

Efficiency of Algorithms

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Objectives

- At the end of this lecture, student will be able to
 - Explain the terms 'efficiency' and 'complexity'
 - Calculate efficiency of an algorithm
 - Express the complexity of algorithms in asymptotic notation
 - Classify algorithms based on their complexity



Contents

- Efficiency and Complexity
- Time Complexity of an algorithm
- Space Complexity of an algorithm
- Measuring complexity of a sequential algorithm
- Measuring complexity of an algorithm with branching
- Measuring complexity of an algorithm with loops



Computer Engineering

Develop Good
Quality Systems

Build *Stable*
Software and
Hardware

Build *Efficient*
Software and
Hardware



Which is Better Algorithm?

```
Algorithm multiply1 (var m, n: Integer): Integer  
var index, temp: Integer;  
begin  
    if (n=0) then  
        begin  
            temp := 0;  
        end  
    else  
        begin  
            temp := m;  
            for index := 2 to n do  
                begin  
                    temp := temp + m;  
                end  
            end  
        end  
    end  
end
```

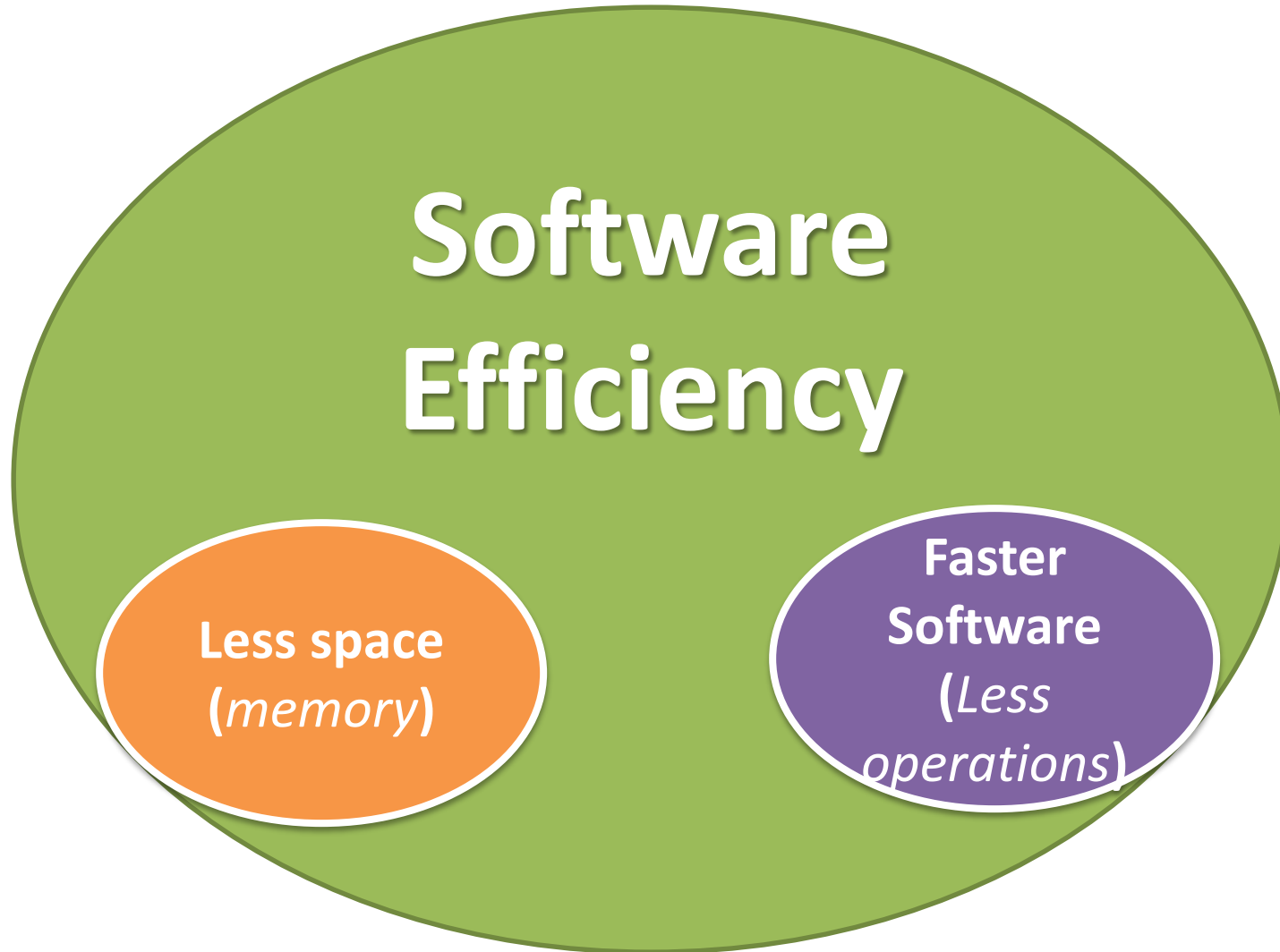
```
Algorithm multiply2 (var m, n :  
Integer): Integer  
var temp: Integer;  
begin  
    temp := m * n;  
end
```

