

Lecture 13 — OpenMP

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So far: Pthreads and automatic parallelization.

Next: “Manual” parallelization using OpenMP (Open Multiprocessing).

You specify parallelization; compiler implements.



“I want 10 turnstiles.”

You use OpenMP by specifying **directives** in the code.

Directives look like `#pragma omp ...`

There are 16 directives; each applies to the immediately-following statement.

Some have a list as an argument (a list of variable names)

All major compilers have OpenMP: (gcc, clang, Solaris, Intel, Microsoft).

Benefits:

- Switch parallelism on and off via compile option
- Make the compiler do hard work
- Makes algorithms easier to read
- Apply them only in parts where needed, incrementally

Simple For Loop Example

```
void calc (double *array1, double *array2, int length) {  
    #pragma omp parallel for  
    for (int i = 0; i < length; i++) {  
        array1[i] += array2[i];  
    }  
}
```

The directive instructs the compiler to parallelize the loop.

Developer's responsibility to make sure this is safe!

You don't need to declare `restrict`, but it's a good idea.

#pragma omp parallel for

Let's look at the parts of this #pragma.



Source: Feuerwehrmagazin.de

- `#pragma omp` indicates an OpenMP directive;
- `parallel` indicates the start of a parallel region.
- `for` tells OpenMP: run the next for loop in parallel.

The runtime library starts a number of threads & assigns a subrange of the range to each of the threads.

```
for (int i = 0; i < length; i++) { ... }
```

Can only parallelize loops which satisfy these conditions:

- form: `for (init expr; test expr; increment expr);`
- loop variable must be integer (signed or unsigned), pointer, or a C++ random access iterator;
- loop variable must be initialized to one end of the range;
- loop increment amount must be loop-invariant (constant with respect to the loop body); and
- test expression must be one of `>`, `>=`, `<`, or `<=`, and the comparison value (bound) must be loop-invariant.

Note: these restrictions thus also apply to automatically parallelized loops.

Get together a team!



- 1 Compiler generates code to spawn a **team** of threads; automatically splits off worker-thread code into a separate procedure.
- 2 Generated code uses fork-join parallelism; when the master thread hits a parallel region, it gives work to the worker threads, which execute and report back.
- 3 Afterwards, the master thread continues running, while the worker threads wait for more work .

You can specify the number of threads by setting the `OMP_NUM_THREADS` environment variable.

(You can also adjust by calling `omp_set_num_threads ()`).

- Solaris compiler tells you what it did if you use the flags `-xopenmp` `-xloopinfo`, or `er_src`.

OpenMP includes three keywords for variable scope and storage:

- private
- shared
- threadprivate

Private: new storage, on a per-thread basis, for the variable.

The scope from the start of the region to the end of that region.

The variable is destroyed afterwards.

Thread equivalent:

```
void* run(void* arg) {  
    int x;  
    // use x  
}
```

Shared: the opposite of a private variable.

All threads have access to the same block of data when accessing such a variable.

Thread equivalent:

```
int x;  
  
void* run(void* arg) {  
    // use x  
}
```

Threadprivate: This is like private, in that each thread makes a copy of the variable.

However, the scope is different.

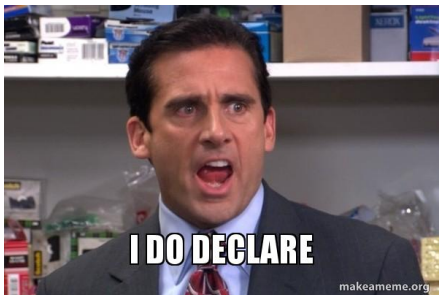
Such variables are accessible to the thread in any parallel region.

This example will make things clearer. The OpenMP code:

```
int x;  
#pragma omp threadprivate(x)
```

maps to this Pthread pseudocode:

```
int x;  
int x[NUM_THREADS];  
  
void* run(void* arg) {  
    // int* id = (int*) arg;  
    // use x[*id]  
}
```



A variable may not appear in **more than one clause** on the same directive.
(There's an exception for `firstprivate` and `lastprivate`)

Default: vars declared in regions are private; those outside regions are shared.

(An exception: anything with dynamic storage is shared).

Let's look at the defaults that OpenMP uses to parallelize the `parallel - for` code:

```
% er_src parallel-for.o
1.  <Function: calc>
```

Source OpenMP region below has tag R1

Private variables in R1: i

Shared variables in R1: array2, length, array1

```
2.      #pragma omp parallel for
```

Source loop below has tag L1

L1 autoparallelized

L1 parallelized by explicit user directive

L1 parallel loop-body code placed in function `_$d1A2.calc`
along with 0 inner loops

L1 multi-versioned for loop-improvement:
dynamic-alias-disambiguation.

