

# Distributed Systems

CS 380D

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# Types of knowledge

- Common knowledge:
  - known by everyone in group
  - each node **can** assume others know this
- Distributed knowledge:
  - known by some members of group
  - a node **cannot** assume others know this
- Simultaneous actions requires common knowledge



# Common Knowledge

- **Impossible** to obtain if communication is over unreliable channels
- Demonstrated in the Coordinated Attack Problem
- Internal Common Knowledge:
  - assume something is common knowledge
  - hope no node encounters state that disproves assumption

# Muddy Children Puzzle

- $n$  children play,  $k$  get muddy
- Each can observe all others, don't know their own state
- Dad says "at least one of you is muddy"
- Dad asks each of them: "do you know if you are muddy"?
- claim: After  $k-1$  rounds, all children will answer yes



# Muddy Children Puzzle

- Children get information from:
  - Observation of other children
  - Hearing what other children say
  - Inferences based on previous rounds
- Common knowledge: father says at start "at least one of you is muddy"

# Proof by induction

- $k = 1$ :

- Muddy child observes all others are clean
- But father said someone is muddy
- Hence child realizes they are muddy, answers yes
- Once other children hear muddy child answer yes, they also answer yes

- $k = 2$ :

- Each muddy child observes one other muddy child
- in first round,  $k = 1$ , all answer no as they are unsure of their own state
- Muddy child realizes they are muddy, since other muddy child answered no in first round (hence other child must see someone muddy)
- In second round, all answer "yes"



# Proof by induction

- $k = 3$
- Say muddy children are  $a, b, c$
- if  $a$  is clean,  $b$  and  $c$  would have answered yes in second round
- Hence  $a$  is not clean;  $b$  and  $c$  do similar reasoning
- All answer yes on third round

# Does father need to provide common knowledge?

- One might think no: for  $k > 1$ , seems like children get the information from direct observation
- However, it is not common knowledge
- For  $k = 2$ , muddy child a observes muddy child b. But **does not know** if b observes a, and therefore knows  $k \geq 1$



# Does father need to provide common knowledge?

- Showing it does not work for  $k = 2$ :
  - Muddy children are A and B
  - In first round, even if A had seen all clean kids, they would have still answered "no" (because they do not know  $k \geq 1$ )
  - In second round, A and B realizing they are muddy depends on muddy child saying yes in round 1
  - A saying "no" in round 1 does not provide B with any information
  - B still thinks  $k = 1$  or  $k = 2$

# Does father need to provide common knowledge?

- Valid sequence if  $k = 1$  from B's viewpoint:
  - A is only muddy child
  - A does not realize  $k \geq 1$ , cannot decide between  $k = 0$  and  $k = 1$
  - A says "no" in first round
- B still cannot decide between  $k = 1$  or  $k = 2$  (both can happen with prior seq)



# Common knowledge

- $k \geq 1$  is distributed knowledge, not common knowledge
- This case clearly shows the difference between the two

# Hierarchy of States of Knowledge

- Agent's knowledge depends on:
  - Starting knowledge
  - Observed history since start
- If agent  $i$  knows  $P$  then  $K_i(P)$
- Agents know only true things



# Hierarchy of States of Knowledge

- $D(G, P)$  = group  $G$  has distributed knowledge of  $P$   
(union of knowledge of  $G$  members =  $P$ )
- $S(G, P)$  = someone in  $G$  knows  $P$
- $E(G, P)$  = everyone in  $G$  knows  $P$
- $E(G, K, P) = E(E(E.. E(G, P))))$   $k$  times
- $E(E(G, P))$  = everyone in  $G$  knows that everyone in  $G$  knows  $P$
- Common knowledge:  $E(G, K, P)$  for all  $K \geq 1$

# Muddy Children Puzzle

- $m$  = "at least one child is muddy"
- Without father speaking,
  - $E(G, K-1, m) = \text{true}$
  - $E(G, K+1, m) = \text{false}$



# Muddy Children Puzzle

- $E(G, m) = \text{true}$
- In  $k = 2$ ,
  - $E(G, 1, m) = \text{true}$ , everyone knows  $m$
  - $E(G, 2, m) = \text{false}$ , everyone does not know everyone knows  $m$
  - Specifically, with muddy children A and B, A does not know B knows  $K \geq 1$
- Father's statement makes  $E(G, 2, m) = \text{true}$

# Knowledge in distributed systems

- Communication in a distributed systems seeks to move up the hierarchy of knowledge:
  - changing  $S(G, P) = E(G, P) = C(G, P)$
- Fact discovery:
  - Changing D to S to E to C
  - Example: finding deadlock in a set of distributed locks
- Fact publication:
  - Changing S to C
  - Example: new protocol for communication



# Common Knowledge

- How does one establish it?
  - By being part of a community
    - Membership procedure imparts common knowledge
    - Example: community of licensed drivers knows what signs mean
  - By being co-present at knowledge creation
    - Example: children being in same room as father when he makes announcement

# Coordinated Attack Problem

- General A sends time in message to General B
- A will not attack without ack from B
- B sends ack to A
- But B will not attack without ack from A
- A sends  $\text{Ack}(\text{Ack}(A))$  to B
- A will not attack without ack of this message
- And so it goes.. A and B cannot agree with finite messages
- Can use induction to prove no set of K messages is enough



# Coordinated Attack Problem

- Generals A and B need common knowledge of the attack time
- After A sends the first message to B,  $E(G) = \text{true}$ , but  $E(E(G))$  is not
- After A sends the ack to B's ack,  $E(E(G)) = \text{true}$ , but  $E(3, G) = \text{false}$
- Coordinated attack requires  $C(G) = E(k, G) = \text{true}$  for any  $k$