Why

- Less space
- Less time
- More stable



Software Efficiency



Complexity of an Algorithm

- Time Complexity
 - Time is a factor in measuring the efficiency of computer program
 - A measure of time taken for an algorithm to execute
 - Number of cycles
- Space Complexity
 - Space is also a factor in measuring the efficiency of computer program
 - A measure of space taken for executing an algorithm
 - Number of words



Complexity of a Sequential Algorithm

- An Example: Swap 2 numbers

Algorithm *swap*

var a,b,temp : **Integer**; {Space complexity: 3 words}

begin

temp := a; {Time complexity: 1 cycle}

a := b; {Time complexity: 1 cycle}

b := temp; {Time complexity: 1 cycle}

end

Total Space complexity: **3 words**

Total Time complexity: **3 cycles**



Complexity of an Algorithm with Branching

Algorithm *isEven* (**var** a:Integer) { Space complexity: 1 word}

var ret: **boolean**; { Space complexity: 1 word}

begin

 {assert a > 0} {Time complexity 1 cycle}

if ((a mod 2) = 0) **then** {Time complexity 1 cycle}

begin

 ret := true; {Time complexity 1 cycle}

end

else

begin

 ret := false;

 {Time complexity 1 cycle}

end

end

Total Space complexity: 2
words

Total Time complexity:
2 cycles + {1 cycle or 1 cycle}
= 3 *cycles*



Complexity of an Algorithm with Looping

Algorithm *factorial* (**var** n:Integer) { Space complexity: 1 word}

var ret, iLoop: Integer; { Space complexity: 2 words}

begin

 {**assert** n >= 0} {Time complexity 1 cycle}

 ret := 1; {Time complexity 1 cycle}

for iLoop **in** 2 to n **do** {Time complexity n cycles+ 1 cycle}

begin

 ret := ret * iLoop; {Time complexity 2 cycle}

end

end

Total Space complexity: 3 **words**

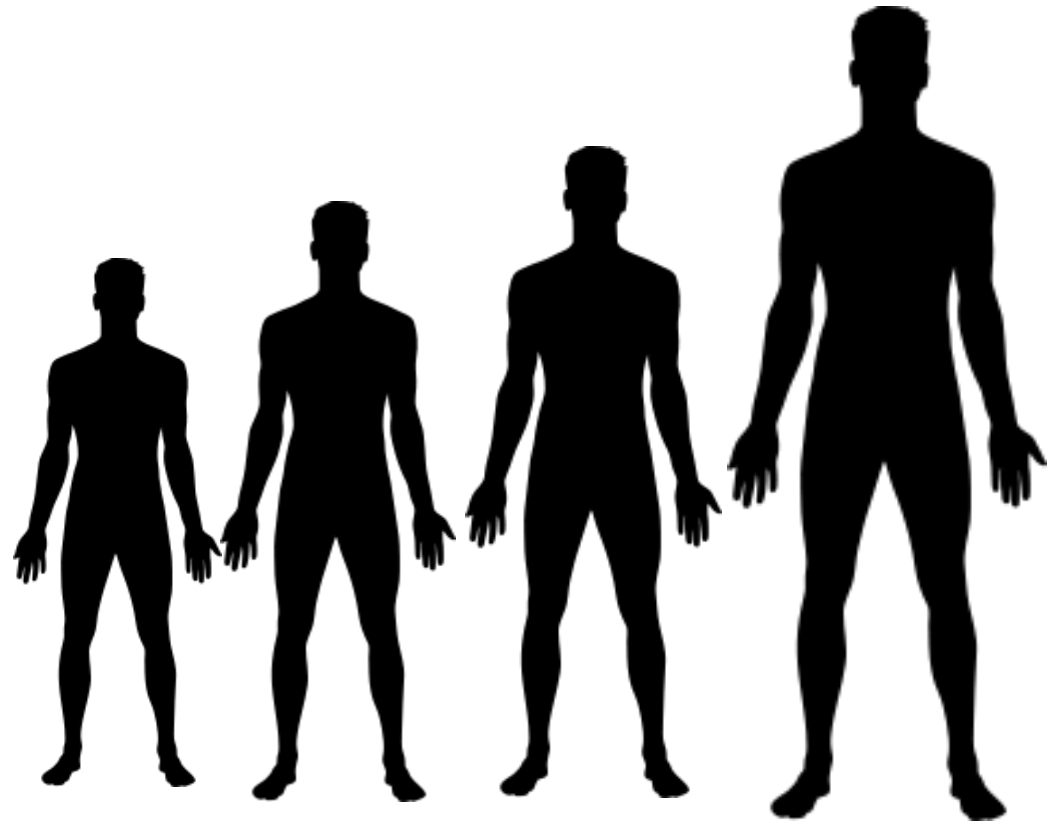
Total Time complexity:

