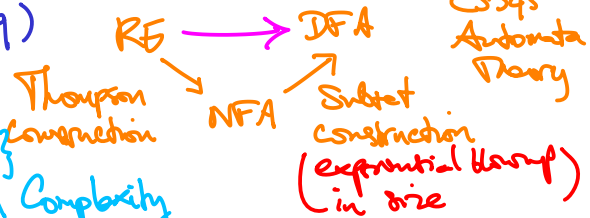


Test 1: 1) Knuth-Morris-Pratt (Ch 32 from Cormen et al.)

Topics 2) Aho-Corasick (see paper + Exercises in Ch 3 Dragon 1)

3) RegExp 2 DFA (Ch 3 Dragon 1, §3.9)  
(no NFA)

4) A proof of why  $L = \{ww : w \in \{0,1\}^*\}$   
is not regular using Communication Complexity



## RegExp to DFA: §3.9, p135

Example:  $R = (a|b)^*abb\# \rightarrow$  DFA?

root 1 2 3 4 5 6

i)  $\text{firstpos}(u)$  = set of positions that can match the 1st symbol of a string generated by the subexpression rooted at  $u$

ii)  $\text{lastpos}(u)$  = ditto but for last symbol

ex:  $\text{firstpos}(F) = \{1, 2\}$

ex:  $\text{lastpos}(F) = \{1, 2\}$

$\text{firstpos}(D) = \{1, 2, 3\}$

$\text{lastpos}(D) = \{3\}$

$\text{firstpos}(A) = \{1, 2, 3\}$

$\text{lastpos}(A) = \{6\}$

iii)  $\text{nullable}(u) = \begin{cases} \text{True} & \text{if } \epsilon \in L(\text{RegExp}_u) \\ \text{False} & \text{else} \end{cases}$

iv)  $\text{followpos}(i) =$  set of positions that can follow  $i$  in syntax tree

$(a|b)^*abb$

ex:

i	Followpos
1	1, 2, 3
2	1, 2, 3
3	4
4	5
5	6

DFA construction: p141 Algorithm 3.5

DStates  $\leftarrow \{\text{firstpos}(\text{root})\}$

While  $\exists$  unmarked state  $T$  in DStates:

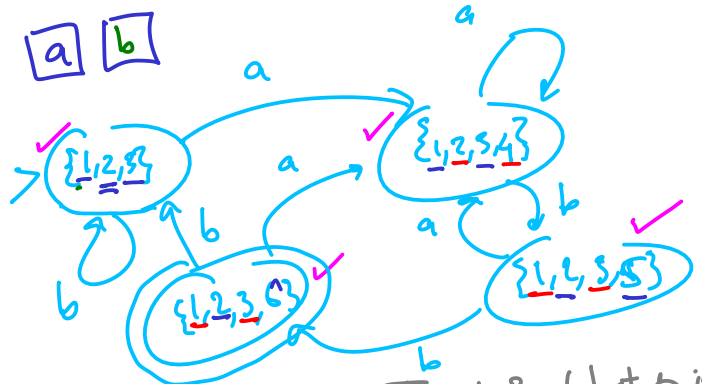
mark  $T$

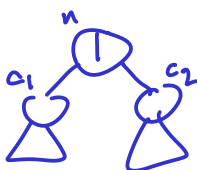
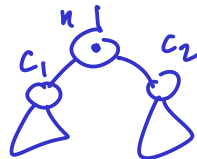

for each  $a \in \Sigma$ :

$u$  is the set of positions in  $\text{followpos}(p)$  for some  $p \in T$  s.t. symbol at  $p$  is  $a$

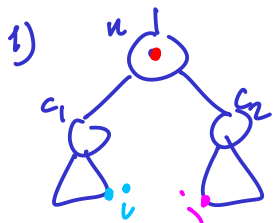
if  $u \neq \emptyset$  and  $u \notin \text{DStates}$ : add  $u$  to DStates

DTrans  $[T, a] = u$



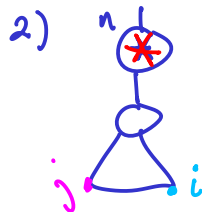
Node $u$	$nullable(u)$	$firstpos(u)$	$lastpos(u)$
<u>P138</u> $a \in \Sigma$ $\begin{array}{c} n \\ \boxed{a} \\ i \end{array}$	true	$\emptyset$	$\emptyset$
$\begin{array}{c} n \\ \boxed{a} \\ i \end{array}$	false	$\{i\}$	$\{i\}$
	$nullable(c_1)$ or $nullable(c_2)$	$firstpos(c_1) \cup firstpos(c_2)$	$lastpos(c_1) \cup lastpos(c_2)$
	$nullable(c_1)$ and $nullable(c_2)$	$\begin{cases} firstpos(c_1) & \text{if } \overline{nullable(c_1)} \\ firstpos(c_1) \cup firstpos(c_2) & \text{if } nullable(c_1) \end{cases}$	$\begin{cases} lastpos(c_2) & \text{if } \overline{nullable(c_2)} \\ lastpos(c_1) \cup lastpos(c_2) & \text{if } nullable(c_2) \end{cases}$
	true	$firstpos(c)$	$lastpos(c)$

$followpos(i)$  :



$i \in lastpos(c_1) \Rightarrow$

$firstpos(c_2) \subseteq followpos(i)$



$i \in lastpos(n) \Rightarrow$

$firstpos(n) \subseteq followpos(i)$