Linear diophantine equations.

$$3x+5y=2$$

 $3x = 2 \mod 5 \iff x = 4 \mod 5$

$$x=4$$
 $y=-2$

ax+by=c, where a,b,c are integers.

All solutions to this eqn.

$$ax + by = C$$

$$\Rightarrow ax \equiv c \mod b$$

$$\Rightarrow x \equiv a^{2}c \mod b$$

$$\Rightarrow x \equiv a^{2}c \mod b$$

$$\Rightarrow how to compute it?$$

$$3x \neq 5y = 26$$

$$3x = 2 \mod 5$$

$$1x = 4 \mod 5$$

$$y = 2 - 3(5t + 4) = [-3t - 2]$$

$$(5t+4,-3t-2)$$
 $t\in\mathbb{Z}$
 $(4,-2)$
 $+5$ -3 $(4,-2)+(5t,-3t)$

$$gcd(a,b)\neq 1$$

$$g|ax g|by \Rightarrow g|c$$

$$4x+6y=5 \times no solutions.$$

$$\frac{a}{9}x+\frac{b}{9}y=\frac{c}{9}$$

$$gcd(a/9,b/9)=1$$

Q) How to find a' mod b? Euclidean algorithm e.g. = a=5, b=7 Fast, #titerations, (min(a,b)).

Idea: To find integers m,n such that 5m+7n=1 $5m=1 \mod 7$ $m=5^{-1} \mod 7$

$$7 = 1 \times 5 + 2$$
 $5 = 2 \times 2 + 1$

$$1 = 5 + (-2) \times 2$$

$$= 5 + (-2)(7 - 5)$$

$$= 3 \times 5 + (-2) \times 7$$

gcd(a,b)=1a' mod b. prime p: Fermat's Theorem $a^{b-1} \equiv 1 \mod b$ $\Rightarrow a \cdot a^{b-2} \equiv 1 \mod b$ $\Rightarrow a^{b-2} = \overline{a}' \mod b.$ exponentiation, fast.

a mod b. Euter's Theorem

Not fast to $a(b) = 1 \mod b$ compute. $a^{-1} = a^{\phi(b)-1} \mod b$

p(b)= number of tve integers in {1,2,...,b} which are coprime to b.

Problem Statement

You are given integers X and Y, which satisfy at least one of $X \neq 0$ and $Y \neq 0$.

Find a pair of integers (A, B) that satisfies all of the following conditions. If no such pair exists, report so.

$$-10^{18} \le A, B \le 10^{18}$$

ullet The area of the triangle with vertices at points (0,0),(X,Y),(A,B) on the xy-plane is 1.

Constraints

- $-10^{17} \le X, Y \le 10^{17}$
- $(X,Y) \neq (0,0)$
- ullet X and Y are integers.

$$\frac{1}{2} | x_1^{0}(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2) | gcd(x,y) | 2$$

$$(x_2,y_2) | (x_3,y_3) = \frac{1}{2} | xB - AY | x = 4 y = 6$$

$$Bx - Ay = 2 \longrightarrow 2B - 3A = 1$$

$$A,B) = (1+2t,2+3t)$$

$$Bx - Ay = -2$$

Euler Totient Function.

 $\phi(n)$: number of integers $a \in \{1,2,...,n\}$ such that $\gcd(a,n)=1$.

$$\phi(6) = 2$$
 (1)2,3,4,5)6
 $\phi(b) = b-1$ (1,2,.., $b-1$)

$$m,n: coprime: \phi(m)\phi(n) = \phi(mn)$$

m=3 n=43t=1 mod 4 a=2 mod 4 1 mod 3 t= 3 mod 4 t=4k+3 a= 10 mod 12 a=12k+10 t=-3t=-1=3 $x = 2 \mod 3$ $x = 1 \mod 4$ $x \equiv 5 \mod 12$ coprime m, n $x \equiv \underline{a} \mod m \implies x \equiv c \mod mn$

 $x = b \mod n$

Chinese Remainder Theorem.

