浙江大学实验报告

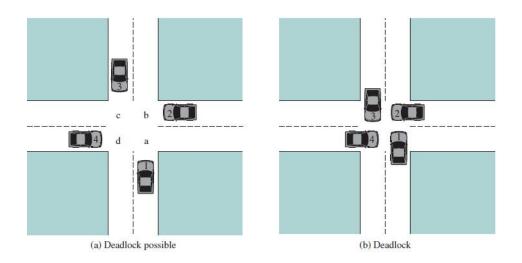
课程名称:	操作系统	实验类型:_	综合型/设计性
实验项目名称:	实验一 同步	互斥和 Linux 内核模	块
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一、实验目的

- 学习使用 Linux 的系统调用和 pthread 线程库编写程序。
- 充分理解对共享变量的访问需要原子操作。
- 进一步理解、掌握操作系统进程和线程概念,进程或线程的同步与互斥。
- 学习编写多线程程序,掌握解决多线程的同步与互斥问题。
- 学习 Linux 模块的实现机理,掌握如何编写 Linux 模块。
- 通过对 Linux 系统中进程的遍历,进一步理解操作系统进程概念和进程结构。

二、实验内容

- 1. 有两条道路双向两个车道,即每条路每个方向只有一个车道,两条道路十字交叉。假设车辆只能向前直行,而不允许转弯和后退。如果有4辆车几乎同时到达这个十字路口,如图 (a) 所示;相互交叉地停下来,如图 (b),此时4辆车都将不能继续向前,这是一个典型的死锁问题。从操作系统原理的资源分配观点,如果4辆车都想驶过十字路口,那么对资源的要求如下:
 - 向北行驶的车 1 需要象限 a 和 b;
 - 向西行驶的车 2 需要象限 b 和 c;
 - 向南行驶的车 3 需要象限 c 和 d;
 - 向东行驶的车 4 需要象限 d 和 a。



我们要实现十字路口交通的车辆同步问题,防止汽车在经过十字路口时产生死锁和饥饿。在我们的系统中,东西南北各个方向不断地有车辆经过十字路口(注意:不只有4辆),同一个方向的车辆依次排队通过十字路口。按照交通规则是右边车辆优先通行,如图(a)中,若只有car1、car2、car3,那么车辆通过十字路口的顺序是car3->car2->car1。车辆通行总的规则:

- 1) 来自同一个方向多个车辆到达十字路口时,车辆靠右行驶,依次顺序通过;
- 2) 有多个方向的车辆同时到达十字路口时,按照右边车辆优先通行规则,除非该车在十字 路口等待时收到一个立即通行的信号;
- 3) 避免产生死锁;
- 4) 避免产生饥饿:
- 5) 任何一个线程(车辆)不得采用单点调度策略;
- 6) 由于使用 AND 型信号量机制会使线程(车辆)并发度降低且引起不公平(部分线程饥饿),本题不得使用 AND 型信号量机制,即在上图中车辆不能要求同时满足两个象限才能顺利通过,如南方车辆不能同时判断 a 和 b 是否有空。

2.编写一个 Linux 的内核模块,其功能是遍历操作系统所有进程。该内核模块输出系统中:每个进程的名字、进程 pid、进程的状态、父进程的名字,以及统计系统中进程个数,包括统计系统中 TASK_RUNNING、TASK_INTERRUPTIBLE、TASK_UNINTERRUPTIBLE、TASK_ZOMBIE、TASK_STOPPED 等(还有其他状态)状态进程的个数。同时还需要编写一个用户态下执行的程序,显示内核模块输出的内容。

三、主要仪器设备(必填)

华硕 X450V, unbantu 操作系统, 虚拟机 2G 内存

四、操作方法和实验步骤

第一部分(具体代码见 deadlock. c 文件): 首先我分析了许多种情况,定义了一些变量和相应的锁程序中一共有两种锁

1: 通行锁(分为四个方向)

- 2: first position lock (即每个路口的第一个位置的占有锁)
- 一开始我会先输入所有需要通行的车辆,之后根据每辆车的方向,在每个方向上 为每辆车排序,比如北方第一辆车到北方第 N 辆车,和西方第一辆车到西方第 N 辆车,每辆车为一个线程

之后我分为以下几种情况:

当某个方向的车(以北方为例)来的时候,他会先等待并判断他的右手边(西边)有没有车,若右手边有车,则等待右手边的车先过,当西边的那辆车通过之后会判断多种情况来发信号:

- 1. 若西边没有车且北边还有车则只给北边发通行信号
- 2. 若西边有车且北边没有车则他会判断南边有没有车,若没有则给西边发通行信号
- 3. 若两边都有车则会判断 westToNorth 这个变量是否为 0, 若为 0 则发出让北边车辆通行的信号, 若为 1 则他会判断南边有没有车, 若没有则给西边发通行信号这样做是为了保证当西边有车通行时, 北边的车能够与西边的车轮流通行, 不导致饥饿

当一辆车通过时,他会给该方向下一辆车发出 first position lock 的信号,使该方向下一辆车占有 first position,当一辆车占有 first position之后,他就会等待放行信号

