

Q.1

a.

$A \twoheadrightarrow B$ does not hold. a_2 has two different values c_1 and c_2 .

b.

$B \twoheadrightarrow D$ holds. b_3 has one value d_3 .

c.

$BCD \twoheadrightarrow A$ does not hold. b_3, c_2, d_3 have two different value a_2, a_3

d.

$F = \{A \twoheadrightarrow D, B \twoheadrightarrow D, AB \twoheadrightarrow D, BC \twoheadrightarrow D, ABC \twoheadrightarrow D\}$

$A^+ = \{A, D\}$

$B^+ = \{B, D\}$

e.

$AB^+ = \{A, B, D\}$ c is missing in AB closure so, it is not a super key or candidate key.

Q.2

a.

$A^+ = \{A, D, E\}$

b.

$CF^+ = \{C, F, A, D, E, B\}$

c.

1. $B \twoheadrightarrow C$

2. $DB \twoheadrightarrow DC$ by augmentation of 1

3. $DC \twoheadrightarrow C$ by reflexivity

4. $DB \twoheadrightarrow C$ by transitivity of 2 and 3

Q.3

a.

$A^+ = \{A, D, B, C\}$ so, A is the candidate key

b.

BCNF:

In every FD $x \twoheadrightarrow y$, x has to be super key

$A \twoheadrightarrow D$ A is a super key

$B \twoheadrightarrow C$ B is not a super key

$A \twoheadrightarrow BD$ A is super key

$D \twoheadrightarrow B$ D is not super key

Does not satisfy BCNF

c.

3NF:

For every non-trivial FD $x \twoheadrightarrow y$:

X is a super key or

Y is a prime attribute

$A \twoheadrightarrow D$ A is a super key and A is prime attribute

$B \twoheadrightarrow C$ B is not a super key, C is not prime attribute so, 3NF fails.

Q.4

a.

$A^+ = \{A, D\}$

$A \twoheadrightarrow D$ A is not a super key so, BCNF fails.

b.

$R_1(A, D), R_2(B, C, E), R_3(A, B, C, F)$

ABCDEF		
AD	ABCEF	
(A \twoheadrightarrow D)	BCE	ABCF
	(BC \twoheadrightarrow E)	(AF \twoheadrightarrow BC)

c.

The relations include the original FD so, decomposition is dependency preserving.

d.

We cannot reach AF \twoheadrightarrow BC from R1 and R2 so, it is not lossless.

e.

AF \twoheadrightarrow BC is lost so, dependency preserving fails.

Q.5

a.

$F = \{A \twoheadrightarrow BC, B \twoheadrightarrow E, BD \twoheadrightarrow C, AD \twoheadrightarrow CE, E \twoheadrightarrow AD\}$

$B^+ = \{B, E, A, D, C\}$

C in B closure so, D is extraneous in $BD \twoheadrightarrow C$

b.

$F' = \{A \twoheadrightarrow B, B \twoheadrightarrow E, BD \twoheadrightarrow C, AD \twoheadrightarrow CE, E \twoheadrightarrow AD\}$

$A^+ = \{A, B, E, D, C\}$

C in A closure so, C is extraneous in $A \twoheadrightarrow BC$

c.

Step1: split RHD

$F = \{A \twoheadrightarrow B, A \twoheadrightarrow C, B \twoheadrightarrow E, BD \twoheadrightarrow C, AD \twoheadrightarrow C, AD \twoheadrightarrow E, E \twoheadrightarrow A, E \twoheadrightarrow D\}$

Step2: remove extraneous attributes (D is extraneous in $AD \twoheadrightarrow E$ because A^+ includes E)

$F = \{A \twoheadrightarrow B, A \twoheadrightarrow C, B \twoheadrightarrow E, B \twoheadrightarrow C, A \twoheadrightarrow E, E \twoheadrightarrow A, E \twoheadrightarrow D\}$

Step3: remove redundant functional dependencies

$A \twoheadrightarrow B$ $A^+ = \{A, E, C, D\}$ not include B so, not redundant

$A \twoheadrightarrow C$ $A^+ = \{A, B, E, C, D\}$ redundant

$B \twoheadrightarrow E$ $B^+ = \{B, C\}$ not redundant

$B \twoheadrightarrow C$ $B^+ = \{B, E, A, D, C\}$ redundant

$A \twoheadrightarrow E$ $A^+ = \{A, B, E, D, C\}$ redundant

$E \twoheadrightarrow A$ $E^+ = \{E, D\}$ not redundant

$E \twoheadrightarrow D$ $E^+ = \{E, A, B, C\}$ not redundant

Canonical form: $F = \{A \twoheadrightarrow B, B \twoheadrightarrow E, E \twoheadrightarrow A, E \twoheadrightarrow D\}$