Specialized version is L2

```
3.      for (int i = 0; i < length; i++) {
4.          array1[i] += array2[i];
5.      }
6.  }
```

You can disable the default rules by specifying `default(none)` on the `parallel`, or you can give explicit scoping:

```
#pragma omp parallel for private(i)  
                        shared(length, array1, array2)
```



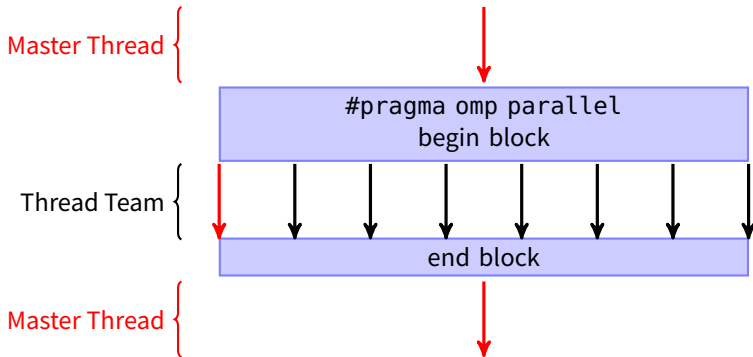
Oh.... Wait... Directives, not Directors.

```
#pragma omp parallel [clause [, clause]*]
```

This is the most basic directive in OpenMP. It tells OpenMP to form a team of threads and start parallel execution.

The thread that enters the region becomes the **master** (thread 0).

Allowed Clauses: **if**, **num_threads**, **default**, **private**, **firstprivate**, **shared**, **copyin**, **reduction**.



```
#pragma omp parallel
{
    printf("Hello!");
}
```

If the number of threads is 4, this produces:

```
Hello!
Hello!
Hello!
Hello!
```

if(*primitive-expression*)

The `if` clause allows you to control at runtime whether or not to run multiple threads or just one thread in its associated parallel section.

If *primitive-expression* evaluates to 0, i.e. `false`, then only one thread will execute in the parallel section.

(It's what you would expect.)

If the parallel section is going to run multiple threads (e.g. `if` expression is true; or if there is no `if` expression), we can then specify how many threads.

num_threads(*integer-expression*)

This spawns at most **num_threads**, depending on the number of threads available.

It can only guarantee the number of threads requested if **dynamic adjustment** for number of threads is off and enough threads aren't busy.

Recall that we introduced the concept of a reduction, e.g.

```
for (int i = 0; i < length; i++) total += array[i];
```

What is the appropriate scope for total?

We want each thread to be able to write to it, but we don't want race conditions.

Fortunately, OpenMP can deal with reductions as a special case:

```
#pragma omp parallel for reduction (+:total)
```


The `total` variable is the accumulator for a reduction over `+`.

OpenMP will create local copies of `total` and combine them at the end of the parallel region.

Operators (and their associated initial value)

<code>+</code>	<code>(0)</code>	<code>-</code>	<code>(0)</code>	<code> </code>	<code>(0)</code>	<code>&&</code>	<code>(1)</code>	<code>max</code>	<code>MAX</code>
<code>*</code>	<code>(1)</code>	<code>&</code>	<code>(~0)</code>	<code>^</code>	<code>(0)</code>	<code> </code>	<code>(0)</code>	<code>min</code>	<code>MIN</code>

Inside a parallel section, we can specify a for loop clause, as follows.

```
#pragma omp for [clause [,] clause]*]
```

Iterations of the loop will be distributed among the current team of threads.

Loop variable is implicitly private; OpenMP sets the correct values.

Allowed Clauses: **private**, **firstprivate**, **lastprivate**, **reduction**, **schedule**, **collapse**, **ordered**, **nowait**.

schedule(*kind*[, *chunk_size*])

The **chunk_size** is the number of iterations a thread should handle at a time.

kind is one of:

- **static**
- **dynamic**
- **guided**
- **auto**
- **runtime**

OpenMP tries to guess how many iterations to distribute to each thread in a team.

The default of equal number of iterations isn't always optimal:

```
double calc(int count) {
    double d = 1.0;
    for (int i = 0; i < count*count; i++) d += d;
    return d;
}

int main() {
    double data[200][200];
    int i, j;
    #pragma omp parallel for private(i, j) shared(data)
    for (int i = 0; i < 200; i++) {
        for (int j = 0; j < 200; j++) {
            data[i][j] = calc(i+j);
        }
    }
    return 0;
}
```

Telling OpenMP to use a *dynamic schedule* can enable better parallelization.