3 cycles + n times 2 cycles

= 2n+3 **cycles**



Growth Rate



Growth Rate

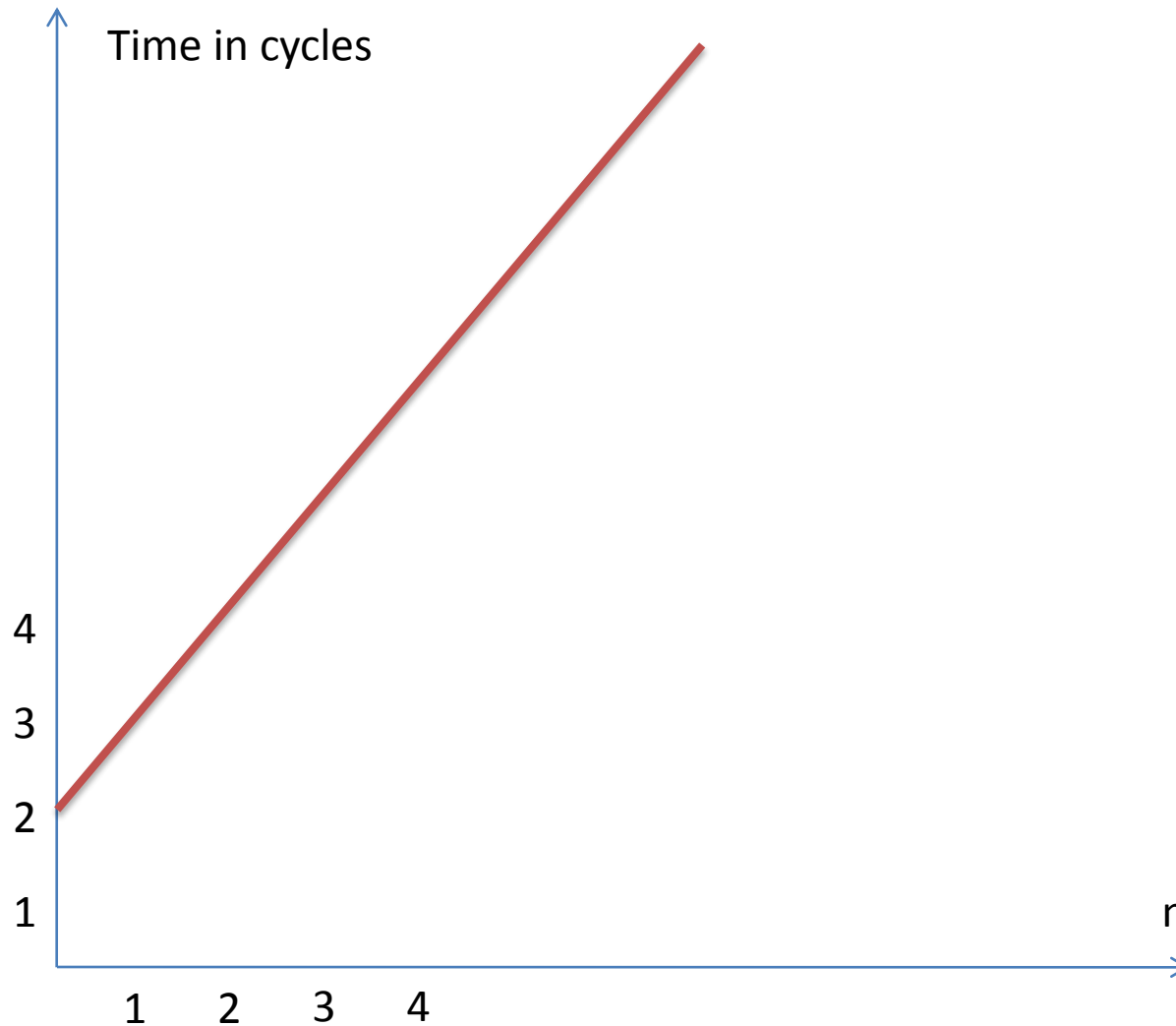
```
Algorithm sumN(array:nIntegerElements, n:Integer):Integer;  
var s,i:Integer; {The partial sum}  
begin  
    s := 0; {1 cycle}  
    for i in 0 to n do {n times + 1 cycle}  
        begin  
            s := s + array[i]; {2 cycle}  
        end  
    end
```

