

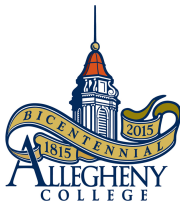
CS202 - Algorithm Analysis

Merge Sort

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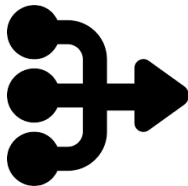
Allegheny College

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Sedgewick 2.2 Merge Sort

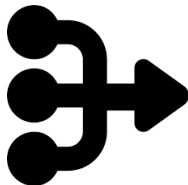
Merge Sort Algorithm



Strategy:

- **Divide:** if S has at least two elements, remove all the elements from S and put them into two sequences S_1 and S_2 , each containing about half of the elements of S . (i.e., S_1 contains the first floor $(n/2)$ elements and S_2 contains the remaining floor $(n/2)$ elements.
- **Conquer:** Sort sequences S_1 and S_2 using Merge Sort.
- **Combine:** Put back the elements into S by merging the sorted sequences S_1 and S_2 into one sorted sequence.

Merge Sort Algorithm



Characteristics:

- sort out of "place", i.e., does require an additional array
- uses divide and conquer principle
- worst case running time is $O(n \times \log(n))$

Merge Sort Algorithm

Merge Procedure (linear)

Algorithm - Merge(A, p, m, r)

Input: an n -element un-sorted array A of integer values, a lower bound p of the array A , and a pivot r in the array A .

Output: an n -element sorted array A of integer values.

Merge Sort Algorithm

Merge Procedure (linear)

```
 $n_1 \leftarrow m - p$   
 $n_2 \leftarrow r - m$   
Initialize Two Arrays L of size  $n_1 + 1$  and R of size  $n_2 + 1$   
for  $i = 0$  to  $n_1$  do  
     $L[i] \leftarrow A[p+i]$   
end for  
for  $j = 0$  to  $n_2$  do  
     $R[j] \leftarrow A[m+j]$   
end for  
 $L[n_1 + 1] \leftarrow \infty$  and  $R[n_2 + 1] \leftarrow \infty$   
Initialize  $i, j \leftarrow 0$   
for  $k = p$  to  $r$  do  
    if  $L[i] \leq R[j]$  then  
         $A[k] \leftarrow L[i]$   
         $i \leftarrow i + 1$   
    else  
         $A[k] \leftarrow R[j]$   
         $j \leftarrow j + 1$   
    end if  
end for
```

Merge Sort Algorithm

MergeSort Procedure (logarithmic)

Algorithm - MergeSort(A, p, r)

Input: an n -element un-sorted array A of integer values, a lower bound p of the array A , and a pivot r in the array A .

Output: an n -element sorted array A of integer values.

if $p < r$ **then**

$m \leftarrow \text{Floor}((p + r)/2)$

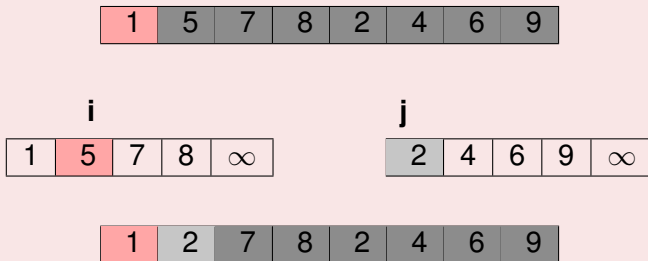
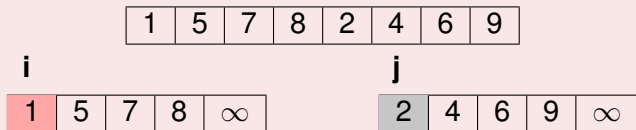
 MergeSort(A, p, m)

 MergeSort($A, m+1, r$)