Let p_n , in be coprime naturals, and let $a,b \in \mathbb{Z}$. Then \exists a unique (mod mn) integer c such that

$$x \equiv a \mod m$$
 $\Rightarrow x \equiv c \mod mn$.
 $x \equiv b \mod n$

mxn possibilities"

-> More general statements hold.

$$\frac{\phi(m)\phi(n)}{\phi(m)} = \frac{\phi(mn)}{\phi(mn)}.$$
Number of integers $a \in \{1, ..., m\}$
 $b \in \{1, ..., n\}$
 $s.t.$

$$\gcd(a, m) = 1$$

$$\gcd(b, n) = 1$$

$$x \equiv a \mod m$$

$$x \equiv a \mod m$$

$$x \equiv b \mod n$$

$$\gcd(x, mn) = 1$$

$$\frac{\phi(m)\phi(n)}{\phi(n)} = \phi(mn).$$
wher of integers $a \in \{1,...,m\}$
 $b \in \{1,...,n\}$
 $\gcd(a,m)=1$
 $d = \gcd(b,n)=1$
 $d = \gcd$

$$\phi(mn) = \phi(m)\phi(n)$$

$$\phi(n) = n \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_k}\right)$$

Q) Find
$$\phi(1), \phi(2), \dots, \phi(n)$$
. Can we do this faster?

Q) Compute
$$\phi(1)$$
, $\phi(2)$,..., $\phi(n)$.

Time: O(n log log n).

$$\phi(n) = n\left(\frac{1-1}{p_1}\right) \cdot - \left(\frac{1-1}{p_R}\right)$$

-> For any prime p/n, we want to multiply 1-1 to its corresponding p.

$$\begin{array}{c}
A[1] = 1 \\
A[2] = 2 \times (1 - \frac{1}{2}) \\
A[3] = 3 \times (1 - \frac{1}{3}) \\
A[4] = 4 \times (1 - \frac{1}{2}) \\
A[5] = 5 \times (1 - \frac{1}{3}) \times (1 - \frac{1}{3}) \\
A[6] = 6 \times (1 - \frac{1}{2}) \times (1 - \frac{1}{3})
\end{array}$$

$$n\left(1-\frac{1}{p_{1}}\right)\cdots\left(1-\frac{1}{p_{K}}\right)$$

$$A[1] = 1$$

$$A[2] = 2 \times (1 - \frac{1}{2})$$

$$A[3] = 3 \times (1 - \frac{1}{3})$$

$$A[4] = 4 \times (1 - \frac{1}{2})$$

$$A[5] = 5 \times (1 - \frac{1}{5})$$

$$A[6] = 6 \times (1 - \frac{1}{2}) \times (1 - \frac{1}{3})$$

$$A[7] = 7 \times (1 - \frac{1}{3})$$

$$A[8] = 8 \times (1 - \frac{1}{2})$$

$$A[8] = 8 \times (1 - \frac{1}{2})$$

$$\phi(n) = n \cdot TT \left(1 - \frac{1}{p}\right)$$

$$Time complexity: p < n$$

$$O(n) + \frac{n}{2} + \frac{n}{3} + \frac{n}{5} + \frac{n}{7} + \cdots + \frac{n}{p}$$

$$= \Theta(n \log \log n).$$

$$n/1 + n/2 + n/3 + \cdots$$

"Sieve" method to find primes

Today Pari and Arya are playing a game called Remainders.

Pari chooses two positive integer x and k, and tells Arya k but not x. Arya have to find the value $x \mod k$. There are n ancient numbers $c_1, c_2, ..., c_n$ and Pari has to tell Arya $x \mod c_i$ if Arya wants Given k and the ancient values, tell us if Arya has a winning strategy independent of value of x or not. Formally, is it true that Arya can understand the value $x \mod k$ for any positive integer x?

Note, that $x \mod y$ means the remainder of x after dividing it by y.

