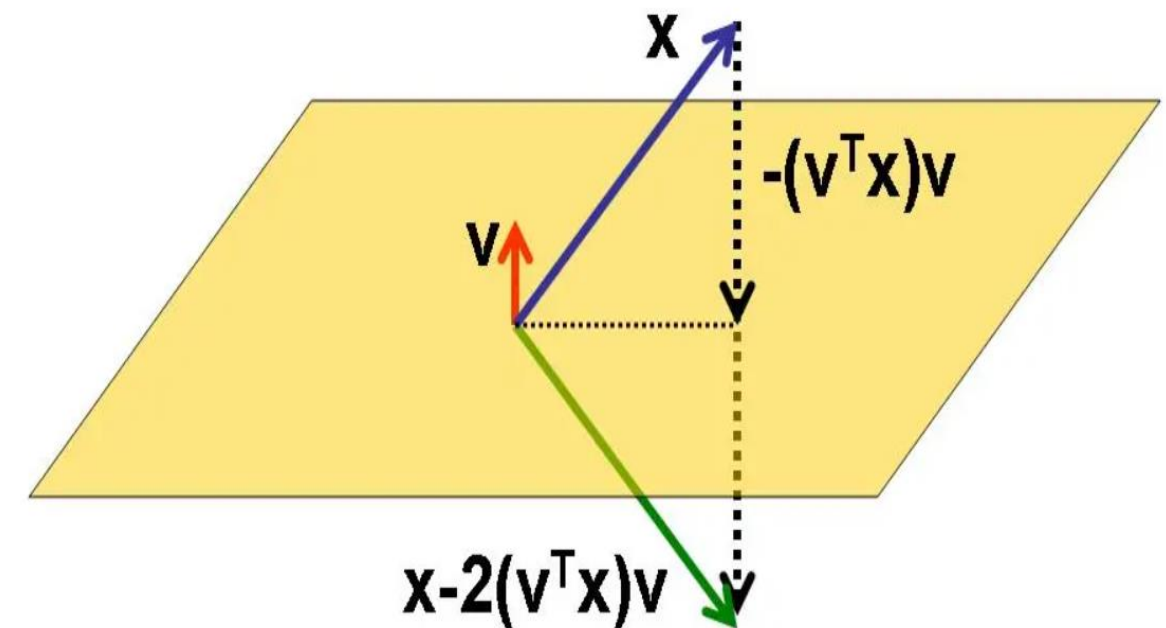


# A block Cholesky-LU-based QR factorization for rectangular matrices

## Team 007

[Link to GitHub](#)

In this presentation, we will explore the powerful technique of QR factorization using Householder method recursively and cutting off the recursion at a small enough block and solving it using the Cholesky-LU method specifically designed for rectangular full rank matrices.



# Motivation

## 1 Numerical Stability

Since the method uses Orthogonal projections , it gets some of the stability as norms are preserved in some of the operations