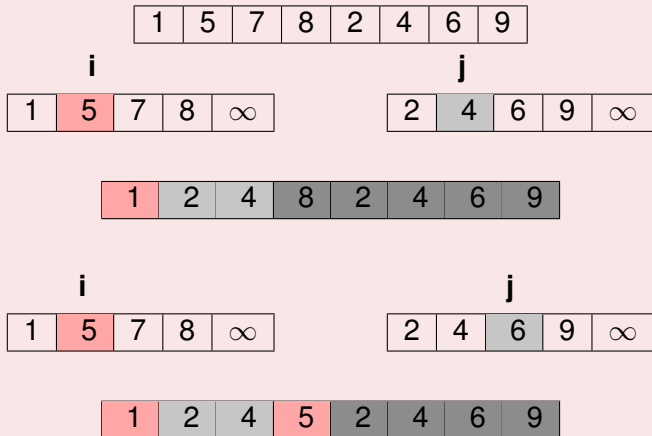
 Merge(A, p, m, r)

end if

Merge Example



Merge Example



Merge Example

1	5	7	8	2	4	6	9
---	---	---	---	---	---	---	---

i

1	5	7	8	∞
---	---	---	---	----------

j

2	4	6	9	∞
---	---	---	---	----------

1	2	4	5	6	4	6	9
---	---	---	---	---	---	---	---

i

1	5	7	8	∞
---	---	---	---	----------

j

2	4	6	9	∞
---	---	---	---	----------

1	2	4	5	6	7	6	9
---	---	---	---	---	---	---	---

Merge Example

1	5	7	8	2	4	6	9
---	---	---	---	---	---	---	---

i

1	5	7	8	∞
---	---	---	---	----------

j

2	4	6	9	∞
---	---	---	---	----------

1	2	4	5	6	7	8	9
---	---	---	---	---	---	---	---

i

1	5	7	8	∞
---	---	---	---	----------

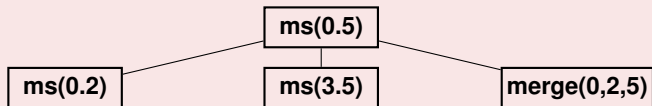
j

2	4	6	9	∞
---	---	---	---	----------

1	2	4	5	6	7	8	9
---	---	---	---	---	---	---	---

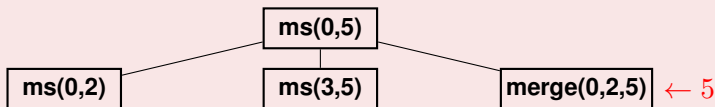
Merge Sort Example

9	6	5	0	8	2
---	---	---	---	---	---



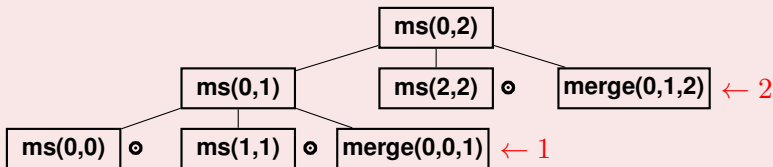
Merge Sort Example

9	6	5	0	8	2
---	---	---	---	---	---



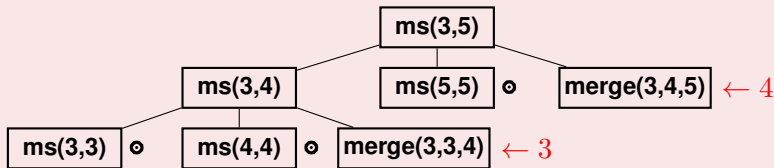
Merge Sort Example

9	6	5	0	8	2
---	---	---	---	---	---



Merge Sort Example

9	6	5	0	8	2
---	---	---	---	---	---



OK done

Merge Sort Example

Merge(0,0,1)

6	9				
---	---	--	--	--	--

9	∞
---	----------

6	∞
---	----------

Merge Sort Example

Merge(0,1,2)

5	6	9			
---	---	---	--	--	--

6	9	∞
---	---	----------

5	∞
---	----------

Merge Sort Example

Merge(3,3,4)

5	6	9	0	8	
---	---	---	---	---	--

0	∞
---	----------

8	∞
---	----------

Merge Sort Example

Merge(3,4,5)

5	6	9	0	2	8
---	---	---	---	---	---

0	8	∞
---	---	----------

2	∞
---	----------

Merge Sort Example

Merge(0,2,5)

0	2	5	6	8	9
---	---	---	---	---	---

5	6	9	∞
---	---	---	----------

0	2	8	∞
---	---	---	----------

Done Sorting!

Merge Sort Algorithm - An analysis

- Space complexity: $O(n)$
- Time complexity: $O(n \times \log(n))$ for Worst, Average, and Best case

Tree Data Structure

Sedgewick 2.2 Merge Sort

Questions?

Please ask if there are any Questions!