

CPSC 313 2012 – Mike Feeley
Assignment #2

Name: Matthew Park
Student Number: 34017103

Problem 1:

Pseudo-code for `rmmovl`

Fetch:

icode:ifun $\leftarrow M_1[PC]$
rA:rB $\leftarrow M_1[PC+1]$
valC $\leftarrow M_4[PC+2]$
valP $\leftarrow PC+6$

Decode:

valA $\leftarrow R[rA]$
valB $\leftarrow R[rB]$

Execute:

valE $\leftarrow valC + (valB * ifun)$

Memory:

$M_4[valE] \leftarrow valA$

Write back:

PC update:

PC $\leftarrow valP$

Pseudo-code for `mrmovl`

Fetch:

icode:ifun $\leftarrow M_1[PC]$
rA:rB $\leftarrow M_1[PC+1]$
valC $\leftarrow M_4[PC+2]$
valP $\leftarrow PC+6$

Decode:

valB $\leftarrow R[rB]$

Execute:

valE $\leftarrow valC + (valB * ifun)$

Memory:

valM $\leftarrow M_4[valE]$

Write back:

$R[rA] \leftarrow valM$

PC update:

PC $\leftarrow valP$

Description of test coverage `movl.s`

According to the comments in the AbstractY86CPU, Y86 CPU can take multiple arguments for the same purpose as proposed in problem 1. They are: 4, 2, 1 and no argument. So the test is split into four parts to test each of the 4 arguments respectively: testx4, testx2, testx1, and testxReg. The workflow is identical for each part; except for the values differ to what is read into memory/register and the constant values to adjust for the offset.

Each part loads a value in a register %ecx that will be inserted into the first two position of an array using the improved rmmovl instruction. It will then take the new value that is stored in memory and place it into a new register using the mrmovl instruction. The workflow will then repeat itself two more times in different positions on memory.

I believe my movl.s has sufficient coverage of the new implementation of rmmovl and mrmovl. During each of the workflows, it inserts different values into different addresses of memory and loads those values into different registers. The workflow is repeated with different values for each argument supported; verifying that each supported argument works and that the intended function works as each position of the array is being overwritten. The test also covers a case of invalid argument for rmmovl and mrmovl, in section "testxInv". When uncommenting those lines of code, Y86 will fail to load movl.s due to an illegal scale.

Test results for movl.s

NOTE: The best way to run the test is to have break points at each of labels.

Label "testxInvalid", test invalid arguments for rmmovl and mrmovl.

Uncommenting line 84 should result an "Assembly error: Illegal scale: must be 1, 2, or 4."

Uncommenting line 85 should result an "Assembly error: Illegal scale: must be 1, 2, or 4."

Pretest:

All values of TheArray should be 7.

Values after running "testx4":

.pos 0x1000

TheArray:	.long	40
	.long	40
	.long	41
	.long	41
	.long	42
	.long	42
%eax		40
%esi		41
%edi		42

Values after running "testx2":

.pos 0x1000

TheArray:	.long	20
	.long	20
	.long	21
	.long	21
	.long	22
	.long	22
%eax		20
%esi		21
%edi		22

Values after running "testx1":

.pos 0x1000

```

TheArray:    .long    10
              .long    10
              .long    11
              .long    11
              .long    12
              .long    12
%eax         10
%esi         11
%edi         12

```

Values after running "testxReg":

```

.pos 0x1000
TheArray:    .long    0
              .long    0
              .long    1
              .long    1
              .long    2
              .long    2
%eax         0
%esi         1
%edi         2

```

Problem 2:

Pseudo-code for I IOPL

Fetch:

```

icode:ifun   ← M1[PC]
rA:rB        ← M1[PC+1]
valC         ← M4[PC+2]
valP         ← PC+6

```

Decode:

```

valB         ← R[rB]

```

Execute:

```

valE         ← valB [+ , - , * , / , and , or , xor] valC

```

Memory:

Write back:

```

R[rB]        ← valE

```

PC Update:

```

PC           ← valP

```

Description of test coverage iopl.s

iopl.s is divided into seven sections, making sure each of the seven arithmetic functions are working correctly. I believe that iopl.s provides sufficient test coverage because each section tests the arithmetic operation at least twice, while testing specific odd cases that mimics i_opl's operation (such as dividing by zero). The arithmetic operation is then checked using I_opl's untouched implementation and errors are then handled to determine which arithmetic operation is not working correctly.

Test results for iopl.s

Simply run the test and check %eax value at the end.

If %eax's value is:

- 0 = test passed
- 1 = iaddl error
- 2 = isubl error
- 3 = imul error
- 4 = idivl error
- 5 = iandl error
- 6 = ixorl error
- 7 = imodl error

Problem 3:

Fetch:

icode:ifun	$\leftarrow M_1[PC]$
rA:rB	$\leftarrow M_1[PC+1]$
valC	$\leftarrow M_4[PC+2]$
valP	$\leftarrow PC+6$

Decode:

valB	$\leftarrow R[rB]$
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Execute:

valE	$\leftarrow valC + (valB * ifun)$
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Memory:

valM	$\leftarrow M_4[valE]$
valE	$\leftarrow valP$

Write back:

R[rA]	$\leftarrow valE$
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PC update:

PC	$\leftarrow valM$
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Description of test coverage call.s

I believe this test is sufficient because, call.s checks to see if both the regression method as well as the new implementation works. By having the successfully coming to the halt statement, this shows that both function works correctly.

Test results for call.s

Run call.s

If programs comes to a halt it runs successfully, and if:

%eax = 7

%edx= 0x117

%ebx= 8

%ecx= 4