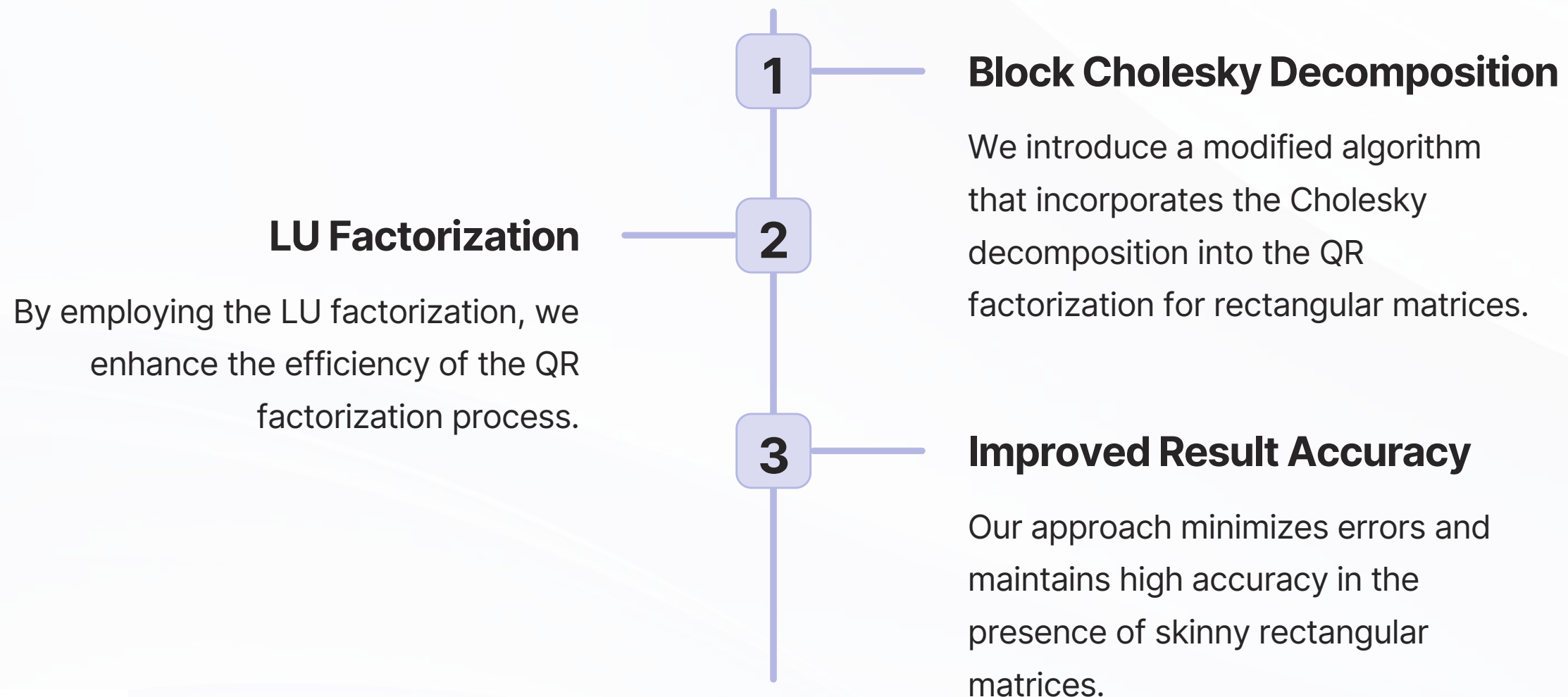
## 2 Computational Efficiency

Block Based computation are more efficient as they maximise cache hit rate in data retrieval

## 3 Best of both Worlds

This method combines the efficient part of the stability of using Householder methods with the efficiency in computing the Block Cholesky LU factorisation

# Cholesky-LU-based QR Factorization for Rectangular Matrices



# Block-based Algorithms for QR Factorization

Prim's Algorithm from a matrix

	1	2	3			
	A	B	C	D	E	F
A	-	8	-	10	-	-
B	8	-	14	12	-	-
C	-	14	-	15	13	-
D	10	12	15	-	12	20
E	-	-	13	12	-	19
F	-	-	-	20	19	-

3. Label column D with a 3. Delete row D. Now choose the smallest uncovered value from columns A or B or D

## Block Householder Transformations

We present a new algorithm that uses block-based Householder transformations to improve the efficiency of QR factorization.

### Partitioned Matrices

- For example, three possible partitions of a 3x4 matrix A:
  - The partition of A into four submatrices  $A_{11}$ ,  $A_{12}$ ,  $A_{21}$ , and  $A_{22}$ 
$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$$
  - The partition of A into its row matrices  $r_1$ ,  $r_2$ , and  $r_3$ 
$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} r_1 \\ r_2 \\ r_3 \end{bmatrix}$$
  - The partition of A into its column matrices  $c_1$ ,  $c_2$ ,  $c_3$ , and  $c_4$ 
$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} c_1 & c_2 & c_3 & c_4 \end{bmatrix}$$

## Matrix Partitioning

The method incorporates block-wise matrix partitioning to optimize the factorization process for larger, rectangular matrices.

### What is Computational Complexity?

- Study of the amount of resources required by an algorithm to solve a problem.
  - Mathematically formalized as Turing Machines

## Computational Complexity

We visualize the empirical time taken in the block-based algorithms and demonstrate their advantages compared to traditional methods.

# Problem Statement

# Algorithm 1

We consider  $m \geq n$

Recursive Householder-based QR factorization of  $A \in \mathbb{R}^{m \times n}$  with orthogonal factor of the form  $Q = I - VSV^T$ .