Input

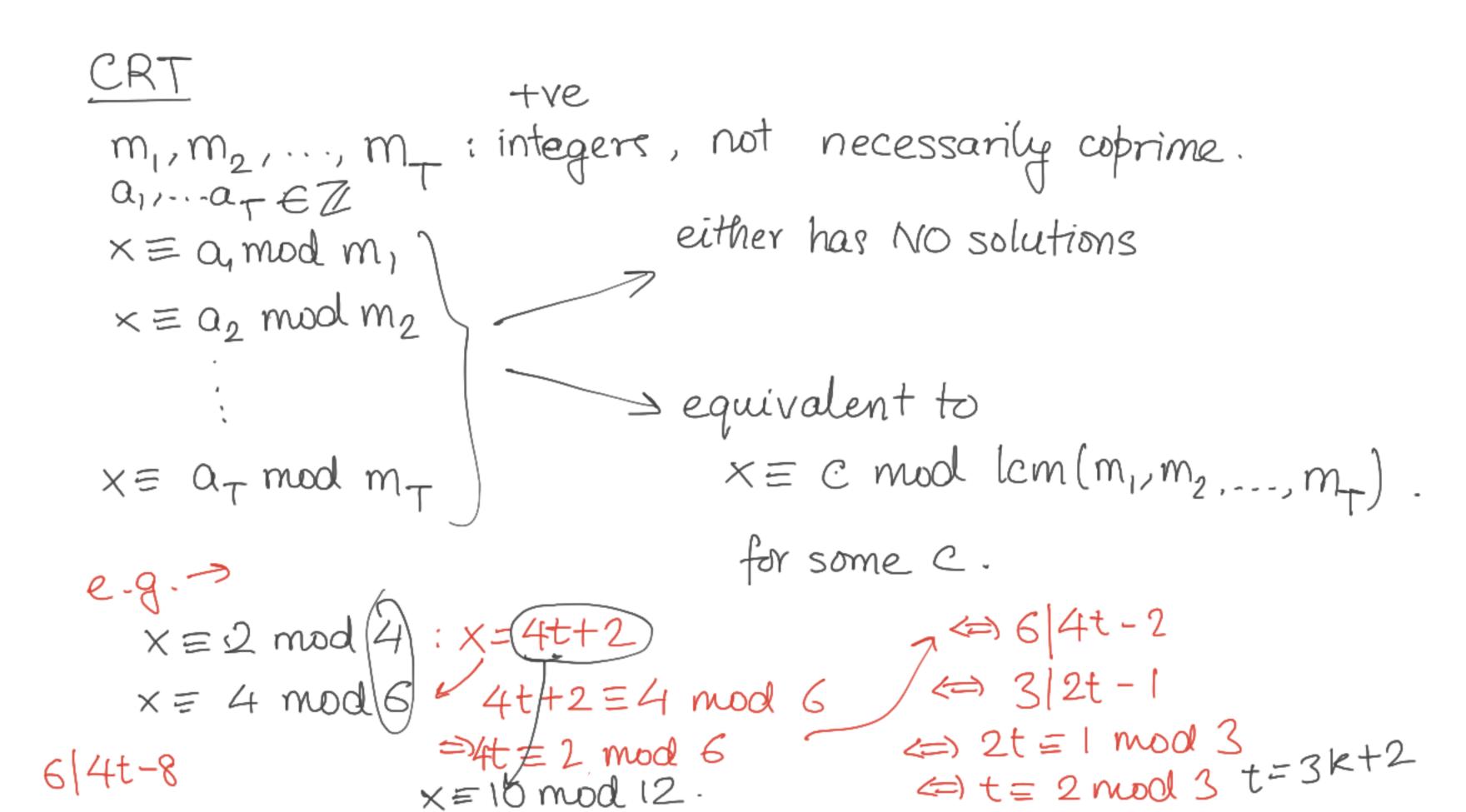
The first line of the input contains two integers n and k ($1 \le n$, $k \le 1000000$) — the number of ancient integers and value k that is chosen by Pari.

The second line contains n integers $c_1, c_2, ..., c_n$ ($1 \le c_i \le 1000000$).

Output

Print "Yes" (without quotes) if Arya has a winning strategy independent of value of x, or "No" (without quotes) otherwise.

what is known? $X \equiv a \mod lc_1 \cdot (C_1, ..., C_n)$ for some a. $X \equiv a \mod c_1$ $X \equiv a \mod c_2$ AND we know $A = a \mod a$ $A = a \mod a$ A



