10; y = 20;  
    test();  
    { int x = 70,y=80;  
        System.out.println("x = " + x + ", y = " + y);  
    }  
    System.out.println("x = " + x + ", y = " + y);  
}  
public static void test() {  
    int x = 50, y = 60;  
    System.out.println("x = " + x + ", y = " + y);  
}
```

Scope of **x,y**

Scope of **x,y**
("hole" in the
scope of **x,y**)

Scope of **x,y**
("hole" in the
scope of **x,y**)

OUTPUT:

x = 50, y = 60

x = 70, y = 80

x = 10, y = 20

What Happens Here? (Java)

```
public static int x,y;  
public static void main(String[] args) {  
    x = 10; y = 20;  
    test1();  
}  
public static void test1() {  
    int x = 50, y = 60;  
    test2();  
}  
public static void test2() {  
    System.out.println("x = " + x + ", y = " + y);  
}
```

Does this print "x = 10, y = 20"
or
"x = 50, y = 60" ?

What Happens in Dynamic Scoping?

```
int x,y;
start() {
    x = 10; y = 20;
    test1();
    test2();
}
test1() {
    int x = 50, y = 60;
    test2();
}
test2() {
    System.out.println("x = " + x + ", y = " + y);
}
```

Outputs "x = 50, y = 60"

Outputs "x = 10, y = 20"

PLP Chapter 03

Questions

Do you have any questions from this class discussion?