Input  $A \in \mathbb{R}^{m \times n}$

Output  $V, R \in \mathbb{R}^{m \times n}$   $S \in \mathbb{R}^{n \times n}$  representing the QR factor of  $A$

$$\begin{pmatrix} A_1 & A_2 \end{pmatrix} = \left( I - \begin{pmatrix} V_1 & V_2 \end{pmatrix} \begin{pmatrix} S_{11} & S_{12} \\ 0 & S_{22} \end{pmatrix} \begin{pmatrix} V_1 & V_2 \end{pmatrix}^T \right) \begin{pmatrix} R_{11} & R_{12} \\ 0 & R_{22} \\ 0 & 0 \end{pmatrix},$$

$$\begin{pmatrix} A_1 \\ \tilde{A}_{22} \end{pmatrix} = \left( \begin{pmatrix} \square & \\ & \square \end{pmatrix} - \begin{pmatrix} V_1 & V_2 \end{pmatrix} \begin{pmatrix} S_{11} & S_{12} \\ & S_{22} \end{pmatrix} \begin{pmatrix} V_1^T & V_2^T \end{pmatrix} \right) \begin{pmatrix} R_{11} & R_{12} \\ & R_{22} \end{pmatrix}$$

# Algorithm 2

We consider  $m \geq n$

**Block Cholesky-LU-based QR factorization of  $A = (A_{\square}^T, A_r^T)^T$  with orthogonal factor of the form  $Q = I - VSV^T$ .**

**Input**  $A \in \mathbb{R}^{m \times n}$

**Output**  $V, R \in \mathbb{R}^{m \times n}$   $S, R \in \mathbb{R}^{n \times n}$  representing the QR factor of  $A$ .

**The block form can be inferred from the equation below**

$$\begin{pmatrix} A \\ A_r \end{pmatrix} = A = QR = (Y|Z) \begin{pmatrix} R \\ 0 \end{pmatrix}$$

# Algorithm 3 (Proposed Method)

We consider  $m \geq n$

Recursive block QR factorization of  $A \in \mathbb{R}^{m \times n}$  with orthogonal factor of the form  $Q = I - VSV^T$  with user specified input  $1 < n < k$

Input  $A \in \mathbb{R}^{m \times n}$

Output  $V, R \in \mathbb{R}^{m \times n}$   $S \in \mathbb{R}^{n \times n}$  representing the QR factor of A.

