

CS201 - PL'S

Subroutines

Aravind Mohan

Allegheny College

October 25, 2021



Arguments, Parameters

The diagram illustrates the difference between arguments and parameters. It shows a function call on the left and a function definition on the right. The call is `x = f(3, a, b+c);` and the definition is `int f(int x, int y, int z){`. Above the call, a red box labeled "Arguments" has three arrows pointing to the values `3`, `a`, and `b+c`. Above the definition, a red box labeled "Parameters" has three arrows pointing to the variables `x`, `y`, and `z`. The background of the slide features a word cloud with various programming and computer science terms.

```
...  
x = f(3, a, b+c);    int f(int x, int y, int z){  
...
```

Arguments are also called “actual parameters” (as opposed to “formal parameters” in the function definition)

Examples of Function Calls (Python)

```
def f(a,b,c):  
...     return 100*a+10*b+c
```

Examples of Function Calls (Python)

```
def f(a=1,b=1,c=1):  
...     return 100*a+10*b+c
```

- Default Parameter Values
- Named Arguments

Examples of Function Calls (Haskell)

```
Prelude> let f a b c = 100 * a + 10 * b + c
Prelude> f 10 20 30
1230
Prelude> let g = f 1
Prelude> g 20 30
330
Prelude> let h = g 1
Prelude> h 30
140
```

g is a function obtained by giving parameter “a” the value “1” in function *f*. This is called “currying”.

More currying; *h x* is equivalent to *g 1 x*, which is equivalent to *f 1 1 x*

Parameter Evaluation

“Applicative Order” evaluation:

arguments evaluated before the function call. “Eager” evaluation.

```
int slow(int n){  
    .../* count to n^2 */  
    return count;  
}  
int f(int a, int b){  
    return a+1;  
}  
int main(){  
    int x;  
    x = f(10,slow(1000000));  
    printf("%d\n",x);  
}
```

```
time ./a.out  
a = 11  
real    0m20.131s  
user    0m20.126s  
sys     0m0.000s
```

Parameter Evaluation

“Normal Order” evaluation:

arguments are not evaluated until they are needed (possibly never)

File `lazy.hs`:

```
slow 0 = 0
```

```
slow n = 1+slow (n-1)
```

```
f a b = a + 1
```

In ghci, try:

```
Prelude> :l lazy
```

```
[1 of 1] Compiling Main ...
```

```
Ok, modules loaded: Main.
```

```
*Main> f 10 (slow 10000000)
```

```
11
```

```
*Main> f (slow 10000000) 10
```

```
10000001
```

“Lazy” evaluation:

- arguments are evaluated at most once (possibly never).
- Even though we used a Haskell example to illustrate “normal order”, it is more accurate to call Haskell’s evaluation order “lazy.”
- In normal order, a parameter could be evaluated more than once (i.e., evaluated each time it appears in the function).

- In languages that support “first-class functions,” a function may be a return value from another function; a function may be assigned to a variable.
- This raises some issues regarding scope.

Closures

JavaScript example: <http://goo.gl/kzUCes>

```
function f(name) {  
    var x = "hi there";  
    function g() {  
        return x+" "+name;  
    }  
    return g;  
}  
var k = f("bob");
```

Function **f** returns the function **g**. Therefore, variable **k** is assigned a *function*.

Once **f** is done, how will **k** (i.e., **g**) know the values of **x** and **name**?

Closures

One solution (NOT the one used by JavaScript!):

use the most recently-declared values of variables “name” and “x”.

– This is called “*shallow binding*.” Common in dynamically-scoped languages.

Another solution (used by JavaScript and most other statically-scoped languages):

bind the variables that are in the environment where the function is defined.

– This is an illustration of “deep binding” and the combination of the function and its defining environment is called a closure.

Closures: Another example (C#)

```
static void Main(string[] args) {  
    Func <int,int> g = returnFunc(7);  
    Console.WriteLine(g(9));  
}  
  
static Func<int,int> returnFunc(int x) {  
    Func <int,int> g =  
        delegate(int z) { int y = 10; return x+y; }  
    return g;  
}
```

See `Closurer.cs` in the repo.

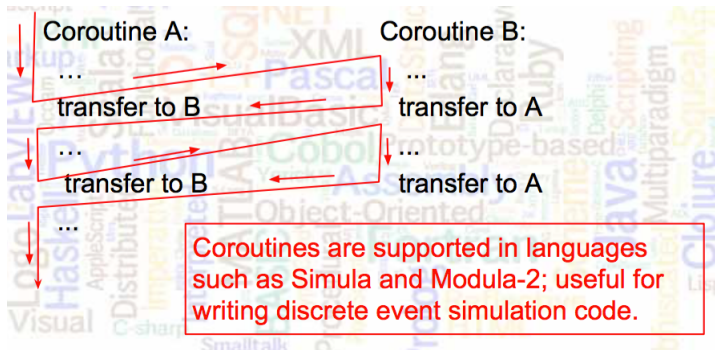
Exceptions

See `Exc.java` in the repo.

Coroutines

- A **coroutine** is a function that can be suspended and resumed.
- Several coroutines can be active at once, transferring control back and forth between them.

Coroutines



Generators in Python

Many other languages have features that allow the implementation of coroutines (even if they are not "built in" to the language).

— Python has generator functions:

```
>>> def gen99():
...     for i in range(100):
...         yield i # NOTE: not "return i"
>>> a = gen99() # call the function just once
>>> next(a)
0
>>> next(a)
1
```


Generators in Python

```
>>> next(a)
2
>>> for i in range(10):
...     print next(a),
...
3 4 5 6 7 8 9 10 11 12
>>> for i in range(10):
...     print next(a),
...
13 14 15 16 17 18 19 20 21 22
```

Generators in Python

- Several generators can be active at the same time; see sample program `gen.py` in the shared repository.
- This isn't precisely a coroutine example (we don't have “call” and “response” directly transferring back and forth).
- See <https://docs.python.org/3/library/asyncio-task.html> (“Tasks and coroutines”, Python 3.5 documentation) intersection.

Functions: Summary

We have considered the following topics:

- Parameter passing (e.g., pass by value, pass by reference)
- Special syntax (default values, named parameters)
- Parameter evaluation (applicative, normal, lazy)
- Closures
- Exceptions
- Coroutines

Functions: Summary

We didn't cover:

- Generics (you have seen a lot of this in CMPSC 101)
- Events (we may revisit this topic in Chapter 13 if schedule permits).

PLP Chapter 08 [8.1 - 8.3; 8.5 - 8.6]

Questions

Do you have any questions from this class discussion?