所以一般情况下一个方向的通信信号由其右手边发出,若其右手边没有车,则由 该方向前一辆车为其发出通行信号

我会维护一个 wait 变量来记录等待的车辆数, 当等待的车辆数到达 4 时, 即判断为死锁, 这时候就让北方的车先走。

第二部分:

我从 init_task 开始用 list_for_each(pos, &(initTask->tasks)) 遍历所有进程

```
list_for_each(pos, &(initTask->tasks))
{
    p=list_entry(pos, struct task_struct, tasks);
    total++;
    printk("%s pid:%d parent: %s parent's pid: %d ",p->comm,p->pid,p->real_parent->comm,p->real_parent->pid);

switch(p->state)
{
    case TASK_RUNNING:printk("state: RUNNING");running++;break;
    case TASK_INITERRUPTIBLE:printk("state: INTERRUPTIBLE");interruptible++;break;
    case TASK_UNINTERRUPTIBLE:printk("state: UNINTERRUPTIBLE");uninterruptible++;break;
    case TASK_TRACED:printk("state: TRACED");traced++;break;
    case TASK_STOPPED:printk("state: STOPPED");stopped++;break;
    case TASK_STOPPED:printk("state: STOPPED");stopped++;break;
    case EXIT_ZOMBIE:zombie++;printk("exit_state: ZOMBIE");break;
    default:printk("state: UNKNOWN");unknown++;break;
}

printk("\n");
}
```

之后再输出到 kern.log 文件中。最后由 printk.c(由我自己编写)输出到控制台

五、实验结果和分析

第一部分:

```
anthony@anthony-virtual-machine:~$ ./deadlock
wesnnswensnwe
the car 1 from North arrives at intersection
the car 1 from South arrives at intersection
the car 1 from East arrives at intersection
the car 1 from West arrives at intersection
DEADLOCK FOUND, the car from north go first
the car 1 from North have left, which is the 1 car
the car 1 from East have left, which is the 2 car
the car 1 from South have left, which is the 3 car
the car 1 from West have left, which is the 4 car
the car 2 from South arrives at intersection
the car 2 from East arrives at intersection
the car 2 from North arrives at intersection
the car 2 from West arrives at intersection
DEADLOCK FOUND, the car from north go first
the car 2 from North have left, which is the 5 car
the car 2 from East have left, which is the 6 car
the car 2 from South have left, which is the 7 car
the car 2 from West have left, which is the 8 car
the car 3 from North arrives at intersection
the car 3 from South arrives at intersection
the car 3 from West arrives at intersection
the car 3 from East arrives at intersection
DEADLOCK FOUND, the car from north go first
the car 3 from North have left, which is the 9 car
the car 3 from East have left, which is the 10 car
the car 3 from South have left, which is the 11 car
the car 3 from West have left, which is the 12 car
the car 4 from North arrives at intersection
the car 4 from North have left, which is the 13 car
anthony@anthony-virtual-machine:~$
```

```
anthony@anthony-virtual-machine:~$ ./deadlock
weswesw
the car 1 from South arrives at intersection
       1 from East arrives at intersection
the car
the car 1 from East have left, which is the 1 car
the car 1 from South have left, which is the 2 car
the car 1 from West arrives at intersection
       2 from South arrives at intersection
the car
the car 2 from East arrives at intersection
the car 2 from East have left, which is the 3 car
       2 from South have left, which is the 4 car
the car 1 from West have left, which is the 5 car
the car 2 from West arrives at intersection
the car 2 from West have left, which is the 6 car
the car 3 from West arrives at intersection
the car 3 from West have left, which is the 7 car
```

第二部分:

```
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875760] update-notifier pid:5618 parent: gnome-session-b parent's pid: 5111 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875763] gvfsd-dnssd pid:5632 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875765] notify-osd pid:5655 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875768] deja-dup-monito pid:5735 parent: gnome-session-b parent's pid: 14 13:03:48 anthony-virtual-machine kernel: [7589.875776] kworker/0:0 pid:6549 parent: kthreadd parent's pid: 2 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875776] dhclient pid:6774 parent: kthreadd parent's pid: 2 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875776] dhclient pid:6774 parent: NetworkManager parent's pid: 4194 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875778] kworker/0:2 pid:6873 parent: kthreadd parent's pid: 2 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875781] update-manager pid:7121 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875783] unity-control-c pid:7121 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875786] unity-scope-loa pid:7200 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875786] unity-scope-loa pid:7210 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875789] unity-scope-loa pid:7210 parent: upstart parent's pid: 4902 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875789] winty-scope-loa pid:7210 parent: kthreadd parent's pid: 2 state: INTERRUPTIBLE
Dec 14 13:03:48 anthony-virtual-machine kernel: [7589.875796] kworker/0:1 pid:7209 parent:
```

六、讨论、心得

在解决饥饿的问题上让我思考了很长时间,从开始到结束一共想了很多种方案,直到最后才得出一个令人满意的答案。一开始采用睡眠的方法,之后发现这样并不是真正的解决问题,之后就采用一个变量来记录两边的通行情况,来使两边能够实现轮流通行,最后才解决了饥饿的问题。

在第二个部分中,我了解了进程的状态,搜了比较久的资料,终于把所有进程的状态和父进程之类的输出了