The Farey sequence of order n is the sequence of completely reduced fractions between 0 and 1 which, when in lowest terms, have denominators less than or equal to n, arranged in ascending order. Farey sequence for different values of n are shown in the figure on the left below:

$$F_1 = \left\{ \frac{0}{1}, \frac{1}{1} \right\}$$

$$F_4 = \left\{ \frac{0}{1}, \frac{1}{4}, \frac{1}{3}, \frac{1}{2}, \frac{2}{3}, \frac{3}{4}, \frac{1}{1} \right\}$$

$$F_7 = \left\{ \frac{0}{1}, \frac{1}{7}, \frac{1}{6}, \frac{1}{5}, \frac{1}{4}, \frac{2}{7}, \frac{1}{3}, \frac{2}{5}, \frac{3}{7}, \frac{1}{2}, \frac{4}{7}, \frac{3}{5}, \frac{2}{3}, \frac{5}{7}, \frac{3}{4}, \frac{4}{5}, \frac{5}{6}, \frac{6}{7}, \frac{1}{1} \right\}$$

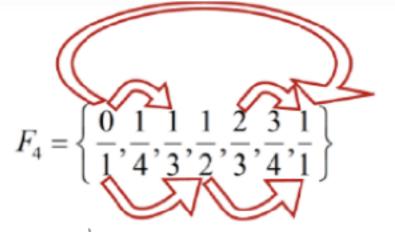


Figure 2: Five desired pairs in F4

It is very well known that if $\frac{m_1}{n_1}$ and $\frac{m_2}{n_2}$ and are two consecutive fractions of a Farey Sequence then $m_2n_1 - m_1n_2 = 1$. But many fractions which are not consecutive also show this property. For example, in F_7 , $\frac{2}{5}$ and $\frac{1}{2}$ also show this property although they are not consecutive fractions in F_7 . Given the value of n, your job is to find number of pair of non-consecutive fractions $\frac{m_i}{n_i}$ and $\frac{m_j}{n_i}$, such that $m_i n_i - m_i n_i = 1$.

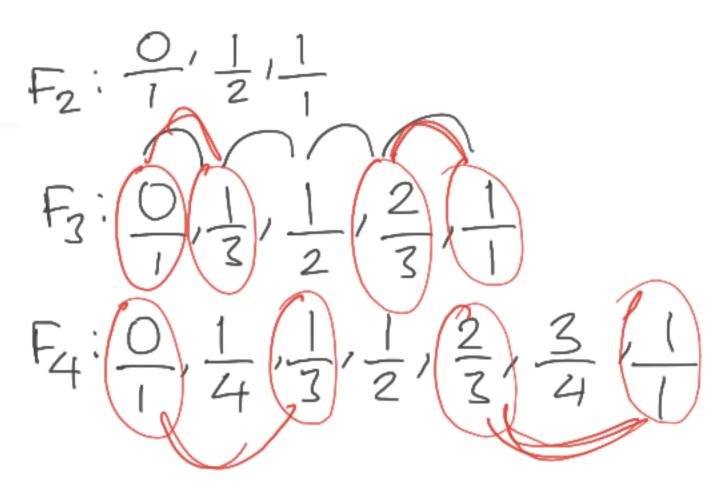
Input

Figure 1:

Input file contains at most 20000 lines of input. Each line contains a positive integer which denotes the value of n (0 < n < 1000001). Input is terminated by a line containing a single zero. This line should not be processed.

Output

For each line of input produce one line of output. This line contains number of pair of non-consecutive fractions $\frac{m_i}{n_i}$ and $\frac{m_j}{n_i}$, (j-i>1) in Farey Series F_n , such that $m_j n_i - m_i n_j = 1$.



Fn+1: $\frac{1}{n}$, ..., $\frac{1}{2}$, $\frac{n-1}{n}$, $\frac{1}{1}$

First fractions with denominator $\frac{a}{n+1}$, $\gcd(a,n+1)=1$

p(n+1) "new" fractions.

each new added fraction creates two non-consec fractions satisfying (*).

$$F_{2}: Q_{1} \xrightarrow{1} : Q_{1} : Q_{1} \xrightarrow{1} : Q_{1} : Q_{1} \xrightarrow{1} : Q_{1} :$$

$$(\phi(2) + \phi(3) + \phi(4) + ... + \phi(n)$$

 $\frac{m}{n}$, $\frac{b}{a}$ $\frac{qm-bn=1}{n}$

1) compute
$$\phi(1), \phi(2), \ldots, \phi(N)$$
 $O(N \log \log N)$
2) compute $\phi(2), \phi(2) + \phi(3), \ldots, \sum_{i=2}^{N} \phi(i)$ $O(N)$
3) For each test case n , output $A[n]$ $O(T)$

O(T+NloglogN)