## 1 Best and Worst Case

For the following functions, provide asymptotic bounds for the best case and worst case runtimes in  $\Theta(\cdot)$  notation.

(a) Give the best and worst case runtimes in terms of M and N. Assume that slam()  $\in \Theta(1)$  and returns a boolean.

```
public void comeon(int M, int N) {
       int j = 0;
2
       for (int i = 0; i < N; i += 1) {</pre>
            for (; j < M; j += 1) {
4
                if (slam(i, j))
                    break;
6
            }
       }
8
       for (int k = 0; k < 1000 * N; k += 1) {
10
            System.out.println("space jam");
11
12
13 }
```

(b) Extra: Give the best case and worst case runtimes for find in terms of N, where N is the length of the input array arr.

```
public static boolean find(int tgt, int[] arr) {
       int N = arr.length;
       return find(tgt, arr, 0, N);
4 }
5 private static boolean find(int tgt, int[] arr, int lo, int hi) {
       if (lo == hi || lo + 1 == hi) {
           return arr[lo] == tgt;
7
8
       int mid = (lo + hi) / 2;
9
       for (int i = 0; i < mid; i += 1) {</pre>
10
11
           System.out.println(arr[i]);
12
       return arr[mid] == tgt || find(tgt, arr, lo, mid)
13
                               || find(tgt, arr, mid, hi);
14
15 }
```

## 2 Best and Worst Case with Recursion

For the following recursive functions, provide asymptotic bounds for the best case and worst case runtimes in  $\Theta(\cdot)$  notation.

(a) Give the runtime in terms of N.

(b) Give the runtime for andwelcome (arr, 0, N) in terms of N, where N is the length of the input array arr. Math.random() returns a double with a value from the range [0,1).

```
public static void andwelcome(int[] arr, int low, int high) {
       System.out.print("[ ");
       for (int i = low; i < high; i += 1) {</pre>
3
            System.out.print("loyal ");
4
5
       System.out.println("]");
6
       if (high - low > 1) {
7
           double coin = Math.random();
9
           if (coin > 0.5) {
               andwelcome(arr, low, low + (high - low) / 2);
10
           } else {
11
               andwelcome(arr, low, low + (high - low) / 2);
               andwelcome(arr, low + (high - low) / 2, high);
13
14
           }
15
       }
```

(c) Give the runtime in terms of N.

```
public int tothe(int N) {
    if (N <= 1) {
        return N;
    }
    return tothe(N - 1) + tothe(N - 1) + tothe(N - 1);
}</pre>
```

(d) *Extra*: Give the runtime in terms of N.

```
public static void spacejam(int N) {
    if (N == 1) {
        return;
}

for (int i = 0; i < N; i += 1) {
        spacejam(N-1);
}

}</pre>
```

## 3 Hey you watchu gon do?

For each example below, there are two algorithms solving the same problem. Given the asymptotic runtimes for each, is one of the algorithms **guaranteed** to be faster? If so, which? And if neither is always faster, explain why. Assume the algorithms have very large input (i.e. *N* is very large).

- (a) Algorithm 1:  $\Theta(N)$ , Algorithm 2:  $\Theta(N^2)$
- (b) Algorithm 1:  $\Omega(N)$ , Algorithm 2:  $\Omega(N^2)$
- (c) Algorithm 1: O(N), Algorithm 2:  $O(N^2)$
- (d) Algorithm 1:  $\Theta(N^2)$ , Algorithm 2:  $O(\log N)$
- (e) Algorithm 1:  $O(N \log N)$ , Algorithm 2:  $\Omega(N \log N)$

Why do we need to assume that N is large?