Total time = $1+n*2+1 = 2n+2$ cycles

How much time will this algorithm take when the value of n is increased from 0 towards infinity in steps of 1?



Growth Rate



Examples

- An Example: Swap 2 numbers

Algorithm *swap* (**var** a,b : **Integer**) { **Space complexity: 2 words**}

var temp : **Integer**;{**Space complexity: 1 word**}

begin

temp := a; {**Time complexity: 1 cycle**}

a := b; {**Time complexity: 1 cycle**}

b := temp; {**Time complexity: 1 cycle**}

end

Type of time complexity: Constant (does not change with different inputs)

Type of space complexity: Constant



Examples

Algorithm *isEven* (**var** a:**Integer**) { Space complexity: 1 word}

var ret: **boolean**; { Space complexity: 1 word}

begin

 {assert a > 0} {Time complexity 1 cycle}

if ((a mod 2) = 0) **then** {Time complexity 1 cycle}

begin

 ret := true; {Time complexity 1 cycle}

end

else

begin

 ret := false; {Time complexity 1 cycle}

end

end

*Type of time
complexity: Constant
Type of space
complexity: Constant*



Examples

Algorithm *factorial* (**var** n:Integer) { Space complexity: 1 word}

var ret, iLoop: Integer; { Space complexity: 2 words}

begin

 {**assert** n >= 0} {Time complexity 1 cycle}

 ret := 1; {Time complexity 1 cycle}

for iLoop **in** 2 to n **do** {Time complexity n cycles}

begin

 ret := ret * iLoop; {Time complexity 1 cycle}

end

end

Type of time complexity:

Linear

Type of space complexity:

Constant



Asymptotic Analysis

- Asymptotic analysis of algorithms
 - Describe behavior of algorithms with bounds
 - A way to group algorithms having similar performance behavior
- Big-Oh (O) notation is the most popular as it provides an upper bound to the behavior of an algorithm
 - $O(f(x))$ tells that the complexity of the algorithm is always limited with $f(x)+c1$ as upper bound, where $c1$ is an arbitrary constant
- For example:
 - 2 cycles and 3 cycles all belong to constant time complexity - $O(1)$
 - n cycles, $n+2$ cycles, $2n+3$ cycles all belong to linear time complexity – $O(n)$



Summary

- Efficiency is the process of achieving maximum productivity with minimum wasted effort or expense
- For algorithms, efficiency is measured in terms of space (memory used) and time (number of operations) complexity
- Space and time complexity are estimated by expressing as growth rate functions
- The upper bound of worst case growth rate is generally used to categorise algorithms and the notation used is Big-Oh notation (O)
- If the steps are in sequence, the time complexity is the sum of the steps
- If there is a branch based on a condition, the time complexity is $1 + \text{worst case branch complexity}$



Summary contd.

- If there is looping for 'n' times on a block, the complexity of the block is multiplied by n times and 1 extra cycle is added for the last check
- Always consider the maximum complexity of the function



References

Toida, S. (2013) Summary of Big-Oh, *Growth Functions*, available at <http://www.cs.odu.edu/~toida/nerzic/content/function/growth.html> (accessed 22 July 2014).

Further Reading

Horowitz, E. And Sahni, S. (1983). Chapter 1: Introduction, *Fundamentals of Data Structures*, Pitman.

