cubical evaluation semantics

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I've been thinking a bit more carefully about the evaluation semantics and the semantic domain. What I think we want to do is take a term and an environment and return a family of values (indexed in restrictions); but crucially, I think it is important that the environment be a vector of *values* and not of value-families. By studying carefully the case for instantiating closures (where the rubber meets the road for variables), I think it becomes clear that this is the right way.

To interpret a variable, we just project from the environment and ignore the restriction. I think the way that we had been thinking of doing it would result in the restriction from inside the closure infecting the argument.

Colors: terms, value families, values, restrictions. I am writing ϕ^*F to mean the reindexing of a family by a restriction.

$$\begin{split}
& \left[\left(\lambda \{ M \} \right) \right]_{\rho}^{\pi} = \left[\phi \right] \lambda \langle M, \rho, \pi, \phi \rangle \\
& \left[\left(\Pi A \{ B \} \right) \right]_{\rho}^{\pi} = \left[\phi \right] \Pi \left(\phi^* \left[A \right]_{\rho}^{\pi}, \langle B, \rho, \pi, \phi \rangle \right) \\
& \left[\left(M N \right) \right]_{\rho}^{\pi} = \underline{\operatorname{app}}(\left[M \right]_{\rho}^{\pi}, \left[N \right]_{\rho}^{\pi}) \\
& \left[x_i \right]_{\rho}^{\pi} = \left[\right] \rho_i
\end{split}$$

We write $\phi \models F \parallel V$ as a shorthand for $F[\phi] \equiv V$.

$$\frac{\phi \models F \Downarrow \lambda\langle M, \rho, \pi, \psi \rangle}{\phi \models \underline{app}(F, N) \Downarrow \llbracket M \rrbracket_{[\rho, N[\phi]]}^{\pi}[\psi]} \qquad \frac{\phi \models F \Downarrow \uparrow^{C}R \qquad \vdash C \Downarrow \Pi(A, \langle B, \rho, \pi, \psi \rangle)}{\phi \models \underline{app}(F, N) \Downarrow \uparrow^{\psi^{*}\llbracket B \rrbracket_{[\rho, N[\phi]]}^{\pi}([\chi] app(\chi^{*}R, \chi \circ \phi^{*}N))}$$

$$\frac{\phi \models F \Downarrow \mathbf{coe}_{(\alpha)C}^{r \to r'} M \qquad \vdash E C \Downarrow \Pi(A, \langle B, \rho, \pi, \psi \rangle)}{\phi \models \underline{app}(F, N) \Downarrow \nu \beta. \nu \gamma. \underline{\mathbf{coe}}_{(\beta) \llbracket B \rrbracket_{[\rho, coe]}^{[r, \beta, \phi]}} \underbrace{\mathbf{app}(M, \underline{\mathbf{coe}}_{(\gamma)(\gamma \alpha) \cdot A}^{r' \to r} M \circ \psi^{*})} \qquad \frac{\phi \models F \Downarrow \mathbf{hcom}_{C}^{r \to r'} M \left[\xi_{i} \hookrightarrow \langle \alpha_{i} \rangle N_{i}\right] \qquad \vdash E C \Downarrow \Pi(A, \langle B, \rho, \pi, \psi \rangle)}{\phi \models \underline{app}(F, N) \Downarrow \nu \beta. \underline{\mathbf{hcom}}_{\psi^{*} \llbracket B \rrbracket_{[\rho, coe]}^{\pi}} \underbrace{\mathbf{app}(M, \phi^{*}O) \left[\xi_{i} \hookrightarrow \langle \beta \rangle \underline{\mathbf{app}}((\alpha_{i} \beta) \cdot N_{i}, \phi^{*}O)\right]}}{\phi \models \underline{\mathbf{app}}(F, N) \Downarrow \nu \beta. \underline{\mathbf{hcom}}_{\psi^{*} \llbracket B \rrbracket_{[\rho, coe]}^{\pi}} \underbrace{\mathbf{app}(M, \phi^{*}O) \left[\xi_{i} \hookrightarrow \langle \beta \rangle \underline{\mathbf{app}}((\alpha_{i} \beta) \cdot N_{i}, \phi^{*}O)\right]}}$$

$$\begin{array}{c} \cdot \vDash r = r' \\ \hline \text{hcom}_{\mathbb{C}}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \equiv M[\cdot] \\ \hline \text{hcom}_{\mathbb{C}}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \equiv N_k [\alpha_k = r'[\cdot]] \\ \hline \text{hcom}_{\mathbb{C}}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \equiv hcom_{\mathbb{C}}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \equiv hcom_{\mathbb{C}}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \not \vdash \overline{\xi}_i \qquad \cdot \vDash \mathbb{C} \stackrel{1}{\downarrow} \text{bool} \\ \hline \text{hcom}_{\mathbb{C}}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \equiv M[\cdot] \\ \hline \phi \vDash hcom_{A}^{r \rightarrow r'} M \left[\overline{\xi}_i \hookrightarrow \langle \alpha_i \rangle N_i \right] \implies \nu \beta, hcom_{\Phi^{+A}}^{\phi^{+}r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r = r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{\phi^{+}r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \rightarrow \Gamma \Rightarrow r \Rightarrow r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \vDash r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \rightarrow \Gamma \Rightarrow gcom_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \Leftrightarrow r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline{\phi}^{*} \xi_i \hookrightarrow \langle \beta \rangle \phi^{*}((\beta | \alpha_i) \circ N_i) \right] \\ \hline \cdot \Leftrightarrow r \neq r' \qquad \cdot \vDash \mathcal{C} \stackrel{1}{\downarrow} \text{fcom}_{\Phi^{+A}}^{r \rightarrow r'} \phi^{*} M \left[\overline$$