`collapse(n)`

This collapses n levels of loops.

Obviously, this only has an effect if $n \geq 2$; otherwise, nothing happens.

Collapsed loop variables are also made private.

```
#pragma omp ordered
```

To use this directive, the containing loop must have an **ordered** clause.

OpenMP will ensure that the ordered directives are executed the same way the sequential loop would (one at a time).

Each iteration of the loop may execute **at most one** ordered directive.

This use of Ordered does not work:

```
void work(int i) {  
    printf("i = %d\n", i);  
}  
  
...  
int i;  
#pragma omp for ordered  
for (i = 0; i < 20; ++i) {  
  
    #pragma omp ordered  
    work(i);  
  
    // Each iteration of the loop has 2 "ordered" clauses!  
    #pragma omp ordered  
    work(i + 100);  
}
```

Here is a valid use of the ordered clause:

```
void work(int i) {
    printf("i = %d\n", i);
}
...
int i;
#pragma omp for ordered
for (i = 0; i < 20; ++i) {
    if (i <= 10) {
        #pragma omp ordered
        work(i);
    }
    if (i > 10) {
        // two ordered clauses are mutually-exclusive
        #pragma omp ordered
        work(i+100);
    }
}
```

Note: if we change $i > 10$ to $i > 9$, the use becomes invalid because the iteration $i = 9$ contains two ordered directives.

```
#include <omp.h>
#include <stdio.h>

int main(int argc, char *argv[])
{
    int j, k, a;
    #pragma omp parallel num_threads(2)
    {
        #pragma omp for collapse(2) ordered private(j,k) \
            schedule(static,3)
        for (k = 1; k <= 3; ++k)
            for (j = 1; j <= 2; ++j) {
                #pragma omp ordered
                printf("t[%d]_k=%d_j=%d\n",
                    omp_get_thread_num(),
                    k, j);
            }
    }
    return 0;
}
```

Output of Larger Example

t [0]	k=1	j =1
t [0]	k=1	j =2
t [0]	k=2	j =1
t [1]	k=2	j =2
t [1]	k=3	j =1
t [1]	k=3	j =2

Note: the output will be deterministic; still, the program will run two threads as long as the thread limit is at least 2.

This directive is shorthand:

```
#pragma omp parallel for [clause [,] clause]*]
```

We could equally well write:

```
#pragma omp parallel  
{  
    #pragma omp for  
    {  
    }  
}
```

Allowed Clauses: everything allowed by `parallel` and `for`, except **nowait**.

```
#pragma omp parallel sections [clause [[, clause]*]
```

As with parallel for, this is basically shorthand for:

```
#pragma omp parallel  
{  
    #pragma omp sections  
    {  
    }  
}
```

Allowed Clauses: everything allowed by parallel and sections, except **nowait**.

When we'd be otherwise running many threads, we can state that some particular code block should be run by only one method.

```
#pragma omp single
```

Only a single thread executes the region following the directive. The thread is not guaranteed to be the master thread.

Allowed Clauses: **private**, **firstprivate**, **copyprivate**, **nowait**. Must not use **copyprivate** with **nowait**

Sometimes we do want to guarantee that only the master thread runs some code.

```
#pragma omp master
```

This is similar to the **single** directive, except that the master thread (and only the master thread) is guaranteed to enter this region. No implied barriers.

Also, no clauses.

```
#pragma omp barrier
```

Waits for all the threads in the team to reach the barrier before continuing.

Loops, sections, and single all have an implicit barrier at the end of their region (unless you use **nowait**).

A barrier inside an `if` statement must be `{ }` separated.

This mechanism is analogous to `pthread_barrier` in `pthread`s.


```
#pragma omp critical [(name)]
```

The enclosed region is guaranteed to only run one thread at a time (on a per-name basis).

This is the same as a block of code in pthreads surrounded by a mutex lock and unlock.

We can also request atomic operations:

```
#pragma omp atomic [read | write | update | capture]  
expression-stmt
```

This ensures that a specific storage location is updated atomically.

Atomics should be more efficient than using critical sections (or else why would they include it?)

Write your code so that simply eliding OpenMP directives gives a valid program. For instance, this is wrong:

```
if (a != 0)
    #pragma omp barrier // wrong!
if (a != 0)
    #pragma omp taskyield // wrong!
```

Use this instead:

```
if (a != 0) {
    #pragma omp barrier
}
if (a != 0) {
    #pragma omp taskyield
}
```
