

Computer Systems B COMS20012

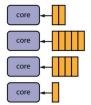
Introduction to Operating Systems and Security

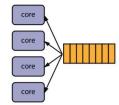


Two possible approaches

- Per-core queues
- Shared queue

Which scales better?
Which offers better performance?





Scalability

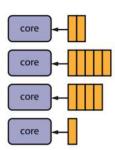
- Contention and scalability
 - Access to the shared queue is a critical section, mutual exclusion is needed
 - As the number of cores grow, contention to access the resource become a problem
- Per-core design scale to many core

Performance

- As a thread run, the data it accesses is loaded in the CPU cache(s)
- If thread changes CPU, the data needs to be loaded again
- Affinity between a thread and one core grows because the cached data
- Per-core approach benefit from this affinity by keeping threads on the same core
- This is not the case for shared queue design

Load balancing

- In per-core design, queues may have different length
- This create load imbalance across cores
 - Cores may be idle why others are busy
 - Threads on lightly loaded core have more CPU time than threads on heavy loaded ones
- Non-issue in shared queue design
- Per-core designs need a mechanism to balance the load across all cores
 - Migrate threads from lightly loaded core to heavy loaded ones





- Pause the video
- Open *kern/thread/thread.c*
- Resume the video





kern/thread/thread.c

- thread_consider_migration defined line 881
- Called from hardclock function (kern/thread/clock.c line 93)
 - Watch Week 7 Video 2 again if confused
- Line 889-901
 - Calculate average load per core (one_share)
 - Load on current core (my_count)
- Line 902-904
 - If bellow average load nothing to do
- Line 906-913

 - Build a list of thread to migrate (victims)
 Remove them from the thread running on the current core (&curcpu->c_runqueue)

kern/thread/thread.c

- Line 915-966
 - Find a core bellow average load (line 921)
 - Migrate victim threads to that core
- Line 973-982
 - Internal management
 - Put back thread that could not be assigned
 - Clean up data structure