**Pseudocode**

If  $1 < n < k$

    Perform Algorithm 2

else

    Perform algorithm 1

endif



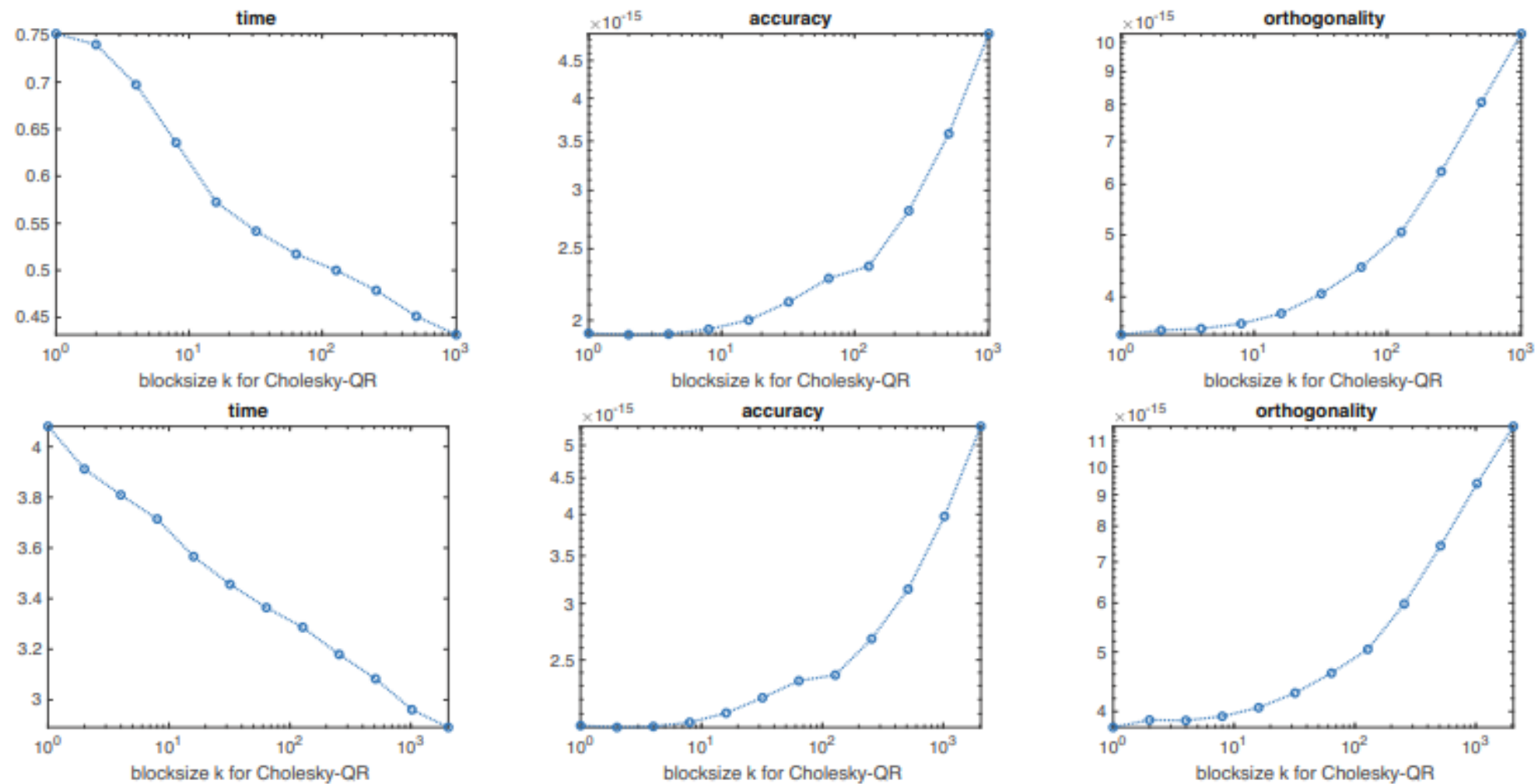
# Problem Statement

**An algorithm that factors tall rectangular matrices into QR orthogonal factors that is computationally efficient.**

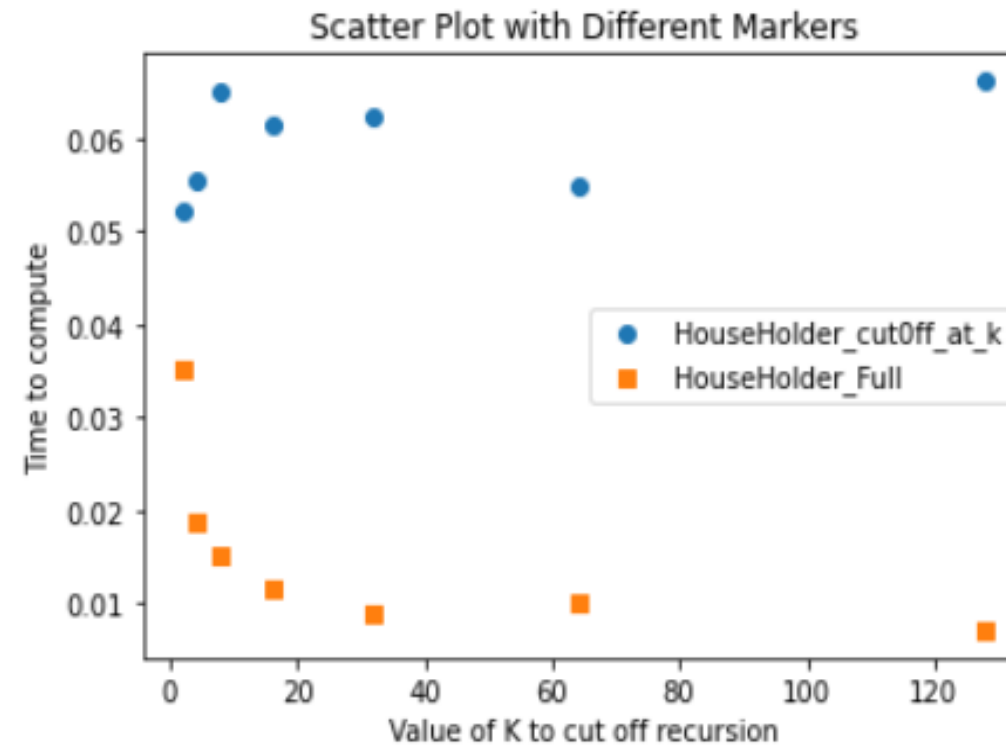
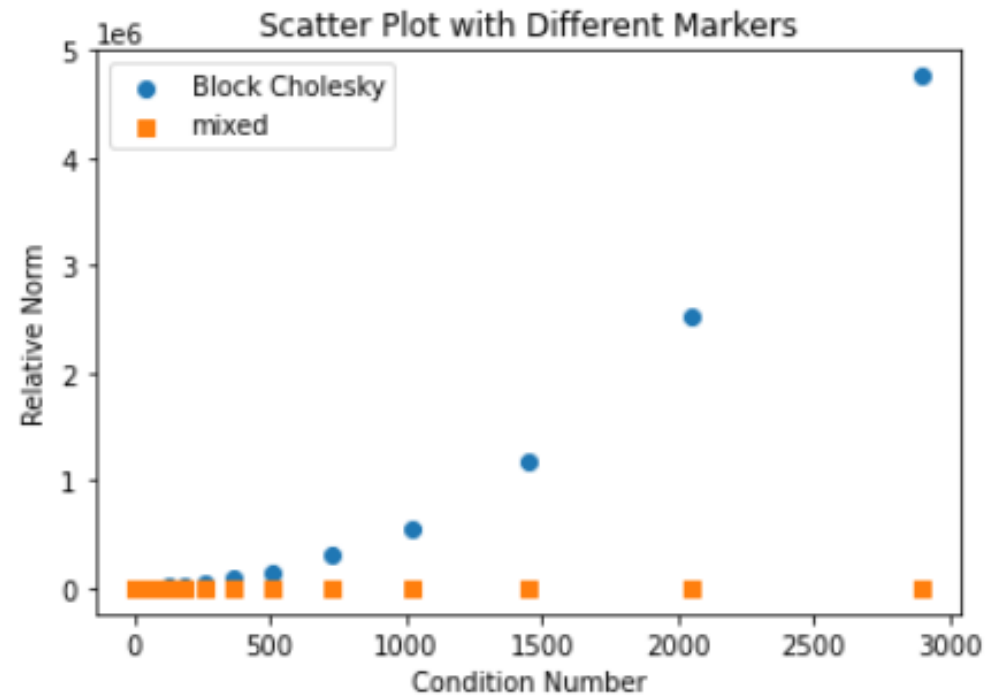
**The known algorithms are either very fast and unstable(eg Cholesky Block LU(Algorithm 2)) or quite slow and stable(eg continuing the HouseHolder method down to a single column (Algorithm 1) )**

**So our task is to find an algorithm that takes advantage of strength of both algorithms**

# Numerical Results



# Numerical Results



Comparison of Algorithm 3 with Algorithm 2 and 1 on stability and computational time respectively for  $m = 2000$  and  $n = 500$

# Team Contribution

- Holdings Ogon - Implemented algorithm 1 and helped with writing the condition number function
- Joshua Udobang - Did the plot graphing and observed the empirical results
