

## Logic Exercises 3

- Which of the following five formulas are in disjunctive normal form (DNF), and which ones are in conjunctive normal form (CNF)?

$$p \wedge \neg(q \vee r), \quad p \vee q \vee \neg r, \quad p \wedge q \wedge \neg r, \quad p \vee (q \wedge \neg r), \quad p \wedge (q \vee \neg r)$$

- Construct formulas in DNF and in CNF that are semantically equivalent to the formula  $\phi$ , based on its truth table:

$p$	$q$	$r$	$\phi$
T	T	T	T
T	T	F	F
T	F	T	F
T	F	F	F
F	T	T	T
F	T	F	F
F	F	T	T
F	F	F	F

- Express the formula  $p \rightarrow q$  using only the Sheffer stroke  $|$  (and  $p$  and  $q$ ).
  - Express the formula  $p \leftrightarrow q$  using only the Sheffer stroke  $|$  (and  $p$  and  $q$ ).
  - Build a tautology and a contradiction using only the Sheffer stroke (and a propositional variable).
- Show that  $\{\neg, \rightarrow\}$  is an adequate system of connectives (meaning that every truth table can be expressed using a formula built from only these connectives and some propositional variables).
  - Argue that  $\{\wedge, \vee, \rightarrow\}$  is not an adequate system of connectives.
  - Is  $\{\neg, \leftrightarrow\}$  an adequate system of connectives?
- Determine a CNF of the formula  $(\neg q \wedge r) \rightarrow (p \wedge \neg q)$  by applying the algorithm CNF.
  - Do the same for the formula  $\neg\neg p \rightarrow \neg(p \rightarrow \neg p)$ .
  - Apply the criterion for determining whether a CNF is a tautology (which was treated in the lecture) to the CNFs you constructed in (a) and (b).
  - Formulate an analogous criterion to determine whether a DNF is a contradiction.

6. Is a DNF a contradiction if and only if each disjunct contains literals  $p$  and  $\neg p$  for some  $p$ ? If so, argue that this is the case. If not, give a counterexample.

7. Apply the DPLL procedure to the following CNFs, to check if they are satisfiable.

(a)  $p \wedge \neg p$

(b)  $p \vee \neg p$

(c)  $(p \vee q) \wedge (\neg p \vee q) \wedge (p \vee \neg q)$

(d)  $(p \vee q \vee \neg r) \wedge (\neg p \vee q \vee \neg r) \wedge (p \vee \neg q \vee \neg r) \wedge r$

(e)  $(p \vee q) \wedge (\neg p \vee q) \wedge (p \vee \neg q) \wedge (\neg p \vee \neg q)$

(f)  $(p \vee q \vee r) \wedge (\neg p \vee q \vee r) \wedge (p \vee \neg q \vee r) \wedge (\neg p \vee \neg q \vee r)$