

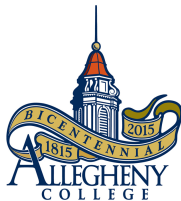
# *CS202 - Algorithm Analysis*

## Selection Sort

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## Sedgewick 1.3, Queues, 2.3 Selection Sort

# Selection Sort Algorithm



## Strategy:

- Place all items in Bag 1.
- Remove the first item from Bag 1 and place it in Bag 2.
- Find the minimum item ( $X$ ) from Bag 1 and switch it with the most recently placed item ( $Y$ ) in Bag 2, **if**  $X < Y$ .
- Repeat steps 2 and 3 until Bag 1 is empty and Bag 2 is full.
- **Note:** Both bags are in one array itself, since this is in-place sorting. The concept of Bag is only logical and not physical.

# Selection Sort Algorithm

**Algorithm** - Selection( $A$ )

**Input:** an  $n$ -element un-sorted array  $A$  of integer values.

**Output:** an  $n$ -element sorted array  $A$  of integer values.

```
1: for  $i = 0$  to  $n - 1$  do
2:    $min \leftarrow i$ 
3:   for  $j = i + 1$  to  $n$  do
4:     if  $A[j] < A[min]$  then
5:        $min \leftarrow j$ 
6:     end if
7:   end for
8:   if  $min \neq i$  then
9:     swap  $A[i]$  &  $A[min]$ 
10:  end if
11: end for
```

# Selection Sort Example

**Input:** [5,4,3,2,1])

```
i=0      j=[1,2,3,4]      min=4      array=[1, 4, 3, 2, 5]
i=1      j=[2,3,4]        min=3      array=[1, 2, 3, 4, 5]
i=2      j=[3,4]          min=2      array=[1, 2, 3, 4, 5]
i=3      j=[4]            min=3      array=[1, 2, 3, 4, 5]
```

# Selection Sort Example

**Input:** [1,2,3,4,5])

```
i=0      j=[1,2,3,4]      min=0      array=[1, 2, 3, 4, 5]
i=1      j=[2,3,4]       min=1      array=[1, 2, 3, 4, 5]
i=2      j=[3,4]        min=2      array=[1, 2, 3, 4, 5]
i=3      j=[4]          min=3      array=[1, 2, 3, 4, 5]
```

# Selection Sort Example

**Input:** [5,6,9,0,8,2,7,1,3])

```
i=0      j=[1,2,3,4,5,6,7,8]    min=3    array=[0, 6, 9, 5, 8, 2, 7, 1, 3]
i=1      j=[2,3,4,5,6,7,8]      min=7      array=[0, 1, 9, 5, 8, 2, 7, 6, 3]
i=2      j=[3,4,5,6,7,8]        min=5      array=[0, 1, 2, 5, 8, 9, 7, 6, 3]
i=3      j=[4,5,6,7,8]          min=8      array=[0, 1, 2, 3, 8, 9, 7, 6, 5]
i=4      j=[5,6,7,8]            min=8      array=[0, 1, 2, 3, 5, 9, 7, 6, 8]
i=5      j=[6,7,8]              min=7      array=[0, 1, 2, 3, 5, 6, 7, 9, 8]
i=6      j=[7,8]                min=6      array=[0, 1, 2, 3, 5, 6, 7, 9, 8]
i=7      j=[8]                  min=8      array=[0, 1, 2, 3, 5, 6, 7, 8, 9]
```

# Worst Case Analysis

- $L_2, L_8$  to  $L_{10}$  executes one time each for every iteration.

$$1 + 1 + 1 + 1 + \dots + n - 1 = \mathbf{n-1}$$

- $L_4$  to  $L_6$  executes in the sequence shown below for all iterations combined.

$$(n-1) + (n-2) + (n-3) + (n-4) + \dots$$

$$+ (n - (n-3)) + (n - (n-2)) + (n - (n-1))$$

$$= (n-1) + (n-2) + (n-3) + (n-4) + \dots + 3 + 2 + 1$$

$$= \frac{n \times (n-1)}{2} = \frac{n^2}{2} = \mathbf{n^2}$$

- Total execution time  $= (n-1) + n^2 = \mathbf{O(n^2)}$



# Best Case Analysis

- $L_2, L_8$  to  $L_{10}$  executes one time each for every iteration.

$$1 + 1 + 1 + 1 + \dots + n - 1 = \mathbf{n-1}$$

- $L_4$  to  $L_6$  executes in the sequence shown below for all iterations combined.

$$(n - 1) + (n - 2) + (n - 3) + (n - 4) + \dots$$

$$+ (n - (n - 3)) + (n - (n - 2)) + (n - (n - 1))$$

$$= (n - 1) + (n - 2) + (n - 3) + (n - 4) + \dots + 3 + 2 + 1$$

$$= \frac{n \times (n - 1)}{2} = \frac{n^2}{2} = \mathbf{n^2}$$

- Total execution time  $= (n - 1) + n^2 = \mathbf{O(n^2)}$

# Average Case Analysis

- $L_2, L_8$  to  $L_{10}$  executes one time each for every iteration.

$$1 + 1 + 1 + 1 + \cdots + n - 1 = \mathbf{n-1}$$

- $L_4$  to  $L_6$  executes in the sequence shown below for all iterations combined.

$$\begin{aligned} & \frac{(n-1) + (n-2) + \cdots + 3 + 2 + 1}{2} \\ &= \frac{n \times (n-1)}{4} = \frac{n^2}{4} = \mathbf{n^2} \end{aligned}$$

- Total execution time =  $(n-1) + n^2 = \mathbf{O(n^2)}$

# Another Sorting Algorithm

**Next: Quick Sort Algorithm**

## Sedgewick 2.3 Selection Sort

# Questions?

**Please ask if there are any Questions!**