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LayerZero Examples

Audit



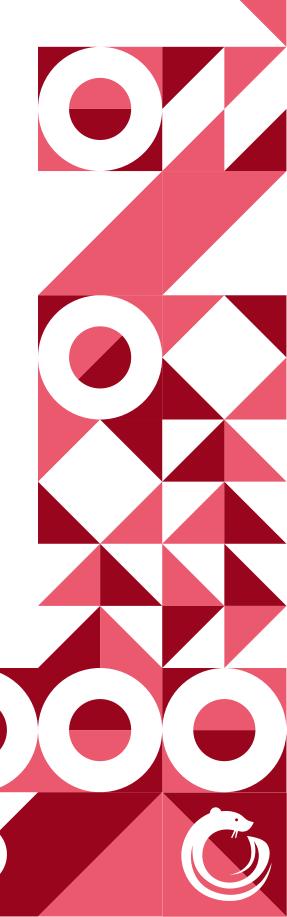
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01 | Executive Summary

Overview

LayerZero engaged OtterSec to perform an assessment of various LayerZero programs under our retainer. This is an ongoing engagement, starting November 11th. For more information on our auditing methodology, see Appendix B.

As part of this engagement, we reviewed LayerZero's OFT v2 upgrade and upgradable contract update.

Key Findings

Over the course of this audit engagement, we produced 5 findings total.

In particular, we found a minor edge case with fee constraint validation (OS-LZR-SUG-01) and a missing initialization check (OS-LZR-SUG-04). Additionally, we made recommendations around reorganizing payload fields (OS-LZR-SUG-00), gas limit ambiguity (OS-LZR-SUG-02), and redundant length check (OS-LZR-SUG-03).

Overall, we commend the LayerZero team for being responsive and knowledgeable throughout the audit.

02 | **Scope**

The source code was delivered to us in a git repository at github.com/LayerZero-Labs/solidity-examples . This audit was performed up to and against commit 2583870.

A brief description of the programs is as follows.

Name	Description
solidity-examples	Solidity code examples building on top of LayerZero.

$03 \mid$ Findings

Overall, we report 5 findings.

We split the findings into **vulnerabilities** and **general findings**. Vulnerabilities have an immediate impact and should be remediated as soon as possible. General findings don't have an immediate impact but will help mitigate future vulnerabilities.

Severity	Count
Critical	0
High	0
Medium	0
Low	0
Informational	5
	High Medium Low

04 | General Findings

Here, we present a discussion of general findings during our audit. While these findings do not present an immediate security impact, they represent antipatterns and could lead to security issues in the future.

ID	Description
OS-LZR-SUG-00	Reorder OFT payload fields to mitigate the impact of variable length decoding.
OS-LZR-SUG-01	Potential edge case for fee constraints
OS-LZR-SUG-02	The gas limit check is ambiguous and can lead to unintended behavior.
OS-LZR-SUG-03	The length check can be safely removed.
OS-LZR-SUG-04	Failure to check for address initialization in isTrustedRemote may lead to unintended behavior.

OS-LZR-SUG-00 | Reorganize Payload Fields

Description

Currently, the OFT payload is encoded via

```
return abi.encodePacked(
PT_SEND_AND_CALL,
_toAddress,
_amountSD,
_addressToBytes32(_from),
_dstGasForCall,
_payload
);tc
```

Note that the security of this encoding depends on the correct size of _toAddress. Otherwise, the critical field _amountSD could get confused with bytes of the _from field.

Remediation

As a defense-in-depth measure, it could make sense to reorder the fields. More specifically, putting _amountSD first in the encoding will likely mitigate future variable length encoding attacks.

OS-LZR-SUG-01 | Fee Constraints Edge Case

Description

In the Fee.sol contract, while setting the defaultFeeBp via setDefaultFeeBp and setFeeBp it should ensure that the fee is less than and **not equal to** the BP_DENOMINATOR. If the defaultFeeBp is equal to the BP_DENOMIMATOR then the transactions will fail.