- Nwachukwu Mmesomachi - Implemented Algorithm 3 and verified the mathematical correctness. Wrote the README file
- Okechukwu Okeke - Implemented algorithm 2 and modularized the code
- All contributed to making the presentation

# References

1. Le Borne S. A block Cholesky-LU-based QR factorization for rectangular matrices. Numer Linear Algebra Appl. 2023;30(5):e2497
2. Barlow J. Block modified Gram-Schmidt algorithms and their analysis. SIAM J Matrix Anal Appl. 2019;40:1257-90.
3. Carson E, Lund K, Rozloznik M, Thomas S. Block Gram-Schmidt algorithms and their stability properties. Linear Algebra Appl. 2022;638:150-95.
4. Terao T, Ozaki K, Ogita T. LU-Cholesky QR algorithms for thin QR decomposition. Parallel Comput. 2020;92:102571. <https://doi.org/10.1016/j.parco.2019.102571>
5. Golub G, Van Loan C. Matrix computations. 3rd. ed. London: John Hopkins; 1996.
6. Kressner D, Susnjara A. Fast QR decomposition of HODLR matrices, 2018. <https://arxiv.org/abs/1809.10585>
7. Le Borne S. Block computation and representation of a sparse nullspace basis of a rectangular matrix. Linear Algebra Appl. 2008;428:2455-67