Remediation

Update the check to ensure that the _feeBp is only less than BP_DENOMINATOR.

```
function setDefaultFeeBp(uint16 _feeBp) public virtual onlyOwner {
    require(_feeBp < BP_DENOMINATOR, "Fee: fee bp must be <
    BP_DENOMINATOR");
    defaultFeeBp = _feeBp;
    emit SetDefaultFeeBp(defaultFeeBp);
}

function setFeeBp(uint16 _dstChainId, bool _enabled, uint16 _feeBp)
    public virtual onlyOwner {
    require(_feeBp < BP_DENOMINATOR, "Fee: fee bp must be <
    BP_DENOMINATOR");
    chainIdToFeeBps[_dstChainId] = FeeConfig(_feeBp, _enabled);
    emit SetFeeBp(_dstChainId, _enabled, _feeBp);
}</pre>
```

OS-LZR-SUG-02 | Gas Limit Ambiguity

Description

The gas limit check has three scenarios.

```
1. minDstGasLookup[_dstChainId][_type] != 0 and _extracGas == 0
```

- 2. minDstGasLookup[_dstChainId][_type] == 0 and _extracGas != 0
- 3. minDstGasLookup[_dstChainId][_type]!=0 and _extracGas!=0

After changing the minimum gas, the admin is unable to revert it back to 0, rendering scenario 2 unusable. This leads to unintended double gas payments.

Remediation

```
require(minDstGasLookup[_dstChainId][_type] > 0)
```

If the intention is to disallow minDstGasLookup[_dstChainId][_type] from being 0, then the gas limit check must be modified to reflect this requirement. Otherwise, the admin must be able to change the minimum gas to 0.

OS-LZR-SUG-03 | Redundant Length Check

```
contracts/contracts-upgradable/lzApp/LzAppUpgradeable.sol

function getTrustedRemoteAddress(uint16 _remoteChainId) external view

    returns (bytes memory) {
    bytes memory path = trustedRemoteLookup[_remoteChainId];
    require(path.length != 0, "LzApp: no trusted path record");
    return path.slice(0, path.length - 20); // the last 20 bytes should

    be address(this)
```

If path.length is zero, subtracting 20 from path.length will result in an arithmetic underflow, causing the program to revert. Hence, it is possible to remove the redundant length check if returning an error message is not needed.

Remediation

Remove the redundant length check.

$OS-LZR-SUG-04 \mid \textbf{Missing Initialization Check}$

Description

The function fails to verify whether srcAddress_ is initialized. Consequently, a function call with a currently uninitialized source chain ID and an empty source address would return true, contrary to expectations.

Remediation

Add an initialization check to isTrustedRemote.

ee rack ert Vulnerability Rating Scale

We rated our findings according to the following scale. Vulnerabilities have immediate security implications. Informational findings can be found in the General Findings section.

Critical

Vulnerabilities that immediately lead to loss of user funds with minimal preconditions

Examples:

- Misconfigured authority or access control validation
- · Improperly designed economic incentives leading to loss of funds

High

Vulnerabilities that could lead to loss of user funds but are potentially difficult to exploit.

Examples:

- Loss of funds requiring specific victim interactions
- Exploitation involving high capital requirement with respect to payout

Medium

Vulnerabilities that could lead to denial of service scenarios or degraded usability.

Examples:

- · Malicious input that causes computational limit exhaustion
- Forced exceptions in normal user flow

Low

Low probability vulnerabilities which could still be exploitable but require extenuating circumstances or undue risk.

Examples:

Oracle manipulation with large capital requirements and multiple transactions

Informational

Best practices to mitigate future security risks. These are classified as general findings.

Examples:

- · Explicit assertion of critical internal invariants
- Improved input validation

B Procedure

As part of our standard auditing procedure, we split our analysis into two main sections: design and implementation.

When auditing the design of a program, we aim to ensure that the overall economic architecture is sound in the context of an on-chain program. In other words, there is no way to steal funds or deny service, ignoring any chain-specific quirks. This usually requires a deep understanding of the program's internal interactions, potential game theory implications, and general on-chain execution primitives.

One example of a design vulnerability would be an on-chain oracle that could be manipulated by flash loans or large deposits. Such a design would generally be unsound regardless of which chain the oracle is deployed on.

On the other hand, auditing the implementation of the program requires a deep understanding of the chain's execution model. While this varies from chain to chain, some common implementation vulnerabilities include reentrancy, account ownership issues, arithmetic overflows, and rounding bugs.

As a general rule of sum, implementation vulnerabilities tend to be more "checklist" style. In contrast, design vulnerabilities require a strong understanding of the underlying system and the various interactions: both with the user and cross-program.

As we approach any new target, we strive to get a comprehensive understanding of the program first. In our audits, we always approach targets with a team of auditors. This allows us to share thoughts and collaborate, picking up on details that the other missed.

While sometimes the line between design and implementation can be blurry, we hope this gives some insight into our auditing procedure and thought process.